

1. Solve the recurrence  $T(n) = T(\sqrt{n}) + 1$ . Give a  $\theta$  bound.

Let  $n = 2^k$ . Thus we have  $T'(k) = T(2^k) = T(\sqrt{2^k}) + 1 = T'(k/2) + 1$ .  
Thus  $T'(k) = \log k$ . Thus  $T(n) = \log \log n$ .

2. Consider the following functions:  $n - \log n$ ,  $\log \log n$ ,  $2^{\log^2 n}$ ,  $\sqrt{n / \log n}$ ,  $n^{\log n}$ .  
Arrange them in order s.t. the  $i$ th function is  $O()$  of the  $i+1$ th function.

You may want to note that if  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} < c$  for any constant  $c$ , then  $f(n) = O(g(n))$ . It is fine if you use this.

$\log \log n, \sqrt{n / \log n}, n - \log n, 2^{\log^2 n} = n^{\log n}$

3. Consider an array  $x[1..n]$ . An element  $x[i]$  is a local minimum if it is no larger than  $x[i-1]$  (if any), as well as  $x[i+1]$  (if any). Give an  $O(\log n)$  algorithm to find a local minimum.

Algorithm:

- (a) If  $n < 10$  then check each element. The smallest number is certainly a local minimum, and this will be found.
  - (b) If  $x[n/2] < x[n/2 + 1]$ , recurse on  $x[1..n/2]$ . If this call returns  $i$  s.t.  $1 \leq i < n/2$  then it will be a local minimum by correctness of the recursion. If it returns  $i = n/2$  then we will have  $x[n/2 - 1] \geq x[n/2]$ . But we already have  $x[n/2] < x[n/2 + 1]$ . Thus  $x[i]$  is indeed a local minimum.
  - (c) Else recurse on  $x[n/2 + 1..n]$ . The rest of the argument is similar.
4. The input to this problem are two sorted arrays, of length  $m, n$  respectively. All the  $m + n$  values in the two arrays are distinct. Another input is an integer  $s$ , satisfying  $s \leq m + n$ . The goal is to find the  $s$ th smallest value from the  $m + n$  values. Show that this can be done using at most  $1 + \lceil \log s \rceil$  comparisons.

Let the arrays be  $A[1..m], B[1..n]$ . Compare  $A[\lceil s/2 \rceil], B[\lceil s/2 \rceil]$ . Say  $A[\lceil s/2 \rceil] < B[\lceil s/2 \rceil]$ . Clearly  $A[1.. \lfloor s/2 \rfloor]$  are all smaller than the  $s$ th

smallest. Thus for the next step we can drop these elements and look for the  $s - \lfloor s/2 \rfloor$ th element. The recursion will end when  $s = 1$ , whereupon we compare  $A[1], B[1]$ .