

CS F364

Design & Analysis of Algorithms

ALGORITHMS - COMPLEXITY

Non-Deterministic Computation - Ideas, Examples

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NON-DETERMINISTIC COMPUTATIONS - DEFINITION

○ In a deterministic TM:

- Computation proceeds by moving from one state to another in a “deterministic way”
 - i.e. given a state, and an input, the next state can be “determined”
 - i.e. the state transitions can be defined by a function (on set of states and set of symbols)

○ By contrast, in a non-deterministic TM:

- Computation proceeds by moving from one state to one of many states in a “non-deterministic” way
 - i.e. given a state, and an input, the next state is *not known a-priori* (before the transition)
 - i.e. the state transitions must be defined by a relation (on set of states and set of symbols)

NON-DETERMINISTIC COMPUTATIONS - DEFINITION [2]

- Non-deterministic computations are defined using abstract machines:
 - There are no inherently non-deterministic computers (if you don't allow Schrodinger's cats!)
- Conceptually, a non-deterministic computation proceeds by making “non-deterministic” choices in real-time i.e.
 - Each choice takes $O(1)$ time irrespective of the number of options available
 - as long as the number of choices is finite

NON-DETERMINISTIC ALGORITHMS - EXAMPLE 1

○ Problem:

- Given an array A of values, and a key k , find whether a value with key k is contained in A .
- Algorithm $\text{NDSearch}(A, k)$
 1. $\text{len} = A.\text{length};$
 2. $\text{ind} = \text{choose}(0, \text{len}-1);$
 3. if $(A[\text{ind}] == k)$ return 1;
 4. else return 0;

NON-DETERMINISTIC ALGORITHMS - EXAMPLE 2

○ Problem SAT(isfiability):

- Given a Boolean expression E with variables x_0, \dots, x_{n-1} find whether there exists an assignment of (Boolean) values that satisfies E .
- NDSAT(E, X)
 1. $len = X.length;$
 2. $for(j=0; j < len; j++) \quad \{ X[j] = choose(0, 1); \}$
 3. $if (evaluate(E, X) == 1) \quad return 1;$
 4. $else return 0;$