

CS F364: Design & Analysis of Algorithm

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Longest Common Subsequence



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<http://ktiwari.in/algo>

Components of Dynamic Programming

1) Optimal substructure, 2) Overlapping subproblems

Let
 $X_m = \langle x_1, x_2, x_3, \dots, x_m \rangle$
 $Y_n = \langle y_1, y_2, \dots, y_n \rangle$
 $Z_k = \langle z_1, \dots, z_k \rangle$ **$z_k > \text{LCS}$**

- 1 If $(x_m = y_n)$ then $Z_k = x_m$ and $Z_{k-1} = \text{LCS}(X_{m-1}, Y_{n-1})$
- 2 If $(x_m \neq y_n)$ then
If $(z_k \neq x_m)$ then $Z_k = \text{LCS}(X_{m-1}, Y_n)$
If $(z_k \neq y_n)$ then $Z_k = \text{LCS}(X_m, Y_{n-1})$

LCS of two sequences contains within LCS of the prefix of the two sequences

$$c[i, j] = \begin{cases} 0 & \text{if } i = 0 \text{ or } j = 0 \\ c[i-1, j-1] + 1 & \text{if } x_i = y_j \\ \max(c[i-1, j], c[i, j-1]) & \text{otherwise} \end{cases}$$

Longest common subsequence

Algorithm 1: LCS-Length(X, Y)

```
1 m = length(X)
2 n = length(Y)
3 for i = 0 to n do
4   c[i, 0] = 0
5 for j = 0 to m do
6   c[0, j] = 0
7 for i = 1 to m do
8   for j = 1 to n do
9     if  $x_i = y_j$  then
10      c[i, j] = c[i-1, j-1] + 1
11      b[i, j] = i-1, j-1
12     else
13      if  $c[i-1, j] > c[i, j-1]$  then
14        c[i, j] = c[i-1, j]
15        b[i, j] = i-1, j
16      else
17        c[i, j] = c[i, j-1]
18        b[i, j] = i, j-1
19 return c and b
```

Complexity? $O(m^2)$

Longest common subsequence

A **subsequence** of a sequence can be obtained by **removing zero or more** elements.

- In the longest-common-subsequence problem, we are given two sequences $X = \langle x_1, x_2, x_3, \dots, x_m \rangle$ and $Y = \langle y_1, y_2, y_3, \dots, y_n \rangle$ and wish to find a maximum-length common subsequence of X and Y .
- Example: DNA sequence $\{A, T, C, G\}$
- Time? exponential
- **Prefix** uses first few items

Solution Sketch

	B	D	C	A	B	A
A						
B						
C						
B						
D						
A						
B						

Longest common subsequence

Algorithm 2: Print-LCS(b, X, i, j)

```
1 if  $i=0$  or  $j=0$  then
2   return
3 if  $b[i, j] = \nwarrow$  then
4   Print-LCS( $b, X, i-1, j-1$ )
5   print  $x_i$ 
6 else
7   if  $b[i, j] = \uparrow$  then
8     Print-LCS( $b, X, i-1, j$ )
9   else
10    Print-LCS( $b, X, i, j-1$ )
```

Example

j	0	1	2	3	4	5	6
i	y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0
1	A	0	↑	↑	↑	↑	←
2	B	0	←	←	←	1	←
3	C	0	↑	↑	←	2	↑
4	B	0	←	↑	2	2	↑
5	D	0	↑	←	2	2	3
6	A	0	↑	↑	2	3	↑
7	B	0	←	↑	2	1	4

Thank You!

Thank you very much for your attention! (Reference¹)

Queries ?