



# Compiler Construction

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# CS F363, Compiler Construction

Lecture topics: Three Address code generation

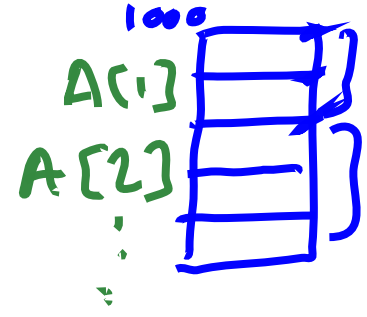
*Code Generation for Arrays Expression*  
*A[100]*

→ Consider an array element  $A[i]$   
 → Start location of element  $A[i]$



low, high.

for  $w=2$



$$\Rightarrow \text{base} + (i - \text{low}) * w$$

$$\Rightarrow \underbrace{(\text{base} - \text{low} * w)}_{\text{constant}} + \underbrace{i * w}_{\text{offset}}$$

e.g.  $A[2]$

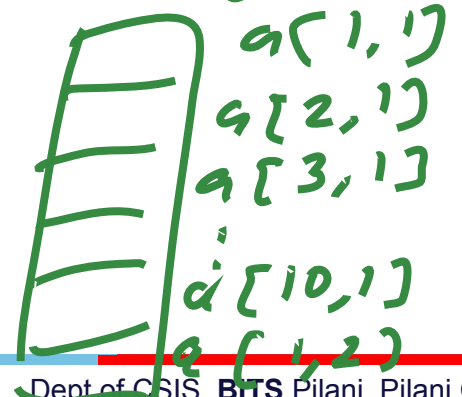
$$(1000 - 1 * 2) + 2 * 2 = \underline{1002}$$

for 2-D Array

i) Row Major Storage ✓



ii) Column major Storage





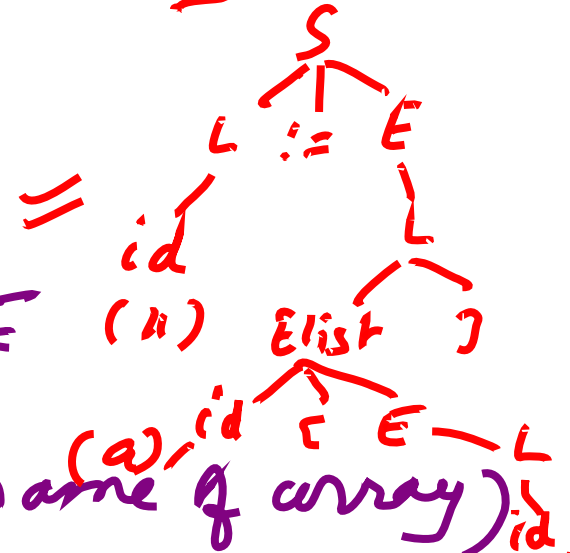
$a[i_1, i_2]$   
 start location of  $a[i_1, i_2]$   $n_2 \rightarrow$  no. of columns.

$$\Rightarrow \text{base} + ((i_1 - \text{low}_1) * n_2 + i_2 - \text{low}_2) * W$$

$$\Rightarrow \underbrace{\text{base} - (\text{low}_1 * n_2 - \text{low}_2) * W}_{\text{constant}} + \underbrace{(i_1 * n_2 + i_2) * W}_{\text{offset}}$$

- Grammar:-
- 1)  $S \rightarrow L := E$
  - 2)  $E \rightarrow E + E \mid (E) \mid L$
  - 3)  $L \rightarrow [E \text{ list}] \mid \text{id}$
  - 4)  $E \text{ list} \rightarrow E \text{ list}, E \mid \text{id} [E$

$$x = a[b]$$



Attributes:-

L.place  $\rightarrow$  holds name of var (or name of array)  
 L.offset  $\rightarrow$  NULL for simple var (offset of element of array)

$E.place \rightarrow$  name of var holding value of  $Exp\ E$

$E.list.place \rightarrow$  name of var holding value of  $index\ emp$

$E.list.dname \rightarrow$  holds the name of array

$E.list.dim \rightarrow$  holds the current dim under consideration.

$S \rightarrow L := E$

if  $L.offset = NULL$

gen ( $L.place := E.place$ )

else

gen ( $L.place[L.offset] = E.place$ )

$E \rightarrow E_1 + E_2$

$E.place = \text{new Temp}()$

$E.code = E_1.code || E_2.code || \text{gen}(E.place =$

$\{ E \rightarrow (E_1)$   
 $E.place = E_1.place$

$\{ E \rightarrow L$   
 if  $L.offset = \text{NULL}$   
 $E.place = L.place$   
 else  $E.place = L.place[L.offset]$

$\{ L \rightarrow id$   
 $L.place = id.place$   
 $L.offset = \text{NULL}$

$\{ L \rightarrow Elist$   
 $L.place = \text{new Temp}$   
 $L.offset = "$   
 $gen (L.place = E(Elist.array))$

$\text{width}(E) = \text{Elist.place} \cdot \text{width}(E.\text{array})$

innovate

achieve

lead

$Elist \rightarrow id \in E$

$Elist.array = id.place$

$Elist.place = E.place$

$Elist.dim = 1$

$Elist \rightarrow Elist.1, E$

$t = \text{new temp}()$

$m = Elist.1.dim + 1$

$gen(t := Elist.1.place * \text{limit}(\frac{Elist.1.array}{m})$

$t = t + E.place$

$Elist.array = Elist1.array \checkmark$

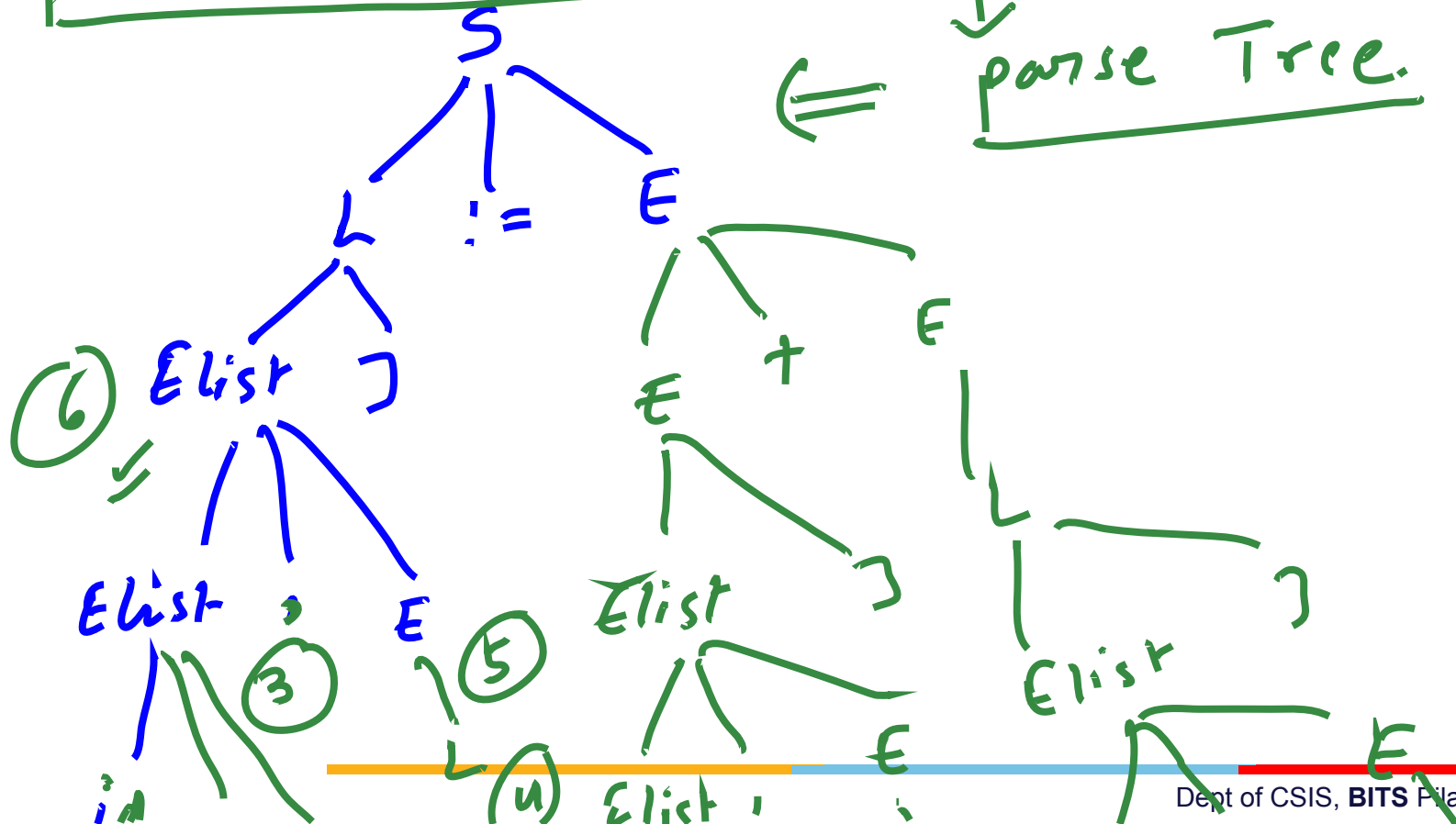
$Elist.place = t \checkmark$

$Elist.dim = m$

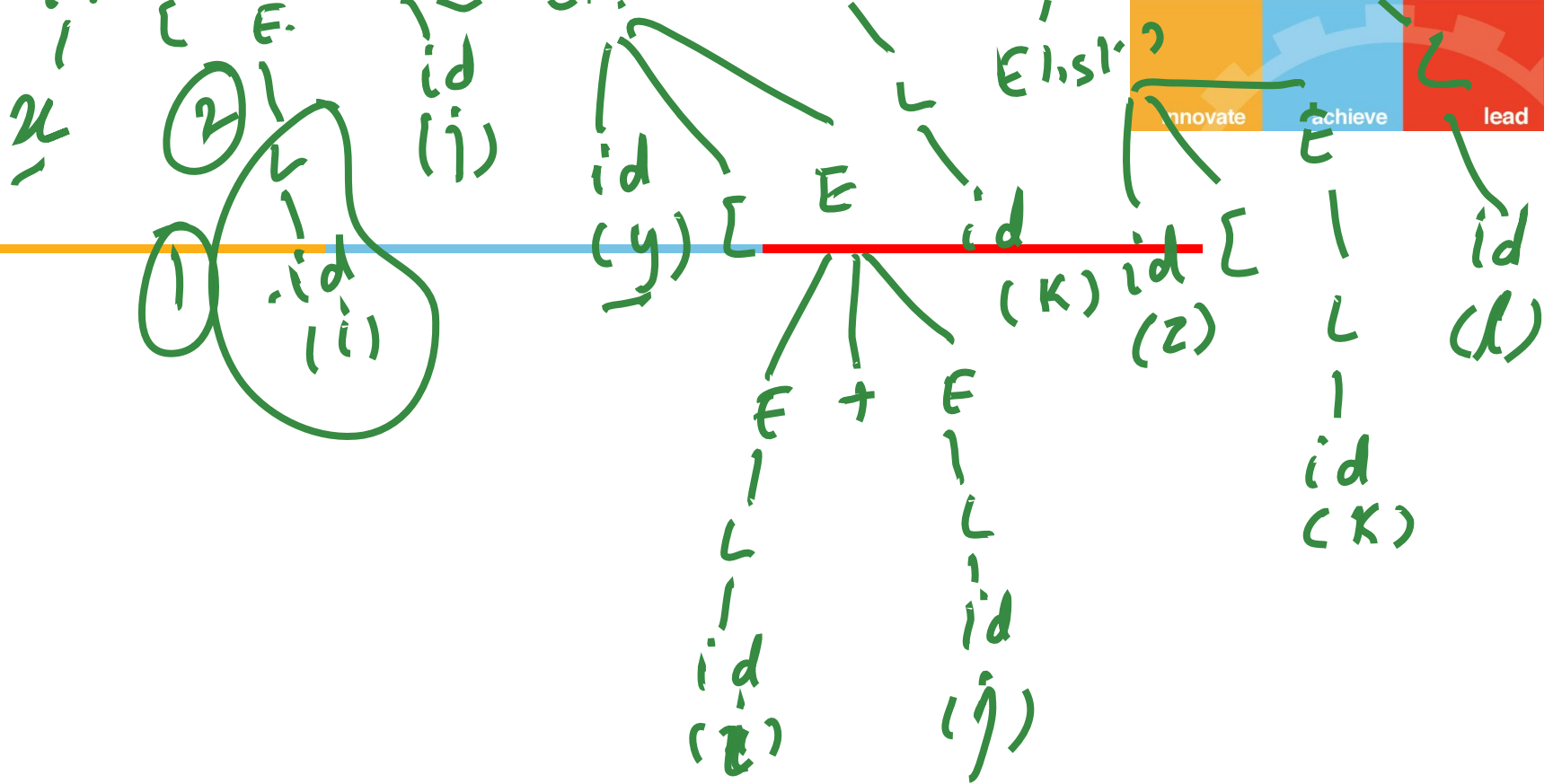
$x(d_1, d_2) \quad y(d_3, d_4) \quad z(d_5, d_6)$

$$x[i, j] = y[i + j, k] + z[k, l] \quad \text{e.g.}$$

$\Leftarrow$  parse Tree.







## ⇒ Semantic Actions.

①  $\begin{cases} h.place = i \\ h.offset = null \end{cases}$

②  $E.place : h.place = i$

③  $E.list = \dots$

- ③
- $\text{list.array} = x$
  - $\text{place} = i$
  - $\text{dim} = 1$  ✓



④  $\text{h.place} = j$  ,  $\text{L.offset} = \text{null}$ .

⑤  $\text{t.place} = j$

⑥  $\text{Elist.array} = x$   
 $\text{elist.place} = t_1 \Rightarrow$   $\begin{cases} t_1 = i * \text{limit}(x, m) \\ t_1 = i * d_2 \\ t = t_1 + j \end{cases}$   
 $\text{elist.dim} = 2$

⑦  $\text{L.place} = t_2$   
 $\text{L.offset} = t_3 \Rightarrow \begin{cases} t_2 = c(x) \\ t_3 = t_1 * w \end{cases}$

⑤ h.place = i  
L.offset = null

④ ..... so on.

final IR code for given 2D array Exp.

$$t_1 = i * d_2$$

$$t_1 = t_1 + j$$

$$t_2 = C(x)$$

$$t_3 = t_1 * w$$

$$t_4 = i + j$$

$$x_5 = t_4 * d_4$$

$$t_5 = x_5 + K$$



$$x_6 = c(r)$$

$$t_7 = t_5 * w$$

$$t_8 = x_6[t_7]$$

$$t_9 = K * d_6$$

$$t_9 = t_9 + l$$

$$t_{10} = c(z)$$

$$t_{11} = t_9 * w$$

$$t_{12} = t_{10}[t_{11}]$$

$$t_{13} = t_{12} + t_{11}$$

$$t_{13} = A_{12} + A_{23}$$

$$t_2[t_3] = t_{13}$$



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# Thank You!