A Companison of Time Complenities of Greedy, Dynamic Programming, and Divide and Longuer: Suppose for a given problem we have threeolgo-ithous by using the abone algorithm design techniques: a Greedy Algorithm, a Dynamic Programming Algorithm, and a Divide and Longuer Algorithm. Divide and longuer of gorithm will be least efficient because it is solving each subpostern even if it is an overlopping subproblem. Dynamic Programming of gorithm will be more efficient than Divide and conquer algorithm because it is avoiding the repeated solutions of subposblems. If we compare the Greedy algorithm with the Dynamic Programming algorithm, the areedy of gorithm will be more efficient than the Pyramic Programming organithm because in Creedy of gorithm we do not solve any subpudsiem I'm general, the time complexities will be: Greedy < Pyramic Programming < Divide & Conquer Example: Consider the Single Source Shortest Poths Pholon Dijkstras Algorithmy Greedy and Los complexity O (1E/6g/V) Bellman-Ford Algorithm is Dynamic Regramming I govithm and has complementy O(IVI(EI) Under the orsumption of $|E| = o(|V|^2)$, we condesign a Divide and Conquer adjouthm by wing the recurrence of Bellman-Ford Algorithm of Complents V2((VI)). The Comparison of three organithms (using E/20(VI)) Greedy: O(|V| 26g |V|) < Dynamic O((V)) < P& CNO((V) W),