



## **Lecture 6**

# **Expression and statement constructs**

# List of constructs in a programming language

- Expressions
- Function definitions and call statements (returning values/ not returning any value)
- Statements
  - Declaration statements
  - Assignment ( $a=b*c+d-f/s$ ; uni-processor view, instruction cycle, semantic)
  - Iterative statements (For, while, do while, repeat until)
  - Conditional statement (If, switch-case values)
  - Goto statements
  - Return, break statements
- Variables
- Datatypes (pointers, records, arrays, objects, lists etc.)
- Separator

# Program elements from language designer's perspective



- **Variables**-characters and words, length and patterns
- **Numbers**-integer, real, precision, form (23.56, 2.56e+2)
- **Operators**- patterns and meaning (=, ==, <=, >=, &&, ||), associativity and precedence of ops
- **Keywords**-patterns, names. why?
- **Grammar**-syntactic structure
- **Constructs for computation** – expressions,
- **Constructs for execution flow**
- **Semantic privileges**-precedence of function def before its call, recursive, overloading etc.
- **Type system**

# Grammar

- A **grammar** is defined as a four valued tuple  $(N, T, S, P)$  where  $N$  is the set of non-terminals,  $T$  is the set of terminals,  $S$  is the start symbol and  $P$  is the set of production rules.
- **Example:** Grammar  $G$  to generate palindrome strings over alphabet  $\{0,1\}$  is defined as below

$$G = (\{S\}, \{0,1\}, S, P)$$

Where  $p$  is

$$S \rightarrow 0 S 0$$

$$S \rightarrow 1 S 1$$

$$S \rightarrow 0$$

$$S \rightarrow 1$$

# Understanding expression construct



- Examples:

**a** *(one term)*

**b\*c** *(one term with a binary operator)*

**a+b\*c** *(one expression with two terms and +)*

**a-d+b\*c** *(one expression with one expression and two terms and +)*

**a+b\*c-d**

# Recursive construction of an expression



- An expression
  - can be a term
  - Can be a combination of two terms with a binary operator
  - A term can also be an expression
- Grammar rules (P) to derive an expression are given below
  1.  $\langle \text{expression} \rangle \rightarrow \langle \text{term} \rangle$
  2.  $\langle \text{expression} \rangle \rightarrow \langle \text{term} \rangle \langle \text{operator} \rangle \langle \text{term} \rangle$
  3.  $\langle \text{term} \rangle \rightarrow \langle \text{expression} \rangle$
  4.  $\langle \text{term} \rangle \rightarrow \text{ID}$  //ID is a token name for a,b,c and d
  5.  $\langle \text{operator} \rangle \rightarrow \text{PLUS}$
  6.  $\langle \text{operator} \rangle \rightarrow \text{MINUS}$
  7.  $\langle \text{operator} \rangle \rightarrow \text{MUL}$
- Grammar  $G = (\{\langle \text{expression} \rangle, \langle \text{term} \rangle, \langle \text{operator} \rangle\}, \{\text{ID}, \text{PLUS}, \text{MINUS}, \text{MUL}\}, \langle \text{expression} \rangle, P)$

# What do we mean by deriving an expression using the underlying grammar?

**a+b\*c-d**

$\langle \text{expression} \rangle \rightarrow \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{expression} \rangle$

$\langle \text{expression} \rangle \rightarrow \langle \text{term}_1 \rangle \langle \text{operator} \rangle \langle \text{term}_2 \rangle$

$\langle \text{term}_1 \rangle \rightarrow a$  //token ID is taken in actual

$\langle \text{operator} \rangle \rightarrow \text{PLUS}$

$\langle \text{term}_2 \rangle \rightarrow \langle \text{expression} \rangle$

$\langle \text{expression} \rangle \rightarrow \langle \text{term}_1 \rangle \langle \text{operator} \rangle \langle \text{term}_2 \rangle$

$\langle \text{term}_1 \rangle \rightarrow \langle \text{expression} \rangle$  // for b\*c

$\langle \text{expression} \rangle \rightarrow \langle \text{term}_1 \rangle \langle \text{operator} \rangle \langle \text{term}_2 \rangle$

$\langle \text{term}_1 \rangle \rightarrow b$

$\langle \text{operator} \rangle \rightarrow \text{MUL}$

$\langle \text{term}_2 \rangle \rightarrow c$

$\langle \text{term}_2 \rangle \rightarrow d$

$\langle \text{operator} \rangle \rightarrow \text{MINUS}$

# Parse tree for $a+b*c-d$

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- Home work



# Redefining grammar for expression

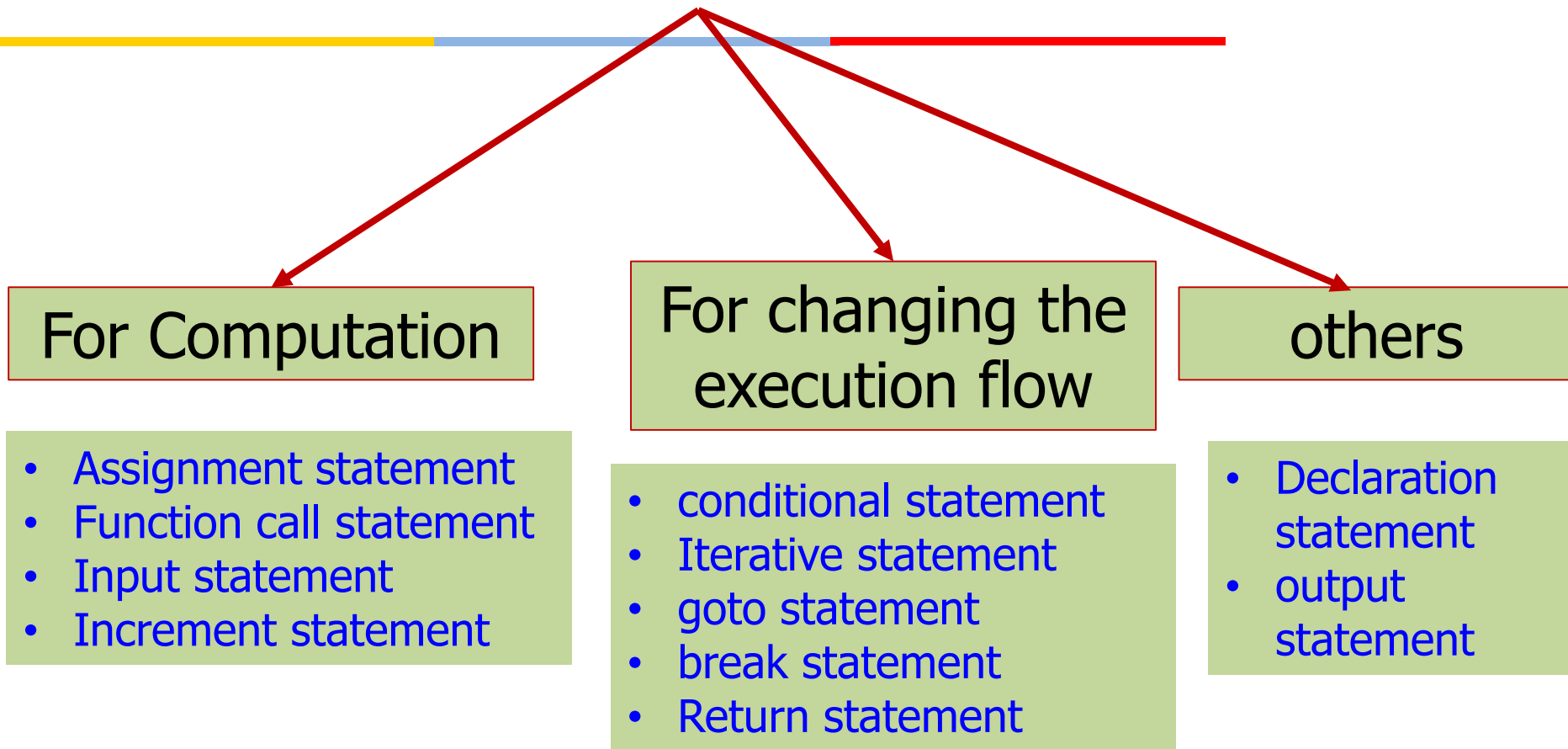
- Grammar rules (P) to derive an expression are given below
  1.  $\langle \text{expression} \rangle \rightarrow \langle \text{expression} \rangle \langle \text{operator1} \rangle \langle \text{term} \rangle$
  2.  $\langle \text{expression} \rangle \rightarrow \langle \text{term} \rangle$
  3.  $\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle \langle \text{operator2} \rangle \langle \text{factor} \rangle$
  4.  $\langle \text{term} \rangle \rightarrow \langle \text{factor} \rangle$
  5.  $\langle \text{factor} \rangle \rightarrow \text{ID}$
  6.  $\langle \text{operator1} \rangle \rightarrow \text{PLUS} \mid \text{MINUS}$
  7.  $\langle \text{operator2} \rangle \rightarrow \text{MUL} \mid \text{DIV}$
- Grammar  $G = (\{\langle \text{expression} \rangle, \langle \text{term} \rangle, \langle \text{operator1} \rangle, \langle \text{operator2} \rangle, \langle \text{factor} \rangle\}, \{\text{ID}, \text{PLUS}, \text{MINUS}, \text{MUL}\}, \langle \text{expression} \rangle, P)$

# Design issues for expression construct



- Which operators are to be included for computations?
- What are the operator associativity rules?
- What are the operator precedence rules?
- Does the language allow operator overloading?
- What will be type rules for operators?  
Whether type conversion will be allowed implicitly?

# 'Statements' construct



# Computation based statements

- The most general form of an assignment statement has at least one variable on the left hand side of the assignment operator.

example `a=b+c*d;`

`a=sum(b,c,d);`

- The RHS of the statement gets the computed value which is then stored in the LHS variable.
- The RHS can be an **expression** or a **function call**.
- Other forms of statements which can change the value of a variable are read statements and increment statements.

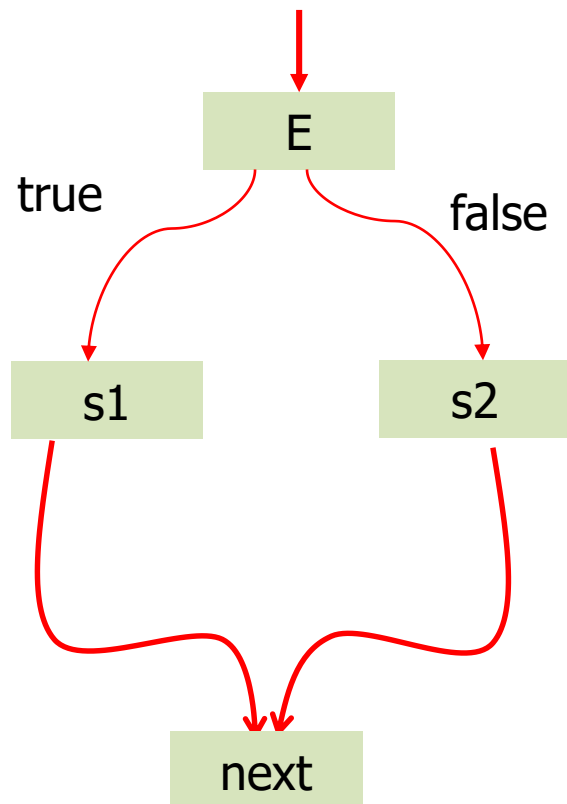
# Statements that handle the execution flow

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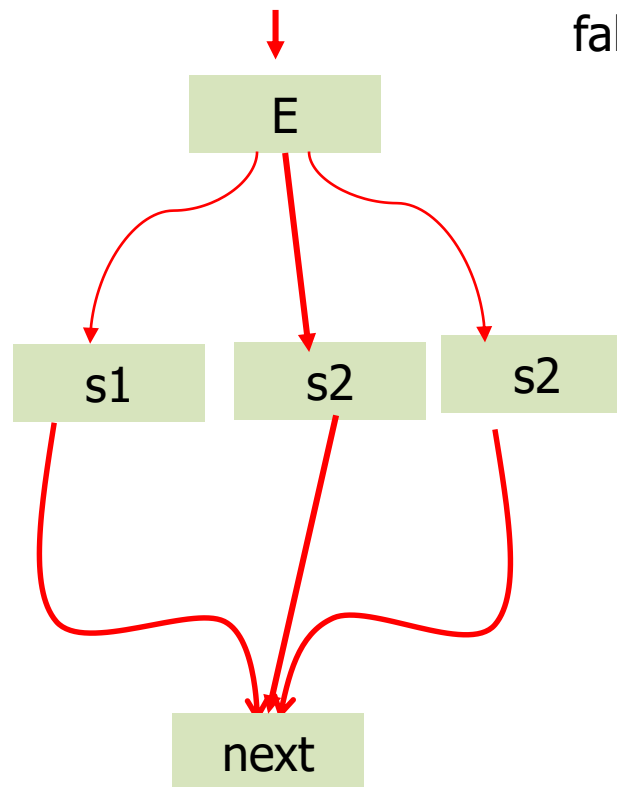


- These statements are called “control statements”
- All such statements exhibit “single entry-single exit” control flow.
- Example statements that change the flow of execution: if-then-else, switch, for, while, break, return etc.

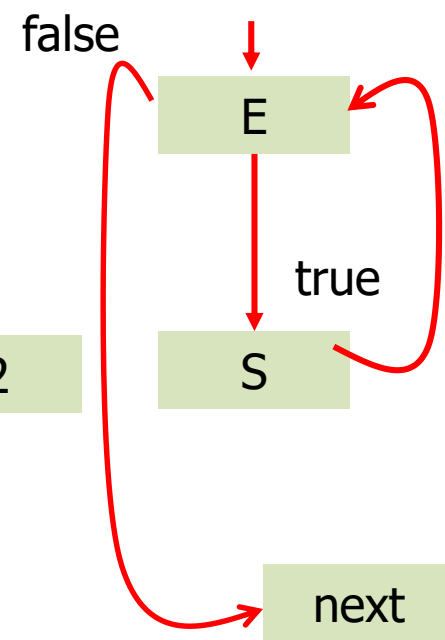
# Flow of execution



If statement



switch  
statement



While  
statement

# Design issues

- The execution flow should be evident from the syntax.
- Selection statements
  - What should be the form and type of expression that controls selection?
  - How should the meaning of nested selector be specified?
- Iterative statements
  - What are the type and scope of loop variables?
  - Can the loop variable be reassigned a value within the loop body?
  - Should the loop parameter be evaluated only once or once for every iteration?
  - Should the loops be counter controlled or logic controlled or both?

# Assignment statement

- Syntax

$\langle \text{assignment\_stmt} \rangle \rightarrow \text{ID ASSIGNOP} \langle \text{expression} \rangle$

- Semantics

Compute the value of  $\langle \text{expression} \rangle$  and copy the value in the memory location bound to the variable identifier (tokenized as ID)

- Example

$\text{abc\_1} = \text{xyz} + \text{qr} * \text{pr} \quad (\text{ID} = \text{ID} + \text{ID} * \text{ID})$

- Nonterminal on the LHS of the grammar rule represents the construct:  $\langle \text{assignment\_stmt} \rangle$



# Function call statement

- Syntax

`<functioncall_stmt> → ID ASSIGNOP FUNCTION_ID <actual para_list>`

- Semantics

Copy the value returned by the function in the memory location bound to the variable identifier (tokenized as ID)

- Example

`abc_1 = sum(a,b)`

- Nonterminal on the LHS of the grammar rule represents the construct: `<functioncall_stmt>`

# Iterative statements (FOR)

- Syntax

**<iterative\_for\_stmt>** → FOR OP <c1> semicolon <c2> semicolon <c3> CP <statements>

**<c1>** → ID ASSIGNOP <expression>

**<c2>** → <expression> <relationalOp> <expression>

**<c3>** → <assignment\_stmt>

**<relationalOp>** → LT|GT|EQ|LE|GE|NE

- Semantics

Initialize <c1> and execute only once

Repeat the execution of <statements> based on <c2>

Execute <c3> everytime before leaving the loop.

- Example

for (x=0; x<a+10; x=x+2) y=x+a\*2;

- Nonterminal on the LHS of the grammar rule represents the construct: **<iterative\_for\_stmt>**