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Equivalence of Regular Expression and NFA

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Designing of Regular Expression (Cont.....)



Construct a regular expression that accepts all strings of 0's and 1's where each string has atleast two 0's

Step 1

- $L = \{ 00, 01011, 11100, \dots \}$

Step 2

- Minimal length String is 00.

Step 3

- $(0+1)^*.0.(0+1)^*.0.(0+1)^*$

Designing of Regular Expression (Cont.....)



Construct a regular expression that accepts all strings of 0's and 1's where each string has '1' at every odd position and the length of the string is odd.

Step 1

- $L = \{ 1, 101, 111, 10111, 10101, \dots \}$

Step 2

- Minimal length String is 1

Step 3

- $1.((0+1).1)^*$

Algebraic Properties of Regular Expression



- $R1 + (R2 + R3) = (R1 + R2) + R3$
- $R1(R2R3) = (R1R2)R3$
- $R1(R2 + R3) = R1R2 + R1R3$
- $(R1 + R2)R3 = R1R3 + R2R3$
- $R1 + R2 = R2 + R1$
- $(R1^*)^* = R1^*$
- $R1.\epsilon = \epsilon.R1 = R1$
- $R1 + \emptyset = R1$

Regular Expressions (REs) and Regular Languages



A language L is regular if and only if $L = L(R)$ for some regular expression R .

REs are equivalent in power to NFAs/DFAs.

Converting RE to NFA

Given a regular expression, we will convert it into a NFA N such that $L(R) = L(N)$.

We will give a case based analysis based on the inductive definition of REs.

Case 1



- Case 1: $R = \phi$



Case 2



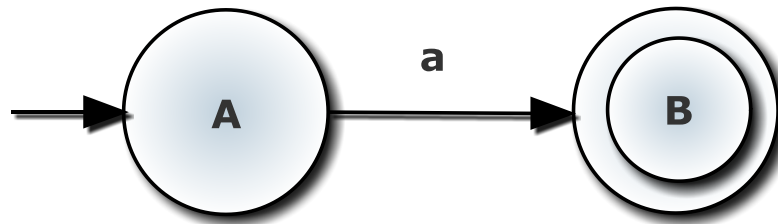
- Case 2: $R = \epsilon$



Case 3



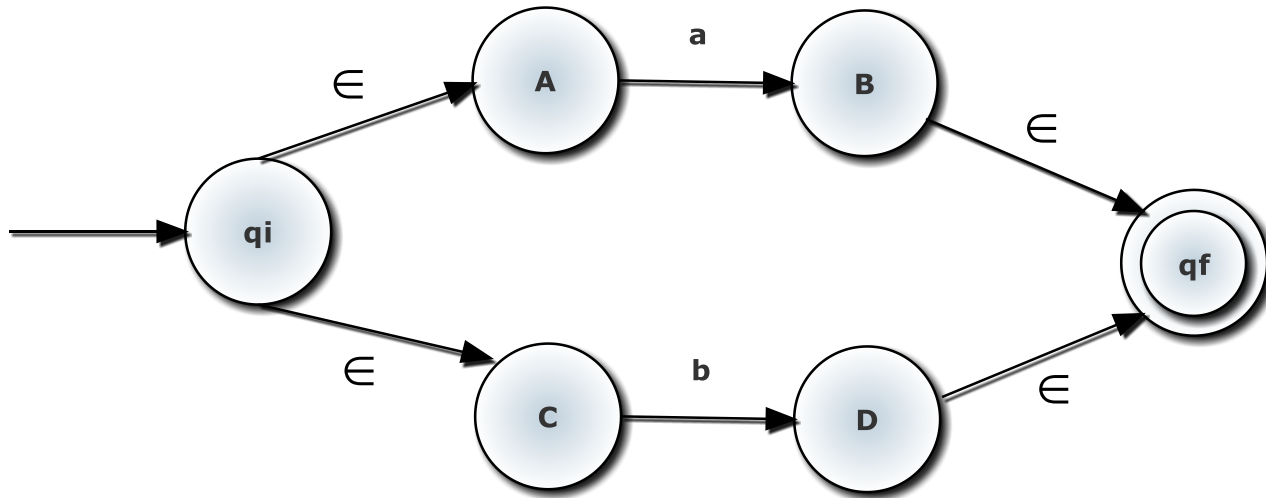
- Case 3: $R = a$



Case 4



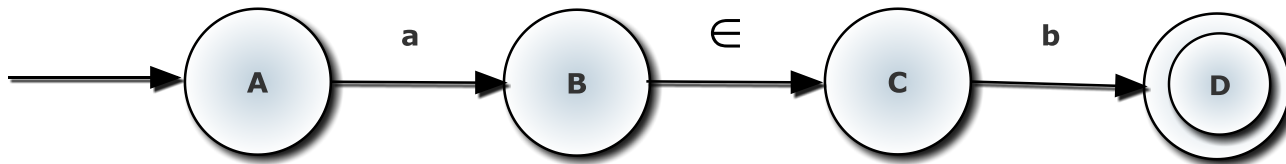
- Case 4: $R = a + b$



Case 5



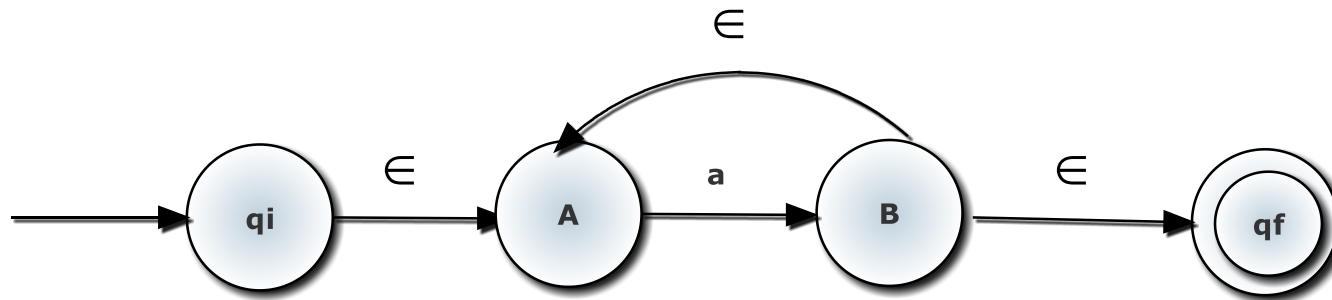
- Case 5: $R = a.b$



Case 6



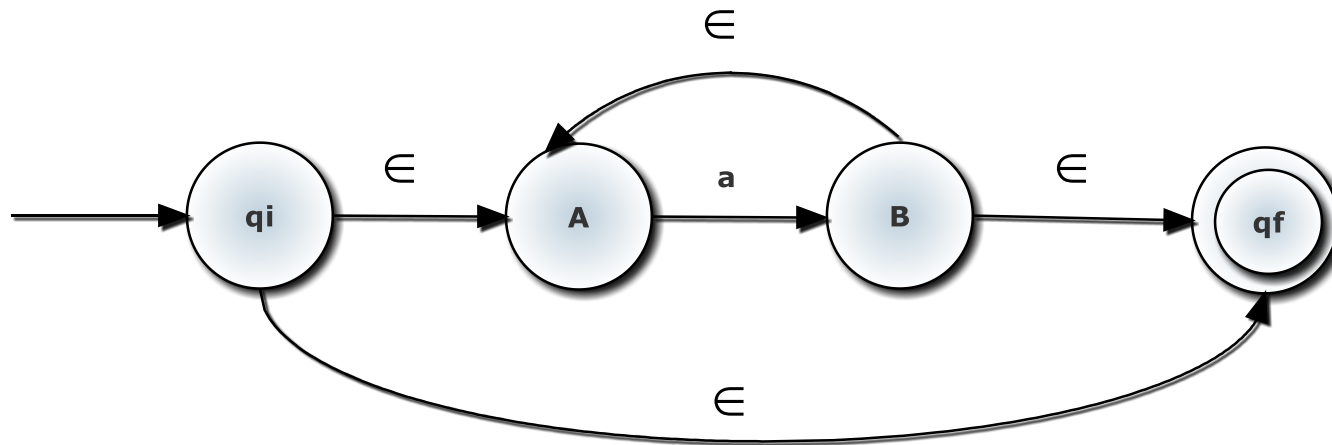
- Case 6: $R = a^+$



Case 7



- Case 7: $R = a^*$



Convert the following regular expression to epsilon NFA

$$(00 + 1)^* 1 (0 + 1)$$