# **OSN Assignment 5**

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# **Report : Question 2 (with bonus)**

- This problem can be broadly broken down into 2 parts, one from the spectators viewpoint and one in terms of the goals in the match.
- So I create <a href="num\_spectators">num\_spectators</a> threads which handle the routine of every spectator via the function <a href="spectator-routine">spectator-routine</a> which takes the spectator structure of its respective spectator as an argument.
- Then there is one thread goals which handles the goals scored in the match as per the goal scoring chances via the function goals\_routine which takes the array of structure goalscoring chance as an argument.
- I have also num\_groups threads called groups\_exit[] which handles the exiting of a group in the simulation ( part of bonus ).
  It uses group\_exit\_routine which takes its respective groups, group element as an argument.

# A. Spectator Routine

This can be further broken down to following sections:

- 1. Arrival of spectator
- 2. Spectator waiting for a seat
- 3. Spectator watching the match
- 4. Spectator at exit gate

#### Arrival of spectator

• For this, I make the thread corresponding to its spectator sleep till the time of arrival of the spectator. Once, sleep finishes, we can assume that the spectator has arrived.

```
sleep(s.time_of_arrival);
printf(ANSI_COLOR_RED "t=%ld : %s has reached the stadium\n" ANSI_COLOR_RESET, time_from_start(), s.name);
```

#### Spectator waiting for a seat

• For this I make use of the function, <code>sem\_timedwait</code> ( refer to man pages ) which takes a semaphore and timespec structure as an argument. I used the timespec structure to store the patience limit of the spectator waiting for their seat.

```
struct timespec ts;
if (clock_gettime(CLOCK_REALTIME, &ts) == -1)
{
    perror("clock gettime");
    exit(EXIT_FAILURE);
}
ts1.tv_sec += s.patience;
```

- sem\_timedwait waits for the counter to go above zero till a specified time ( set via the timespec struct ). If the function returns 0, that means it received sem\_post before timeout. But if the function returns -1 and errno == ETIMEDOUT, that means it wasn't able to receive sem\_post by the specified time.
- As we know, Home Fan → Home, Neutral seats
- Neutral Fan → Home, Neutral, Away seats
- Away Fan → Away seats

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- Since H can access from a pool of H,N seats, N from H,N,A seats and A from only A seats, I created three semaphores:
  - · HN for Home fan
  - o HNA for Neutral fan
  - o A for Away fan

```
sem_init(&HN, 0, Hzone.capacity + Nzone.capacity);
sem_init(&HNA, 0, Hzone.capacity + Nzone.capacity + Azone.capacity);
sem_init(&A, 0, Azone.capacity);
```

- · Now for a Home fan,
  - I use sem\_timedwait on HN and wait till its patience limit to receive a post on it.
  - If a timeout occurs, I send the fan to the exit gate.
  - If it receives a post before timeout, I lock mutexes H\_lock and N\_lock as I will now access the capacity of Zones H and N.
  - o If both have seats available, I'll chose randomly which seat to occupy.
  - o If only one seat is available, I'll choose that one only.
  - After selecting the seat, I increment the value of currently filled seats in that zone and also do sem\_wait on another semaphore.
    - This is because if I occupy seat H from the pool H and N (available to a Home fan), I also have to decrease the counter for the pool H,N,A as it also has seats for H zone. Thus, I do sem\_wait(&HNA)
  - Then I unlock the mutexes and move to the next stage.
- Similar steps are done for Neutral and Away fans (only different pool of seats).

#### Spectator watching the match

- Here I use <a href="https://pthread\_cond\_timedwait">pthread\_cond\_timedwait</a> which is another function that makes use of struct timespec to set timeout time to <a href="mailto:spectating\_time">spectating\_time</a> of the spectators. Returns 0 if successful, else <a href="mailto:error">error</a>
- I define two condition variables Agoal\_cond and Hgoal\_cond, locked by Agoal\_lock and Hgoal\_lock for away and home goals respectively.
- · For Spectator Type H,

```
pthread_mutex_lock(&Agoal_lock);
int to = time_from_start();
while (A_goals < s.enrage_limit && j != ETIMEDOUT)
    j = pthread_cond_timedwait(&Agoal_cond, &Agoal_lock, &ts2); // see if away goal is scored
int t1 = time_from_start();
pthread_mutex_unlock(&Agoal_lock);</pre>
```

- I wait for a signal on the condition variable from thread <code>goals</code>, and check in a while loop if it has timedout or the goals by opponent team is still less that enrage\_limit of the spectator.
- Once it breaks out of the loop, I also take an account of the time it stayed in the loop and see if it is greater than the timeout limit.
  - This is done because at times I will get a signal on my cond\_variable but still it might not break out of the loop if goals
    are less than enrage limit of the spectator. Still a timeout, just not shown by j == ETIMEDOUT

```
if ((t1 - t0) >= spectating_time || j == ETIMEDOUT)
{
    // Timed out
    // stands up from his/her seat
    s.is_seated = 0;
    printf(ANSI_COLOR_GREEN "t=%ld : %s watched the match for %d seconds and is leaving\n" ANSI_COLOR_RESET, time_from_start(), s.na
} else
```

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```
{
  // Gets Enraged
printf(ANSI_COLOR_GREEN "t=%ld : %s is leaving due to bad performance of his team\n" ANSI_COLOR_RESET, time_from_start(), s.name
  s.is_seated = 0;
}
```

- Similarly it is done for Away spectator. But for a neutral spectator ( no enrage limit ), I just sleep till its spectating time and send him to exit gate.
- For vacating his/her seat,

```
if (s.seated_zone == 'H')
{ \  \  \, //\  \,  is part of pools HN and HNA
    pthread_mutex_lock(&H_lock);
    Hzone.currently_filled-
    pthread_mutex_unlock(&H_lock);
    sem_post(&HN);
    sem_post(&HNA);
else if (s.seated_zone == 'N')
{ // is part of pools HN and HNA
    pthread_mutex_lock(&N_lock);
    Nzone.currently_filled--;
    pthread_mutex_unlock(&N_lock);
    sem_post(&HN);
    sem_post(&HNA);
else
{ // is part of pool A
    pthread_mutex_lock(&A_lock);
    Azone.currently_filled--;
    pthread_mutex_unlock(&A_lock);
    sem_post(&A);
    sem_post(&HNA);
```

#### Spectator at exit gate

• I make use of mutex locks group\_lock[s.group\_id] for each group of spectators. And also conditional variables
group\_cond[s.group\_id] to signal an update in the spectators at exit for a particular group to the groups\_exit thread.

```
// Exit the stadium
s.is_at_exit = 1;
printf(ANSI_COLOR_BLUE "t=%ld : %s is waiting for their friends at the exit\n" ANSI_COLOR_RESET, time_from_start(), s.name);
pthread_mutex_lock(&group_lock[s.group_id]);
group_members_at_exit[s.group_id]++;
// Signals the exit_routine
pthread_cond_broadcast(&group_cond[s.group_id]);
pthread_mutex_unlock(&group_lock[s.group_id]);
```

#### **B.** Goals Routine

• In this I iterate through the goalscoring\_chance array and sleep till the time for the next goalscoring chance has come.

```
int current_time = time_from_start();
int time_to_wait = g[1].time_since_start - current_time;
if (time_to_wait > 0)
    sleep(time_to_wait);
```

- Now when the chance has come, I generate a random float number between 0 and 1 (inclusive) and use it as probability. If it is less than or equal to the probability of goal to occur, a goal is given.
- I check the team which has the chance and give it a goal if it satisfies the probability. Then I also broadcast a signal to its respective condition variable.

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```
if (probability <= g[i].probability)
{
    if (g[i].type == 'H')
    {
        pthread_mutex_lock(&Hgoal_lock);
        H_goals++;
        pthread_cond_broadcast(&Hgoal_cond);
        pthread_mutex_unlock(&Hgoal_lock);
    }
    else
    {
        pthread_mutex_lock(&Agoal_lock);
        A_goals++;
        pthread_cond_broadcast(&Agoal_cond);
        pthread_mutex_unlock(&Agoal_lock);
    }
    printf("t=%ld : Team %c has scored their %s goal\n", time_from_start(), g[i].type, key);
} else
    {
        printf("t=%ld : Team %c missed the chance to score their %s goal\n", time_from_start(), g[i].type, key);
}</pre>
```

• I keep iterating till all chances are over.

### C. Group Exit Routine

- I make use of mutex locks and conditional variables for each group to check and wait for all the members of a particular group to reach the exit gate.
- Once all group members are at exit, I break out of the loop and make the group leave for dinner.

```
// Wait for all spectators of a group to leave
pthread_mutex_lock(&group_lock[g->group_id]);
while (group_members_at_exit[g->group_id] < g->group_size)
    pthread_cond_wait(&group_cond[g->group_id], &group_lock[g->group_id]);
pthread_mutex_unlock(&group_lock[g->group_id]);

// Now the time has come
printf(ANSI_COLOR_YELLOW "t=%ld : Group %d is leaving for dinner\n" ANSI_COLOR_RESET, time_from_start(), (g->group_id + 1));
```

# **Main function**

- In the main function I take inputs, initialize threads, locks, condition variables and semaphores.
- I create threads and join them.
- $\bullet\,$  I also free the pointers ( used for passing args to threads) allocated via malloc.

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