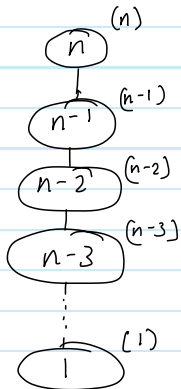


# Sk10101\_hw4\_q1\_q2

Tuesday, October 25, 2022 9:51 PM

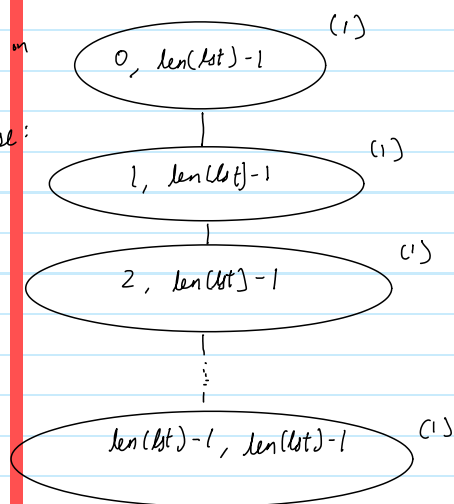
Q1)

2) (i) For  $\text{sum\_lst1}(\text{lst})$ :

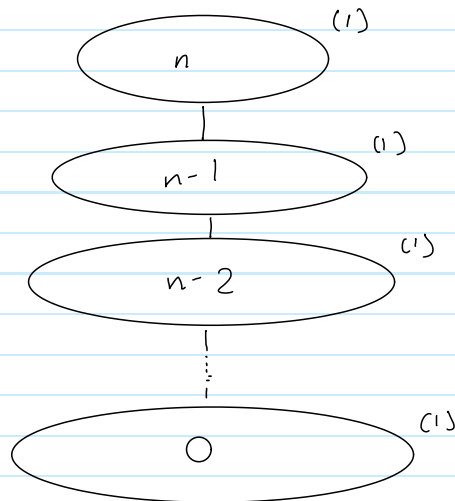


For  $\text{sum\_lst2}(\text{lst}, \text{low high})$ :

Based on  
TA's  
response:



Based on professor's approach:



(ii) For  $\text{sum\_lst1}(\text{lst})$ :

For each recursion of the function over  $n$  items,  
we slice the list.

$$\therefore \text{Total cost} = n + n-1 + n-2 + \dots + 1 = \frac{n(n+1)}{2} = \frac{n^2}{2} + \frac{n}{2} = \Theta\left(\frac{n^2}{2} + \frac{n}{2}\right)$$

$$= \Theta(n^2)$$

For  $\text{sum\_lst2}(\text{lst})$ :

For each recursion of the function, we have  $\Theta(1)$  implementations over  
 $n$  times

$$\therefore \text{Total cost} = \boxed{\Theta(n)}$$

3) Since  $n^2 > n$ ,  $\text{sum\_lst2}(\text{lst})$  will be asymptotically faster