# **Distributed Computing**

# Experiment – 7

Prog	gram to demonstrate Deadlock Management in Distributed Systems			
Learning Objective:	To explore mutual exclusion algorithms and deadlock handling in the distributed system.			
Learning Outcome:	Ability to demonstrate Deadlock Management in Distributed Systems			
Course Outcome:	CSL801.4			
Program Outcome:	(PO1) Engineering knowledge: Apply the knowledge of mathematics, scient engineering fundamentals, and an engineering specialization to the solution complex engineering problems.			
	(PO2) Problem analysis: Identify, formulate, review research literature, and analyze complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.			
	(PO3) Design/development of solutions: Design solutions for complex			
	gineering problems and design system components or processes that meet the			
	specified needs with appropriate consideration for the public health and safety, and			
	the cultural, societal, and environmental considerations.			
	<b>(PO5) Modern tool usage</b> : Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools including prediction and modeling to complex engineering activities with an understanding of the limitations.			
Bloom's Taxonomy Level:	Analysis, Create			
Theory:	Deadlock Management in Distributed Systems in details.			
	Assigned to Resource 1 Waiting for			
	Process 1  Process 2  Resource 2  Assigned to			
	Resource 2			

A deadlock is a situation that arises in concurrent systems, including distributed systems, where two or more processes or threads are waiting indefinitely for one another to release resources that they need to proceed. This leads to a complete standstill, where none of the processes can make progress, resulting in a system that is essentially hung or frozen.

## **Types of Deadlocks**

- 1. **Resource Deadlock**: This type of deadlock occurs when two or more processes are waiting for resources that are held by other processes, creating a circular wait condition.
- Communication Deadlock: In distributed systems, processes often
  communicate with each other by sending and receiving messages. A
  communication deadlock can occur when processes are waiting for
  messages that will never arrive due to a cyclic dependency in the
  communication pattern.
- 3. **Deadlock by No Indefinite Delay**: This type of deadlock occurs when a process or thread is waiting for a resource that will never become available due to a programming error or a system failure.

# **Deadlock Management Procedure**

The deadlock management procedure typically involves the following steps:

- 1. **Deadlock Detection**: The first step is to detect the presence of a deadlock in the system. This can be achieved through various algorithms, such as resource allocation graphs or wait-for graphs.
- 2. **Deadlock Resolution**: Once a deadlock is detected, the system must resolve the deadlock by breaking the circular wait condition. This can be done through several strategies, such as process termination, resource preemption, or rollback.
- Deadlock Prevention: To prevent future deadlocks, the system can employ various techniques, such as careful resource allocation, avoidance of circular wait conditions, or enforcing ordering among resource acquisition.
- 4. **Deadlock Avoidance:** This involves dynamically monitoring resource allocation and denying requests that could potentially lead to a deadlock situation.

# Types of Deadlock Management

- 1. **Deadlock Prevention**: This approach aims to prevent deadlocks from occurring in the first place by enforcing certain rules or protocols for resource allocation. Examples include the "no preemption" policy, the "hold and wait" policy, and the "circular wait" policy.
- 2. **Deadlock Avoidance**: In this approach, the system dynamically examines

resource allocation requests and grants or denies them based on an analysis of whether granting the request could lead to a deadlock. The Banker's Algorithm is a well-known example of a deadlock avoidance strategy.

- 3. **Deadlock Detection and Resolution**: This approach allows deadlocks to occur but detects them as soon as possible and resolves them through techniques like process termination, resource preemption, or rollback.
- 4. **Deadlock Ignorance**: In some systems, deadlocks are considered so rare or their impact so minimal that the system simply ignores them. This approach is typically used in systems where deadlocks are not a significant concern or where the cost of prevention or detection is deemed too high.

### Algorithm:

- 1. Let Work and Finish be vectors of length m and n respectively. Initialize Work= Available. For i=0, 1, ...., n-1, if Requesti = 0, then Finish[i] = true; otherwise, Finish[i] = false.
- Find an index i such that both
   a) Finish[i] == false
   b) Requesti <= Work</li>
  - If no such i exists go to step 4.
- Work= Work+ Allocationi
   Finish[i] = true
   Go to Step 2.
- 4. If Finish[i] == false for some i, 0 <= i < n, then the system is in a deadlocked state. Moreover, if Finish[i] == false the process Pi is deadlocked.

### Outcome:

#### **Program Code -**

## 1. Example for Deadlock -

import java.util.concurrent.locks.Lock;

import java.util.concurrent.locks.ReentrantLock;

public class DeadlockExample {

private Lock lock1 = new ReentrantLock();

```
private Lock lock2 = new ReentrantLock();
  public static void main(String[] args) {
    DeadlockExample deadlock = new DeadlockExample();
    new Thread(deadlock::operation1, "T1").start();
    new Thread(deadlock::operation2, "T2").start();
  }
  public void operation1() {
    lock1.lock();
     System.out.println(Thread.currentThread().getName() + ": lock1 acquired,
waiting to acquire lock2.");
    sleep(50);
    lock2.lock();
     System.out.println(Thread.currentThread().getName() + ": lock2
acquired");
     System.out.println(Thread.currentThread().getName() + ": executing first
operation.");
    lock2.unlock();
    lock1.unlock();
  }
  public void operation2() {
```

```
lock2.lock();
    System.out.println(Thread.currentThread().getName() + ": lock2 acquired,
waiting to acquire lock1.");
    sleep(50);
    lock1.lock();
    System.out.println(Thread.currentThread().getName() + ": lock1
acquired");
    System.out.println(Thread.currentThread().getName() + ": executing
second operation.");
    lock1.unlock();
    lock2.unlock();
  }
  // helper methods
  private void sleep(int milliseconds) {
    try {
       Thread.sleep(milliseconds);
     } catch (InterruptedException e) {
       e.printStackTrace();
    }
  }
```

```
2. Example for Deadlock Resolution -
```

```
import java.util.Random;
import java.util.concurrent.locks.Lock;
import java.util.concurrent.locks.ReentrantLock;
public class LiveLockExample {
  private Lock lock1 = new ReentrantLock();
  private Lock lock2 = new ReentrantLock();
  public static void main(String[] args) {
    LiveLockExample livelock = new LiveLockExample();
    new Thread(livelock::operation1, "T1").start();
    new Thread(livelock::operation2, "T2").start();
  }
  public void operation1() {
    while (true) {
       tryLock(lock1, 10000); // Increase timeout to 10 seconds
       System.out.println(Thread.currentThread().getName() + ": lock1
acquired, trying to acquire lock2.");
       sleep(50);
       if (tryLock(lock2, 10000)) { // Increase timeout to 10 seconds
```

```
System.out.println(Thread.currentThread().getName() + ": lock2
acquired.");
          break;
       } else {
         System.out.println(Thread.currentThread().getName() + ": cannot
acquire lock2, releasing lock1.");
          lock1.unlock();
         randomDelay();
       }
     }
     System.out.println(Thread.currentThread().getName() + ": executing first
operation.");
    lock2.unlock();
    lock1.unlock();
  }
  public void operation2() {
    while (true) {
       tryLock(lock2, 10000); // Increase timeout to 10 seconds
       System.out.println(Thread.currentThread().getName() + ": lock2
acquired, trying to acquire lock1.");
       sleep(50);
       if (tryLock(lock1, 10000)) { // Increase timeout to 10 seconds
          System.out.println(Thread.currentThread().getName() + ": lock1
acquired.");
          break;
```

```
} else {
          System.out.println(Thread.currentThread().getName() + ": cannot
acquire lock1, releasing lock2.");
         lock2.unlock();
         randomDelay();
       }
    }
    System.out.println(Thread.currentThread().getName() + ": executing
second operation.");
    lock1.unlock();
    lock2.unlock();
  }
  // helper methods
  private boolean tryLock(Lock lock, long timeout) {
    try {
       return lock.tryLock(timeout,
java.util.concurrent.TimeUnit.MILLISECONDS);
    } catch (InterruptedException e) {
       return false;
    }
  }
  private void sleep(int milliseconds) {
    try {
       Thread.sleep(milliseconds);
```

```
} catch (InterruptedException e) {
    e.printStackTrace();
}

private void randomDelay() {
    try {
        Thread.sleep(new Random().nextInt(100)); // Introduce a random delay to break symmetry
    } catch (InterruptedException e) {
        e.printStackTrace();
    }
}
```

### Output -

1. Example for Deadlock -

```
Windows PowerShell
Copyright (C) Microsoft Corporation. All rights reserved.

Install the latest PowerShell for new features and improvements! https://aka.ms/PSWindows

PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> javac DeadlockExample.java
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> java DeadlockExample
11: lockl acquired, waiting to acquire lock2.

T2: lock2 acquired, waiting to acquire lock1.
```

2. Example for Deadlock Resolution -

For first 10 seconds, the 2 Operations are under deadlock -

```
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> javac LiveLockExample.java
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> java LiveLockExample
T2: lock2 acquired, trying to acquire lock1.
T1: lock1 acquired, trying to acquire lock2.
```

After 10 seconds, timeout releases the deadlock -

```
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> javac LiveLockExample.java
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> javac LiveLockExample.java
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7> java LiveLockExample
T2: lock2 acquired, trying to acquire lock1.
T1: lock1 acquired lock2, releasing lock1.
T2: cannot acquire lock1, releasing lock2.
T2: lock2 acquired, trying to acquire lock2.
T2: cannot acquire lock1, releasing lock2.
T1: lock2 acquired.
T1: lock2 acquired.
T1: exceuting first operation.
T2: lock2 acquired, trying to acquire lock1.
T2: lock1 acquired, trying to acquire lock1.
T2: lock1 acquired.
T2: lock2 acquired.
T2: lock3 acquired.
T2: lock4 acquired.
T2: lock4 bcquired.
T2: exceuting second operation.
PS C:\Desktop\Don Bosco\College Sem 8\Sem 8 Practicals\Distributed Computing\Code\Exp 7>
```

#### Conclusion:

Thus, we have studied and implemented deadlock management in distributed computing using Java to ensure system stability and prevent resource contention issues. We have implemented the methodology to detect, prevent, and resolve deadlocks, thereby enhancing the reliability and performance of distributed systems.

### References:

- <a href="https://www.baeldung.com/java-deadlock-livelock">https://www.baeldung.com/java-deadlock-livelock</a>
- <a href="https://www.geeksforgeeks.org/introduction-of-deadlock-in-operating-system/">https://www.geeksforgeeks.org/introduction-of-deadlock-in-operating-system/</a>
- <a href="https://www.geeksforgeeks.org/deadlock-handling-strategies-in-distributed-system/">https://www.geeksforgeeks.org/deadlock-handling-strategies-in-distributed-system/</a>
- <a href="https://en.wikipedia.org/wiki/Deadlock">https://en.wikipedia.org/wiki/Deadlock</a>

# **Rubrics for Assessment**

Timely Sub- mission	Late	Submitted just after deadline	On time Submission
	0 points	1 points	2 points
Under- stand- ing	Student is confused about the concept	Students has justifiably understood the concept	Students is very clear about the concepts
	0 points	2 points	3 points
Per- form- ance	Students has not per- formed the Experiment	Student has performed with help	Student has independently performed the experiment
	0 points	2 points	3 points
Devel- op- ment	Student struggles to write code	Student can write code with the requirement stated	Student can write exceptional code with his own ideas
	0 points	2 points	3 points
Overall Discip- line in	No Discipline maintained	Follows Rules	Sincerely performs the practical is always on time
Lab	0 points	1 points	4 points