

- 1) A bit string, 0111101111101111110, needs to be transmitted at the data link layer. What is the string actually transmitted after bit stuffing?

Insert a 0 after every consecutive 1s: 011110111110011111010

- 2) What is the remainder obtained by dividing  $x^7 + x^5 + 1$  by the generator polynomial  $x^3 + 1$ ? (give your answer as bit string)

Get remainder  $-x^2 + x + 1$ , correspond to bit string 111

- 3) A channel has a bit rate of 4 kbps and a propagation delay of 20 msec. For what range of frame sizes does stop-and-wait give an efficiency of at least 50 percent?

$$50\% \leq \frac{\frac{\text{FrameSize}}{4 \times 10^3 \text{ bps}}}{\frac{\text{FrameSize}}{4 \times 10^3 \text{ bps}} + 2 \times 20 \times 10^{-3} \text{ s}}$$

Solve it and get FrameSize  $\geq 160$  bits

- 4) Consider an error-free 64-kbps satellite channel used to send 512-byte data frames in one direction, with very short acknowledgements coming back the other way. What is the maximum throughput for window sizes of 1, 7, 15? The earth-satellite propagation time is 270 msec. (give your answer as an integer)

4.a) for window size=1

- Total time per frame:  $T = \text{TransmissionDelay} + 2 \times \text{PropagationDelay} = \frac{512 \times 8 \text{ bits}}{64 \times 10^3 \text{ bps}} + 2 \times 270 \times 10^{-3} \text{ s} = 0.604 \text{ s}$
- $TP = \frac{\text{WindowSize} \times \text{FrameSize}}{\text{TotalTime}} = \frac{1 \times 4096 \text{ bits}}{0.604 \text{ s}} = 6781 \text{ bps}$

4.b) for window size=7

- $TP = \frac{\text{WindowSize} \times \text{FrameSize}}{\text{TotalTime}} = \frac{7 \times 4096 \text{ bits}}{0.604 \text{ s}} = 47470 \text{ bps}$

4.c) for window size=15

- $TP = \frac{\text{WindowSize} \times \text{FrameSize}}{\text{TotalTime}} = \frac{15 \times 4096 \text{ bits}}{0.604 \text{ s}} = 101722 \text{ bps}$ , but it transcends the channel capacity, so the maximum throughput is 64000 bps

- 5) A 100-km-long cable runs at the T1 data rate. The propagation speed in the cable is  $\frac{2}{3}$  the speed of light in vacuum. How many bits fit in the cable?

- $T_1 = 1.554 \text{ Mbps}$
- $\text{TransmissionDelay} = \frac{100 \times 10^3 \text{ m}}{2 \times 10^8 \text{ m/s}} = 5 \times 10^{-4} \text{ s}$
- “Bits fit” = bandwidth-delay product  $= 5 \times 10^{-4} \text{ s} \times 1.554 \times 10^6 \text{ bps} = 777 \text{ bits}$

- 6) A CRC generator polynomial is  $G(X) = X^{16} + X^{15} + X^2 + 1$ . How many bits will the checksum(发送方进行计算后得到余数的位数) be?

- The degree of the generator polynomial is 16, so the checksum will be 16 bits.

- 7) Assume the sequence number has 3 bits. What is the maximum number of outstanding sending frames for a go back N protocol?

- $W_{T(\max)} = 2^3 - 1 = 7$
- The maximum number of outstanding sending frames is 7.

- 8) Assume the sequence number has 5 bits. What is the maximum number of outstanding sending frames for a selective repeat protocol?

- $W_{T(\max)} = 2^{5-1} = 16$
- The maximum number of outstanding sending frames is 16.

9) A CSMA/CD Ethernet network has a length of 10 km, a data rate of 10 Mb/s, and a signal propagation speed of 200 m/μs. What's the minimum frame length for this network?

- $\text{minFrameLength} = \text{DataRate} \times \text{ContentionPeriod} = 10 \times 10^6 \text{ bps} \times 2 \times \frac{10 \times 10^3 \text{ m}}{200 \times 10^6 \text{ m/s}} = 1000 \text{ bits}$

10) After the sender first sends frames from 0 to 6 and at the end of timeout receives the acknowledgements for frame 1, 3, and 5, the next frame it will re-transmit is which frame? (assume the protocol is go-back-N)

- Frame 6, from which timeout starts.