

5 $q^{a+ib} \leq \frac{q \pm \sqrt{q^2 - 4}}{2}$ ⑤ April Saturday 05-04-2014

$$q^{a+ib} \leq \frac{\sqrt{q}(\sqrt{q}) \pm \sqrt{q(q-4)}}{2}$$

$$q^{a+ib} \leq \frac{\sqrt{q}(\sqrt{q} \pm \sqrt{q-4})}{2}$$

$$\frac{q^{a+ib}}{q^{1/2}} \leq \frac{\sqrt{q} \pm \sqrt{q-4}}{2}$$

$$q^{a-\frac{1}{2}+ib} \leq \frac{\sqrt{q} \pm \sqrt{q-4}}{2}$$
 6 Sunday 06-04-2014

$$(a - \frac{1}{2} + ib) \log q \leq \log \left(\frac{\sqrt{q} \pm \sqrt{q-4}}{2} \right)$$

$$(a - \frac{1}{2} + ib) \leq \frac{\log(\sqrt{q} \pm \sqrt{q-4})}{\log q}$$

RHS (July) $\Rightarrow a \in \mathbb{R}$ For $q < 4$ ($2, 3$)
above has imaginary part

if $q = 2$: $a - \frac{1}{2} + ib \leq \log \left(\frac{\sqrt{2} \pm i\sqrt{2}}{2} \right)$

if $q = 3$: $a - \frac{1}{2} + ib \leq \log \left(\frac{1 \pm i}{\sqrt{2}} \right)$

May 2014

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8 For $q \geq 5$ April Tuesday 08-04-2014

$$a - \frac{1}{2} + ib \leq \log \left(\frac{\sqrt{q} \pm \sqrt{q-4}}{2} \right)$$

$$a - \frac{1}{2} \leq \log \left(\frac{\sqrt{q} \pm \sqrt{q-4}}{2} \right)$$

$$a \leq \frac{\log \left(\frac{\sqrt{q} \pm \sqrt{q-4}}{2} \right)}{\log q} + \frac{1}{2}$$

$$q^s + q^{1-s} \leq q$$

$$q^s + \frac{q}{q^s} \leq q$$

$$q^{2s} + q \leq q^{s+1}$$

March 2014

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09-04-2014 $q^{2a+2ib} + q - q^{a+ib} \leq 0$ April Wednesday 9

$e^{(2a+2ib)\log q} + q - e^{(a+ib)\log q} \leq 0$

$e^{2a\log q} \cdot e^{2ib\log q} + q - e^{(a+ib)\log q} \cdot e^{ib\log q} \leq 0$

$e^{2a\log q} \left(\cos(2b\log q) + i\sin(2b\log q) \right) + q + e^{(a+ib)\log q} \left(\cos(b\log q) + i\sin(b\log q) \right) \leq 0$

$e^{2a\log q} \cos(2b\log q) + q + e^{(a+ib)\log q} \cos b\log q \leq 0$ ①

$e^{2a\log q} \sin(2b\log q) + e^{(a+ib)\log q} \sin b\log q \leq 0$ ②

Reminder for 2: $\cos x + i\sin x$

Re(z) $q^{2a} \cos(2b\log q) + q + q^{(a+ib)\log q} \cos b\log q \leq 0$

Im(z) $q^{2a} \sin(2b\log q) + q^{(a+ib)\log q} \sin b\log q \leq 0$

May 2014

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| March 2014 | | | | | | |
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| May 2014 | | | | | | |
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12 April Saturday

$$\frac{2 \cos^2(b \log_2) \times 1}{2 \cos^2(b \log_2)} = 1$$

12-04-2014

$$\Rightarrow a = \frac{1}{2} \text{ for all } s = a + ib$$

$$\text{from } \det(I - A_q^{-s} + q^{1-2s} I) = 0$$

and using $q^s + q^{1-s} \leq q$

$$(or) \quad q^{25} + q - q^{51} \leq 0$$

Assumption of $A = \frac{1}{2}$ Kan Smizraasan
 13 Sunday Mahaveer Jayanthi, 13-04-2014

Assam
12 Sunday

13 Sunday for all $\{ = a + ib$ Manavee Jayanthi, 13-04-2017 Non-Ramanujan

doesn't give any contradiction. graphs)

Ramanujan graphs have RZF as zeros

$$\text{for } a \neq k_2 \quad \frac{q^{2a} \sin(2b\log z)}{q^{2a} \cos(2b\log z) + q^a}$$

March 2014

would not be equal to RHS

as of would not be common in
numeral denom

For rummicator product = 0. (12)
14-04-2014 (Page 1 29/3/2014) April
Tamil New Year's Day, Dr. Ambedkar Jayanthi, Pournami Monday 14

14-04-2014 Tamil New Year's Day, Dr.Ambedkar Jayanthi, Pournami O

April 14
Monday

$1 + q^{-s} = 0$ for some $s = a + ib$
 where RZ in RHS is zero.

$$9 - 5 = -1$$

$$q^s = -1$$

$$g^{a+ib} = -1$$

$$(a+ib) \log r = -1$$

$$e^{a \log z} \cdot e^{b \log z} = -1$$

$$e^{a \log r} (\cos b \log r + i \sin b \log r) = -1$$

$$e^{a \log r} \cos b \log r = -1$$

$$e^{\alpha \log r} \sin \beta \log r = 0$$

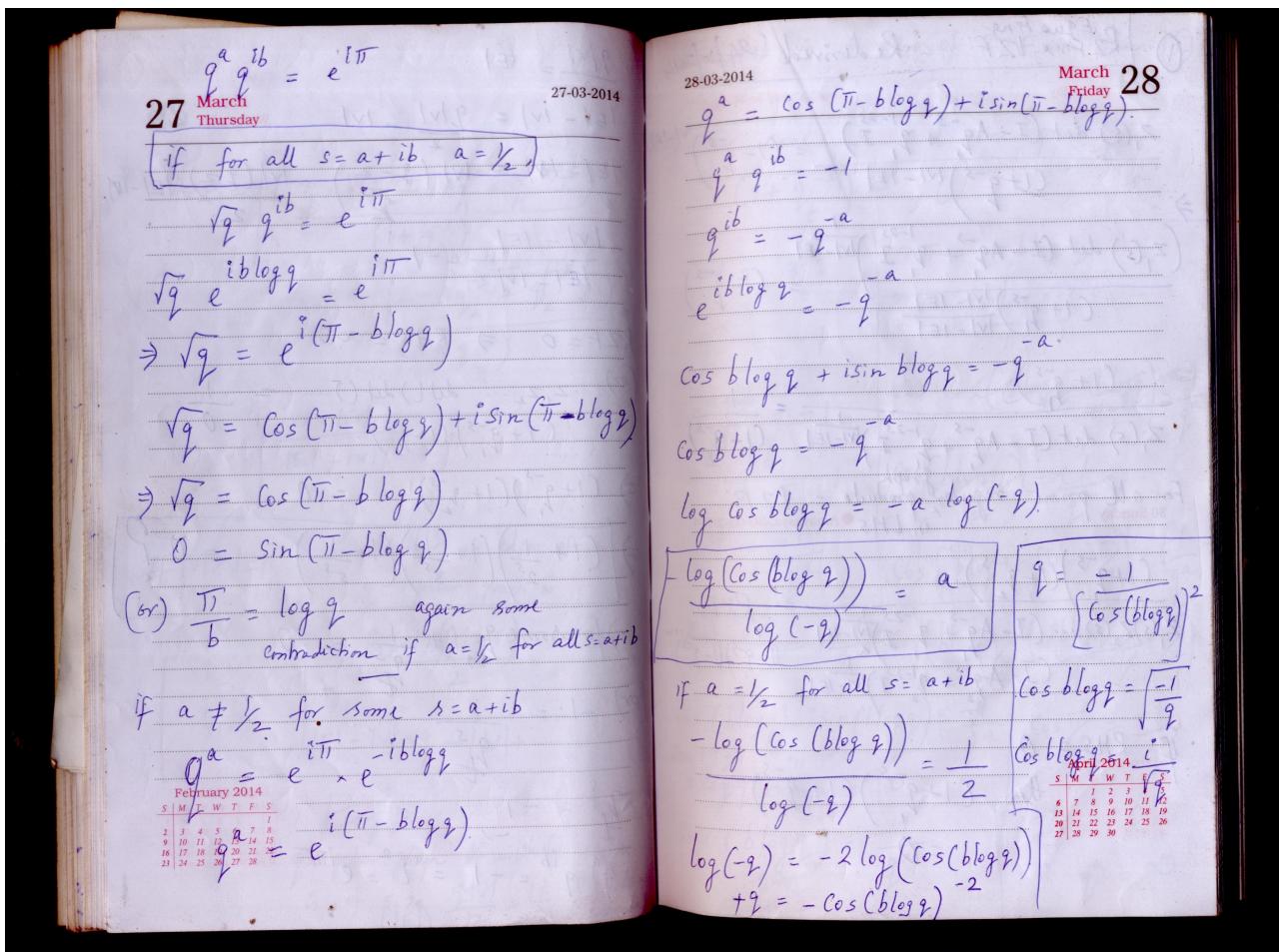
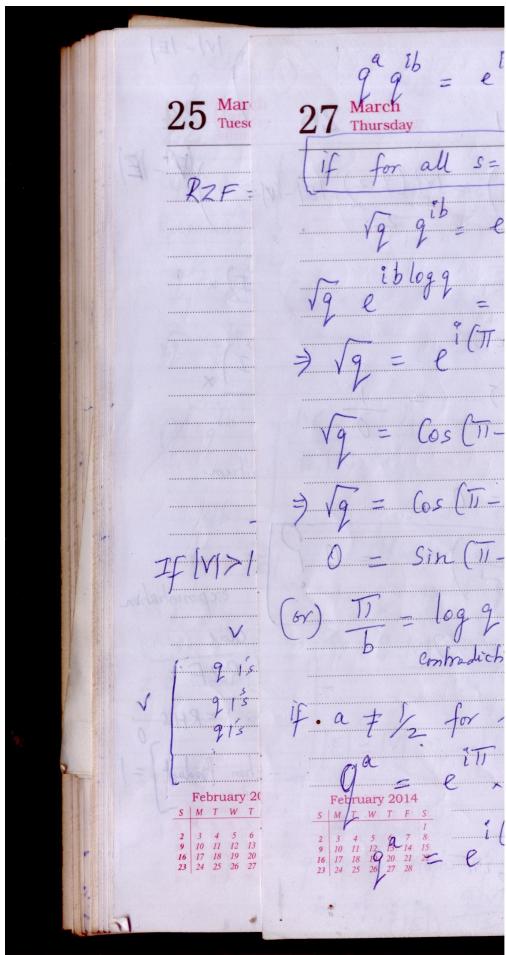
for all

$\Rightarrow b = 0$ or $\log q = 0$ be the case
 incorrect. Thus numerator never vanishes.

| May 2014 | | | | | | |
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|-----------------------|---|--------------------------|
| 15 April Tuesday | Thus denominator product of these identities for all prime regular graphs $\Rightarrow \infty$ and only $a = \frac{1}{2}$ satisfies it for non-ramanujan graphs. Thus Riemann Hypothesis is TRUE (or is it?) <i>Ka Shrinivasan</i> 25/04/2014 1pm. | 15-04-2014 O Pournami |
| 16 April Wednesday | | April Wednesday |

| | | |
|---|--|--------------------|
| 25 March Tuesday | 25-03-2014 | March Wednesday |
| $RZF = \left[\frac{(Z_1 \times \det(I - Aq_1^{-s} + q_1^{1-2s} I))}{(1 + q_1^{-s})} \times \frac{(Z_2 \times \det(I - Aq_2^{-s} + q_2^{1-2s} I))}{(1 + q_2^{-s})} \right] \frac{ V - E }{ E - V }$ | $ V = E $ $ E - V = \frac{ V }{2} - V = -\frac{ V }{2}$ $ E - E = 0$ $RZF = 0 \Rightarrow RHS = \infty$ | March Wednesday |

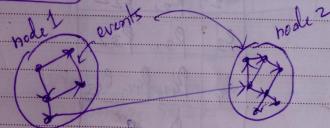


Text files EventNet Edges.txt and EventNet Vertices.txt are read by Python script on 08-05-2014.

Event Net is needed in AsFar, VIRGO, Linux and KingCobra.

AsFar can be python and c/c++ in VIRGO, Linux and KingCobra (gboost::graph).

4/12/2014



partaker Id format:

<ip address>:<threadid>

Id(s) of the events in node 2 are > the cause event in node 1.

Each conversation requires a log message being sent to EventNet service which updates the input edges & vertices text input files. Python script in a loop recomputes the ordering a graph.

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10/12/2014 25/4/170 4B7 N8/41750 May 9

① Circuit Value Problem is P-Complete
 (whether given input satisfies a circuit) - only if input is known (verifying a certificate - not finding it)
 (continued from page 16/1/2014). (Certificate complexity)

SAT circuit for each voter is just

CNF formulae which is NP Complete. It is like each voter searching solution space is superpolynomial or exponential (or) finding a certificate assignment.

Philosophically, candidate tried hard to convince voters - each voter to be precise. Number of voters is infinite. Theoretically is infinite majority decidable? This is like a streaming algorithm to find heavy hitter in any online voting - finding majority in a stream of binary digits.

The uniform distribution in page 16/1/2014 cannot be ruled out - for example if the voters are equally "rational" or intelligent.

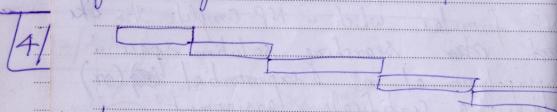
As in the case of nodes in a cloud which have same software installed.

② P can be equal to NP when above converges & in uniform distribution & undecidable if majority is undecidable in infinite.

Text 6 10 May Saturday 470/BN BNF 10-05-2014

470/BN 15

② Discrete Hyperbolic Factorization:
 Backman-Schubert-Vishkin & others algorithms using PRAM algorithms are NC algorithms by PRAM \equiv NC equivalence.



$O(\log n)$ time and $n^{O(1)}$ processors (PRAMs).

reduced to $\log n$ depth a

$n^{O(1)}$ size NC circuit. Each PRAM cell is of size $\leq \log n$. Thus factorization EN

the (3) graphs are ramanujan or non-ramanujan.

E RZF is represented as infinite set of $(\text{prime}+1)$ -regular graphs - product of their Ihara idenitities.

$\text{Re}(\text{atib}) = \frac{1}{2} (a = k_2)$

for both ramanujan & non-ramanujan graphs (thus RZF)

April 2014

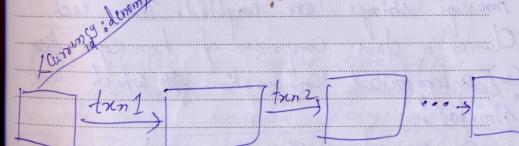
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All 3 results seem to be big and if correct (which they look) are unbelievable.

La Srinivasan.

Continued from 19/3/2013 page of Diary (2)

EventNet unique id for KingCobra MAC currency (MINT).



each transaction creates eventlog message.

An example transaction for MAC

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Sender $\xrightarrow{\text{MAC}}$ Receiver 2-Way or Add [MAC-CR] 3-Way handshake with MAC-ACK-ACK

Some expansion in notes from UBI diary

15 May Thursday 15-05-2014 11/4/2014

$nC_0 = \frac{n!}{0! n!}$

$nC_1 = \frac{n!}{1! (n-1)!}$

$nC_x = \frac{n!}{x! (n-x)!}$

$x! (n-x)! = (n-x)! x!$

$nC_{\frac{n}{2}} = \frac{n!}{\frac{n}{2}! \frac{n}{2}!}$ though n is even

$\frac{n!}{(\frac{n}{2})! (\frac{n}{2})!} = nC_{\frac{n}{2}}$

$\frac{n!}{(\frac{n}{2}+1)! (\frac{n}{2}-1)!} = \frac{n!}{(\frac{n}{2}+1)! (\frac{n}{2}-1)!}$

$4C_0 + 4C_1 + 4C_2 + 4C_3 + 4C_4 = 2^4$

$5C_0 + 5C_1 + 5C_2 + 5C_3 + 5C_4 + 5C_5 = 2^5$

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$\frac{5!}{2! 3!} = \frac{5!}{3! 2!}$

16 May Friday 16-05-2014

Special case of $\sum nC_{\frac{n}{2}}$ for even n :

$2x + nC_{\frac{n}{2}} = 2^n$

$2x = 2^n - nC_{\frac{n}{2}}$

$x = \frac{2^n - nC_{\frac{n}{2}}}{2}$

$x + nC_{\frac{n}{2}}$ to get $4C_3 + 4C_4$

$2^n - nC_{\frac{n}{2}} + nC_{\frac{n}{2}}$

$= \frac{2^n + nC_{\frac{n}{2}}}{2} = 2^{\frac{n}{2}} (4C_3 + 4C_4)$

$2^n + nC_{\frac{n}{2}} = 2^{\frac{n}{2}} (n + nC_{\frac{n}{2}})$

$\frac{2^n + nC_{\frac{n}{2}}}{2 \cdot 2^n} = \frac{1}{2} + \frac{nC_{\frac{n}{2}}}{2^n}$

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17 May Saturday 17-05-2014

For even population, $P(\text{good}) = \frac{1}{2} + \epsilon$

where $\epsilon = \frac{n!}{2^n}$

$\frac{(n!)^2}{(2^n)!} = \frac{n!}{2^n}$

Stirling approximation:

$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$

$\left(\frac{n}{2}\right)! \approx \sqrt{\pi n} \left(\frac{n}{2e}\right)^{\frac{n}{2}}$

$\epsilon = \frac{\sqrt{2\pi n} \left(\frac{n}{e}\right)^n}{(\sqrt{\pi n})^2 \left(\frac{n}{2e}\right)^{\frac{n}{2}}}$

$\epsilon = \frac{\sqrt{2\pi n} (n^{\frac{n}{2}} \times (2e)^n)}{(\sqrt{\pi n})^2 e^n (\pi n)^{\frac{n}{2}}} = \frac{\sqrt{2} \times 2^{\frac{n}{2}}}{\sqrt{\pi n}}$

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18 Sunday 18-05-2014

$\epsilon = \frac{\sqrt{2} \times 2^{\frac{n}{2}}}{\sqrt{\pi n}} = \frac{\sqrt{2} \times 2^{\frac{n}{2}}}{\sqrt{12.5}} \approx \frac{\sqrt{2} \times 2^{\frac{n}{2}}}{\sqrt{12.5}} \approx 0.4$

19 May Monday 19-05-2014

$\frac{\sqrt{2} \times 2^{\frac{n}{2}}}{\sqrt{\pi n}} = \epsilon$

$\epsilon = \frac{\sqrt{2}}{\sqrt{\pi n}} \quad \text{for } n=1, nC_{\frac{n}{2}} \text{ is not valid}$

$\epsilon = \frac{\sqrt{2}}{\sqrt{\pi n}} \quad \text{(must be ignored)} \quad \sqrt{\frac{2}{\pi n}} = \frac{1.41}{\sqrt{1.77}}$

$\epsilon = \sqrt{\frac{2}{\pi n}} \approx \sqrt{0.6} \approx 0.77$

for large n $\epsilon \approx \frac{\sqrt{2}}{\sqrt{\pi n}} \Rightarrow 0$

ϵ tends to zero for large n

for $n > 2$ (even n)

$\frac{\sqrt{2}}{\sqrt{\pi n}} \approx \frac{\sqrt{2}}{\sqrt{12.5}} \approx \frac{\sqrt{2}}{\sqrt{12.5}} \approx 0.4$

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This implies that for even
20 May Tuesday population 20-05-2014

20-05-2014

But for large n (∞).

21-05-2014 Both were odd May Wednesday 21

May 21
Wednesday

$$\begin{aligned}
 & \frac{4C_0 + 4C_1 + 4C_2 + 4C_3 + 4C_4}{2^4} \\
 & = \frac{2^4 + 4C_2}{2} = \frac{16 + \frac{24}{2^2}}{2} \\
 & = \frac{16 + 6}{2} = 11 \\
 & \frac{11}{2^4} = \frac{11}{16} \\
 & 4C_0 = \frac{4!}{0!4!} = 1 \quad 4C_4 = 1 \\
 & 4C_1 = \frac{4!}{1!3!} = 4 \quad \sum_{i=0}^4 4C_i = 16 \\
 & \begin{array}{c} \text{April 2014} \\ \hline \text{S} & \text{M} & \text{T} & \text{W} & \text{T} & \text{F} & \text>S \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ 8 & 9 & 10 & 11 & 12 & 13 & 14 \\ 15 & 16 & 17 & 18 & 19 & 20 & 21 \\ 22 & 23 & 24 & 25 & 26 & 27 & 28 \\ 29 & 30 & & & & & \end{array} \\
 & 4C_1 = \frac{4!}{1!3!} = 4 \\
 & 4C_3 = \frac{4!}{3!1!} = 4
 \end{aligned}$$

$$2^n - \frac{2^n}{2} = \frac{nC_n}{2}$$

$$\therefore 2^n - \frac{nC_n}{2}$$

$$\frac{2^n - 2^{n-1}}{2} = \frac{10}{2 \times 16} = \frac{5}{16}$$

$$\left. \begin{aligned}
 & \begin{array}{c} 1 \\ + \\ 2 \end{array} \xrightarrow{2^n} \begin{array}{c} n \\ \gamma_2 \end{array} \\
 & \begin{array}{c} 2 \\ \xrightarrow{2 \times 2^n} \end{array}
 \end{aligned} \right\} = P(\text{good}) \text{ for even } \gamma_2$$

22/12/2014 for even n (e.g. $n=4$) Special case for even electrode
 27 May Tuesday $P_c(\text{prod})$ 27-05-2014 P.M.

$\sum C_i = 1$ $4C_0 + 4C_1 + 4C_2 + 4C_3 + 4C_4 = 2^4$

$4C_1 = 4 \quad hC_x = hC_{n-x}$

$4C_2 = 6$

$4C_3 = 4 \quad 2x + 4C_2 = 2^4$

$4C_4 = 1 \quad x = \frac{2^4 - 4C_2}{2}$

$(+) \quad 16 = 2^4$

$4C_3 + 4C_4 = -(x + 4C_2) + 2^4$

$= \left(\frac{2^4 - 4C_2}{2} + 4C_2 \right) + 2^4$

$4C_3 + 4C_4 = -\left(\frac{2^4 + 4C_2}{2} \right) + 2^4$

for any even n .

$$\begin{aligned}
 nC_0 + nC_1 + \dots + nC_n &= - \left(\frac{2^n + nC_{n/2}}{2} \right) + 2^n \\
 &= 2^n - \frac{2^n}{2} - \frac{nC_{n/2}}{2}
 \end{aligned}$$

28-05-2014
Kritikai ★

May 28
Wednesday

$$\begin{aligned}
 &= \frac{2^n}{2} - \frac{nC_{n-2}}{2} \\
 &= \frac{2^n}{2} - nC_{n-2} = nC_{n-1} + nC_{n-2} + \dots + nC_n
 \end{aligned}$$

RHS Pr₁(good) for democracy circuit, above 1/2?

$$\frac{2^n - nC_{n/2}}{2 \times 2^n} = \frac{1}{2} - \frac{nC_{n/2}}{2^n}$$

$$\frac{n!}{\frac{n^c}{2^n}} = e = \frac{n!}{\overline{\left(\frac{n}{2}\right)!\left(\frac{n}{2}\right)!}}$$

$$\approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n = \frac{\sqrt{2\pi n}}{e^n} \cdot n^n$$

$$2 \left[\sqrt{\frac{2\pi n}{e}} \times \left(\frac{n}{e}\right)^{\frac{n}{2}} \times \frac{1}{2^{\frac{n}{2}}} \right]^2 = \frac{2^{2n}}{\sqrt{2\pi n} \cdot \frac{n^n}{e^n}}$$

29 May Thursday

29-05-2014

$$\begin{aligned}
 & \frac{\sqrt{2\pi n}}{e^n} = \frac{\sqrt{2\pi n}}{e^n} \times e^n \\
 \approx & \frac{\frac{2\pi n}{2^n} \frac{n^{2n}}{e^{2n}}}{e^{2n}} = \frac{\frac{2}{\sqrt{\pi n}}}{e^{2n}} = \frac{2}{\sqrt{\pi n}} \\
 \approx & \frac{\sqrt{2\pi n}}{e^n} \times \frac{n}{e^n} \times 2 \times e^n \\
 \approx & \frac{\sqrt{2\pi n}}{e^n} \times \frac{n}{e^n} \times 2 \times e^n \\
 \approx & \frac{\sqrt{2\pi n}}{e^n} \times \frac{n}{e^n} \times 2 \times e^n \\
 \approx & \frac{\sqrt{2\pi n}}{e^n} \times \frac{n}{e^n} \times 2 \times e^n \\
 \approx & \frac{\sqrt{2\pi n}}{e^n} \times \frac{n}{e^n} \times 2 \times e^n
 \end{aligned}$$

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$$\text{30-05-2014} \quad \frac{1}{h} \left(1 + \frac{1}{n}\right)^{2n} = e$$

May Friday 30

$$\frac{h}{n} \frac{\sqrt{2}}{\sqrt{\pi n}} \times \left(\frac{2e}{n}\right)^n = 0 = e$$

For finite even electorate, e does not vanish and always less than $\frac{1}{2}$ (?) and thus pseudorandomness is better than majority voting for non-uniform distribution.

But for infinite RHS, e is still zero.

Is infinite majority circuit decidable?

for $n=4$

$$4C_3 + 4C_4 = \frac{2^4}{2} - 4C_2 = 16 - \frac{4 \times 3 \times 2}{2 \times 2} = 16 - 6 = 10$$

$$\text{Thus } \Pr(\text{good}) \text{ RHS} = \frac{5}{16} = \frac{5}{2^4}$$

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Thus $\Pr(\text{good})$ RHS is

$$\frac{2^4}{2^5} = \frac{1}{2}$$

in uniform distribution.

Thus Odd electorate has no issues at all.

For even n uniform distribution:

$$\text{RHS } \Pr(\text{good}) = \frac{1}{2} - \frac{nC_{n/2}}{2^n}$$

$$= \frac{1}{2} - \left(\frac{n!}{(\frac{n}{2})! (\frac{n}{2})!} \right) \frac{2^n}{2^n}$$

$$\frac{n!}{(\frac{n}{2})! (\frac{n}{2})!} \approx \frac{\sqrt{2\pi n} (\frac{n}{e})^n}{\sqrt{\pi n} \frac{n}{2} (\frac{n}{2e})^{\frac{n}{2}} \times \sqrt{\pi n} (\frac{n}{2e})^{\frac{n}{2}}}$$

$$\approx \frac{\sqrt{2\pi n} (\frac{n}{e})^n \times (2e)^n}{\sqrt{\pi n} \sqrt{\pi n} n^n} = \frac{\sqrt{2\pi n} (\frac{n}{e})^n \times 2^n}{\sqrt{\pi n} n^n}$$

$$\approx \frac{\sqrt{2\pi n} \frac{n}{e} \times 2^n}{\sqrt{\pi n} \times n^n} \approx \frac{2}{\sqrt{\pi n}} \times 2^n$$

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31 May Saturday

31-05-2014

| | | |
|----|------|-------|
| 0 | 0000 | (1+1) |
| 1 | 0001 | |
| 2 | 0010 | |
| 3 | 0011 | |
| 4 | 0100 | |
| 5 | 0101 | |
| 6 | 0110 | |
| 7 | 0111 | |
| 8 | 1000 | |
| 9 | 1001 | |
| 10 | 1010 | |
| 11 | 1011 | |
| 12 | 1100 | |
| 13 | 1101 | |
| 14 | 1110 | |
| 15 | 1111 | |

Even electorate, when finite is puzzling

For odd electorate (e.g. $n=5$)

$$5C_0 + 5C_1 + 5C_2 + 5C_3 + 5C_4 + 5C_5$$

$$\underbrace{5C_0 + 5C_1 + 5C_2}_{x} + \underbrace{5C_3 + 5C_4 + 5C_5}_{1} = 2^5 - x$$

$$2^5 - x = 2^5$$

$$x = 2^4$$

$$2x = 2^5$$

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$$x = 2^4$$

$$2$$

3 Thus RHS $P(\text{good})$

June
Tuesday

03-06-2014

$$\frac{1}{2} - \frac{\sqrt{\frac{2}{n!}} \times 2^{\frac{n}{2}}}{2^{\frac{n}{2}}} = \left[\frac{1}{2} - \frac{\sqrt{\frac{2}{n!}}}{\sqrt{\frac{n}{2}}} \right] \text{ if } n \rightarrow \infty$$

Special case of even number of voters:

for even electorate democracy.

Choice can be worse than pseudorandom choice (uniform distribution) when finite and thus democracies with even number of electors are failures (?)

K. Shrinivasan.

$n=2$

$$\frac{2^2 - 2C_1}{2^2} = \frac{4 - 2}{8} = \frac{2}{8} = \frac{1}{4}$$

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$$\frac{1}{2} - \frac{\sqrt{\frac{2}{2!}}}{2} = \frac{1}{2}$$

Infinite majority is undecidable

04-06-2014

June
Wednesday

4

Proof:

Hypothetical turing machine computing majority of infinite bit stream would be looping always without an 'yes', 'no' output. Thus TM is not recursive but recursively enumerable language. As RE is recursive are undecidable, infinite majority is undecidable. Thus Arrow's theorem for infinite number of candidates should also be undecidable.

23/12/2014

Arrow's theorem for 3-candidate Condorcet election proves that there is non-zero error in voting. Does this imply that 2-candidate Condorcet election has error also?

If yes, the RHS of $P(\text{good})$ binomial coefficient summation for democracy circuit with finite SAT inputs for voters (i.e. finite number of electors) can never be 1 by Arrow's theorem.

For 2 candidates of the ranking of Condorcet in trivial

The counter example in there exists a

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democracy circuit ($\text{Maj} + \text{SAT}$) with odd number of voters (sat count per voter) with uniform distribution

which converges to $\frac{1}{2}$ (Same as LHS pseudorandom choice).

This counterexample is tantamount to saying that there exists a PTIME algorithm using a PRGenerator for an NP-complete democracy problem.

$\text{CLHS} = \text{RHS}$ in $P(\text{good})$ binomial coefficient summation.

Infinite version seems to be undecidable.

K. Shrinivasan.

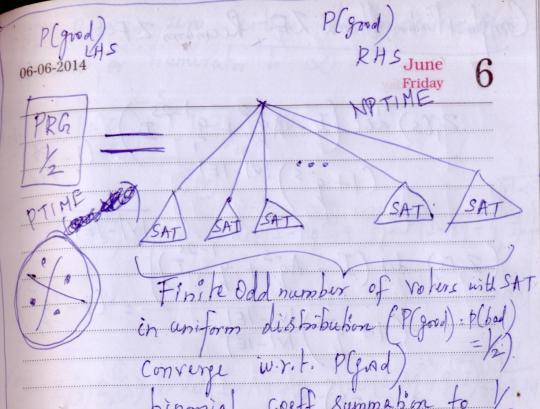
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SAT + Maj with odd number of voter SAT inputs to majority circuit, when finite, with uniform distribution converges to $\frac{1}{2}$ Same as LHS for pseudorandom choice.

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This is the counter example when $P = NP$



Counter example for $P \neq NP$ Also if $P = 1$

Philosophically this implies that there are scenarios where hard problems are easy due to "lucky guesses".

As mentioned above:

If there is no error in both LHS and RHS $P = 1$ and hence,

LHS is 1

RHS is $nC_n (s) = 1$

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Complement function, Ihara ZF, Riemann ZF

7 June Saturday

07-06-2014

$$\frac{Z_1(s) \det(I - Aq_i^{-s} + q_i^{1-2s} I)}{(1+q_i^{-s})^{1/2(1-2s)}} = (1-q_i^{-s})^{\frac{1}{M-1}}$$

$$\frac{(Z_1(s) \det(I - Aq_i^{-s} + q_i^{1-2s} I))}{(1+q_i^{-s})^{1/2(1-2s)}} = (1-q_i^{-s})^{\frac{1}{M-1}}$$

$$\frac{(1+q_i^{-s})}{\left(Z_1(s) \det(I - Aq_i^{-s} + q_i^{1-2s} I) \right)^{\frac{1}{M-1}}} = (1-q_i^{-s})^{\frac{1}{M-1}}$$

$$\frac{(1+q_i^{-s})(1+q_{i_2}^{-s}) \cdots}{(Z_1(s) Z_2(s) \cdots \det(I - Aq_{i_1}^{-s} + q_{i_1}^{1-2s} I) \det(I - Aq_{i_2}^{-s} + q_{i_2}^{1-2s} I) \cdots)^{\frac{1}{M-1}}} = (1-q_{i_1}^{-s})^{\frac{1}{M-1}} (1-q_{i_2}^{-s})^{\frac{1}{M-1}} \cdots$$

$$= RZF$$

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If RHS is zero, either denominator is zero
or numerator is ∞

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$$\frac{(1+q_{i_1}^{-s})(1+q_{i_2}^{-s}) \cdots}{(1-q_{i_1}^{-s})^{\frac{1}{M-1}}} = 0$$

$$\left(\frac{(1+q_{i_1}^{-s})(1+q_{i_2}^{-s}) \cdots \det(I - Aq_{i_1}^{-s} + q_{i_1}^{1-2s} I) \det(I - Aq_{i_2}^{-s} + q_{i_2}^{1-2s} I)}{(1-q_{i_1}^{-s})^{\frac{1}{M-1}} (1-q_{i_2}^{-s})^{\frac{1}{M-1}}} \right) = \infty$$

$$(b) \left[Z_1(s) Z_2(s) \cdots \det(I - Aq_{i_1}^{-s} + q_{i_1}^{1-2s} I) \det(I - Aq_{i_2}^{-s} + q_{i_2}^{1-2s} I) \cdots \right]^{\frac{1}{M-1}} = 0$$

$$\Rightarrow \text{either } Z_1(s) Z_2(s) \cdots = 0$$

which implies that zeros of RZF in RHS are zeros of Ihara ZF also. Thus the graph must be Ramanujan (a known result).

$$\det(I - Aq_{i_1}^{-s} + q_{i_1}^{1-2s} I) \det(I - Aq_{i_2}^{-s} + q_{i_2}^{1-2s} I) \cdots = 0$$

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Thus for some q_i

$$\det(I - Aq_{i_1}^{-s} + I + q_{i_1}^{1-2s} I) = 0$$

multiplying by q_i^s

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10-06-2014

$$\det\left[-A + Iq_i^s + q_i^{1-s} I\right] = 0$$

$$\det\left[-A - Iq_i^s - q_i^{1-s} I\right] = 0$$

$$\det\left[-A - I(q_i^s + q_i^{1-s})\right] = 0$$

$$\text{eigen value is } q_i^s + q_i^{1-s} = \nu$$

$$q_i^{2s} + q_i^{-s} = \nu q_i^s$$

$$q_i^{2s} - \nu q_i^s + q_i^{-s} = 0 \quad \begin{matrix} \text{(from this onwards} \\ \text{derivation is} \\ \text{same as} \\ \text{page 17/4/2014)} \end{matrix}$$

$$q_i^{2a+2ib} - \nu q_i^a q_i^{ib} + q_i^{-s} = 0$$

$$q_i^{2a} q_i^{2ib} - \nu q_i^a q_i^{ib} + q_i^{-s} = 0$$

$$q_i^{2a} e^{2ib \log q_i} - \nu q_i^a e^{ib \log q_i} + q_i^{-s} = 0$$

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$$q^{2a} (\cos(2b \log \nu) + i \sin(2b \log \nu))$$

$$- \nu q^a (\cos(b \log \nu) + i \sin(b \log \nu))$$

$$+ q^{-s} = 0$$

equating to $0 + i0$

$$q^{2a} \cos(2b \log \nu) - \nu q^a \cos(b \log \nu) + q^{-s} = 0$$

$$q^{2a} \sin(2b \log \nu) - \nu q^a \sin(b \log \nu) = 0$$

$$\frac{q^{2a} \cos(2b \log \nu) + q^{-s}}{q^{2a} \sin(2b \log \nu)} = \frac{\nu q^a \cos(b \log \nu)}{\nu q^a \sin(b \log \nu)}$$

$$\frac{q^{2a} \sin(2b \log \nu)}{q^{2a} \cos(2b \log \nu) + q^{-s}} = \tan(b \log \nu)$$

$$\frac{q^{2a} (\cos(2b \log \nu) + q^{-s})}{q^{2a} \sin(2b \log \nu)} = \frac{2 \sin(b \log \nu) \cos(b \log \nu)}{q^{2a} \cos^2(b \log \nu) - q^{2a} \sin^2(b \log \nu) + q^{-s}}$$

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Thursday

12-06-2014
O Pournami

$$q^{2a} \cos^2 \theta \text{cosec}^2 \theta =$$

$$q^{2a} \cos^2 \theta \text{cosec}^2 \theta - q^{2a} \sin^2 \theta \text{cosec}^2 \theta + q^{2a}$$

$$q^{2a} \cos^2 \theta \text{cosec}^2 \theta = -q^{2a} \sin^2 \theta \text{cosec}^2 \theta + q^{2a}$$

$$q^{2a} (\cos^2 \theta \text{cosec}^2 \theta + \sin^2 \theta \text{cosec}^2 \theta) = q^{2a}$$

$$q^{2a} = q$$

$a = \frac{1}{2}$ for any eigenvalue
which implies for
non-ramanujan graphs
also

Thus infinitely many (prime + 1)
regular graphs (both ramanujan &
non-ramanujan).

May 2014
represent Riemann Zeta function
as infinite graph. Since above

derivation is irrespective of eigen values
of graph adj matrix, both ramanujan &
non-ramanujan graphs (prime +
regular) have $\text{Re}(\zeta) = \frac{1}{2}$.

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Ka Shrinivasan

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