parta) (Allog (log(n))

void fl (int n)

of

int 1=2

while (i < n) &

/* do something that takes O(1) time */

7

3

The number of times the while loop runs varies depending on the value of n. For example, if n=2 the while runs zero times. If n=4, the while loop runs once. If n=16, the while loop runs twice. If n=256, the while loop runs three times. If we create a variable K, and set it equal to the number of times the white loop runs, we can determine a relationship between n and K. This relationship is $n=2^{2^K}$. From, this we can determine that the loop runs $K=\log(\log n)$ times. Thus, my final runtime is $\theta(\log(\log n))$.

K	0	1	2	3	Albiray K	stop when
i	2	4	16	256	1=22x	x = n
/V	u55	124	n>16	N ≥256		

Since the code segment stops when i=n, the wast case run time is $n=2^{2^k}$, where $k=\log(\log n)$, thus, the final runtine is $\Theta(\log n)$.

part b) (0 (n" void f2(int n) for (intiel; i <= n; i++) & if ((1,00 cint) sqrt(n)) ==0) { for (int K=0; K<pow(i, 3); K++) of / do something that takes oci) time */ 6 6 Jo In the worst case scenario, the inner for loop runs n3 times. This is the case in which i= no. The code inside the if Statement itself will run it times. For example, if n=100, the code inside the is statement will on when i = 10, 20, 30, 40, 50, Thus, the total runtime for the code is: Thus, the total runtine for the code is: $\theta(n+\sqrt{n}\cdot n^3) + \theta(n-\sqrt{n}) = 0$ this accounts this accounts for the worst rave for the case assuming the if where the if statement is enlaced Stalement isn't $\sum_{i=1}^{n} \theta(i) + \sum_{i=1}^{n} \sum_{k=1}^{n} \theta(i) + \sum_{i=1}^{n} \theta(i) - \sum_{i=1}^{n} \theta(i)$ $\theta(n) + \sum_{i=1}^{\sqrt{n}} (n^3) + \theta(n) = \theta(n + \sqrt{n} \cdot n^3) + \theta(n - \sqrt{n}) = \theta(n)$

worst case whon i=n

port c) $(\theta(n^2)$ for (int i=1; i <=n; i+1) { for (int K=1) K <= n; K++) of if (A[h] == i) of for (int m=1; m <=n; m=m+m) of I do something that takes O(1) time // Assume the contents of the A[] array me not changed The inner for loop runs log (n) times since it runs until m=n and m is doubted after each iteration. The it statement will evaluate to true a maximum of 1 times so the code inside the it statement will run a total of 1 times. Thus the run time for the code including the it statement or the inner for loop is $\theta(nlogn)$. The outer for loops will run n times each so the runtime for the outerfor loops is O(n2). Thus the runtime for the code is: $O(n^2 + n \log n) + O(n^2 - n)$. Thus, my final runtime for the code is $\theta(n^2)$. $\theta(n^2 + n \cdot \log n)$ This accounts for This accounts the worst case assuming for the the if statement is worst case Mfered assuming the if statement isn $\sum_{i=1}^{n} \sum_{k=1}^{n} \theta(i) + \sum_{j=1}^{n} \theta(\log n) + \sum_{i=1}^{n} \sum_{k=1}^{n} \theta(i) - \sum_{j=1}^{n} \theta(i) = \theta(n^{2} + n \log n) + \theta(n^{2} - n)$

dummy voriable

if statement is

tive

; for # of times

during voriable for ;= = (n2) & \tall(n2)

 $=(\theta(n^2))$

for # of times it.

Statement is true

```
part d) (O(n)
int f (int n)
  int or = new int [10]
 int size = 10;
  for (int i=0, i < n; i++)
        if (i==size)
           int newsize = 3 * size /2;
           int "b = new int [newsize];
          for (int j=0; j<size; j++) b[j] = a[i];
           delete []a;
           Size = newsize;
     a[i] = i*i;
The outer for loop runs in times. The number of times the inner for loop
run is dependent on the value of n. For example, if n is 20, the
inner for loop will run 10 times when i = 10 and 15 times when
i is equal to 15. If n is 30, the inner for loop will run 10 times when
is equal to 10, 16 times when it is equal to 15 and 22 times when
i is equal to 22. This means that the inner for loop runs ((3) *.10) times each
time the size is reached. Defending on the value of n, the value of x energes.
The first time the if statement is entered, x is O, and the inner for loop runs 10 times.
Then, the second time the it statement is entered, x is I, and the inner for loop runs 15 times.
When n is 30, the first time the its lalement is entered, x is 0, and the inner for loop runs
```

10 times. The second time the if statement is run, x is 1, and the inner for loop runs 15 times. The third time the if statement is entered, x is 2 and the inner for loop run 22 times.

port of (continued)

Since the size variables are integers, we must floor the vale of $(\frac{3}{2})^2 \cdot 10$, which is why the inner took 100p runs 22 times even though $(\frac{3}{2})^2 \cdot 10$ evaluates to 22.5. From are previous alculations through example values of 0, we can conclude that the if statement is entered $\log_{\frac{3}{2}}(\frac{10}{10})$ times. This means the total runtime for the inner for loop is $\log_{\frac{3}{2}}(\frac{10}{10})$. The total runtime for the outer for loop is $\log_{\frac{3}{2}}(\frac{10}{10})$.

Thus, our warst case runtime for this code segment will evaluate to

$$\left(n + \sum_{x=0}^{\lfloor \frac{n}{2} \rfloor} |0 \cdot \left(\frac{3}{2}\right)^{x}\right) + \Theta\left(n - \log_{\frac{3}{2}}(n)\right)$$

Outer and inner for loop and for loop while accounting for when the if statement statement doesn't evaluate to the doesn't evaluate to the

does evaluate to

we con evaluate the summation using the geometric series formula from lecture:

$$\sum_{x=0}^{9/3} \frac{\binom{3}{2}}{\binom{3}{2}} \cdot 10 = \underbrace{10 \cdot \binom{3}{2}}{\frac{3}{2}} \cdot \frac{\log_3(\frac{n}{10})}{\frac{3}{2}} - 1 = \underbrace{10 \cdot (\frac{n}{10})}{\frac{3}{2}} - 1 = \underbrace{n-10}_{\frac{1}{2}} = \theta(n)$$

Additionally, if we plug in the most possible walves to the summation equation, we have $10 \cdot \binom{3}{2} \log_{\frac{3}{2}} \binom{1}{18}$ which evaluates to n. Each number before n in the summation must be less than this value since this is the max, so the other numbers in the summation can be represented as $n-c_1$, $n-c_2$... where c_1 and c_2 are ansteads. This preams that the summation will be $c_3(n)$ where c_3 is some anstate, thus, for this vacuum and the geometric one above, the summation is $\theta(n)$.

Thus the total runtine evaluates to $\theta(n+n) + \theta(n) = \theta(n)$