

Constraint Satisfaction Problems I

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50.021 Artificial Intelligence

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Outline & Objectives

- Understand the differences between standard search problems and constraint satisfaction problems
- Able to formulate a constraint satisfaction problem
- Understand the workings behind backtracking search and the various heuristics used to enhance its efficiency
- Able to use backtracking search to solve a CSP

Recap: Standard Search Problem Formulation

- State space, e.g. At(Arad), At(Bucharest)
- Initial state, e.g. At(Arad)
- Actions, set of actions given a specific state
 - Transition model e.g., Result(At(Arad), Go(Zerind)) → At(Zerind)
 - Path cost (additive), e.g., sum of distances, number of actions, etc
- Goal test, can be
 - Explicit, e.g. At(Bucharest)
 - Implicit, e.g. checkmate(x)

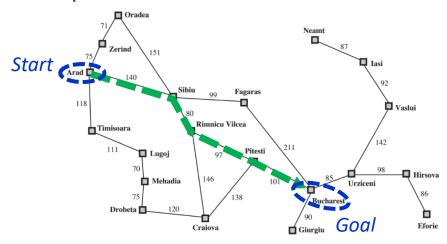


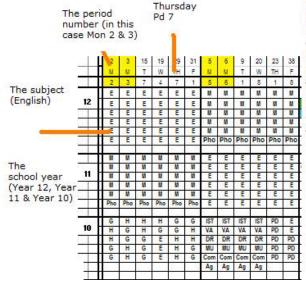
Recap: Standard Search Problem Solution

- A solution is a sequence of actions from the initial state to a goal state
 - E.g., Arad → Sibiu → Fagarus → Bucharest
- An optimal solution is a solution with the lowest path cost

Compare and Contrast

What are the differences among these problems?





Key:
E=English
M=Maths
Pho=Photography
G=Geography
H-History
IST=Info Tech
PD=PDHPE
VA=Art
Dr=Drama
Mu=Music
Com=Commerce
Ag=Agriculture

Initial State



Goal State

	1	2
3	4	5
6	7	8

Initial State

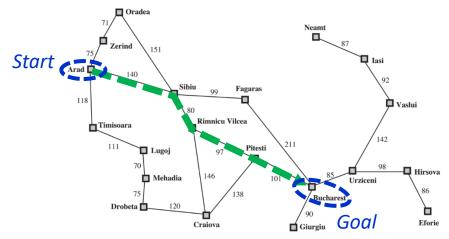
	en,					9		
		6						
			2	4	1		3	
			9			7		
					2			4
	80			7			2	
8	5							
	9		7		4			
					6			1

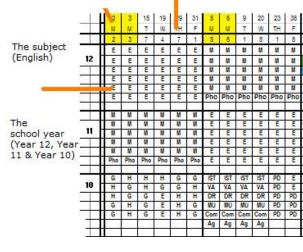
Goal State

1	3	2	5	6	7	9	4	8
5	4	6	3	8	9	2	1	7
9	7	8	2	4	1	6	3	5
2	6	4	9	1	8	7	5	3
7	1	5	6	3	2	8	9	4
3	8	9	4	7	5	1	2	6
8	5	7	1	2	3	4	6	9
6	9	1	7	5	4	3	8	2
4	2	3	80	9	6	5	7	1

Compare and Contrast

 What are the differences among these problems? Two main types





Thursday

The period

number (in this

case Mon 2 & 3)

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Initial State

8		6
5	4	7
2	3	1

Goal State

	1	2
3	4	5
6	7	8

Initial State

	3					9		
		6						
			2	4	1		3	
			9			7		
					2			4
	80			7			2	
8	5							
	9		7		4			
					6			1

Goal State

1	3	2	5	6	7	9	4	8
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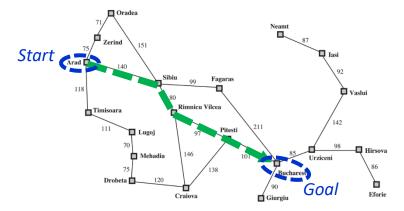
What is the Solution?

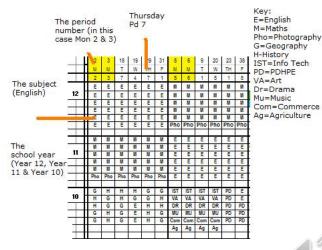
Standard Search Problems

- More interested in the sequence of actions (path) to the goal
- Paths have various costs, depths
- Heuristics give problem-specific guidance

Constraint Satisfaction Problems

- More interested in the goal itself, not the sequence of actions (path) there
- All paths at the same depth (for some formulations)
- CSPs are specialized for this type of task





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Constraint Satisfaction Problems

State

Defined by variables X_i that take on values from domain D_i

Goal Test

- A set of constraints C_i specifying allowable combinations of values for subsets of variables
- In contrast to standard search problems
 - State is a "black box" any old data structure that supports goal test, evaluation, successor

Constraint Satisfaction Problems

State

Defined by variables X_i that take on values from domain D_i

Goal Test

- A set of constraints C_i specifying allowable combinations of values for subsets of variables
- Simple example of a formal representation language
- Allows useful general-purpose algorithms with more power than standard search algorithms

Constraint Satisfaction Problems Formulation

- Finite set of *variables* $X = \{X_1, X_2, ..., X_n\}$
- Non-empty **domain D** of k possible values for each variable D_i where $D_i = \{v_1, ..., v_k\}$
- Finite set of *constraints C* = { C_1 , C_2 , ..., C_m }
 - Each constraint C_i limits the values that variables can take, e.g., $X_1 \neq X_2$



Constraint Satisfaction Problems Formulation

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- Finite set of *constraints C* = { C_1 , C_2 , ..., C_m }
 - Each constraint C_i limits the values that variables can take, e.g., $V_1 \neq V_2$
- In relation to a search problem, we have:
 - State is defined by variables X_i that take on values from domain D_i
 - Goal Test is a set of constraints C_i specifying allowable combinations of values for subsets of variables

Constraint Satisfaction Problems

- A complete assignment is where every variable is assigned a value
- A consistent assignment does not violate any constraint
- A CSP solution is a complete and consistent assignment for all variables

Advantages of CSPs

- Formal representation language that can be used to formalize many problems types
 - Represent problem as a CSP and solve with general-purpose solver
- Able to use general-purpose solver, which are generally more efficient than standard search
 - Constraints allow us to focus the search to valid branches
 - Branches that violate constraints are removed
 - Non-trivial to do this for standard search (need manual selection of actions)

Exercise: Map Colouring

Variables:

• WA, NT, Q, NSW, V , SA, T

Openains:

∘ D_i = {red, green, blue}

Constraints:

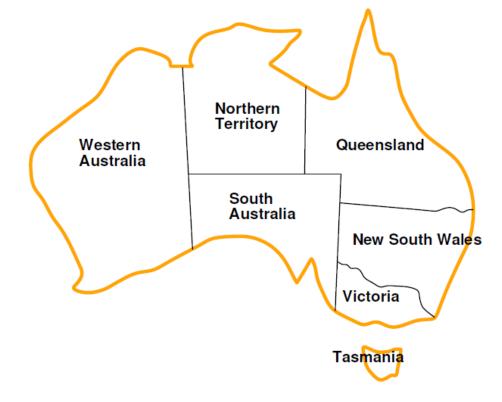
- Adjacent regions must have different colors
- E.g., WA ≠ NT (if the language allows this), or
- E.g., (WA,NT) ∈ {(red, green), (red, blue), (green, red), (green, blue), . . .}





Exercise: Map Colouring

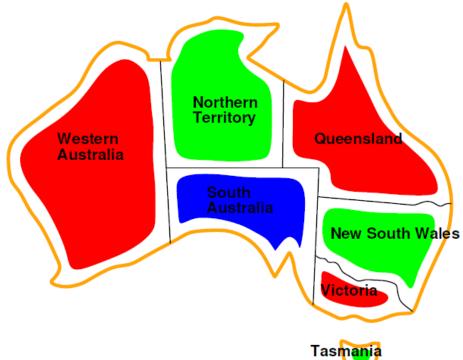
- Now, try to complete this map colouring exercise
- What strategies did you use to:
 - Select the state to start?
 - Assign colours to states?
 - Select the next state to continue?





Exercise: Map Colouring

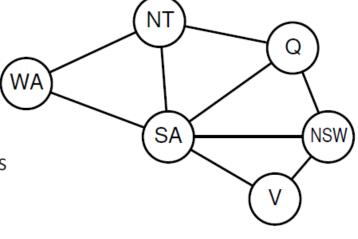
- Solutions are complete and consistent assignment for all variables
- That is, all variables are assigned and all constraints are satisfied
 - E.g., {WA=red, NT = green, Q=red, NSW=green, V=red, SA=blue, T=green}





Constraint graph

- Similar to search problems, we can represent CSPs using graphs
- Constraint graphs
 - Nodes = variables
 - Edges = constraints
- Binary CSPs
 - Each constraint related to at most two variables
- General-purpose CSP algorithms use the graph structure to speed up the search
 - E.g., Tasmania is an independent subproblem!





Varieties of CSPs

Discrete variables

- Finite domains: O(dⁿ) complete assignments for n variables, domain size d
 - e.g., map-colouring, scheduling with time limits
- Infinite domains: Integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - Need a constraint language, e.g., StartJob₁ + 5 ≤ StartJob₃

Continuous variables

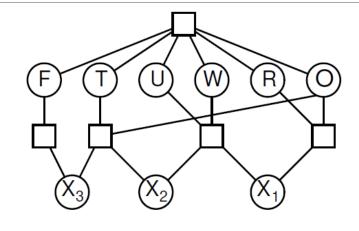
e.g., start/end times for Hubble Space Telescope observations



Varieties of Constraints

- Unary constraints involve a single variable,
 - e.g., SA ≠ green
- Binary constraints involve pairs of variables,
 - e.g., SA ≠ WA
- Higher-order constraints involve 3 or more variables
 - e.g., cryptarithmetic column constraints
- Preference (soft constraints)
 - e.g. *red is better than green* can be represented by a cost for each variable assignment, aka constrained optimization problems.

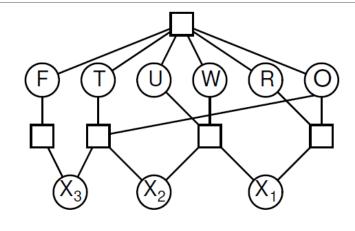
Exercise: Cryptarithmetic



- Variables: FTUWROX1X2X3
- o Domains: {0, 1, 2, 3, 4, 5, 6, 7, 8, 9}
- Constraints
 - alldiff(F, T,U,W,R,O)
 - $O + O = R + 10 \cdot X1$, etc.

Exercise: Cryptarithmetic

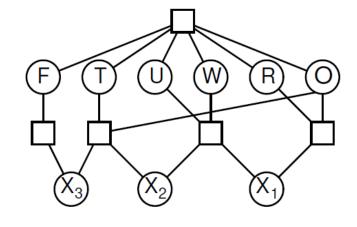
Task: Work out the remaining constraints



- Variables: FTUWROX1X2X3
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Exercise: Cryptarithmetic

Task: Work out the remaining constraints



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- o Domains: {0, 1, 2, 3, 4, 5, 6, 7, 8, 9}
- Constraints
 - alldiff(F, T,U,W,R,O)
 - $O + O = R + 10 \cdot X1$, etc.
 - \circ W +W + X₁ = U + 10 · X₂
 - $T + T + X_2 = O + 10 \cdot X_3$
 - ∘ *X*₃ = *F*

Real-world CSPs

- Assignment problems
 - e.g., who teaches what class
- Timetabling problems
 - e.g., which class is offered when and where?
- Hardware configuration
- Spreadsheets
- Transportation scheduling
- Factory scheduling
- Floor planning
- Notice that many real-world problems involve real-valued variables



Solving CSPs

- One idea is to solve CSPs as though it were standard search
 - Then we can apply the standard search algorithms
- First need to formulate this search problem



- CSPs can be easily formulated/represented as standard search problems
 - Initial State: An empty assignment {}
 - Successor function (actions): Assign a value to an unassigned variable
 - Path cost: A constant cost for every step
 - Goal test: The current assignment is complete and consistent



- For CSPs, the sequence of actions or path is not important, only the final goal state
 - E.g., for map colouring, it does not matter the sequence in which you colour the states, as long as all states are coloured with no conflicting colours
- Thus, the solution will appear at depth *n* for *n* variables
 - To find the solution, we can use depth-first search

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- Thus, the solution will appear at depth *n* for *n* variables
 - To find the solution, we can use depth-first search
 - There are potentially n!dⁿ leaves in the search tree
 - How did we get this?



- There are potentially n!dⁿ leaves in the search tree
 - How did we get this?
- For a CSP with n variables with d domains, we have:
 - Depth 1: branching factor of nd
 - Depth 2: branching factor of (n-1)d
 - Depth 3: branching factor of (n-2)d
 - 0
 - Depth n: branching factor of d



- There are potentially n!dⁿ leaves in the search tree
 - How did we get this?
- o n!dⁿ leaves is too many to search through, can we reduce it?
 - Next: Backtracking search with various enhancements

