

TRANSACTION THEORY

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Applied class : A01 Wednesday 8:00 am G16 room

a) **Write through database:**

From the diagram we can observe the following details :

- 1) Transactions T3 and T6 were initiated before the checkpoint and successfully committed before the checkpoint occurred. Therefore, no further action is necessary for these transactions.
- 2) Transactions T1, T4, and T7 were initiated prior to the checkpoint but did not reach the commit. Therefore, these transactions require an undo or rollback process.
- 3) Transactions T9 and T11 were initiated after the checkpoint and never reached the commit . In such cases, an undo or rollback operation is needed.
- 4) Transactions T2, T5, and T8 started before the checkpoint and were successfully committed before the power outage or failure. Therefore, a roll-forward or redo procedure is essential for these transactions.
- 5) Transaction T10 began after the checkpoint and was successfully committed before the power outage or failure. As a result, a roll-forward or redo process is necessary for this transaction.

The Three stages of the recovery process are as follows:

Stage 1: Compiling REDO and UNDO Lists

During this stage, we identify the last checkpoint before the crash in the log file.

Two lists are constructed:

REDO list: Contains the transaction IDs of transactions that were committed.

UNDO list: Contains the transaction IDs of transactions that never committed.

1) **Transactions T1, T4, and T7:**

These transactions were initiated before the checkpoint but did not reach the commit point. They are included in the **UNDO** list.

2) **Transactions T9 and T11:**

These transactions were initiated after the checkpoint and never reached the commit point. They are also included in the **UNDO** list.

3) **Transactions T2, T5, and T8:**

These transactions started before the checkpoint and were successfully committed before the power fail. They are included in the **REDO** list.

4) Transaction T10:

This transaction began after the checkpoint and was successfully committed before the power fail. It is also included in the **REDO** list.

Hence,

REDO List :

| | | | |
|----|----|----|-----|
| T2 | T5 | T8 | T10 |
|----|----|----|-----|

UNDO List :

| | | | | |
|----|----|----|----|-----|
| T1 | T4 | T7 | T9 | T11 |
|----|----|----|----|-----|

Stage 2: UNDO Operations (Rollback)

During this stage, **UNDO** operations are applied to transactions in the **UNDO** list, starting from the newest transactions.

UNDO List :

| | | | | |
|-----|----|----|----|----|
| T11 | T9 | T7 | T4 | T1 |
|-----|----|----|----|----|

Stage 3: REDO Operations (Rollforward)

During this stage, REDO operations to transactions in the REDO list needs to be done , starting from the oldest transactions.

REDO List :

| | | | |
|----|----|----|-----|
| T2 | T5 | T8 | T10 |
|----|----|----|-----|

After completing these three stages, the database should be restored to a consistent state based on the checkpoint, with all committed transactions reapplied (REDO) and incomplete or rolled-back transactions undone (UNDO).

b)

| TIME | TRANS | ACTION | A | B | C | D | E | F | G | H |
|------|-------|----------|--------------|-------|--------------|-------------|------------|------------|------------|------------|
| 0 | T1 | Read A | S(T1) | | | | | | | |
| 1 | T2 | Read B | | S(T2) | | | | | | |
| 2 | T1 | Read C | | | S(T1) | | | | | |
| 3 | T4 | Read D | | | | S(T4) | | | | |
| 4 | T5 | Read A | S(T5) | | | | | | | |
| 5 | T2 | Read E | | | | | S(T2) | | | |
| 6 | T2 | Update E | | | | | X(T2) | | | |
| 7 | T3 | Read F | | | | | | S(T3) | | |
| 8 | T2 | Read F | | | | | | S(T2) | | |
| 9 | T5 | Update A | T5 wait T1 | | | | | | | |
| 10 | T1 | Commit | X(T5) | | | | | | | |
| 11 | T6 | Read A | T6 wait T5 | | | | | | | |
| 12 | T5 | Rollback | S(T6) | | | | | | | |
| 13 | T6 | Read C | | | S(T6) | | | | | |
| 14 | T6 | Update C | | | X(T6) | | | | | |
| 15 | T7 | Read G | | | | | | | S(T7) | |
| 16 | T8 | Read H | | | | | | | | S(T8) |
| 17 | T9 | Read G | | | | | | | S(T9) | |
| 18 | T9 | Update G | | | | | | | T9 wait T7 | |
| 19 | T8 | Read E | | | | | T8 wait T2 | | | |
| 20 | T7 | Commit | | | | | | | X(T9) | |
| 21 | T9 | Read H | | | | | | | | S(T9) |
| 22 | T3 | Read G | | | | | | | T3 wait T9 | |
| 23 | T10 | Read A | S(T10) | | | | | | | |
| 24 | T9 | Update H | | | | | | | | T9 wait T8 |
| 25 | T6 | Commit | | | | | | | | |
| 26 | T11 | Read C | | | S(T11) | | | | | |
| 27 | T12 | Read D | | | | S(T12) | | | | |
| 28 | T12 | Read C | | | S(T12) | | | | | |
| 29 | T2 | Update F | | | | | | T2 wait T3 | | |
| 30 | T11 | Update C | | | T11 wait T12 | | | | | |
| 31 | T12 | Read A | S(T12) | | | | | | | |
| 32 | T10 | Update A | T10 wait T12 | | | | | | | |
| 33 | T12 | Update D | | | | T12 wait T4 | | | | |
| 34 | T4 | Read G | | | | | | | T4 wait T3 | |

i) Item A : T10 wait T12

Item F : T2 wait T3

Item C : T11 wait T12

Item G :T3 wait T9

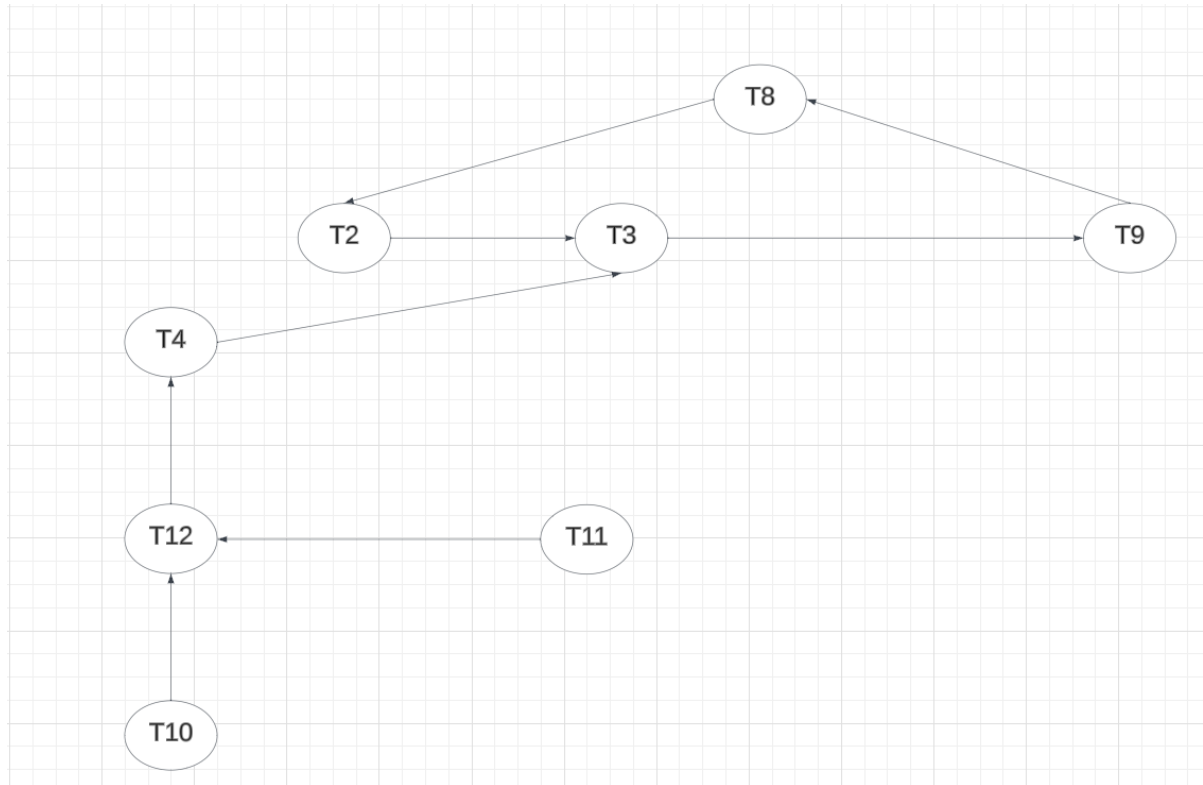
Item D: T12 wait T4

Item G :T4 wait T3

Item E : T8 wait T2

Item H : T9 wait T8

ii) Wait for graph indicating the state of waiting locks at time 34 is as follows:



iii) Deadlock exists as T2 waits for T3 ,T3 wait for T9, T9 waits for T8 and T8 wait for T2 which results in a deadlock.
The transactions involved in Deadlock are **T9 , T8 , T3 and T2**.