

## ✳ 1) Introduction to Transport Layer

- It is **Layer 4 of OSI model (Layer 3 in TCP/IP)**.
  - Provides **end-to-end communication between applications** on different computers.
  - Ensures data is **delivered accurately, in order, and without errors**.
  - Uses **port numbers** to deliver data to the correct application (e.g. web uses port 80).
  - **Example protocols: TCP (reliable), UDP (unreliable)**
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## ✳ 2) Multiplexing and Demultiplexing

- **Multiplexing:** Combining messages from different applications into segments for sending.
  - **Demultiplexing:** Delivering received segments to correct applications.
  - Uses **IP addresses + port numbers** to identify correct destination.
  - **UDP:** Uses only destination port for demux.
  - **TCP:** Uses 4-tuple → source IP, source port, destination IP, destination port.
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## 📦 3) Connectionless Transport – UDP

- UDP = User Datagram Protocol
  - **No connection setup** between sender and receiver.
  - Adds **source/destination port numbers** and passes to network layer.
  - **Unreliable & unordered:** Packets may get lost, duplicated, or arrive out of order.
  - **Used in:** DNS, SNMP, online gaming, live video/audio streaming.
  - **Header fields:** Source port, Destination port, Length, Checksum
  - **Checksum:** Used to detect errors in the segment.
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## ✅ 4) Principles of Reliable Data Transfer (RDT)

- Reliable protocols are built over unreliable channels.
- Use **ACK (Acknowledgement), NAK (Negative ACK), sequence numbers**, and **timers** to ensure reliability.
- **rdt 1.0:** Assumes perfect channel (no errors).
- **rdt 2.0:** Handles bit errors using checksum + ACK/NAK (stop-and-wait).
- **rdt 2.1:** Adds sequence numbers (0,1) to detect duplicate packets.
- **rdt 2.2:** Uses only ACKs, no NAK. Duplicate ACK acts like NAK.

- **rdt 3.0:** Handles errors + packet loss using **timeout + retransmission**.
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## 5) Pipelined Protocols

- Allows **multiple unacknowledged packets in transit at the same time**.
- Improves efficiency by using network bandwidth fully.
- Two types:

### a) Go-Back-N (GBN)

- Sender sends up to **N unACKed packets**.
- If one packet is lost, **all after it are retransmitted**.
- Receiver only accepts packets in order and sends **cumulative ACK**.
- Uses **one timer for oldest unACKed packet**.

### b) Selective Repeat (SR)

- Sender sends up to **N unACKed packets**.
  - Receiver **accepts out-of-order packets** and buffers them.
  - Sender **retransmits only lost packets**.
  - Uses **separate timer for each packet**.
  - More efficient than GBN.
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## 6) Connection-Oriented Transport – TCP

- TCP = Transmission Control Protocol.
- **Reliable, connection-oriented, byte-stream protocol.**
- Breaks data into segments and ensures **ordered, error-free delivery**.
- **Example:** Like courier service with tracking number.
- **Features:**
  - Reliable delivery
  - Ordered data transfer
  - Error checking (checksum)
  - Retransmission of lost data
- **3-Way Handshake:**
  1. SYN – client says "I want to connect"
  2. SYN+ACK – server replies "I'm ready"

3. ACK – client confirms "Let's start"
- **TCP Segment:** includes source port, dest port, sequence number, ack number, flags (SYN, ACK, FIN), checksum, window size.
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## 7) Flow Control

- Prevents **fast sender from overwhelming slow receiver**.
  - Receiver advertises a **window size (rwnd)** showing how much data it can handle.
  - Sender sends only within this window until it gets ACKs.
  - Avoids receiver buffer overflow.
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## 8) Congestion Control

- Prevents **too much data entering the network** which causes packet loss/delay.
  - Causes:
    - Too many senders sending fast
    - Full router buffers
    - Retransmissions increasing load
  - Two approaches:
    - **End-to-End (TCP):** Detect congestion via packet loss/delay.
    - **Network-assisted:** Routers give feedback to senders (rare).
  - **TCP Slow Start:**
    - Starts with small congestion window.
    - Gradually increases window size until packet loss occurs.
    - Helps find network's capacity safely.
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## 9) RPC (Remote Procedure Call)

- Allows a program to **call a function on another computer as if it is local**.
- Programmer just writes `result = add(5,7)` even if `add()` is on remote system.
- Uses transport layer to:
  - Send function call request as a message.
  - Wait for result and return it.
- Uses **TCP for reliable RPC** (e.g. gRPC) or **UDP for fast RPC** (e.g. ONC RPC).

- RPC itself is **not a protocol**, it **uses transport protocols**.