Unit - 5

Data Link Layer – Introduction

- Nodes → Hosts & Routers (devices).
- **Links** → Channels connecting adjacent nodes (Wired, Wireless, LANs).
- At this layer, data is called a **Frame** (datagram + header & trailer).
- **Job** → Transfer a datagram from one node to its physically adjacent node.

Link Layer Services

1. Framing

o Encapsulates datagram into a frame (adds header + trailer).

2. Link Access

- o Uses MAC addresses in frame headers to identify source & destination.
- o MAC ≠ IP (different addressing).

3. Reliable Delivery

- o Guarantees each datagram moves across the link without error.
- Done using Acknowledgment + Retransmission.

4. Flow Control

o Ensures proper pacing between sender & receiver.

5. Error Detection & Correction

- o Errors occur due to signal attenuation & noise.
- o Receiver detects/corrects errors or requests retransmission.

Framing Techniques

1. Bit Stuffing

- Used in bit-oriented protocols.
- Frames are marked with an **8-bit flag: 01111110** at beginning & end.
- Problem → Same flag pattern may appear in data part.
- Solution → After **five consecutive 1s**, sender inserts an extra **0**.
- Receiver removes the stuffed 0.

***** Example (from PPT):

If data has sequence 01111110 (flag), sender changes it to 011111010. Receiver removes 0 before passing data.

2. Byte Stuffing

- Uses a **special flag byte** at start & end of frame.
- Problem → If flag byte appears inside data, confusion occurs.
- Solution → Sender inserts an **ESC (Escape) byte** before it.
- Receiver removes ESC before passing data up.

Example (from PPT):

If data has accidental flag 01111110, sender sends ESC + flag. Receiver removes ESC and keeps flag as data.

K Error Detection & Correction Techniques

1. Parity Check

- Add 1 extra bit → keeps total number of 1s even (even parity) or odd (odd parity).
- Receiver counts 1s → If parity matches → data accepted.
- Detects single-bit errors, but not multiple-bit errors.

***** Example:

Data: 1011 (3 ones = odd). With even parity \rightarrow add 1 \rightarrow 10111. Receiver checks \rightarrow total 4 ones (even) \rightarrow correct.

2. Checksum

• Sender side:

- o Data divided into k segments of m bits.
- Segments added using 1's complement arithmetic.
- Complement of sum = checksum \rightarrow sent with data.

Receiver side:

- o Adds all segments + checksum.
- If result = $0 \rightarrow$ correct. Else \rightarrow error.

* Example (from PPT):

Two 16-bit words are added.

If wraparound occurs, carry is added back.

If final result = 111111111111111, data is correct.

If any 0 appears, error detected.

3. Cyclic Redundancy Check (CRC)

• Most powerful error detection. Based on binary division.

• Sender side:

- 1. Append (k-1) zeros (where k = degree of polynomial + 1).
- 2. Divide data by generator polynomial.
- 3. Remainder is added to data \rightarrow transmitted.

Receiver side:

- Divides received data by same polynomial.
- If remainder = $0 \rightarrow$ no error. Else \rightarrow error.

* Example (from PPT):

Data = 100100, Polynomial = 1101 (x^3+x^2+1) .

Sender appends 3 zeros \rightarrow 100100000.

Remainder = 001.

Transmitted data = 100100001.

Receiver divides \rightarrow remainder 000 \rightarrow data correct.

Multiple Access Links

• Point-to-Point Link

- One sender ← One receiver.
- o Example: Direct cable connection between two computers.

Broadcast Link

- o Multiple senders & receivers share one channel.
- When one sends, all receive a copy.
- \circ Example: Classic Ethernet LAN \rightarrow all devices connected via same bus cable.