

1. Decoding Logical bit 1 and 0 by the best station.
- (2) Suppose station A (code a) station B (code b) station C (code c) station D (code d) station E (code e) sends following logical bits:
 A : Logical 1, B : Logical 0, C : logical 1, D : logical 0 and E does not send any thing
- (3) Base station shall give the tick signal and A, B, C, D transmit their data
- (4) Base station then calculate the sum signal S ~~by~~ chip by chip

$$S = a + \bar{b} + c + \bar{d}$$
- Now base station knows codes allocated to A, B, C, D, E are as a, b, c, d, e.
- Base station shall parallelly

iii) If the any of the above value = 1 : station has transmitted logical 1
= -1 : Logical 0
= 0 : station has not transmitted anything

iv) Reading
Let us take the above example for A: $S.a = (a + \bar{b} + c + \bar{d}) \cdot a$
 $= a \cdot a + \bar{b} \cdot a + c \cdot a + \bar{d} \cdot a = 1 + 0 + 0 + 0 = 1 \Rightarrow A \text{ has sent logical 1}$

$S.b = -1 \Rightarrow \text{Logical 0}$
 $S.c = 1 \Rightarrow \text{Logical 1}$

$S.e = 0$
 $S.f = 0$
 $S.m = 0$ } A has not sent anything

02. Control channels are also created in code: Similar (not exactly same) to TDMA

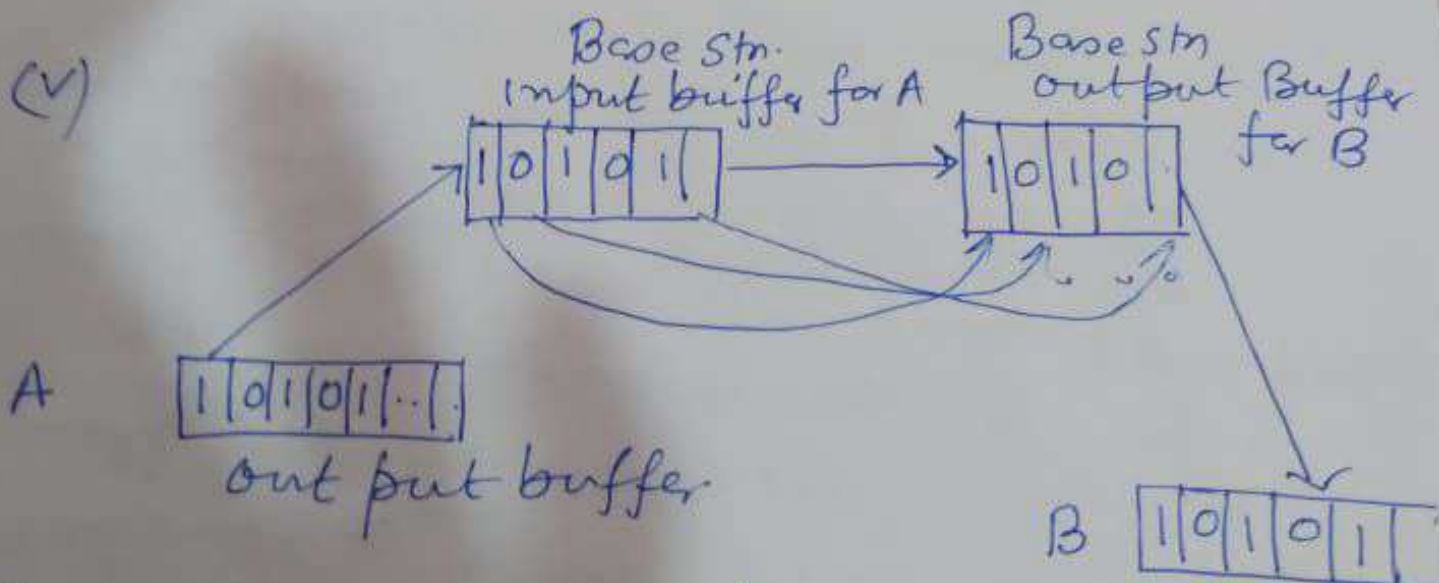
03. (i) To make a call per mobile A send send a call request to base station giving his and B's number through request channel

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(i) Base shall check if A has sufficient Balance, B is not busy and has one uplink and down link code free than it will connect the call allocating uplink and down link code (= channel to A and B. Base stn shall register

(ii) Station shall not a Connection $A \rightarrow B$. Base stn.

(iv) Actual Route A to B to B



A and B shall have their own output buffer to transmit and input buffer to receive respectively.

(vi) on making connection ~~to~~ from A to B, Base stn other than channel allocation to A, & B like above page 3(ii), shall also allocate an input buffer for A and output buffer for B in Base stn's data base.

(vii) - Now Suppose A sends logical 1 bit (first bit in his input buffer)

- Base stn shall decode it like page (2) vii and find it as logical 1 and store it in A's input buffer in his database. As base stn shall note that

(viii) for A, the connection is to B, it will transfer that bit from input buffer of A to output buffer of B.

(ix)

Now Assume there are other connections also.

$A \rightarrow B$

$C \rightarrow E$

$F \rightarrow G$

$K \Rightarrow L$

in buffer of A to B transfer; out buffer; all first

(x) So in similar manner, first bits will be transferred from input buffer C to output buffer E, then from input buffer F to output buffer G, then from input buffer K to output buffer L.

(xi) Base station shall note the bit content of current bits of B, E, G, L.

Suppose $B \Rightarrow 1$, $E \Rightarrow 0$, $G \Rightarrow 1$, $L \Rightarrow 0$

Now base station shall make a sum $S = B + \bar{E} + G + \bar{L}$ (bit chip by chip) and send it in the downlink.

(x) All stations B, E, G, L and also others will receive S through

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their corresponding downlink channels.

(xi) After receiving $\sum S$, each station who are in connection shall decode their received by as follows

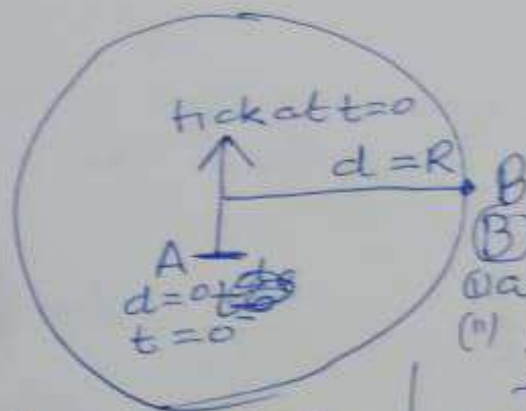
B: b	}	b, e, g, e are
E: e		the downlink
G: g		codes assigned
K: e		to b, e, g, l.

04. Synchronization problems in uplink channels.

(i) The assumption was that after giving the tick signal at $t=0$, all stations shall receive it at $t=0$ and send their codes in uplink channel at $t=0$ and base station gets all codes at $t=0$ so that the first chip of all codes as well reach at $t=0$ and other chips will also reach simultaneously.

for all

(ii) But above is not true in Practical situation ~~for~~ due to propagation delay

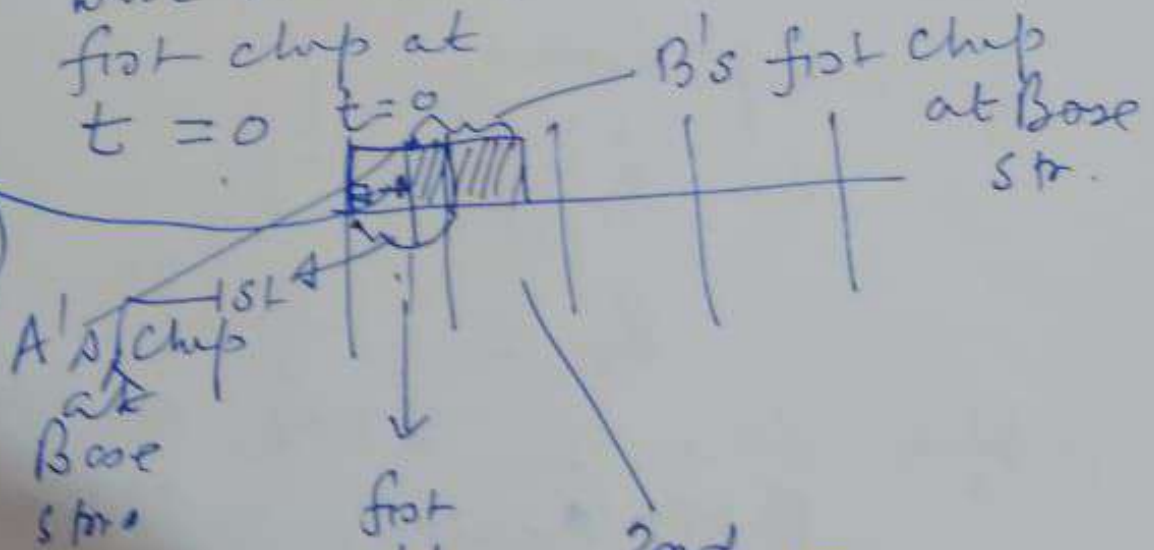


A receives tick at $t=0$
 sends first chip at $t=0$
 Base stn Receives first chip at $t=0$

(i) A starts sending 2nd first chip
 (ii) first chip received by base stn at $t = 2t_p = \frac{2R}{c}$
 ($t_p + t_p$)

(iii)

This delay
 $= 2 \frac{R}{c}$



(iii) while Base stn calculates sum of first

chip at $t = 0$, ^(Far 08) it does not get the 1st chip of B. So the calculation of sum signal shall go wrong.

(iv) Soln.

Base sta instead of calculating the sum signal chip by chip at the beginning of chip time, it calculates the chip at the middle of chip time.

If $\frac{2R}{v} = \frac{t_c}{2}$ then the chip of B (farthest station) shall also reach at time $\frac{t_c}{2}$, and all other chip value is available at the middle of chip time.

$$\text{So if } \frac{2R}{v} = \frac{t_c}{2}$$

$$R = \frac{t_c v}{2}$$

$v = \text{speed of signal}$

Suppose $BW_u = 2 \text{ MHz}$, Modulation = 4

Baud rate (Bandwidth) = 2 M baud/sec.

Chip rate = $2 \times 2 = 4 \text{ M chips/sec}$

$$\text{chip time } t_p = \frac{1}{4 \times 10^6} \text{ sec}$$

$R = \text{known if } v = 3 \times 10^8 \text{ m/sec}$