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e. i. CSMA/CD Protocol.

(i) DIFS carrier sensing ~~the~~
— A station before transmition
starts sensing the carrier for
continuous DIFS duration.
If the channel ~~become busy~~
during the DIFS for continuous
DIFS duration Then The
station goes for Back-off
otherwise keep on sensing
the carrier free for continuous
DIFS duration.

(ii) For back-off following
Definitions are required.
— (a) senders initial window
size W_0 which is normally
 $= 04$.

(b) senders current window
variable W_i

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(c) Initially $w_2 = w_0 = 4$.

(d)(i) The sender picks up a

random num $r \in \{0, \dots, w_i - 1\}$

(d)(ii) If $r = 0$ Then sender immediately sends ~~RTS frame~~ ^{RTS frame}. 3.

(e) If $r > 0$ sender starts decreasing

r after every slot time, if

the channel is free. ~~sensed~~

(f) If the channel ~~is busy~~ busy, then

The station starts sensing the corner again ~~for~~ for continuous DIFS duration. When it finds

the channel free for continuous DIFS duration it starts back-off with the value of r , with which it left back-off in (f).

(g)(i) After decrementing r as mentioned in step (e) if value of r becomes 0, Then the sender sends a RTS frame.

(g)(ii) If more than one sender started DIFS corner sensing and completed that at the same time and picked up the

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The same random no 'r' Then they on decreasing r slot by slot as mentioned in (e), will find $r = 0$ at the same time Then sends RTS frame simultaneously Then there shall be collision of RTS frame.

Here the senders shall detect the collision ~~of RTS frame~~ in a definite mechanism as mentioned later and mitigate post collision activity.

- (i) If only one station transmitted RTS frame as mentioned in q(ii)
then there will not be collision of RTS frame Assume A transmit RTS for B.
(ii) If ~~only~~ Then all the non-hidden terminals shall be alerted not to transmit for RTS duration (duration mentioned in RTS frame)
Assume receiver B ~~sends~~ gets the RTS frame correctly (No hidden terminal transmitting)

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- (k) Then the B shell sends the CTS frame to A.
- (l) If the CTS frame comes back to A correctly Then the ~~near~~ hidden terminals shall be alerted and (A) shall get the channel with full certainty. But CTS frame can collide also because of the hidden terminal E.
- (m) If no collision of CTS frame Then (A) shall transmit Data frame. Now data frame will not collide, but can be corrupted by transmission here.
- (n) If Receiver B receives the frame correctly, Then it will send acknowledgement to A.
- (o) After A gets the acknowledgement Then one frame transmission sequence is successfully completed.

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- (p) If RTS frame collided ~~does~~
as mentioned in q(2i) or
RTS frame is corrupted by
transmission error or CTS frame
collided or CTS' frame was
in transmission error Then
B shall not send CTS frame.
If A does not receive the
CTS frame ^{within} a time known
as $CTS_{TIMEOUT} = DIFS + RCTS$
frame transmission time Then
A will assume that the collision
has taken place (Although no
collision in case of transmission
error) & it will start the post
collision activity.
- (q) Also if after sending frame
does not come back to A [?] within
a time period known as ACKTIME
 $- OUT = DIFS + \text{Acknowledgment}$
frame transmission time
Then also it will start the
post collision Activity.
Because data frame can be corrupted and
B does not get the data frame and does

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not send the acknowledgement frame or B sends the Acknowledgement frame receiving data frame correctly but acknowledgement is corrupted by error.

- (r) If both above case (p) & (q)
(g) The post collision activity starts started by any sender here A as example.
- (s) Post collision Activity P is done by running an Algorithm known as binary back off exponential algorithm.
Here the sender double their Wi variable i.e. $W_i = 2 W_i$ (Here W_i variable i.e. $W_i = 2 W_i$)
on first collision W_i becomes 8, 2nd $W_i = 16 \dots$
Then repeat the same steps from 1(i) - 7 page(1).
- (t) ~~But~~ The Above process goes on till 10th collision.
(u) From 11 to 16 collision the Above process goes on

but the window value \uparrow is not incremented.

(v) On 17th collision, assuming that something serious has happened (Cable break up etc) the interface network Interface card shell detect it and through SNMP protocol the network shell be delivered as broken down.

v(ii) How Binary exponential Backoff algorithm reduces the probability of collision.
on the probability first attempt probability of collision (that two stations picking up same random no for 0 to 3) is $\frac{1}{4} \times \frac{1}{4}$ (given probability) i.e. probability of collis. = $\frac{1}{16}$

- 2nd attempt probability

$$\text{of collision} = \frac{1}{8} \times \frac{1}{8} = \frac{1}{64}$$

In the 10th to 16th attempt $\frac{1}{1024} \times \frac{1}{1024} \dots$