

## **Dynamic Memory Allocation Problem**

For the solution of the Dynamic memory allocation problem, I have used the certain data structures to the ease the process of memory allocation and deallocation.

The following structures will be used in the program :

- **struct mem\_space** : This structure is assigned to each memory block whether used or free. It consists two variables and four pointers.

Two variables are:

- ➔ size : Size of the memory block
- ➔ used : It shows whether the block is allocated or free.

Pointers are :

- ➔ prev : Points to the previous allocated block of memory
- ➔ next : Points to the next allocated block of memory
- ➔ prev1 : Points to the previous free block of memory ( Not NULL if this structure is of free memory block otherwise NULL)
- ➔ next1 : Points to the next free block of memory ( Not NULL if this structure is of free memory block otherwise NULL)

Size	Used	Prev	Next	Prev1	next1
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- **Memory pool (void \*mp)** : This is an space which is initialized in the **MemInit()** function which will be used as the memory pool of which the memory will be allocated and deallocated.
- **Hash Table** : This hash table is used to store the free block of memory according to a hash funtion based on its size. The hash function used in this algorithm is greatest integer ( $\log_2(\text{size})$ ) . Each entry of this Hash-Table points to list of free spaces. These free spaces have been added using the above mentioned hash-function.  
Eg : A free space of size 16 to 31 . It will be mapped to position  $x = 4$  in the Hash-Table. Hence we can refine our search to the space of  $2^n$  sizes.

$2^0-2^1$	$2^1-2^2$	$2^2-2^3$	$2^3-2^4$	.....	$2^{n-1}-2^n$
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## Algorithm :

The basic approach for solving dynamic memory allocation problem is to form a list of free blocks and search through this list for the first free block of big enough size to the requested size. This is called the first-fit algorithm. I have implemented a hash-table data structure to optimize this algorithm.

**void MemInit()** : Allocate a 160 KB of memory to the memory pool of which the memory block will be allocated and deallocated. It also initializes the hash-table with all entries pointing to NULL except the last entry pointing to a list containing one element of a free block of memory of size 160 KB.

**MyMalloc(int bytes)** : This function would allocate a block of memory to a request of memory of particular size. It looks at the particular entry in the hash table if it finds a block of memory of that size it allocates it or searches further down the list for a free block of larger size and allocates it. It also adds an entry in the hash table for the left-over free block which is  $(\text{size} - \text{size\_requested} - \text{sizeof}(\text{struct mem\_space}))$ . Whenever it finds a free block of memory it also assigns to it a **mem\_space struct** to store the information of that block. This function returns a pointer to the starting address of free memory block.

**MyFree(void \*p)** : This function would deallocate the allocated block of memory with a pointer p pointing to the starting address of the allocated memory. It checks in the **struct mem\_space** assigned to the block and frees that block of memory. If that block of memory that is to be freed is beside a free block then these block are coalesced into one and an entry for the combined block is added to the hash-table according to the hash-function.

## Order analysis :

In memory allocation, it would take  $O(1)$  ( amortized cost ) time.

In memory deallocation, it takes  $O(1)$  ( amortized cost ) with some overhead in coalescing.