



Subject : Data Structures and Algorithms

Assignment: 03

Problem Definition:

Create an array inordered threaded binary tree and perform inorder & preorder traversal. Analyze time & space complexity of the algorithm

Objectives:

- To understand threaded binary trees
- To understand & implement various operations on threaded binary tree - creation, traversal, deletion.
- To avoid overhead of stack data structure while performing preorder & inorder traversal.

Outcomes:

- To implement threaded binary trees & its operations
- To use classes in c++ & pointers.

Hardware Requirements:

Manufactures: Acer Inc.

System type: X-64 type PC

Processor & RAM: Intel (R) Core i5-8265U @ 1.60 GHz & 8 GB of RAM

Software Requirements:

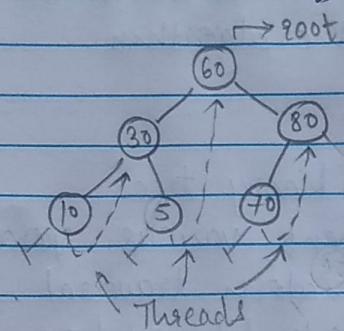
Operating System: MS Windows 10

IDE: Eclipse (2020 Edition)

Theory:

A threaded binary tree is used to make improvements in binary tree traversal by eliminating stack requirements. It is faster.

A binary tree is made threaded by making all right child pointers point to the inorder successor of the node (if exists)



ADT:

```
class Node {
private:
    int data;
    Node *lchild, *rchild;
    bool lTh, rTh;
};
```

```
class TBST {
private:
    Node *root, *head;
public:
    Node* getRoot();
    void InsertNode(int x);
    void ThInorder();
    void ThPreorder();
    bool Search();
    void delete();
};
```



Pseudocode:

insert(key) { // inserting key in threaded BST

if (root == NULL) {

root = new Node(x);

root->lchild = root->rchild = head;

}

else {

Node* curr = root, *prev = NULL;

while (curr != head) {

prev = curr;

if (x < curr->data) {

if (curr->lTh == false)

curr = curr->lchild;

else break;

}

else if (x > curr->data) {

if (curr->rTh == false)

curr = curr->rchild;

else break;

}

else if (x == curr->data)

return;

}

curr = new Node(x);

if (x < prev->data) {

curr->rchild = prev;

curr->lchild = prev->lchild;

prev->lchild = curr;

prev->lTh = false;

}



else if ($x > \text{prev} \rightarrow \text{data}$) {

 curr \rightarrow lchild = prev;

 curr \rightarrow rchild = prev \rightarrow rchild;

 prev \rightarrow rchild = curr;

 prev \rightarrow lth = false;

}

Th Inorder() { // Inorder traversal

 if ($\text{root} == \text{NULL}$) {

 print("EMPTY TREE")

 return;

}

 Node* curr = root;

 while ($\text{curr} \rightarrow \text{lth} == \text{false}$)

 curr = curr \rightarrow lchild;

 while ($\text{curr} != \text{head}$) {

 print($\text{curr} \rightarrow \text{data}$);

 if ($\text{curr} \rightarrow \text{rth} == \text{false}$) {

 curr = curr \rightarrow rchild;

 while ($\text{curr} \rightarrow \text{lth} == \text{false}$)

 curr = curr \rightarrow lchild;

}

 else curr = curr \rightarrow rchild;

}

}

Th Preorder() { // Preorder traversal

 if ($\text{root} == \text{NULL}$) {

 print("EMPTY TREE")

 return;

}

(3)



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Node* curr = root;

while (curr != head) {

printf(curr->data)

if (curr->lTh == false)

curr = curr->lchild;

else {

if (curr->rTh == false)

curr = curr->rchild;

else {

while ((curr == head) || curr->rTh == true)

curr = curr->rchild;

if (curr == head)

break;

curr = curr->rchild;

{

{

deleteNode (int x);

Node* curr = NULL, *parent = NULL;

if (!search (root, x, curr, parent)) {

printf ("NOT FOUND");

return;

{

 if (curr->lTh == false && curr->rTh == false) (1) Intermediate
 Node

Node* temp = curr->rchild;

parent = curr;

while (temp->lTh == false) {

parent = temp;

temp = temp->lchild;

{



$\text{curr} \rightarrow \text{data} = \text{temp} \rightarrow \text{data};$
 $x = \text{temp} \rightarrow \text{data};$
 $\text{curr} = \text{temp}.$

{

else

If ($\text{curr} \rightarrow \text{LTh} == \text{true}$ || $\text{curr} \rightarrow \text{RTh} == \text{true}$) { // Leaf Node.

If ($\text{curr} == \text{parent} \rightarrow \text{lchild}$) {

$\text{parent} \rightarrow \text{lchild} = \text{curr} \rightarrow \text{lchild};$

$\text{parent} \rightarrow \text{LTh} = \text{true};$

{

else if ($\text{curr} == \text{parent} \rightarrow \text{rchild}$) {

$\text{parent} \rightarrow \text{rchild} = \text{curr} \rightarrow \text{rchild};$

$\text{parent} \rightarrow \text{RTh} = \text{true};$

{

delete curr;

{

else if ($\text{curr} \rightarrow \text{LTh} == \text{false}$ & $\text{curr} \rightarrow \text{RTh} == \text{true}$) {

Node^{*} temp = $\text{curr} \rightarrow \text{lchild};$

If ($\text{parent} \rightarrow \text{lchild} == \text{curr}$)

$\text{parent} \rightarrow \text{lchild} = \text{temp};$

else $\text{parent} \rightarrow \text{rchild} = \text{temp};$

while ($\text{temp} \rightarrow \text{RTh} == \text{false}$)

$\text{temp} = \text{temp} \rightarrow \text{rchild};$

$\text{temp} \rightarrow \text{rchild} = \text{curr} \rightarrow \text{rchild};$

delete curr;

{

else if ($\text{curr} \rightarrow \text{RTh} == \text{true}$ & $\text{curr} \rightarrow \text{LTh} == \text{false}$) {

Node^{*} temp = $\text{curr} \rightarrow \text{rchild};$

If ($\text{parent} \rightarrow \text{lchild} == \text{curr}$)

$\text{parent} \rightarrow \text{lchild} = \text{temp};$

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else parent->lchild = temp;

while (temp->lTh == false)

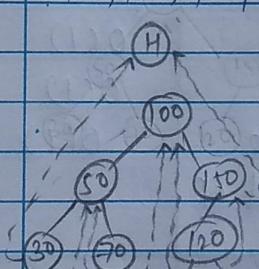
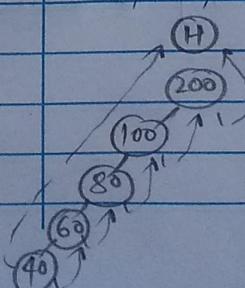
temp = temp->lchild;

temp->lchild = curr->lchild;

delete curr;

{ }

Test Cases:

No	Testcase Description	Expected O/p	Actual O/p.	Result:
1>	Insert nodes 100, 150, 120, 50, 30, 70	Threaded BST should be Created	Threaded BST is Created	Passed
				
	Inorder Traversal:	30, 50, 70, 100, 120, 150	30, 50, 70, 100, 120, 150	Passed
	Preorder Traversal:	100, 50, 30, 70, 120, 150	100, 50, 30, 70, 120, 150	Passed
2>	Insert Nodes 200, 100, 80, 60, 40	Threaded BST should be Created	Threaded BST is Created	Passed
				



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Inorder Traversal	40, 60, 80, 100, 200	40, 60, 80, 100, 200	Passed
Preorder Traversal	200, 100, 80, 60, 40	200, 100, 80, 60, 40	Passed

Algorithm Analysis (Time & Space Complexity)

Method.	Time complexity	Space complexity
1) InsertNode()	O(n) $O(h)$ where h is height of tree	No extra space required hence $O(1)$
2) InorderTraversal()	$O(n)$	$O(1)$
3) PreorderTraversal()	$O(n)$	$O(1)$
4) DeleteNode()	$O(h)$ where h is height of tree	$O(1)$

Applications of Threaded Binary Tree:

- 1) Threaded binary trees are used for various searching algorithms.
- 2) They are used for avoiding auxiliary stack for traversal of binary trees.
- 3) They are also useful for maintaining a sorted stream of data.



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Conclusion:

Through this assignment, we learned about threaded binary trees & their usage, along with their traversals.