Heterogeneous Donor Circles for Fair Liver Transplant Allocation*

Shubham Akshat

The Robert H. Smith School of Business, University of Maryland, College Park, MD 20742, sakshat@umd.edu

Sommer E. Gentry

Department of Mathematics, United States Naval Academy, Annapolis, MD 21402; and Johns Hopkins University School of Medicine, Baltimore, MD 21287, gentry@usna.edu

S. Raghavan

The Robert H. Smith School of Business and Institute for Systems Research University of Maryland, College Park, MD 20742, raghavan@umd.edu

The United States (U.S.) Department of Health and Human Services is interested in increasing geographical equity in access to liver transplant. The geographical disparity in the U.S. is fundamentally an outcome of variation in organ supply to patient demand (s/d) ratios across the country (which cannot be treated as a single unit due to its size). To design a fairer system, we develop a nonlinear integer programming model that allocates organ supply in order to maximize the minimum s/d ratios across all transplant centers. We design circular donation regions, that are able to address issues raised in legal challenges to earlier organ distribution frameworks. This allows us to reformulate our model as a set-partitioning problem. Our policy can be viewed as a heterogeneous donor circle policy, where the integer program optimizes the radius of the circle around each donation location. Compared to the current policy that has fixed radius circles around donation locations, the heterogeneous donor circle policy greatly improves both the worst s/d ratio, and the range between the maximum and minimum s/d ratio. We found that with the fixed radius policy of 500 nautical miles (NM) the s/d ratio ranges from 0.37 to 0.84 at transplant centers, while with the heterogeneous circle policy capped at a maximum radius of 500 NM the s/d ratio ranges from 0.55 to 0.60, closely matching the national s/d ratio average of 0.5983. Our model matches supply and demand in a more equitable fashion than existing policies, and has a significant potential to improve the liver transplantation landscape.

Key words: health care policy, liver transplant, geographical disparity, optimization

^{*}The data reported here have been supplied by the Hennepin Healthcare Research Institute as the contractor for the Scientific Registry of Transplant Recipients (SRTR). The interpretation and reporting of these data are the responsibility of the authors and in no way should be seen as an official policy or interpretation by the SRTR or the U.S. Government.

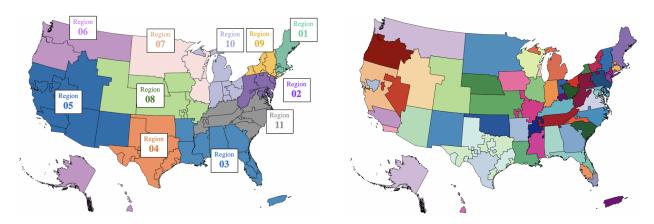


Figure 1 Pre February 4, 2020 policy divided the U.S. in to 11 regions (left), which comprised of 58 DSAs (right).

1 Introduction

In 2019, 8896 liver transplants took place in the United States (U.S.) while 12,941 patients were added to the waiting list. Unfortunately, in 2019, on average three people in the U.S. died every day awaiting a liver transplant for a total of 1,202 lives lost. Since demand for liver transplant outstrips supply, allocating deceased donor livers judiciously and justly is extremely important. For thirty years (from 1989 to Feb. 4, 2020) transplant allocation policy divided the U.S. into 58 Donation Service Areas (DSAs) that were grouped into 11 geographical regions (Figure 1). Livers were offered to the candidates in a DSA in decreasing order of medical urgency, quantitatively measured by Model for End-stage Liver Disease (MELD) score. (Pediatric End-Stage Liver Disease (PELD) severity score, a measure calculated slightly differently, is used for patients \leq 12 years old.) MELD score reflects the probability of death within a 3-month period and ranges from 6 to 40; a higher score indicates greater mortality risk (Freeman et al. 2002). More serious patients are assigned Status 1A and 1B, and their number is fewer than 50 nationwide at any time.

By law, the deceased donor organs are national resources in the U.S. The U.S. government created the Organ Procurement and Transplantation Network (OPTN) in 1984 to coordinate a nationwide transplant system and optimize the usage of the limited resource of donor organs for transplants. Since 1986, the United Network for Organ Sharing (UNOS), a non-profit private organization has overseen the operations of the OPTN. A key regulatory framework that guides organ transplantation is the "Final Rule" that was adopted in 1998 by the Department of Health and Human Services (HHS) to establish a more detailed framework for the structure and operations of the OPTN (HHS 1998). The Final Rule requires that policies shall not be based on the candidate's place of residence or place of listing, except to the extent required by the other requirements of the Rule.

Geographic inequity in access to liver transplantation across DSAs is well documented in the literature (see Yeh et al. 2011). Indeed, as early as 2008 a HHS Advisory Committee on Transplantation recommended that organ allocation be evidence-based and not on the arbitrary boundaries of the DSAs; and in 2012 the OPTN board adopted a strategic plan that included reducing geographic disparities in access to transplantation. Despite implementation in 2013 of broader sharing of organs in a region for candidates with MELD scores ≥ 35 ; geographic inequities remained in the system. The U.S. Scientific Registry of Transplant Recipient's (SRTR) Liver Transplant Waiting List Outcomes Tool¹ (that is built on historical data from 2017 to 2019) shows that for waitlist candidates in Los Angeles with MELD score in the range 25-29, only 15% received a transplant within 90 days, while for candidates in Indianapolis (with MELD scores in the range 25-29), 72% received a transplant within 90 days. The DSA/Region allocation policy resulted in significant disparities even for candidates on transplant lists in close proximity. For example, SRTR's Liver Transplant Waiting List Outcomes Tool shows that for waitlist candidates in New York City with MELD score in the range 25-29, only 15% received a transplant within 90 days, while for similar candidates in Newark, New Jersey, just 15 miles away, 41% received a transplant within 90 days. Since MELD scores directly correlate with the probability of death in the absence of an organ transplant in the next 90 days, different transplant wait times for candidates with the same MELD score across DSAs implies (i) significantly different mortality rates for candidates with the same MELD score in different DSAs, and (ii) a significant variation in the median MELD score at transplant (MMaT).² Indeed variance of MMaT has typically been used by UNOS as a key metric in evaluating the effectiveness of a proposal in mitigating geographic disparity (i.e., a lower value of the variance of MMaT indicates less disparity).

In November 2017, a New York City resident Miriam Holman (a patient who had a rare form of pulmonary hypertension for which there is no medical therapy, and which is rapidly fatal without lung transplantation) filed a lawsuit (hereafter, "lung lawsuit") against HHS.³ Due to the particular lung allocation policy in place at that time, a donor lung could become available across the river in New Jersey (less than four miles away), but because the location of the donor lung is in a different geographical DSA, it would be offered to every candidate waiting for lungs in that New Jersey DSA (even to candidates who are much further away and far less medically critical) before it could be offered to Miriam (Glazier 2018). In July 2018, six liver transplant waiting list patients in New York, California, and Massachusetts filed a lawsuit (hereafter, "liver lawsuit") against HHS.⁴ The liver lawsuit pointed out the wide geographical variability in the median MELD scores in recipients for deceased donor transplants, arguing that place of residence largely determines the chances of one's survival in the policy in place.

To address these issues, in June 2018, the UNOS board (based on the recommendations of a Geography Committee formed in December 2017) adopted the following set of principles to guide future organ transplant policy relating to geographic aspects of organ distribution (that were also identified to be consistent with the final rule).

- 1. Reduce inherent differences in the ratio of donor supply and demand across the country.
- 2. Reduce travel time expected to have a clinically significant effect on ischemic time and organ quality.
 - 3. Increase organ utilization and prevent organ wastage.
 - 4. Increase efficiencies of donation and transplant system resources.

The Geography Committee identified three potential distribution frameworks that fit with these four principles: (1) fixed distance from the donor hospital, (2) mathematically optimized boundaries, and (3) continuous scoring (candidates to be ranked on the offer list on a combination of their clinical characteristics and proximity to a donor).⁵

Following public comment, on December 3, 2018, the UNOS board adopted an Acuity Circles Policy (an implementation of the fixed distance from the donor hospital framework that we discuss in Section 2.2). Although there were legal challenges and political pressures from several quarters to maintain the existing system, the new Acuity Circles policy was implemented on May 14, 2019. However, within a day, on May 15, 2019; a federal court issued an injunction and UNOS was required to revert to the prior system while the legal challenges to the policy were pending. On January 16, 2020, the federal court reversed itself and decided not to keep the injunction in place while the case is pending. Subsequently, the Acuity Circles policy was implemented again on February 4, 2020.

In this paper, we use UNOS's stated principle of reducing inherent differences in the ratio of supply to demand (s/d) as our objective explicitly within a mathematical optimization framework to design heterogeneous sized areas around the donation locations. One approach to reduce inequity is through the central distributive principle, proposed by Rawls (1971): the least well-off group in the society should be made as well off as possible. We use this *maximin* principle to design heterogeneous sized areas that maximize the minimum value of the s/d ratio across all transplant centers (or DSAs). We then apply a secondary optimization to minimize the disparity between the transplant centers (or DSAs) with the highest and lowest s/d ratios.

Our mathematical optimization model can be applied using zip codes or DSAs as the geographical units. When using zip codes as the geographical units the model may be viewed as a heterogeneous circle policy (as compared to a fixed circle policy). When using DSAs as the geographical units the model may be viewed as a type of neighborhood model (Kilambi and Mehrotra 2017), where the neighborhood around a DSA is somewhat circular in shape.

Without the sharing of organs between DSAs, we found that the s/d ratio ranges from 0.31 to 1.98. With 500 nautical miles (NM) fixed circles, the s/d ratio improves and ranges from 0.37 to 0.84. We show that using heterogeneous circles around the donation zip codes, the s/d ratio ranges from 0.55 to 0.60, meaning that there is a much lower disparity in access to organs between the transplant centers. Further, when we examine the s/d ratio disparity for transplant centers that are close to one another (specifically, within 150 NM of each other) the heterogeneous circle policy reduces the s/d ratio disparity to one-fourth compared to the fixed 500 NM circle policy.

We ran simulations with SRTR's Liver Simulated Allocation Model (LSAM, version 2014) using historical patient and organ donor data. The version of the tool available to us was based on DSAs. Hence, we compared our optimized geographical neighborhoods using DSAs. The results show that in comparison to the prior OPTN 11 region policy (in place till February 4, 2020), an allocation policy based on our optimized heterogeneous circular neighborhoods (around DSAs), with a maximum radius of 500 NM and full regional sharing of all organs with MELD \geq 15, drastically reduces the variance of MMaT across DSAs (from 13.66 to 2.00) and average annual deaths (from 3745 to 3568), for a modest increase in average travel distance (from 199 NM to 258 NM).

A key policy insight is that the one-size-fits-all framework (i.e., the currently proposed Acuity Circles Policy) approach taken by UNOS does not adequately address the problem of reducing the differences in the ratio of donor supply to demand across the country. Rather, a customized approach that accounts for where organ supply and demand occur, and adjusts radii of the circles, more effectively addresses UNOS stated goal of equalizing s/d ratios. The remainder of the paper is organized as follows. In the next section, we give a brief overview of the liver allocation system in the U.S., and review proposals and related research. Section 3 presents our optimization methodology. Section 4 describes our findings and the projected outcomes. Section 5 summarizes and provides concluding remarks.

2 Liver Allocation Policy and Literature Review

UNOS supervises the transplantation network in the U.S. Its primary responsibilities are to manage the national transplant waiting list, match organs from deceased donors to candidates, establish medical criteria for allocating organs, facilitate organ distribution, frame equitable policies, etc. Some of the main UNOS members are the 142 liver transplant centers and Organ Procurement Organizations (OPOs) in the 58 DSAs. The OPO coordinates the local procurement of deceased donor organs and allocation in a DSA.

Each transplant center evaluates patients and adds candidates to the waitlist. The medical data about the candidates is shared with UNOS. This pooled data of candidates across all transplant

Sequence #	Candidates that are within:	And are:
1	OPO's region	Adult status 1A or pediatric status 1A/1B
2	OPO's region	$MELD/PELD \ge 35$
3	OPO's DSA	$MELD/PELD \ge 15$
4	OPO's region	$MELD/PELD \ge 15$
5	Nation	Adult status 1A or pediatric status 1A/1B
6	Nation	$MELD/PELD \ge 15$
7	OPO's DSA	$MELD/PELD \le 15$
8	OPO's region	$MELD/PELD \le 15$
9	Nation	$MELD/PELD \le 15$

Table 1 Liver Allocation Policy (a.k.a. Share 35) prior to February 4, 2020

hospitals is constantly updated when new candidates get added and existing candidates are either removed or their medical conditions (e.g. MELD scores) are updated. When a deceased donor organ becomes available, the OPO sends medical data about the organ donor to UNOS. Subsequently, the UNOS matching system compares the donor information with the candidate pool to rank the candidates for the organ offer as per the allocation policy. Upon receiving an offer, the transplant surgeon or physician, in consultation with the candidate decides whether to accept the offer.

2.1 Share 35 Policy

In the previous allocation policy (a.k.a. Share 35) in place from June 18, 2013, until February 4, 2020, deceased donor livers were offered hierarchically to candidates, in decreasing order of MELD scores within each hierarchy, according to the priority list in Table 1. First, an organ was offered to Status 1A and 1B candidates in the region, followed by candidates with MELD \geq 35. After that, candidates with MELD score between 15 and 35 in the OPO's DSA were preferred over candidates outside the DSA. Next in the hierarchy were candidates with MELD score between 15 and 35 in the OPO's region, followed by candidates with MELD score between 15 and 35 outside the OPO's region. Finally, candidates with MELD \leq 15 in the DSA were preferred over candidates with MELD \leq 15 outside the DSA but in the region, who in turn were preferred over candidates with MELD \leq 15 outside the region.

Due to differences in demographics, disease incidence, and mortality leading to organ donations among the DSAs, there was a huge disparity in the s/d ratios across the DSAs. Figure 2 shows the wide variability in the s/d ratio (left), and an inverse relationship of this variability with observed MMaT scores (right). The s/d ratios (at DSAs) varied from 0.31 in NYRT (a DSA in New York) to 1.98 in FLWC (a DSA in Florida). This disparity primarily drove the differences in MMaT among the DSAs. In a study by Wey et al. (2018), the s/d ratios in a DSA were found to be associated with MMaT in DSAs (r = -0.56; P < 0.001).

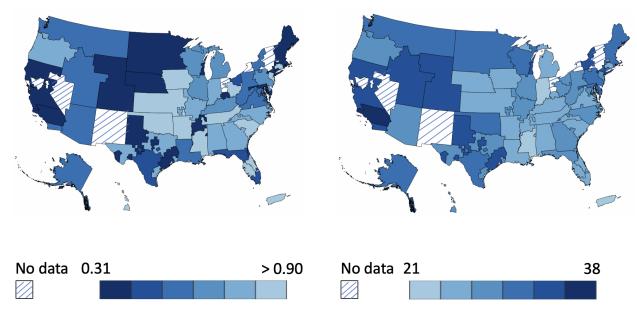


Figure 2 Lower supply to demand (s/d) ratios at a DSA (left) correspond to a higher MMaT at the DSA (right). The time period of analysis is from July, 2013 to June, 2017.

2.2 Current Policy: Acuity Circles

This policy progressively shares organs in circles of radius 150 NM, 250 NM, and 500 NM around the donor hospital, with the following hierarchy: (1) Status 1 candidates within 500 NM; (2) candidates with MELD score \geq 37 within 150 NM, then 250 NM, then 500 NM; (3) candidates with MELD score \geq 33 within 150 NM, then 250 NM, then 500 NM; (4) candidates with MELD score \geq 29 within 150 NM, then 250 NM, then 500 NM; (5) candidates with MELD score \geq 15 within 150 NM, then 250 NM, then 500 NM, and then nationally.⁶ This is a "one-size-fits-all" policy because it does not account for the organ arrival rate, the candidate waiting list, nor the distances of the transplant centers from a donor hospital.

2.3 Related Research

Redistricting is a problem that occurs frequently in multiple domains (e.g., political redistricting, school redistricting, and sales territory assignment) where a finite, denumerable set of non-overlapping geographical units are aggregated into regions/districts subject to some criteria. Hess et al. (1965) and Garfinkel and Nemhauser (1970) introduced the use of optimization techniques for political redistricting. Zoltners and Sinha (1983) discuss an application of redistricting in sales territory assignment, and Caro et al. (2004) discuss school redistricting using integer programming. Much of the redistricting literature is focused on political redistricting (see Gopalan et al. 2013, Kim and Xiao 2017, Ricca et al. 2013). Two important considerations in redistricting problems are the contiguity and compactness of the districts. In this regard, Shirabe (2009) proposed a

flow-based model for contiguity constraints, that has been typically used in subsequent integer programming approaches. However, contiguity constraints make redistricting problems notoriously hard to solve exactly (see Kim and Xiao 2017, Ricca et al. 2013).

Focusing on transplant, and disregarding geographical equity for the moment, Kong et al. (2010) studied the problem of maximizing the efficiency, by maximizing total intraregional transplants through the redesign of the liver allocation regions. They formulate the problem as a set-partitioning problem and use a branch-and-price algorithm to approximate solutions. Stahl et al. (2005) consider geographical equity as measured using intraregional transplant rates in their objective function along with efficiency (measured by total intraregional transplants) but they restrict their regions to contain up to eight DSAs due to computational challenges. Extending their work, Demirci et al. (2012) developed a branch-and-price algorithm to incorporate a larger set of potential regions and explored the efficient frontier in a trade-off between efficiency and geographical equity. Their metric of geographical equity maximizes the minimum in-district viability-adjusted transplant rates per waiting list candidate, which is sensitive to the number of waiting list patients added by the transplant centers and thus problematic. For low MELD patients, the survival benefit of transplantation is minimal (Merion et al. 2005) and the chances of receiving an organ vary across geographies, and consequently, the transplant centers differ in their practices of adding low MELD patients to the waiting list.

Gentry et al. (2015) used optimization to reorganize DSAs into regions/districts to reduce geographical disparity. Their objective was to minimize the sum of absolute differences between the number of deceased-donor livers recovered in each district and the ideal number of livers that would be offered in each district if each liver was given to the medically most urgent candidate in the country. Working closely with the liver committee of UNOS, they proposed eight-district and four-district (reorganized DSA) maps. The proposed maps were under active consideration by UNOS from 2015 to 2017, but ultimately after significant debate and public comment were not adopted.

Kilambi and Mehrotra (2017) introduced the neighborhood framework in organ allocation as a way to provide for broader sharing and improve geographic equity. Each DSA has its own neighborhood, which consists of a unique set of other DSAs (or neighbors) to which it shares its organs. A DSA can be part of multiple neighborhoods, therefore the neighborhoods can be overlapping, which makes the representation of all neighborhoods on a single map difficult. Interconnectivity and overlap among neighborhoods provide resilience to supply and demand uncertainty. The neighborhood framework reduces to redistricting when all the DSAs in a neighborhood have the identical neighborhood. Thus, the redistricting framework can be viewed as a special case of a neighborhood

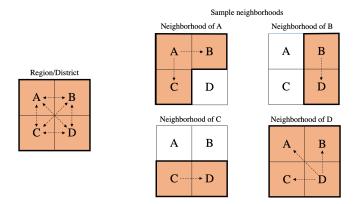


Figure 3 Illustration of difference between Regions/Districts and neighborhoods. Let DSAs A, B, C and D form a region or district. They all share with each other. However, the neighborhood of A consists of DSAs A, B and C, therefore A shares with only A, B and C. Similarly, B shares with B and D, and so on.

framework. Figure 3 illustrates the difference between regions/districts and neighborhood framework. Using the neighborhood framework Kilambi and Mehrotra (2017) developed an optimization model to design DSA neighborhoods to minimize the absolute deviation of the s/d ratios across the neighborhoods from the national average.

Ata et al. (2017) used fluid approximation and game theory to show that multiple listings (a patient lists at more than one transplant center (potentially in other DSA or region) so that he/she can get organ offers from multiple places) can reduce geographical disparity in kidney allocation. However, fewer than 2% of patients (On April 14, 2021, the OPTN website shows 181 out of 11868 candidates on the waiting list are multiple listed) waiting for liver transplant multiple list. Bertsimas et al. (2020) suggest the use of tradeoff curves for the assessment of the three organ distribution frameworks identified by the Geography Committee. Running a large number of simulations for the three distribution frameworks, they plot tradeoff curves of efficiency (measured as average travel distance) versus fairness (measured as deaths or variance of MMaT). For a given value of the efficiency metric, the tradeoff curve then identifies the policy with the greatest fairness. However, they did not consider the neighborhood and heterogeneous circles distribution frameworks in their study.

There are two methods of organ donations: (1) living donation, and (2) deceased donation. Alagoz et al. (2004) study the optimal timing of living-donor liver transplantation when the patient is either ineligible or has decided not to receive organs from deceased donors. They ignore the risk to living donors in their model. Ergin et al. (2020) model liver exchange as a market-design problem where they account for risk to donors and compatibility issues. Using data on South Korean population, they show that their proposed mechanism can increase the number of living-donor transplants by 30%. However, deceased donation has been contributing to greater than 95%

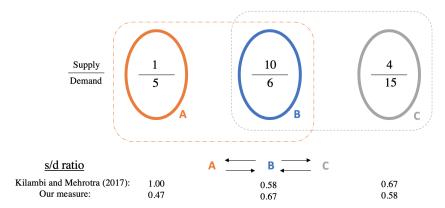


Figure 4 Comparing our s/d ratio measure with that of Kilambi and Mehrotra (2017). Their measure artificially inflates the s/d ratio.

of liver donations in the last 15 years in the U.S. Unlike living donation that can be arranged privately between a patient-donor pair, deceased donor organs are considered national resources (whose allocation is determined by government policy). We focus on deceased donation in this study, and the parameters used in our model and their policy implications are likely to remain unaffected with recent promising developments in living donation.

3 Model Formulation

Consistent with UNOS's stated principles, our approach is to design an organ distribution policy that equalizes s/d ratios across the transplant centers and thus mitigates geographical disparities. We start by aggregating the historical supply and demand of organs by geographical location for the period of study. We assume that the distribution of organ quality (Appendix A compares transplant organ quality on a four year data set used in our study and finds that there are no significant differences in the distribution of organ quality at recovery across the different regions) and patient's health characteristics are similar across the donor hospitals and transplant centers. While there are certainly differences currently in the patient health characteristics from state to state (e.g., presently California has a higher proportion of high MELD candidates than Tennessee), this is largely a function of accumulated disparity over the years; in steady state, the distribution of MELD scores should be similar.

We formulate an Integer Programming model (IP) that uses a neighborhood framework. Each supply location (DSA, zip code, or donor hospital) is assigned a unique set of demand locations (DSA or transplant center), which is referred to as its neighborhood. In a setting where geographical units of supply and demand are DSAs, a neighborhood of a DSA consists of other DSAs (including itself) to which it shares its organs.

Kilambi and Mehrotra (2017) pioneered the idea of defining neighborhoods for DSAs. However, their definition of supply to demand ratio at a DSA is somewhat problematic. They model the

Notation	Description
$i \in \mathcal{I} = \{1,, N_{sup}\}$ $j \in \mathcal{J} = \{1,, N_{dem}\}$	Supply locations
$j \in \mathcal{J} = \{1,, N_{dem}\}$	Demand locations
Parameters:	
s_i	Number of livers from deceased donors recovered (or supply) at i
d_{j}	Number of incident waiting list additions (or demand) at j
$ au_{ij}^-$	Distance between location i and j
$ au_{max}$	Maximum permissible distance from a supply location to a demand
	location
c_{j}	Number of transplant centers in demand location j
$c_i^{(r)}$	Number of transplant centers that are $\leq r$ distance units away from
c_{min}	supply location i Minimum number of transplant centers a supply location must
	share its organs with
$\lambda_{{\lceil S-1 ceil}}^*$	Minimum s/d ratio value to be used in Stage 2 optimization
$\lambda_{[S-1]\atop s_{ij}^{(r)}}^*$	Apportioned share of organs from i to j when the farthest demand
ij	location in $i's$ neighborhood is r units away
Decision variables:	location in v o holghoothood is v ainto away
x_{ij} (General model)	1 if i shares its organs with j , and 0 otherwise
x_{ir} (Set-partitioning model)	1 if the farthest member in the neighborhood of i is r units away
and (200 per oronomias inforcer)	from i , and 0 otherwise
λ	Minimum s/d ratio for an allocation
β	Maximum s/d ratio for an allocation
<i>P</i>	Wide find the first of the control o

Table 2 Model Notation

s/d ratio of a DSA as the ratio of total supply to total demand in the DSA's neighborhood. In other words, they treat all the DSAs in that neighborhood as a single unit, but a DSA can also be part of another neighborhood, which results in the artificial inflation of the s/d ratio. To illustrate, consider three DSAs A (Supply: 1, Demand: 5), B (Supply: 10, Demand: 6), and C (Supply: 4, Demand: 15) as shown in Figure 4. A shares with B and receives from B, B shares with and receives from both A and C, and C shares with B and receives from B. The neighborhood of A consists of A and B, neighborhood of B consists of A, B and C, and neighborhood of C consists of B and C. Kilambi and Mehrotra (2017) compute the s/d ratios of A, B and C as 1.00 (11/11), 0.58 (15/26), and 0.67 (14/21), respectively. But, in aggregate the s/d ratio for this three region system is only 0.58! Further, their objective function is to minimize the absolute deviation of the s/d ratios from a target value (the national average), which effectively treats deviations below the average identically to deviations above the average. Unfortunately, locations with deviations below the average (i.e., lower s/d ratios and higher MMaT scores) have poorer outcomes (greater chances of dying while waiting for a transplant) than locations with deviations above the average. Thus, in a setting where the desire is to minimize the disparities it does not seem appropriate to treat these two deviations identically. By maximizing the worst s/d ratio, our primary focus is on minimizing the deviation below the national average. Finally, we note that our model does not require symmetric organ sharing (which they enforce) giving more flexibility in optimization.

3.1 Supply-Demand Ratio Calculation

First, we define our s/d ratio measure. Recall that we assumed the MELD scores of candidates across the geographies are independent and identically distributed (i.i.d), and when an organ is recovered, all locations in the neighborhood are treated alike. For a given demand location j that is in the neighborhood of supply location i, we model the expected supply received (by j) from i to be proportional to j's demand over the total demand competing for i's supply. Using this expected allocation of supply in the example in Figure 4, we find that 5/11 units of the supply from A are allotted to A, and 6/11 units of the supply from A are allotted to B. Similarly, $(5/26)\times10$, $(6/26)\times10$, and $(15/26)\times10$ units of the supply from B are allotted to A, B, and C, respectively. Finally, $(6/21)\times4$ units of the supply from C are allotted to B, and $(15/21)\times4$ units of the supply from C are allotted to each location by its demand we find s/d ratios of 0.47, 0.67 and 0.58 for A, B, and C, respectively, with our measure. Using the notations described in Table 2 we formally calculate,

Expected supply from
$$i$$
 to $j = \frac{d_j}{\sum_{k=1}^{N_{dem}} d_k x_{ik}} s_i x_{ij}$

To determine overall supply to demand ratio, we first sum the expected supply over all supply locations and then divide by j's demand, d_j giving:

s/d ratio at
$$j = \sum_{i=1}^{N_{sup}} \frac{1}{\sum_{k=1}^{N_{dem}} d_k x_{ik}} s_i x_{ij}$$

We note that the way we calculate the expected s/d ratio does not account for organs that a DSA may receive only due to national sharing. However, these organs are generally a very small fraction (less than 4% in a four year data set used in our study) and should not significantly impact the s/d ratios realized in practice.

3.2 General Model

We now describe our model, which solves the problem in two stages. In Stage 1, we apply the maximin equity principle to maximize the performance of the worst demand location (i.e., we maximize the value of the lowest s/d ratio across all demand locations). In Stage 2, we reduce the disparity among the different demand locations. To do this, we minimize the disparity between the best and worst demand locations, while ensuring the s/d ratio of the worst demand location remains at the optimum value obtained from the Stage 1 optimization. We now present Mixed-Integer Linear Programs (MIPs) for the different stages.

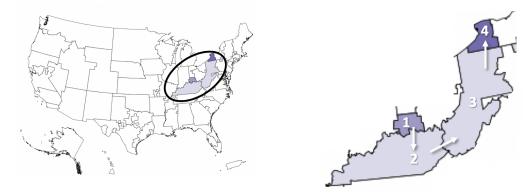


Figure 5 Illustration of sharing and receiving contiguity. If $x_{14} = 1$, with sharing contiguity, $x_{12} = x_{13} = 1$; and with receiving contiguity $x_{24} = x_{34} = 1$.

3.2.1 Stage 1 Formulation:

In Stage 1, we seek to maximize the s/d ratio of the worst demand location.

[S-1] Maximize
$$\lambda$$
 (1)

subject to:
$$\lambda \leq \sum_{i=1}^{N_{sup}} \frac{1}{\sum_{k=1}^{N_{dem}} d_k x_{ik}} s_i x_{ij} \quad \forall j \in \mathcal{J}$$

$$x_{ij} \tau_{ij} \leq \tau_{max} \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$

$$(3)$$

$$x_{ij} \tau_{ij} \le \tau_{max}$$
 $\forall i \in \mathcal{I}, j \in \mathcal{J}$ (3)

$$x_{ij} = 1$$
 $\forall i = j, i \in \mathcal{I}, j \in \mathcal{J}$ (4)

$$\sum_{j=1}^{N_{dem}} c_j x_{ij} \ge c_{min} \qquad \forall i \in \mathcal{I}$$
 (5)

$$x_{ij} \in \{0,1\}$$
 $\forall i \in \mathcal{I}, j \in \mathcal{J}$ (7)

Constraint (2) models λ as the lower bound of the s/d ratios across all the demand locations, and the objective is to maximize this lower bound. Constraint (3) limits the size of the neighborhood (by limiting how far an organ can be transported for transplant), constraint (4) implies that if a supply and demand location coincide (e.g., a DSA or zip code that has both a donor hospital and a transplant center) it must share with itself, and constraint (5) ensures that there are at least c_{min} transplant centers in a neighborhood. We also include contiguity constraints to ensure that the designed neighborhoods are contiguous and somewhat compact in shape. This is enforced by an adjacency matrix that describes locations that are geographically adjacent to each other, and two types of contiguity constraints. Sharing contiguity ensures that if location r supplies organs to location t that is not adjacent to it, then all locations between r and t also receive organs from location r. Receiving contiguity ensures that if location r supplies organs to location t that is not adjacent to it, then all locations between r and t also supply organs to location t. Figure 5 illustrates receiving and sharing contiguity, ensuring that if location 1 shares its organs with location 4, locations 2 and location 3 also share their organs with location 4; and locations 2 and 3 also receive organs from location 1. Appendix B describes flow-based mathematical constraints applying Shirabe (2009)'s approach, that can be used to enforce sharing and receiving contiguity with any geographical shapes, as well as a linearization of constraint (2) in the non-linear integer programming model [S-1].

3.2.2 Stage 2 Formulation:

In Stage 2, we minimize the maximum absolute difference of s/d ratios among demand locations. This is achieved by constraining the lowest s/d ratio value to be greater than or equal to the Stage 1 objective $\lambda_{[S-1]}^*$, and minimizing the maximum s/d ratio value across all the demand locations.

[S-2] Minimize
$$\beta$$

subject to: $\beta \ge \sum_{i=1}^{N_{sup}} \frac{1}{\sum_{k=1}^{N_{dem}} d_k x_{ik}} s_i x_{ij} \quad \forall j \in \mathcal{J}$ (8)

$$\lambda \ge \lambda_{[S-1]}^* \tag{9}$$

All constraints from
$$[S-1]$$
 (10)

The optimal values of x_{ij} obtained by optimizing [S-1] followed by [S-2] are used to construct the new optimized geographical scheme.

3.3 Circular Contiguity and a Set-Partitioning Model

One of the chief complaints in the liver and lung lawsuits was that a candidate receiving the transplant organ may be geographically further away from the donated organ than another sicker candidate. In other words, neighborhood boundaries that allow an organ to be transported further away to a less sick candidate than a closer sicker candidate (because the sicker candidate is outside the neighborhood) goes against generally accepted perceptions of fairness. This suggests that we consider (roughly) circular contiguity for neighborhoods. If the radius of a neighborhood is r units around the supply location, then all demand locations that are within r units away be in the neighborhood (Figure 6).

Circular contiguity allows for a more computationally tractable reformulation of the previous model. For a neighborhood of a given radius r, one can easily calculate (a priori) the amount of supply allocated to each demand location in the neighborhood. This enables us to reformulate [S-1] and [S-2] linearly as Set-Partitioning Problems, which also makes them scalable. In the set-partitioning formulation, x_{ir} is a binary decision variable that takes value 1 if the radius of the neighborhood of i is r units (all demand locations $\leq r$ units from i are part of the neighborhood) and is 0 otherwise.

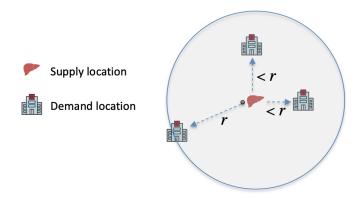


Figure 6 Illustration of circular contiguity: If a neighborhood is r units in radius around the supply location, then all demand locations that are within r units must be in the neighborhood.

Stage 1 Formulation: 3.3.1

[SP-1] Maximize
$$\lambda$$
 (11)

[SP-1] Maximize
$$\lambda$$
 (11)
subject to:
$$\lambda \leq \sum_{i=1}^{N_{sup}} \sum_{r \in \mathcal{R}_i} \frac{x_{ir} s_{ij}^{(r)}}{d_j} \quad \forall j \in \mathcal{J}$$
 (12)

$$\sum_{r \in \mathcal{R}_i} x_{ir} = 1 \qquad \forall i \in \mathcal{I}$$
 (13)

$$\sum_{r \in \mathcal{R}_i} c_i^{(r)} x_{ir} \geq c_{min} \qquad \forall i \in \mathcal{I}$$
 (14)

$$x_{ir} \in \{0,1\} \qquad \forall i \in \mathcal{I}, \forall r \in \mathcal{R}_i$$
 (15)

$$\sum_{r \in \mathcal{R}_i} x_{ir} = 1 \qquad \forall i \in \mathcal{I}$$
 (13)

$$\sum_{r \in \mathcal{R}_i} c_i^{(r)} x_{ir} \ge c_{min} \qquad \forall i \in \mathcal{I}$$
 (14)

$$x_{ir} \in \{0, 1\}$$
 $\forall i \in \mathcal{I}, \forall r \in \mathcal{R}_i$ (15)

For a given radius r, $s_{ij}^{(r)}$ denotes the apportioned share of i's organs that are expected to be offered to location j. In other words, $s_{ij}^{(r)} = \frac{d_j}{\sum_{k:\tau_{ik} \leq r} d_k} s_i$, which can be precomputed for a given radius r. Note that for a given supply location i, we do not need to consider a continuum of possible neighborhood radii. Rather (because this apportionment of organs will only change when a new demand location is added to the neighborhood) we only need to consider a finite set of values of r that correspond to the distance from i to each of the other demand locations that are within τ_{max} . In [SP-1], the set \mathcal{R}_i contains the possible values of r created accordingly. Constraint (12) models λ as the lower bound of the s/d ratios across all the demand locations, and the objective is to maximize this lower bound. Constraint (13) allows one assignment of r to each supply location, and constraint (14) ensures a minimum number of transplant centers in the neighborhood.

Stage 2 Formulation: 3.3.2

Once the optimal solution $\lambda_{[SP-1]}^*$ to [SP-1] is obtained, we can solve [SP-2] to minimize the maximum s/d ratio while ensuring that the minimum s/d ratio remains at least $\lambda_{[SP-1]}^*$.

[SP-2] Minimize
$$\beta$$

subject to:
$$\beta \ge \sum_{i=1}^{N_{sup}} \sum_{r \in \mathcal{R}_i} \frac{x_{ir} s_{ij}^{(r)}}{d_j} \qquad \forall j \in \mathcal{J}$$
 (16)

$$\lambda \ge \lambda_{[SP-1]}^* \tag{17}$$

4 Data and Results

This study used data from SRTR. The SRTR data system includes data on all donors, wait-listed candidates, and transplant recipients in the U.S., submitted by the members of the OPTN. The Health Resources and Services Administration (HRSA), U.S. Department of Health and Human Services provides oversight to the activities of the OPTN and SRTR contractors.

In the data, encompassing the four years starting July 2013 and ending June 2017, the supply or the total number of livers (from deceased donors) donated from all donor hospitals in the U.S. is 26,899. The patient pool is dynamic: new patients enlist, waiting candidates die or become too sick for transplant and are removed, and the MELD scores get updated periodically. We measure demand (44,959) as the total incident⁸ adult patients whose MELD scores became at least 15 during the four years; which gives a national s/d ratio of 0.5983. There are two reasons for excluding low MELD patients from the demand: (1) patients with MELD <15 have no survival benefit from transplantation (Merion et al. 2005), therefore our demand measure is less sensitive to the number of low-MELD score patients added to the waiting list, and (2) transplant centers differ in their practices of listing low MELD score patients across the country (which would create an artificial increase in demand for a transplant center that lists low MELD score patients compared to a transplant center that does not). In practice, the fraction of transplants to low MELD patients is relatively very low—about 1.08% (in the four years encompassing our study), supporting the decision to exclude them.

We apply the set-partitioning optimization model to two versions of the data: a zip-code cluster version where supply locations are zip-code clusters (clustered by the first three-digits and first four-digits) and demand locations are the 142 transplant centers, and a DSA version where supply and demand locations are the DSAs. We restrict r (radius around the supply locations) within the range 150 NM to τ_{max} for every \mathcal{R}_i , constraining the minimum and maximum size of the neighborhoods. We set $c_{min} = 3$, ensuring that at least three transplant centers are present in a neighborhood. We used R 3.5.1 and a commercial solver Gurobi 8.1.1 to solve the set-partitioning optimization models on a 3.2 GHz 6-Core Intel Core i7 iMac with 32 GB RAM.

4.1 Zip-code Cluster Version

The location of zip codes and transplant centers are indicated by their latitude and longitude values. To calculate the distance between a three-digit (four-digit) zip-code cluster and a transplant

center, we first find the centroid of the zip-codes in the cluster having the same first three digits (four digits) and then use the "geosphere" package in R to calculate the shortest distance between two points (centroid of the zip cluster and transplant center) according to the "Vincenty (ellipsoid)" method.

There are a total of 641 three-digit and 1380 four-digit zip-code clusters with supply in our data. We vary τ_{max} from 350 NM to 700 NM in steps of 50 NM. We do not include the zip codes in Hawaii and Puerto Rico in our analysis, since they are more than 1000 miles from transplant centers in mainland U.S. Consistent with the current policy zip codes in Alaska are considered to be situated at the Seattle Tacoma Airport in Washington State. We require that the minimum radius of a neighborhood be 150 NM (to try and keep parity with the radius of the innermost concentric circle in the acuity circles policy). Since a transplant hospital may not necessarily be exactly at 150 NM from a zip-code cluster, this is enforced by ensuring the closest transplant center greater than or equal to 150 NM away is included in the neighborhood, unless it is greater than τ_{max} miles away. Appendix C provides computational details—problem size, running times, cutting planes, simplex iterations etc.—for the set partitioning model on the four-digit zip-clusters.

Table 3 provides a comparison of the s/d ratios. To compare against the fixed radius type of policy in place currently (i.e., acuity circles), we also computed the s/d ratio for homogeneous radii circles by fixing the radius of each zip-code cluster to τ_{max} . Compared to the heterogeneous radius circle policy, the "one-size-fits-all" fixed radius policy does a poor job in equalizing s/d ratios across transplant centers. The heterogeneous circle policy at $\tau_{max} = 500$ NM is able to keep the ratio at transplant centers between 0.55 and 0.60 (compared to the national s/d ratio of 0.5983), while the fixed 500 NM radius circle policy has an s/d ratio variation between 0.37 and 0.84.

We also examine the difference in the s/d ratio of nearby transplant centers (defined as being within 150 NM). Table 3 provides both the maximum and the median values of this difference. As is evident, in the heterogeneous circles policy the value of the s/d ratio at nearby transplant centers is very similar—which can hopefully lead to more equitable transplant outcomes in nearby transplant centers. For most of the transplant centers, the difference in s/d ratio is in the scale of 10^{-4} , as implied by the median values.

As we increase τ_{max} from 350 NM to 700 NM, the minimum s/d ratio increases, and the range of s/d ratio decreases. When $\tau_{max} = 400$ NM the s/d ratio range is already quite narrow at 0.53-0.61, and once $\tau_{max} = 500$ NM, the s/d ratio range stays steady at 0.55-0.60. Figure 7 shows the quartiles of radii when using four-digit zip-code clusters. When τ_{max} is 500 NM, the first, second, and third quartiles of radii are 211, 305, and 415 NM, respectively. Compared to the fixed radii circles policy, the heterogeneous radii circles policy achieves an equalization in the s/d ratio (near the national average) while keeping transport distances lower. This has an added benefit. Since the radii of the

Allocation Policy		d ratio	Maximum (Median) s/d ratio
· ·	Range	Std. deviation	difference within 150 NM from TC
	0.20.1.00	0.123	0.59920 (0.04611)
Fixed radius circles (Three-digit zip) Fixed radius circles (Four-digit zip)	0.39-1.09 0.38-1.09	$0.123 \\ 0.123$	$0.69920 (0.04011) \\ 0.60235 (0.04559)$
Three-digit zip-code cluster model	0.50-1.09 0.51-0.88	0.123 0.098	0.33043 (0.05047)
Four-digit zip-code cluster model	0.51 - 0.88	0.103	0.33813 (0.07085)
Ü .	0.01 0.00	0.100	0.00010 (0.01000)
$\tau_{max} = 400 \text{ NM}$			
Fixed radius circles (Three-digit zip)	0.37 - 0.85	0.112	$0.23255 \; (0.03690)$
Fixed radius circles (Four-digit zip)	0.37 - 0.84	0.112	$0.22818 \ (0.03633)$
Three-digit zip-code cluster model	0.53-0.62	0.033	$0.08571 \ (0.00042)$
Four-digit zip-code cluster model	0.53 - 0.61	0.030	$0.07763 \; (0.00028)$
$\tau_{max} = 450 \text{ NM}$			
Fixed radius circles (Three-digit zip)	0.38 - 0.88	0.124	$0.20629 \ (0.02900)$
Fixed radius circles (Four-digit zip)	0.38 - 0.87	0.124	$0.19770\ (0.02048)$
Three-digit zip-code cluster model	0.54 - 0.61	0.023	0.05277 (0.00108)
Four-digit zip-code cluster model	0.54 - 0.61	0.024	0.06125 (0.00043)
500 NM			
$ \frac{\tau_{max} = 500 \text{ NM}}{\text{Fixed radius circles (Three-digit zip)}} $	0.27.0.94	0.137	0.20941 (0.03632)
Fixed radius circles (Three-digit zip) Fixed radius circles (Four-digit zip)	0.37 - 0.84 $0.37 - 0.84$	$0.137 \\ 0.137$	$0.20941 (0.03032) \\ 0.20851 (0.04489)$
Three-digit zip-code cluster model	0.57 - 0.64 0.55 - 0.60	0.137 0.022	0.20831 (0.04489) 0.04621 (0.00009)
Four-digit zip-code cluster model	0.55 - 0.60	0.022 0.021	0.04921 (0.00003) 0.04922 (0.00025)
-	0.00 0.00	0.021	0.01022 (0.00020)
$ \underline{\tau_{max}} = 550 \text{ NM} $			
Fixed radius circles (Three-digit zip)	0.37 - 0.91	0.145	$0.17331 \ (0.03808)$
Fixed radius circles (Four-digit zip)	0.36 - 0.91	0.146	$0.17213 \; (0.03882)$
Three-digit zip-code cluster model	0.55 - 0.60	0.020	$0.05070 \; (0.00029)$
Four-digit zip-code cluster model	0.55 - 0.60	0.019	$0.04387 \; (0.00025)$
$\tau_{max} = 600 \text{ NM}$			
Fixed radius circles (Three-digit zip)	0.34-0.97	0.152	0.17767 (0.04866)
Fixed radius circles (Four-digit zip)	0.34-0.96	$0.152 \\ 0.152$	0.17819 (0.04473)
Three-digit zip-code cluster model	0.55-0.60	0.018	0.05407 (0.00113)
Four-digit zip-code cluster model	0.55-0.60	0.018	0.03613 (0.00015)
_			,
$ \underline{\tau_{max} = 650 \text{ NM}} $	0.00.004	0.450	0.40=10.(0.00110)
Fixed radius circles (Three-digit zip)	0.33-0.94	0.152	$0.16743 \ (0.02449)$
Fixed radius circles (Four-digit zip)	0.33-0.93	0.152	$0.17091 \ (0.02457)$
Three-digit zip-code cluster model	0.55-0.60	0.017	$0.05049 \ (0.00016)$
Four-digit zip-code cluster model	0.55 - 0.60	0.018	$0.03336 \ (0.00012)$
$\tau_{max} = 700 \text{ NM}$			
Fixed radius circles (Three-digit zip)	0.32 - 0.94	0.145	$0.17881 \ (0.04773)$
Fixed radius circles (Four-digit zip)	0.32 - 0.94	0.145	0.18275 (0.04654)
Three-digit zip-code cluster model	0.55 - 0.60	0.016	$0.05100 \ (0.00010)$
Four-digit zip-code cluster model	0.55 - 0.60	0.017	$0.03402\ (0.00007)$

Table 3 Comparison of s/d ratios between fixed and heterogeneous circles (supply and demand locations are zip-code clusters and transplant centers (TCs), respectively).

circles are smaller, each donor zip-code cluster on average has 24 (median 20) transplant centers, as compared to the fixed radii circles that have on average 39 transplant centers (median 43). The logistics of a donor hospital (zip-code cluster) coordinating with a smaller number of transplant centers can be much simpler. One may wonder whether fixed population circles (i.e., the radius of the circle around each transplant center is set so that they all cover the same number of people) would reduce disparity. Using the s/d metric defined and introduced in this paper, Haugen et al.

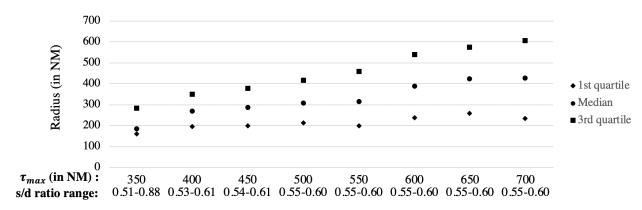


Figure 7 Quartiles of radii in the four-digit zip-code cluster models.

(2019) analyze the disparity in s/d ratios across the transplant centers with fixed population circles. They find circles covering a population of 12 million individuals provides s/d values ranging from 0.27 to 2.14. Increasing the size of the circles to cover 50 million individuals decreases the s/d variation to 0.43–1.01.

To check if our solution is robust to variations in supply and demand across time, using the optimal radii obtained using the four years data, we recalculate the s/d ratio range skipping one year's (supply and demand) data at a time. We find that, on average, the minimum (absolute) s/d ratio changes by 0.016 points, and maximum (absolute) changes by 0.018 points (based on $\tau_{max} = 500, 550, 600, 650$ and 700 NM); indicating the results are fairly robust to variations in the data.

Since the current implementation of LSAM does not support schemes based on zip-code clusters, we could not evaluate our zip code based allocation policy via the LSAM simulation model. Instead, we use the results of the DSA version described in the next section and run the LSAM simulation on the neighborhoods it generates to evaluate the effectiveness of our allocation policy in reducing geographical disparity.

4.2 DSA Version

Using DSAs as the geographical unit preserves the existing important relationships between donor hospitals and the OPO in each DSA. If indeed, the court rules in a manner that reinstates DSAs as a geographical unit, then our method shows how they could share organs to achieve equitable outcomes with regards to the s/d ratio.

The distance between any two DSAs i and j, τ_{ij} , is calculated as the mean of the transplant-volume-weighted distance between donor hospitals in DSA i and transplant center in DSA j, and the reverse. Since six DSAs do not have a transplant hospital, there are 58 DSAs with supply and 52 DSAs with demand. Consistent with Gentry et al. (2015) and Kilambi and Mehrotra (2017), we

Allocation Policy	S/	s/d ratio		Max. (Median) s/d ratio
Anocation 1 oney	Range	Std. deviation	$(\inf^{\tau_{max}} NM)$	difference among adjacent DSAs
OPTN 11 regions	0.42 - 0.76	0.109	843, 401	0.228 (0.117)
Gentry et al. (2015)	0.52 - 0.69	0.054	975, 569	$0.120 \ (0.036)$
Kilambi and Mehrotra (2017)	0.35 - 0.99	0.157	1380, 666	0.615~(0.246)
[SP-2], $\tau_{max} = 500 \text{ NM}$	0.50 - 0.65	0.054	500, 349	0.151 (0.086)
[SP-2], $\tau_{max} = 600 \text{ NM}$	0.52 - 0.65	0.051	600, 409	$0.132\ (0.077)$
[SP-2], $\tau_{max} = 600 \text{ NM}$ [SP-2], $\tau_{max} = 700 \text{ NM}$	0.53-0.63	0.033	700, 422	0.096 (0.036)

Table 4 Comparison of s/d ratios among different allocation policies in the DSA version (supply and demand locations are 58 DSAs and 52 DSAs, respectively). τ_{max} and $\bar{\tau}$ represent the maximum and average distance, respectively, of the farthest DSA in a neighborhood/region/district in each allocation policy.

allow (as exceptions to τ_{max}) the DSAs located in Hawaii, and Puerto Rico, to share and receive organs from other DSAs located in California and Oregon, and Florida, respectively.

Table 4 summarizes the results for τ_{max} set to 500 NM, 600 NM and 700 NM, and compares it with the prior 11 region system and other proposed geographical allocation policies. As is evident, our model produces a neighborhood that results in the narrowest range of s/d ratios across DSAs: 0.15 when $\tau_{max} = 500$ NM, 0.13 when $\tau_{max} = 600$ NM and 0.10 when $\tau_{max} = 700$ NM, as compared to 0.34 (OPTN 11 regions), 0.17 (Gentry et al. 2015, 8 districts) and 0.64 (Kilambi and Mehrotra 2017). Our model also produces relatively more uniform and smaller size neighborhoods. It does not contain any unusually large neighborhoods (as evidenced by the value of τ_{max}). Further, the average distance of the farthest DSAs in the neighborhoods ($\bar{\tau}$) is much smaller than Gentry et al. (2015), Kilambi and Mehrotra (2017) and is comparable with OPTN 11 regions. The maximum s/d ratio difference among adjacent DSAs is also reduced significantly. For example, with $\tau_{max} = 700$ NM, the maximum difference of s/d ratio among adjacent DSAs is 0.096, much smaller compared to OPTN 11 regions (0.228).

Table 5 presents the s/d ratios for each DSA in the different proposals. This allows a deeper examination of how each DSA is affected by the proposed reallocations. The maximum and minimum s/d ratio values in every proposal are highlighted in bold. Appendix D describes the DSA neighborhoods obtained by our models for $\tau_{max} = 500$, 600, and 700 NM, respectively.

The computational benefit of [SP-1] over [S-1] is easily seen on the DSA version. For example, when $\tau_{max} = 500$ NM, the size of [S-1] using only sharing contiguity was 16286 rows and 18883 columns, and the MIP gap (MIP gap = $\frac{|Objective\ bound-Objective\ value|}{|Objective\ value|}$) was 1.19% after 2 hours of running time. Whereas, the size of [SP-1] was 110 rows and 742 columns, and took only 0.66 seconds to reach optimality.

4.2.1 Liver Simulated Allocation Model (LSAM) Results

Next, we wanted to see how the proposed (DSA based) allocation policies perform on metrics that policy makers have traditionally examined to evaluate policies, like the variance of MMaT

				s/d ratio				
\mathbf{DSA}	Local, or	OPTN 11 regions	Gentry et al. (2015)	Kilambi and Mehrotra (2017)			P-2]	
	no sharing	$\tau_{max} = 843$ NM	$\tau_{max} = 975 \text{ NM}$	$\tau_{max} = 1380 \text{ NM}$	τ_{max} :	500 NM	600 NM	700 NM
ALOB	0.72	0.76	0.61	0.56		0.62	0.65	0.63
AROR	0.97	0.76	0.61	0.96		0.65	0.52	0.58
AZOB	0.55	0.52	0.54	0.88		0.53	0.53	0.59
CADN	0.38	0.52	0.52	0.45		0.51	0.52	0.53
CAOP CASD	$0.39 \\ 0.55$	$0.52 \\ 0.52$	$\begin{array}{c} 0.52 \\ 0.52 \end{array}$	0.5 0.35		0.54 0.5	0.53	0.53 0.59
CORS	$0.35 \\ 0.37$	0.64	0.52	0.35		0.51	$0.53 \\ 0.54$	0.59 0.53
CTOP	0.95	$0.04 \\ 0.42$	0.57	0.77		0.56	0.62	0.59
DCTC	0.58	0.57	0.57	0.46		0.64	0.64	0.63
FLFH	1.3	0.76	0.61	0.65		0.54	0.52	0.62
FLMP	0.5	0.76	0.61	0.65		0.61	0.65	0.61
FLUF	0.47	0.76	0.61	0.81		0.65	0.58	0.61
FLWC	1.98	0.76	0.61	0.65		0.64	0.52	0.62
GALL	0.72	0.76	0.57	0.99		0.65	0.65	0.62
HIOP	0.97	0.66	0.52	0.37		0.63	0.64	0.54
IAOP	1.23	0.64	0.64	0.58		0.62	0.64	0.61
ILIP	0.69	0.55	0.69	0.62		0.65	0.62	0.62
INOP KYDA	0.78	0.66 0.76	0.69	0.67 0.69		0.62 0.65	0.63	$0.59 \\ 0.6$
LAOP	$0.66 \\ 0.55$	0.76	0.69 0.61	0.69		0.63	0.64 0.65	0.63
MAOB	0.39	$0.76 \\ 0.42$	0.61	0.4		0.54	0.61	0.56
MDPC	0.34	0.57	0.57	0.4		0.64	0.65	0.63
MIOP	0.68	0.66	0.69	0.49		0.54	0.64	0.63
MNOP	0.4	0.55	0.64	0.53		0.51	0.56	0.56
MOMA	0.71	0.64	0.61	0.73		0.65	0.63	0.63
MSOP	1.49	0.76	0.61	0.56		0.58	0.55	0.63
MWOB	1.04	0.64	0.64	0.7		0.5	0.52	0.56
NCCM	0.73	0.76	0.57	0.44		0.65	0.64	0.62
NCNC	0.77	0.76	0.57	0.63		0.65	0.6	0.62
NEOR	0.41	0.64	0.64	0.44		0.51	0.54	0.63
$ \begin{array}{c} \text{NJTO} \\ \text{NYFL} \end{array} $	$\frac{1.19}{0.56}$	$0.57 \\ 0.42$	0.57 0.69	$0.47 \\ 0.59$		0.65 0.53	0.65 0.53	0.63 0.62
NYRT	0.30	$0.42 \\ 0.42$	0.69	$0.39 \\ 0.47$		0.65	0.65	0.62
OHLB	0.47	0.66	0.69	0.67		0.65	0.65	0.62
ÖHLP	0.9	0.66	0.69	0.83		0.65	0.61	0.62
OHOV	0.33	0.66	0.69	0.51		0.65	0.65	0.61
OKOP	0.91	0.53	0.64	0.81		0.58	0.52	0.63
ORUO	0.71	0.66	0.52	0.62		0.62	0.65	0.58
PADV	0.62	0.57	0.57	0.6		0.62	0.65	0.63
PATF	0.58	0.57	0.69	0.83		0.64	0.59	0.63
PRLL	1.69	0.76	0.57	0.56		0.54	0.6	0.53
$\frac{\text{SCOP}}{\text{TNDS}}$	$\frac{1.02}{1.17}$	$0.76 \\ 0.76$	0.57 0.69	$0.38 \\ 0.77$		$\substack{0.65\\0.65}$	$0.62 \\ 0.64$	0.62 0.63
TNMS	0.36	0.76	0.69	0.77		0.65	$0.64 \\ 0.64$	0.62
TXGC	0.36	0.78	0.61	$0.83 \\ 0.52$		0.64	0.58	0.62
TXSA	0.5	0.53	0.61	$0.32 \\ 0.44$		0.53	0.52	0.53
TXSB	0.77	0.53	0.61	0.5		0.64	0.55	0.62
UTOP	0.53	0.52	0.54	0.47		0.54	0.55	0.56
VATB	0.6	0.76	0.57	0.85		0.63	0.65	0.62
WALC	0.6	0.66	0.52	0.6		0.62	0.62	0.55
WIDN	0.4	0.55	0.69	0.5		0.5	0.54	0.62
WIUW	0.61	0.55	0.69	0.72		0.62	0.63	0.63

Table 5 Comparison of s/d ratios among different DSA based allocation policies (supply and demand locations are 58 DSAs and 52 DSAs, respectively).

across geographies, distance traveled, and the number of deaths. To that end, we use LSAM to simulate our neighborhoods [SP-2], OPTN 11 regions, Gentry et al. (2015) (8 districts), and Kilambi and Mehrotra (2017)'s neighborhoods. There are two main inputs to LSAM: (1) patient and organ arrival processes, and (2) allocation policy that includes geographical schemes and offer prioritization rules.

LSAM uses historical data of donors and patients to simulate waiting list patient's health state transitions, organ acceptance behavior, and post-transplant survival outcomes. When an organ becomes available, candidates on the waiting list are prioritized for the organ offer as per the allocation policy. When a candidate receives a transplant, the simulation determines the survival time of the transplanted organ, and uses this information to determine when in the future the candidate may die or relist. Using LSAM in its current form does have some limitations. It uses a probability acceptance function that is built on past data where distance is more strongly correlated with acceptance of an organ because of the lack of broader sharing. It also does not account for organ availability in determining organ acceptance. These limitations may underestimate the effects of broader sharing and the equalization of the s/d ratios. Despite these limitations, it is instructive

	Avg. (Quartiles)	Waitlist	Total	I	Across DSAs
Allocation Policy	travel distance	deaths	deaths	Variance	Std. deviation of avg
	(in NM)	(annual)	(annual)	of MMaT	travel distance (NM)
OPTN 11 regions	258 (75, 194, 347)	1411.6	3658.8	7.26	109
Gentry et al. (2015)	309 (101, 226, 429)	1376.1	3600.0	5.22	124
Kilambi and Mehrotra (2017)	305 (124, 240, 395)	1348.2	3555.4	2.68	142
[SP-2], $\tau_{max} = 500 \text{ NM}$	258 (112, 220, 341)	1356.4	3567.7	2.00	56
[SP-2], $\tau_{max} = 600 \text{ NM}$	283 (125, 251, 384)	1343.6	3551.4	1.98	55
$ \begin{array}{l} [\text{SP-2}], \; \tau_{max} = 600 \; \text{NM} \\ [\text{SP-2}], \; \tau_{max} = 700 \; \text{NM} \end{array} $	293 (125, 250, 399)	1343.4	3544.6	1.61	64

Table 6 Comparison of LSAM simulation results for DSA-based allocation policies under Enhanced Share 15.

to use LSAM as a first step in evaluating the potential benefit of the heterogeneous radii circles policy.

In the simulation study (to model broader sharing within a circle) we allow for full sharing of organs to Status 1A/1B and $MELD \geq 15$ candidates in the neighborhood or region/district in which the organ is recovered in the first level of allocation. In the next allocation level, the organ is offered nationally to Status 1A/1B, then nationally to candidates with $MELD \geq 15$. Next, it is offered to candidates with MELD < 15 locally (DSA in which the organ is recovered), then in the neighborhood or region/district, and then nationally, before being discarded after 100 offers. The above policy, that we refer to as "Enhanced Share 15", skips sequences # 2 and 3 of the Share 35 policy described in Table 1. For benchmarking, we also compared using the prior Share 35 policy. We simulated the different DSA-based geographical allocation policies using the organ and patient arrival data, consisting of 3 years (July 2013 to June 2016). We ran the simulation ten times (the maximum allowed by LSAM) by resampling the input files.

Table 6 compares the simulation results under Enhanced Share 15. The average number of annual waitlist deaths and total deaths is smallest for [SP-2], $\tau_{max} = 700$ NM, with a projected savings of 114 lives annually compared to OPTN 11 regions. The average travel distance, although slightly higher in our allocation policy compared to OPTN 11 regions, is smaller than the other policies. To measure the differences between DSAs, we consider the variance of MMaT and the standard deviation of average organ travel distance, across DSAs. The variance of MMaT across the DSAs is smallest in our allocation policies (2.00 when $\tau_{max} = 500$ NM, 1.98 when $\tau_{max} = 600$ NM and 1.61 when $\tau_{max} = 700$ NM), compared to 7.26, 5.22 and 2.68 in OPTN 11 regions, Gentry et al. (2015) and Kilambi and Mehrotra (2017) respectively. Since the different proposals vary in their efficiency (travel distance) and fairness (MMaT) metrics, it is instructive to compare the fairness of proposals with similar levels of efficiency. To that end comparing [SP-2], $\tau_{max} = 500$ NM against OPTN 11 regions, shows a significant reduction in both total deaths and variance of MMaT. Similarly, comparing [SP-2], $\tau_{max} = 700$ NM against Gentry et al. (2015) and Kilambi and Mehrotra (2017), shows a significant reduction in the variance of MMaT. Overall, we see that greater fairness can be achieved by DSA-based geographical allocation policies that equalize s/d ratios. The standard

	Avg. (Quartiles)	Waitlist	Total		Across DSAs
Allocation Policy	travel distance	deaths	deaths	Variance	Std. deviation of avg.
	(in NM)	(annual)	(annual)	of MMaT	travel distance (NM)
OPTN 11 regions	199 (20, 105, 258)	1455.5	3744.9	13.66	88
Gentry et al. (2015)	231 (25, 130, 314)	1419.5	3696.4	10.49	102
Kilambi and Mehrotra (2017)	230 (32, 150, 309)	1389.0	3656.3	11.87	104
[SP-2], $\tau_{max} = 500 \text{ NM}$	203 (29, 142, 291)	1399.9	3664.8	10.30	57
[SP-2], $\tau_{max} = 600 \text{ NM}$	221 (32, 157, 326)	1384.8	3645.2	8.80	57
[SP-2], $\tau_{max} = 600 \text{ NM}$ [SP-2], $\tau_{max} = 700 \text{ NM}$	233 (36, 162, 344)	1397.3	3636.4	10.04	64

Table 7 Comparison of LSAM simulation results for DSA-based allocation policies under Share 35.

deviation of average travel distance across the DSAs (Hawaii and Puerto Rico are excluded from our distance analysis) in our allocation policies is less than half that of the others. This indicates that there is less disparity in travel distance between DSAs because our neighborhoods have relatively similar sizes.

Table 7 compares the LSAM simulation results under Share 35. We note that our neighborhoods are optimized under the assumption of full sharing, which is closer to Enhanced Share 15 than Share 35, and thus the full benefits of improved MMaT is less likely to be seen. Since there is less sharing under Share 35 (organs offers are restricted to within DSA patients ($15 \le \text{MELD} < 35$) before being offered broadly (neighborhood or region/district and nationally)) the average travel distance significantly decreased, and the number of the waitlist and total deaths increased for all policies. Even so, comparing [SP-2], $\tau_{max} = 500 \text{ NM}$ against OPTN 11 regions, shows a significant reduction in both total deaths and variance of MMaT. Similarly, comparing [SP-2], $\tau_{max} = 600 \text{ NM}$ against Gentry et al. (2015) and Kilambi and Mehrotra (2017), shows a significant reduction in the variance of MMaT. Like Enhanced Share 15, we observe again the standard deviation of average travel distances (across DSAs) is much lower for our allocation policies.

Ultimately comparing our allocation policy $\tau_{max} = 500$ NM under Enhanced Share 15 against OPTN 11 regions under Share 35 shows that a drastic reduction in the variance of MMaT across DSAs (from 13.66 to 2.00) and deaths (from 3745 to 3568) can be achieved, for a modest increase in the average travel distance (from 199 NM to 258 NM).

5 Conclusions

We use the Rawlsian maximin principle to minimize the variability in deceased donor liver access across geographies. In contrast to the current fixed radius policy, we propose heterogeneous radii circles. The benefit of heterogeneous radii circles is that they account for where organ supply and demand occur, and adjust radii of the circles so that each transplant center's s/d ratio can be close to the national average. Moreover, equalizing the s/d ratios at the transplant centers is achieved without a significant increase in anticipated travel distance. In fact, the median radius is approximately 305 NM. In other words, the optimization model only increases the radii of donor circles when necessary.

By using DSA as the geographical unit, we demonstrate that low geographical variation in s/d ratio can be achieved while maintaining DSA boundaries by judiciously creating neighborhoods for each DSA. An LSAM evaluation of our DSA neighborhoods predicts significant reduction in the number of deaths, and the overall variation in MMaT and average travel distance across DSAs.

As noted earlier, there are limitations of our analysis since LSAM's organ acceptance function may not accurately reflect the change in candidate/transplant center behaviors when organ accessibility and availability changes. For instance candidates at organ rich locations might behave more selectively in accepting organs, than at locations with low s/d ratios. Additionally, with broader sharing the correlation between organ acceptance probability and distance may decrease within the radius of a circle (or neighborhood). Studying changes in organ acceptance behavior with organ accessibility can be an interesting future research work.

In terms of a logistical implementation of the heterogeneous circles policy, we have a few suggestions. First, we believe the circles should be defined around the donor location rather than the transplant location (note that in a fixed radius policy there is no difference between circles defined around donor locations and transplant locations, but with heterogeneous circles there is a difference), else the issues raised in the lawsuits (i.e., organs being offered to a less sicker candidate who is further away) would not be addressed. Second, we expect small variations in supply and demand over time. Hence, we suggest that the optimization model be run occasionally to account for changes in demographics.

Our approach can be viewed as a combination of the fixed distance from a donor hospital and a mathematical optimized boundaries framework identified by the Geography Committee. There is considerable debate in the transplant community about using continuous scoring (the third distribution framework identified by the Geography Committee). Rather than a one-size-fits-all framework for continuous scoring, that we do not believe will address geographical inequities adequately, we would recommend a mathematically optimized continuous scoring function that accounts for regional differences in supply and demand.

Clearly the optimization concepts applied to mitigate geographical disparities in the liver transplantation setting could also be applied to other organs. We hope this research will spur similar work in the other organ transplantation settings, and thus reduce/mitigate the geographical disparities that are inherent to all of these systems!

Endnotes

- 1. https://www.srtr.org/reports-tools/waiting-list-calculator/
- 2. This is seen for example in 2016 data (Kim et al. 2018) where the highest MMaT is 39 and the lowest MMaT is 20.

- 3. Miriam Holman v. HHS, (S.D.N.Y 17-CV-09041)
- 4. Cruz et al. v. HHS, (S.D.N.Y 18-CV-06371)
- 5. https://optn.transplant.hrsa.gov/media/2565/geography_publiccomment_201808.pdf
- 6. https://unos.org/policy/liver-distribution/
- 7. Deceased-donor livers vary in quality, and marginal livers are more likely to be used and less likely to be discarded when more competition exists among transplant centers (Halldorson et al. 2013, Garonzik-Wang et al. 2013). Thus UNOS requires that a minimum number of transplant centers be in contention for organs from a supply location.
- 8. We consider incident patients so that the model parameters are not biased because of accumulated disparity and thus are exogenous to the geographical scheme.
- 9. For every DSA with demand, there are close to 3 (142/52=2.73) transplant centers.

References

- Alagoz, O., L. M. Maillart, A. J. Schaefer, M. S. Roberts. 2004. The Optimal Timing of Living-Donor Liver Transplantation. *Management Science*, 50 (10), 1420-1430.
- Ata, B., A. Skaro, S. Tayur. 2017. Organjet: Overcoming geographical disparities in access to deceased donor kidneys in the United States. *Management Science*, 63 (9), 2776-2794.
- Bertsimas, D., T. Papalexopoulos, N. Trichakis, Y. Wang, R. Hirose, P. A. Vagefi. 2020. Balancing efficiency and fairness in liver transplant access: Tradeoff curves for the assessment of organ distribution policies. *Transplantation*, 104 (5), 981-987.
- Caro, F., T. Shirable, M. Guignard, A. Weintraub. 2004. School redistricting: embedding GIS tools with integer programming. *Journal of the Operational Research Society*, 55 (8), 836-849.
- Demirci, M. C., A. J. Schaefer, H. E. Romeijn, M. S. Roberts. 2012. An exact method for balancing efficiency and equity in the liver allocation hierarchy. *INFORMS Journal on Computing*, 24 (2), 260-275.
- Ergin, H., T. Sönmez, M. U. Ünver. 2020. Efficient and Incentive-Compatible Liver Exchange. *Econometrica*, 88 (3), 965-1005.
- Feng, S., N. P. Goodrich, J. L. Bragg-Gresham, D. M. Dykstra, J. D. Punch, M. A. DebRoy, S. M. Greenstein, R. M. Merion. 2006. Characteristics associated with liver graft failure: The concept of a donor risk index. American Journal of Transplantation, 6 783790.
- Freeman, R. B., R. H. Wiesner, A. Harper, S. V. McDiarmid, J. Lake, E. Edwards, R. Merion, R. Wolfe, J. Turcotte, L. Teperman. 2002. The new liver allocation system: Moving toward evidence based transplantation policy. *Liver Transplantation*, 8 (9), 851-858.
- Garfinkel, R. S., G. L. Nemhauser. 1970. Optimal political districting by implicit enumeration techniques. *Management Science*, 16 (8), B495-B508.
- Garonzik-Wang, J. M., N. T. James, K. J. Van Arendonk, N. Gupta, B. J. Orandi, E. C. Hall, A. B. Massie, R. A. Montgomery, N. N. Dagher, A. L. Singer, A. M. Cameron, D. L. Segev. 2013. The aggressive

- phenotype revisited: Utilization of higher-risk liver allografts. American Journal of Transplantation, 13 (4), 936-942.
- Gentry, S., E. Chow, A. B. Massie, D. L. Segev. 2015. Gerrymandering for justice: Redistricting U.S. liver allocation. *Interfaces*, 45 (5), 462-480.
- Glazier, A. K. 2018. The lung lawsuit: A case study in organ allocation policy and administrative law. Journal of Health & Biomedical Law, XIV 139-148.
- Gopalan, R., S. O. Kimbrough, F. H. Murphy, N. Quintus. 2013. The Philadelphia districting contest: Designing territories for city council based upon the 2010 census. *Interfaces*, 43 (5), 477-489.
- Halldorson, J. B., H. J. Paarsch, J. L. Dodge, A. M. Segre, J. Lai, J. P. Roberts. 2013. Center competition and outcomes following liver transplantation. *Liver Transplantation*, 19 (1), 96-104.
- Haugen, C. E., T. Ishaque, A. Sapirstein, A. Cauneac, D. L. Segev, S. Gentry. 2019. Geographic disparities in liver supply/demand ratio within fixed-distance and fixed-population circles. *American Journal of Transplantation*, 19 (7), 2044-2052.
- Hess, S. W., J. B. Weaver, J. N. Whelan, P. A. Zitlau. 1965. Nonpartisan political redistricting by computer. *Operations Research*, 13 (6), 998-1006.
- HHS. 1998. Organ Procurement and Transplantation Network; Final Rule (42 CFR Part 121). Federal Register, 63 (63), 16296-16338.
- Kilambi, V., S. Mehrotra. 2017. Improving liver allocation using optimized neighborhoods. *Transplantation*, 101 350-359.
- Kim, M., N. Xiao. 2017. Contiguity-based optimization models for political redistricting problems. *International Journal of Applied Geospatial Research (IJAGR)*, 8 (4), 1-18.
- Kim, W. R., J. R. Lake, J. M. Smith, D. P. Schladt, M. A. Skeans, A. M. Harper, J. L. Wainright, J. J. Snyder, A. K. Israni, B. L. Kasiske. 2018. OPTN/SRTR 2016 Annual data report: Liver. American Journal of Transplantation, 18, Suppl 1 172-253.
- Kong, N., A. J. Schaefer, B. Hunsaker, M. S. Roberts. 2010. Maximizing the efficiency of the U.S. liver allocation system through region design. *Management Science*, 56 (12), 2111-2122.
- Merion, R. M., D. E. Schaubel, D. M. Dykstra, R. B. Freeman, F. K. Port, R. A. Wolfe. 2005. The survival benefit of liver transplantation. *American Journal of Transplantation*, 5 307-313.
- Rawls, J. 1971. A Theory of Justice.. Harvard University Press, Cambridge, MA.
- Ricca, F., A. Scozzari, B. Simeone. 2013. Political districting: from classical models to recent approaches.

 Annals of Operations Research, 204 (1), 271-299.
- Shirabe, T. 2009. Districting modeling with exact contiguity constraints. *Environment and Planning B: Planning and Design*, 36 1053-1066.
- Stahl, J. E., N. Kong, S. M. Shechter, A. J. Schaefer, M. S. Roberts. 2005. A methodological framework for optimally reorganizing liver transplant regions. *Medical Decision Making*, 25 (1), 35-46.

Wey, A., J. Pyke, D. P. Schladt, S. E. Gentry, T. Weaver, N. Salkowski, B. L. Kasiske, A. K. Israni, J. J. Snyder. 2018. Offer acceptance practices and geographic variability in allocation model for end-stage liver disease at transplant. *Liver Transplantation*, 24 (4), 478-487.

Yeh, H., E. Smoot, D. A. Schoenfeld, J. F. Markmann. 2011. Geographic inequity in access to livers for transplantation. *Transplantation*, 91 (4), 479-486.

Zoltners, A. A., P. Sinha. 1983. Sales territory alignment: A review and model. *Management Science*, 29 (11), 1237-1256.

Appendix

A Comparing Organ Quality

We use the metric, donor risk index (DRI), proposed by Feng et al. (2006) to evaluate the quality of organs in our data set. This index measures the quality of an organ using demographic factors (age, race, height), cause and type of donor death, sharing type (local/regional/national), and cold ischemia time. A higher DRI is associated with a greater risk of graft failure. Since we want to assess the quality of organs at the time of recovery, we exclude cold ischemia time that depends on the transplant locations and assume local sharing for adequate comparison. In Figure 8, we compare the box plots of DRI across the regions. We see that there are no significant differences in the distributions of organ quality across the regions.

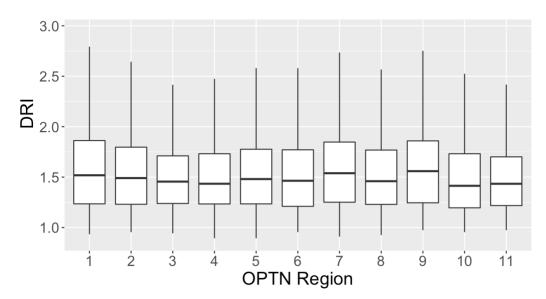


Figure 8 Comparison of organ quality across the regions using donor risk index (DRI).

Linearization of s/d Ratio and Contiguity Constraints

To linearize the right hand side of constraint (2) in [S-1], i.e., $\sum_{i=1}^{N_{sup}} \frac{1}{\sum_{k=1}^{N_{den}} d_k x_{ik}} s_i x_{ij}$, we introduce auxiliary variables: $y_{ij} \ge 0$ and $t_i \ge 0$ that are defined as follows.

$$t_{i} = \frac{1}{\sum_{k=1}^{N_{dem}} d_{k} x_{ik}} \qquad \forall i \in \mathcal{I}$$

$$y_{ij} = t_{i} x_{ij} \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$

$$(19)$$

$$y_{ij} = t_i x_{ij}$$
 $\forall i \in \mathcal{I}, j \in \mathcal{J}$ (20)

Together they imply:

$$\lambda \leq \sum_{i=1}^{N_{sup}} s_i y_{ij} \qquad \forall j \in \mathcal{J}$$

$$\sum_{k=1}^{N_{dem}} d_k y_{ik} = 1 \qquad \forall i \in \mathcal{I}$$
(21)

$$\sum_{k=1}^{N_{dem}} d_k y_{ik} = 1 \qquad \forall i \in \mathcal{I}$$
 (22)

A set of linear constraints (22) model equation (19). We note that $t_i \le 1$ and, $y_{ij} = \begin{cases} 0 & \text{if } x_{ij} = 0 \\ t_i & \text{if } x_{ij} = 1 \end{cases}$. Since y_{ij} is a product of two variables and therefore non-linear, the following linear constraints model $y_{ij} = t_i x_{ij}$:

$$y_{ij} \le t_i \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$
 (23)

$$y_{ij} \le x_{ij} \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$
 (24)

$$(1 - x_{ij}) + y_{ij} \ge t_i \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$
 (25)

$$y_{ij}, t_i \ge 0 \qquad \forall i \in \mathcal{I}, j \in \mathcal{J}$$
 (26)

Therefore, constraint (2) in [S-1] can be replaced by constraints (21)-(26). Constraint (8) in [S-2] can be linearized identically.

Shirabe (2009) describes flow-based contiguity constraints for districting problems. We adapt those constraints to model receiving contiguity and sharing contiguity in our neighborhood framework through equations (27)-(29) and (30)-(32), respectively. With receiving (sharing) contiguity, the suppliers (recipients) assigned to a recipient (supplier) form a continuous geography on the map. Let m_1 (m_2) be the maximum number of supply (demand) locations that can be assigned to a demand (supply) location. Parameter $a_{ik} = 1$, if supply locations i and k are geographically adjacent, and 0 otherwise. We use flow variables f_{ik}^{j} to model receiving contiguity and flow variables g_{jk}^i to model sharing contiguity. Flow variable f_{ik}^j denotes the flow from i to k (only defined when $a_{ik} = 1$) destined for demand location j, while flow variable g_{jk}^i denotes the flow from j to k (only defined when $a_{jk} = 1$) destined for supply location i. The first three constraints involving the flow variables f_{ik}^j ensure that if $x_{ij} = 1$ for a supply location i and a demand location j that are

non-adjacent, then every supply location on the path from i to j also supplies demand location j. The next set of three constraints involving the flow variables g_{jk}^i ensure that if $x_{ij} = 1$ for a supply location i and a demand location j that are non-adjacent, then every demand location on the path from i to j is also supplied by i.

$$\sum_{k=1}^{N_{sup}} f_{ik}^{j} a_{ik} - \sum_{k=1}^{N_{sup}} f_{ki}^{j} a_{ki} = x_{ij} \qquad \forall i \neq j, i \in \mathcal{I}, j \in \mathcal{J}$$

$$(27)$$

$$\sum_{k=1}^{N_{sup}} f_{ki}^{j} a_{ki} \le (m_{1} - 1) x_{ij} \ \forall i \in \mathcal{I}, j \in \mathcal{J}$$
 (28)

$$\sum_{k=1}^{N_{sup}} f_{jk}^{j} a_{jk} = 0 \qquad \forall j \in \mathcal{J}$$
 (29)

$$\sum_{k=1}^{N_{dem}} g_{jk}^{i} a_{jk} - \sum_{k=1}^{N_{dem}} g_{kj}^{i} a_{kj} = x_{ij} \qquad \forall i \neq j, i \in \mathcal{I}, j \in \mathcal{J}$$
(30)

$$\sum_{k=1}^{N_{dem}} g_{kj}^{i} a_{kj} \le (m_2 - 1) x_{ij} \ \forall i \in \mathcal{I}, j \in \mathcal{J}$$
(31)

$$\sum_{k=1}^{N_{dem}} g_{ik}^{i} a_{ik} = 0 \qquad \forall i \in \mathcal{I}$$

$$f_{ik}^{j} \geq 0 \qquad \forall i, k \in \mathcal{I}, j \in \mathcal{J}$$

$$(32)$$

$$f_{ik}^j \ge 0 \qquad \forall i, k \in \mathcal{I}, j \in \mathcal{J}$$
 (33)

$$g_{jk}^{i} \ge 0$$
 $\forall i \in \mathcal{I}, k, j \in \mathcal{J}$ (34)

Together constraints (27)-(34) refer to constraint (6) in [S-1].

$ au_{max}$			Size	Cutting	Run time	Nodes	Simplex	Best	Best	MIP
(in NM)		Rows	Columns	Planes	(in secs)	explored	iterations	obj.	bound	gap
	[SP-1]	1525	19061	2	7200	3444938	7290609	0.519	0.519	0.01%
350	SP-2	1668	19062	22	6	1	2047	0.884	0.884	0.00%
	[81 2]	1000	10002		O	-	2011	0.001	0.001	0.0070
400	[SP-1]	1527	24776	0	7200	2785686	9013065	0.537	0.537	0.01%
400	SP-2	1670	$\frac{24777}{24777}$	ŏ	7200	262484	20797966	0.611	0.610	0.12%
	[01 -2]	1010	24111	U	1200	202404	20131300	0.011	0.010	0.12/0
	[SP-1]	1528	30362	6	45	1	10922	0.543	0.543	0.01%
450	SP-2	1671	30363	16	7200	72601	18971000	0.606	0.606	0.10%
	[51 -2]	1071	30303	10	1200	72001	10971000	0.000	0.000	0.1070
	[SP-1]	1528	35990	2420	7200	499437	6067640	0.551	0.551	0.02%
500	SP-1	1671								
	[SP-2]	1071	35991	0	7200	85424	30217420	0.605	0.604	0.03%
	[SP-1]	1528	42175	3005	7200	220154	6747918	0.554	0.554	0.01%
550	SP-2									, -
	[SP-2]	1671	42176	495	7200	44898	11247871	0.604	0.604	0.03%
	[SP-1]	1528	48381	1936	7200	201431	1759417	0.555	0.555	0.01%
600										
	[SP-2]	1671	48382	93	7200	26765	10007894	0.602	0.602	0.08%
	[CD 1]	1500	F 4202	1000	7000	200106	05 42027	0.554	0.554	0.0107
650	[SP-1]	1528	54323	1908	7200	209106	2543837	0.554	0.554	0.01%
	[SP-2]	1671	54324	39	7200	28708	6409075	0.602	0.602	0.06%
	[CD 1]	1500	FOOTO	1 5	7000	027750	4501000	0.550	0.550	0.0107
700	[SP-1]	1528	59952	15	7200	837756	4521290	0.556	0.556	0.01%
. 00	[SP-2]	1671	59953	20	7200	22964	5031591	0.602	0.601	0.05%

Computational details of the set-partitioning model run on four-digit zip-clusters.

C Set Partitioning Model Computational Details

Table 8 contains computational details of the set-partitioning model run on four-digit zip-clusters. We report the problem size, total number of cutting planes used by the solver, run time until termination, nodes explored, simplex iterations, best objective, best bound and MIP gap. We observe that by two hours of running time, MIP gap of [SP-1] and [SP-2] typically reaches below 0.02% and 0.12% respectively, which corresponds to difference in s/d ratio at fourth or higher decimal places between the best objective and best bound.

D DSA Neighborhoods

Table 9, 10, and 11, present the neighborhoods obtained by our model based on the DSA version when the maximum distance to any DSA in the neighborhood is constrained to 500 NM, 600

	D . 1'	
DSA	Radius (in NM)	Neighbors
ALOB AROR	336 305	KYDA, ALOB, NCCM, TNDS, MSOP, AROR, SCOP, TNMS, FLUF, GALL, LAOP
AZOB	499	MSOP, AROR, TXSB, TNMS, MOMÁ, MWOB, TXGC, LAOP, OKÓP AZOB, CORS, UTOP, CASD, CAOP
CADN	366	HIOP, CADN, CASD, CAOP
CAGS CAOP	318 499	HIOP, CADN, CAOP HIOP, AZOB, UTOP, CADN, CASD, CAOP
CASD	366	HIOP, AZOB, CADN, CASD, CAOP CORS, UTOP
$\begin{array}{c} \text{CORS} \\ \text{CTOP} \end{array}$	$\frac{316}{203}$	CORS, UTOP MAOB, NYFL, CTOP, PADV, NJTO, NYRT
DCTC	339	OHOV, NCNC, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, NYRT
$_{ m FLFH}$	492	NCNC, ALOB, PRLL, NCCM, TNDS, FLMP, MSOP, SCOP, FLFH, FLUF, FLWC, GALL
FLUF	$\frac{472}{422}$	PRLL, FLMP, SCOP, FLFH, FLUF, FLWC, GALL NCNC, ALOB, PRLL, NCCM, TNDS, FLMP, MSOP, SCOP, FLFH, FLUF, FLWC, GALL, LAOP NCNC, ALOB, PRLL, NCCM, FLMP, MSOP, SCOP, FLFH, FLUF, FLWC, GALL
FLWC	488	NCNC, ALOB, PRLL, NCCM, FLMP, MSOP, SCOP, FLFH, FLUF, FLWC, GALL
GALL	484	OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, FLMP, MSOP, AROR, SCOP, TNMS, VATB, MOMA, OHID, FILEH OHID, FILE INOD, FILMC, CALL, LAOP.
HIOP	NA	OHLP, FLFH, OHLB, FLUF, INOP, FLWC, GALL, LAOP HIOP, CADN, CASD, CAOP
IAOP ILIP	$\frac{210}{447}$	MNOP, WIUW, IAOP, NEOR OHOV, MNOP, KYDA, PATF, TNDS, WIUW, ILIP, AROR, IAOP, TNMS, MOMA, MWOB, OHLP, WIDN, NEOR,
ILIF	447	OHUS, INOP, MIOA, PAIF, INDS, WIOW, ILIF, AROK, IAOF, INMS, MOMA, MWOB, OHLF, WIDN, NEOK, OHLB, INOP, MIOP
INOP	224	OHOV, KYDA, ILIP, OHLP, WIDN, OHLB, INOP, MIOP
KYDA	497	OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, MSOP, AROR, IAOP, SCOP,
LAOP	396	TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, OHLB, FLUF, INOP, GALL, MIOP TXSA ALOB MSOP AROR TXSB TNMS TXCC LAOP OKOP
MAOB	170	TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, OHLB, FLUF, INOP, GALL, MIOP TXSA, ALOB, MSOP, AROR, TXSB, TNMS, TXGC, LAOP, OKOP MAOB, CTOP, NJTO, NYRT NCNC, MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, NYRT
MDPC MIOP	$\frac{308}{482}$	NCNC, MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, NYRT OHOV, MNOP, KYDA, NYFL, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, IAOP, VATB, MDPC, MOMA,
	102	OHLP, NJTO, WIDN, OHLB, INOP, NYRT, MIOP
MNOP	270	MNOP, WIUW, IAOP, NEOR
MOMA MSOP	233 406	ILIP, AROR, TNMS, MOMA, MWOB ALOB, TNDS, MSOP, AROR, TXSB, TNMS, MOMA, TXGC, FLUF, GALL, LAOP
MWOB	477	MNOP, KYDÁ, WIUW, ILIP, AROR, IAOP, TXSB, TNMS, MOMÁ, MWÓB, WIDN, CORS, NEOR, TXGC, INOP,
NCCM	403	OKOP OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, SCOP, VATB, MDPC, OHLP, FLFH, OHLB,
	100	FLUF, INOP, GALL
NCNC	489	OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, SCOP, VATB, MDPC, OHLP,
NEOR	390	NJTO, FLFH, OHLB, FLUF, INOP, FLWC, NYRT, GALL, MIOP MNOP, WIUW, IAOP, MOMA, MWOB, CORS, NEOR, OKOP
NJTO	188	MAOB, NYFL, CTOP, DCTC, PADV, MDPC, NJTO, NYRT
NMOP NVLV	$\frac{282}{308}$	AZOB AZOB, CADN, CASD, CAOP
NYAP	182	MAOB, NYFL, CTOP, PADV, NJTO, NYRT
$\begin{array}{c} \mathrm{NYFL} \\ \mathrm{NYRT} \end{array}$	458	OHOV, MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, INOP, NYRT, MIOP
NYWN	$\frac{193}{234}$	MAOB, NYFL, CTOP, PADV, NJTO, NYRT OHOV, MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, INOP, NYRT, MIOP MAOB, NYFL, CTOP, PADV, MDPC, NJTO, NYRT NYFL, PATF, PADV, OHLB, MIOP OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB, MDPC,
OHLB	459	OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB, MDPC,
OHLC	329	MOMA, OHLP, NJTO, WIDN, OHLB, INOP, NYRT, MIOP OHOV, KYDA, PATE, DCTC, TNDS, WIIIW II.IP, OHLP, WIDN, OHLB, INOP, MIOP
OHLP	307	OHOV, KYDA, PATF, DCTC, TNDS, WIÚW, ILIP, OHLP, WIDN, OHLB, INOP, MIOP OHOV, KYDA, PATF, NCCM, DCTC, TNDS, ILIP, VATB, MDPC, OHLP, WIDN, OHLB, INOP, MIOP
OHOV OKOP	$\frac{284}{475}$	OHOV, KYDA, PATF, NCCM, TNDS, ILIP, OHLP, WIDN, OHLB, INOP, MIOP TXSA, MSOP, AROR, IAOP, TXSB, TNMS, MOMA, MWOB, CORS, NEOR, TXGC, LAOP, OKOP
ORUO	234	WALC, HIOP, ORUO
PADV PATF	$\frac{235}{278}$	MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, NYRT OHOV, NCNC, NYFL, PATF, DCTC, PADV, VATB, MDPC, OHLP, NJTO, OHLB, NYRT, MIOP
PRLL	NA	PRLL, FLMP, FLWC
$\frac{\text{SCOP}}{\text{TNDS}}$	326 493	NCNĆ, KYDÁ, ALOB, NCCM, TNDS, SCOP, VATB, FLFH, FLUF, GALL OHOV, NCNC, KYDÁ, ALOB, PATÉ, NCCM, DCTC, TNDS, ILIP, MSOP, AROR, SCOP, TNMS, VATB, MDPC,
INDS	490	MOMA, MWOB, OHLP, WIDN, FLFH, OHLB, FLUF, INDP, GALL, MIOP, LAOP
TNMS	492	OHOV, KYDA, ALOB, NCCM, TNDS, ILIP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS, MOMA, MWOB, OHLP,
TVCC	276	WIDN, TXGC, FLUF, INOP, GALL, LAOP, OKOP
$_{\mathrm{TXSA}}^{\mathrm{TXGC}}$	$\frac{276}{363}$	TXSA, TXSB, TXGC, LAOP, OKOP TXSA, TXSB, TXGC, LAOP
TXSB	218	TXSA, TXSB, TXGC, OKOP
UTOP VATB	$\frac{480}{196}$	AZOB, CORS, UTOP, CADN NCNC, DCTC, PADV, VATB, MDPC
WALC	234	WALC, ORUO
WIDN WIUW	$\frac{284}{426}$	OHOV, WIUW, ILIP, IAOP, WIDN, INOP, MIOP OHOV, MNOP, WIUW, ILIP, IAOP, MOMA, MWOB, OHLP, WIDN, NEOR, OHLB, INOP, MIOP
		,,, mor, morni, m, ob, omb, mbon, omb, mor, mior

Table 9 Geographical allocation policy for DSAs when the maximum permitted distance to a neighboring DSA is

NM, and 700 NM, respectively. The column, "Radius" provides the radius (in terms of the transplant volume weighted distance as discussed in Section 4.2) of each neighborhood. The column, "Neighbors" contains the DSAs to which the DSA in the first column will share its organs to.

DSA	Radius (in NM)	Neighbors
ALOB	517	OHOV, NCNC, KYDA, ALOB, NCCM, TNDS, MSOP, AROR, SCOP, TNMS, VATB, MOMA, OHLP, FLFH, TXGC,
AROR	569	FLUF, INOP, FLWC, GALL, LAOP OHOV, TXSA, KYDA, ALOB, TNDS, ILIP, MSOP, AROR, IAOP, TXSB, TNMS, MOMA, MWOB, WIDN, NEOR, TXGC FLUF INOP GALL LAOP OKOP
AZOB CADN	546	TXGC, FLUF, INOP, GALL, LAOP, OKOP AZOB, CORS, UTOP, CADN, CASD, CAOP HIOP, CADN, CASD, CAOP HIOP, CADN, CASD, CAOP
$\begin{array}{c} \text{CADN} \\ \text{CAGS} \end{array}$	366 446	HIOP, CADN, CASD, CAOP
CAOP	499	HIÓP, ÓRUÓ, ÚTOP, CADN, CASD, CAOP HIOP, AZOB, UTOP, CADN, CASD, CAOP
CASD	265	HIOP, AZOB, CASD, CAOP CORS, UTOP MAOB, NYFL, CTOP, PADV, MDPC, NJTO, NYRT
$\begin{array}{c} \text{CORS} \\ \text{CTOP} \end{array}$	$\frac{316}{244}$	CORS, UTOP MAOR NYFI, CTOP PADV MDPC NITO NYRT
DCTC	346	OHOV, NCNC, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, PADV, SCOP, VATB, MDPC, OHLP, NJTO, OHLB,
		NYRT
$_{ m FLFH}$	$\frac{578}{573}$	NCNC, ALOB, PRLL, NCCM, TNDS, FLMP, MSOP, SCOP, TNMS, VATB, FLFH, FLUF, FLWC, GALL, LAOP NCNC, ALOB, PRLL, NCCM, FLMP, SCOP, FLFH, FLUF, FLWC, GALL OHOV, NCNC, KYDA, ALOB, PRLL, NCCM, TNDS, FLMP, MSOP, AROR, SCOP, TNMS, VATB, OHLP, FLFH, FLUF,
FLUF	592	OHOV, NCNC, KYDA, ALOB, PRLL, NCCM, TNDS, FLMP, MSOP, AROR, SCOP, TNMS, VATB, OHLP, FLFH, FLUF,
		FLWC, GALL, LAOP
FLWC	330	FLWC, GALL, LAOP PRLL, FLMP, FLFH, FLUF, FLWC, GALL OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, FLMP, ILIP, MSOP, AROR, SCOP, TNMS, VATB,
GALL	565	MDDC MOMA OUI B FIFE OUI B FIFE INOD FIWE CALL MIOD LACE
HIOP	NA	MDPC, MOMA, OHLP, FLFH, OHLB, FLUF, INOP, FLWC, GALL, MIOP, LAOP HIOP, CADN, CASD, CAOP
IAOP	335	MNOP, WIUW, ILIP, IAOP, MOMA, MWOB, WIDN, NEOR, INOP OHOV, MNOP, KYDA, TNDS, WIUW, ILIP, IAOP, MOMA, MWOB, OHLP, WIDN, OHLB, INOP, MIOP OHOV, MNOP, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, MSOP, AROR, IAOP,
ILIP INOP	365 571	OHOV, MNOP, KYDA, TNDS, WIUW, ILIP, IAOP, MOMA, MWOB, OHLP, WIDN, OHLB, INOP, MIOP OHOV MNOP NENC KYDA ALOR NYFL PATE NECM DETE TNDS PADV WIIW ILIP MSOP AROR IAOP
11101	011	SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, OHLB, INOP, NYRT, GALL, MIOP,
		OKOP
KYDA	497	OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, MSOP, AROR, IAOP, SCOP,
LAOD	200	TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, OHLB, FLUF, INOP, GALL, MIOP TXSA, ALOB, MSOP, AROR, TXSB, TNMS, TXGC, LAOP, OKOP MAOB, CTOP, NJTO, NYRT ONEY, MACE, NYRL, DATE, NGCM, CTCD, DATE, DATE, NGCM, CTCD, DATE, DATE, NGCM, CTCD, DATE, DAT
$_{ m MAOB}^{ m LAOP}$	$\frac{396}{170}$	TASA, ALOB, MSOP, AROR, TASB, TNMS, TAGC, LAOP, OROP MAOB, CTOP, NJTO, NYRT
MDPC	407	OHOV, NCNC, KYDA, MAOB, NYFE, PAIF, NCCM, CTOP, DCTC, PADV, SCOP, VAIB, MDPC, OHEP, NJTO,
		OHLB, NYRT, MIOP OHOV, MNOP, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, IAOP, SCOP,
MIOP	594	OHOV, MNOP, NCNC, KYDA, MAOB, NYFL, PATF, NCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, IAOP, SCOP,
MNOP	270	TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, OHLB, INOP, NYRT, GALL, MIOP MNOP, WIUW, IAOP, NEOR
MOMA	600	MNOP, WIUW, IAOP, NEOR OHOV, MNOP, NCNC, KYDA, ALOB, PATF, NCCM, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS,
11000	400	MOMA, MWOB, OHLP, WIDN, NEOR, OHLB, TXGC, FLUF, INOP, GALL, MIOP, LAOP, OKOP TXSA, KYDA, ALOB, NCCM, TNDS, MSOP, AROR, SCOP, TXSB, TNMS, MOMA, MWOB, FLFH, TXGC, FLUF,
MSOP	482	TXSA, KYDA, ALOB, NCCM, TNDS, MSOP, AROR, SCOP, TXSB, TNMS, MOMA, MWOB, FLFH, TXGC, FLUF,
MWOB	590	FLWC, GALL, LAOP, OKOP OHOV, TXSA, MNOP, KYDA, ALOB, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, TXSB, TNMS, MOMA, MWOB,
		WIDN, CORS, NEOR, TXGC, INOP, MIOP, LAOP, OKOP OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, ILIP, MSOP, SCOP, TNMS, VATB, MDPC, MOMA,
NCCM	498	OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, ILIP, MSOP, SCOP, TNMS, VATB, MDPC, MOMA,
NCNC	570	OHLP, NJTO, FLFH, OHLB, FLUF, INOP, FLWC, NYRT, GALL, MIOP OHOV, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, ILIP, MSOP, SCOP, TNMS,
110110	570	VATE MDPC OHLP NITO FIFH OHLE FILIT INCOM, OT NYRT GALL MIOP
NEOR	377	VATB, MDPC, OHLP, NJTO, FLFH, OHLB, FLUF, INOP, FLWC, NYRT, GALL, MIOP MNOP, IAOP, MOMA, MWOB, CORS, NEOR, OKOP MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, NYRT
NJTO NMOP	$\frac{254}{282}$	MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, NYRT AZOB
NVLV	308	AZOB, CADN, CASD, CAOP
NYAP	182	ÄZÖĞ, CADN, CASD, CAOP MAOB, NYFL, CTOP, PADV, NJTO, NYRT OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, PADV, WIUW, ILIP, SCOP, VATB, MDPC, OHLP,
NYFL	582	OHOV, NUNC, KYDA, MAOB, NYFL, PATF, NCCM, GTOP, DCTC, PADV, WIOW, ILIP, SCOP, VATB, MDPC, OHLP, NJTO, WIDN, OHLB, INOP, NYRT, MIOP
NYRT	195	MAOB, NYFL, CTOP, DCTC, PADV, MDPC, NJTO, NYRT
NYWN	326	MAOB, NYFL, CTOP, DCTC, PADV, MDPC, NJTO, NYRT MAOB, NYFL, PATF, CTOP, DCTC, PADV, MDPC, OHLP, NJTO, OHLB, NYRT, MIOP OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB, MDPC,
OHLB	457	OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB, MDPC,
OHLC	346	OHLP, NJTO, WIDN, OHLB, INOP, NYRT, MIOP OHOV. KYDA, PATF. NCCM. DCTC. TNDS. WIUW. ILIP. MDPC. MOMA. OHLP. WIDN. OHLB. INOP. MIOP
OHLP	366	OHOV, KYDA, PATF, NCCM, DCTC, TNDS, WIUW, ILIP, MDPC, MOMA, OHLP, WIDN, OHLB, INOP, MIOP OHOV, NCNC, KYDA, NYFL, PATF, NCCM, DCTC, TNDS, PADV, ILIP, SCOP, VATB, MDPC, OHLP, WIDN, OHLB,
OHOV	599	INOP, MIOP OHOV, MNOP, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, MSOP,
OHOV	599	OHOV, MINOT, NEIDA, ALOB, NIFL, FAIF, NECM, CIOF, DELC, INDS, FADV, WILW, ILIF, MSOF, AROR, IAOP, SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, OHLB, FLUF, INOP,
		NYRT GALL MIOP
OKOP	397	TXSA, AROR, TXSB, TNMS, MOMA, MWOB, NEOR, TXGC, LAOP, OKOP WALC, HIOP, ORUO
ORUO PADV	$\frac{234}{275}$	WALC, HIOP, ORUO
PATF	406	MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, OHLB, NYRT OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, ILIP, SCOP, VATB, MDPC, OHLP,
		NJTO, WIDN, OHLB, INOP, NYRT, MIOP
$\begin{array}{c} \text{PRLL} \\ \text{SCOP} \end{array}$	NA 450	PRLL, FLMP OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, SCOP, TNMS, VATB, MDPC, OHLP, FLFH,
3001	450	OHLE FLIF INOP FINC GALL.
TNDS	525	OHLB, FLUF, INOP, FLWC, GALL OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, ILIP, MSOP, AROR, IAOP, SCOP, TNMS, VATB,
m		MDPC, MOMA, MWOB, OHLP, WIDN, FLFH, OHLB, FLUF, INOP, FLWC, GALL, MIOP, LAOP OHOV, TXSA, NCNC, KYDA, ALOB, PATF, NCCM, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS,
TNMS	594	OHOV, TASA, NUNC, KYDA, ALOB, PATF, NUCM, TNDS, WILW, LILP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS, VATB, MOMA, MWOD, OHID, WIDN, NEOP, STEIL OHID, TYGG, PLUE, WOD, FRING CALL, WOD, TAGE CYCLE
TXGC	276	VATB, MOMA, MWOB, OHLP, WIDN, NEOR, FLFH, OHLB, TXGC, FLUF, INOP, FLWC, GALL, MIOP, LAOP, OKOP TXSA, TXSB, TXGC, LAOP, OKOP TXSA, TXSB, TXGC
TXSA	179	TXSA, TXSE, TXGC, LLCC, CALLERY, CALLERY, CALLERY, TXSE, TXGC, LLCC, CALLERY, CALLER
$\begin{array}{c} { m TXSB} \\ { m UTOP} \end{array}$	598	TXSA, ALOB, MSOP, AROR, TXSB, TNMS, MOMA, MWOB, CORS, NEOR, TXGC, LAOP, OKOP AZOB, CORS, UTOP, CADN
VATB	$\frac{480}{437}$	AZOB, CORS, UTOP, CADIN OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, SCOP, VATB, MDPC, OHLP,
		NJTO, OHLB, INOP, NYRT, GALL, MIOP
WALC	234	WALC, ORUÓ WIUW, ILIP, WIDN, INOP, MIOP
WIDN WIUW	$\frac{200}{248}$	WIUW, ILIP, WIDN, INOP, MIOP MNOP, WIUW, ILIP, IAOP, WIDN, MIOP
		- , , , ,

Table 10 Geographical allocation policy for DSAs when the maximum permitted distance to a neighboring DSA is 600 NM.

DSA	Radius	Neighbors
ALOB	(in NM) 684	OHOV, TXSA, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, FLMP, ILIP, MSOP, AROR, IAOP, SCOP,
ALOB	084	TXSB, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, FLFH, OHLB, TXGC, FLUF, INOP, FLWC, GALL, MIOP, LAOP, OKOP
AROR	297	MSOP, AROR, TXSB, TNMS, MOMA, MWOB, LAOP, OKOP
AZOB	546	MSOP, AROR, TXSB, TNMS, MOMA, MWOB, LAOP, OKOP AZOB, CORS, UTOP, CADN, CASD, CAOP WALC, HIOP, AZOB, ORUO, UTOP, CADN, CASD, CAOP
$_{\mathrm{CAGS}}^{\mathrm{CADN}}$	$\frac{655}{318}$	WALC, HIOP, AZOB, ORUO, UTOP, CADN, CASD, CAOP HIOP, CADN, CAOP
CAOP	316	HIOP, AZOB, CADN, CASD, CAOP
CASD	537	HIOP, AZOB, UTOP, CADN, CASD, CAOP
CORS	316	CORS, UTOP
$\begin{array}{c} \text{CTOP} \\ \text{DCTC} \end{array}$	$\frac{343}{421}$	MÃOB, NYFL, PATF, CTOP, DCTC, PADV, MDPC, NJTO, NYRT OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, PADV, SCOP, VATB, MDPC, OHLP, NJTO,
FLFH	687	OHLB, INOP, NYRT, MIOP
		OHOV, NCNC, KYDA, ALOB, PRLL, NCCM, DCTC, TNDS, FLMP, MSOP, SCOP, TNMS, VATB, MDPC, OHLP, FLUF, FLUE, GLLL, LAOP, DDLL, FLOR, GLLL, LAOP
$_{ m FLMP}$	$\frac{287}{674}$	PRLL, FLMP, FLFH, FLUF, FLWC OHOV, NCNC, KYDA, ALOB, PATF, PRLL, NCCM, DCTC, TNDS, FLMP, MSOP, AROR, SCOP, TNMS, VATB,
FLWC	488	MDPC, MOMA, OHLP, FLFH, OHLB, FLUF, INOP, FLWC, GALL, LAOP NCNC, ALOB, PRLL, NCCM, FLMP, MSOP, SCOP, FLFH, FLUF, FLWC, GALL OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, FLMP, ILIP, MSOP, AROR, SCOP, TNMS,
GALL	671	OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, FLMP, ILIP, MSOP, AROR, SCOP, TNMS,
		VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, FLFH, OHLB, TXGC, FLUF, INOP, FLWC, NYRT, GALL,
		MIOP, LAOP, OKOP
HIOP	NA	HIOP, CASD
IAOP	235	MNOP, WIUW, ILIP, IAOP, MWOB, WIDN, NEOR
ILIP	447	OHOV, MNOP, KYDA, PATF, TNDS, WIUW, ILIP, AROR, IAOP, TNMS, MOMA, MWOB, OHLP, WIDN, NEOR, OHLB, INOP, MIOP
INOP	510	OHOV, MNOP, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, MSOP, AROR, IAOP,
TANKE A	000	SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, NEOR, OHLB, INOP, GALL, MIOP OHOV, MNOP, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP,
KYDA	696	OHOV, MINOP, NENC, KYDA, ALOB, MAOB, NYFL, PAIT, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP,
		MSOP, AROR, IAOP, SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, FLFH, OHLB,
LAOP	396	TXGC, FLUF, INOP, FLWC, NYRT, GALL, MIOP, LAOP, OKOP
MAOB	170	TXSA, ALOB, MSOP, AROR, TXSB, TNMS, TXGC, LAOP, OKOP MAOB, CTOP, NJTO, NYRT MAOB, CTOP, NJTO, NYRT MAOB, CTOP, NJTO, NYRT MAOB, NYRT, NGCM, CTOP, DCTG, TNDG, PADV, WILLIAM, H. D. GGOD, VATER MONOY, NGCC, WARD, MAOB, NYRT, DATE, NGCM, CTOP, DCTG, TNDG, PADV, WILLIAM, H. D. GGOD, VATER
MDPC	645	OHOV, NCNC, KYDA, ALOB, MAOB, NYFE, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB,
		MDPC, OHLP, NJTO, WIDN, OHLB, FLUF, INOP, NYRT, GALL, MIOP
MIOP	$\frac{216}{270}$	OHOV, ILIP, OHLP, WIDN, OHLB, INOP, MIOP MNOP, WIUW, IAOP, NEOR
MNOP MOMA	694	MNOP, WILW, LAOP, NEOK OHOV, TXSA, MNOP, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, SCOP,
MOMIN	004	TXSB, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, WIDN, CORS, NEOR, OHLB, TXGC, FLUF, INOP, GALL,
		MIOP, LAOP, OKOP
MSOP	418	KYDA, ALOB, TNDS, MSOP, AROR, TXSB, TNMS, MOMA, TXGC, FLUF, GALL, LAOP, OKOP
MWOB	590	OHOV, TXSA, MNOP, KYDA, ALOB, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, TXSB, TNMS, MOMA, MWOB,
NOOM	020	WIDN, CORS, NEOR, TXGC, INOP, MIOP, LAOP, OKOP NCNC, KYDA, NCCM, TNDS, SCOP, VATB, GALL NCNC, KYDA, PATF, NCCM, DCTC, TNDS, PADV, SCOP, VATB, MDPC, OHLP, GALL MNOP, IAOP, MOMA, MWOB, CORS, NEOR, OKOP WAGE, MADE, MADE, CORD, DCTC, DADY, WDPC, NUTC, NAPT, WAGE, MADE, CORD, CORG, DADY, WDPC, NUTC, NAPT,
NCCM NCNC	$\frac{238}{320}$	NCNC, KYDA, NCCM, INDS, SCOP, VATB, GALL NCNC KYDA PATE NCCM DCTC TNDS PADV SCOP VATB MDPC OHLP GALL
NEOR	377	MNOP, IAOP, MOMA, MWOB, CORS, NEOR, OKOP
NJTO	254	MAOB, NIFE, FAIF, CIOF, DCIC, FADV, MDFC, NJIO, NIKI
NMOP NVLV	$\frac{282}{311}$	AZOB AZOB, UTOP, CADN, CASD, CAOP
NYAP	182	MAOB, NYFL, CAOP, PADV, NJTO, NYRT
NYFL	186	NYFL, PADV
NYRT	195	MAOB, NYFL, CTOP, DCTC, PADV, MDPC, NJTO, NYRT
NYWN OHLB	$\frac{153}{484}$	NYFL, OHLB OHOV, NCNC, KYDA, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB, MDPC,
OHLD	404	MOMA OHLP NITO WIDN OHLB INOP NYRT GALL MIOP
OHLC	678	MOMA, OHLP, NJTO, WIDN, OHLB, INOP, NYRT, GALL, MIOP OHOV, MNOP, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP,
		MSOP, AROR, IAOP, SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, OHLB, FLUF,
		INOP, NYRT, GALL, MIOP, OKOP
OHLP	624	OHOV, MNOP, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP,
		MSOP, AROR, IAOP, SCOP, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, OHLB, FLUF, INOP,
OHOU	410	NYRT, GALL, MIOP
OHOV	410	OHOV, NCNC, KYDA, ALOB, PATF, NCCM, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, TNMS, VATB, MDPC,
OKOP	418	MOMA, OHLP, WIDN, OHLB, INOP, GALL, MIOP TXSA, MSOP, AROR, TXSB, TNMS, MOMA, MWOB, NEOR, TXGC, LAOP, OKOP WALC, HIOP, ORUO, UTOP, CADN, CAOP MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, NYRT OHOV, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB,
ORUO	662	WALC, HIOP, ORUO, UTOP, CADN, CAOP
PADV	235	MAOB, NYFL, PATF, CTOP, DCTC, PADV, VATB, MDPC, NJTO, NYRT
PATF	544	OHOV, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, ILIP, SCOP, VATB,
PRII	NΑ	MDPC, MOMA, OHLP, NJTO, WIDN, OHLB, INOP, NYRT, GALL, MIOP PRLL, FLMP, FLFH, FLWC
$\begin{array}{c} \mathrm{PRLL} \\ \mathrm{SCOP} \end{array}$	NA 263	NCNC, NCCM, TNDS, SCOP, VATB, FLUF, GALL
TNDS	694	NCNC, NCCM, TNDS, SCOP, VATB, FLUF, GALL OHOV, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, FLMP, ILIP, MSOP, AROR, NCNC, KYDA, ALOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, WIUW, FLMP, ILIP, MSOP, AROR, NCNC, WILLIAM, NAME, NAME
		TAOP, SCOP, TXSB, TNMS, VATB, MDPC, MOMA, MWOB, OHLP, NJTO, WIDN, NEOR, FLFH, OHLB, TXGC,
		FLUF, INOP, FLWC, NYRT, GALL, MIOP, LAOP, OKOP OHOV, TXSA, NCNC, KYDA, ALOB, NCCM, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS, MOMA,
TNMS	561	OHOV, TXSA, NCNC, KYDA, ALOB, NCCM, TNDS, WIUW, ILIP, MSOP, AROR, IAOP, SCOP, TXSB, TNMS, MOMA,
TVCC	276	MWOB, OHLP, WIDN, NEOR, OHLB, TXGC, FLUF, INOP, GALL, MIOP, LAOP, OKOP
TXSA	$\frac{276}{179}$	1A3A, 1A3B, 1AGC, LAOP, UKOP TXSA TXSR TXGC
TXSB	598	TXSA, ALOB, MSOP, AROR, TXSB, TNMS, MOMA, MWOB. CORS. NEOR. TXGC. LAOP. OKOP
TXGC TXSA TXSB UTOP	537	MWOB, OHLP, WIDN, NEOR, OHLB, TXGC, FLUF, INOP, GALL, MIOP, LAOP, OKOP TXSA, TXSB, TXGC, LAOP, OKOP TXSA, TXSB, TXGC TXSA, ALOB, MSOP, AROR, TXSB, TNMS, MOMA, MWOB, CORS, NEOR, TXGC, LAOP, OKOP AZOB, ORUO, CORS, UTOP, CADN, CASD, CAOP OHOV, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, SCOP, VATB, MDPC,
VATB	517	OHOV, NCNC, KYDA, ALOB, MAOB, NYFL, PATF, NCCM, CTOP, DCTC, TNDS, PADV, SCOP, VATB, MDPC,
WALC	234	OHLP, NJTO, OHLB, FLUF, INOP, NYRT, GALL, MIOP WALC, ORUO OHOV, MNOP, KYDA, PATF, WIUW, ILIP, IAOP, MOMA, MWOB, OHLP, WIDN, OHLB, INOP, MIOP MNOP, WIUW, ILIP, IAOP, WIDN, MIOP
WIDN	$\frac{234}{423}$	OHOY, MNOY, KYDA, PATF, WIUW, ILIP, IAOP, MOMA. MWOB. OHLP. WIDN, OHLB. INOP. MIOP
WIUW	248	MNOP, WIUW, ILIP, IAOP, WIDN, MIOP

Table 11 Geographical allocation policy for DSAs when the maximum permitted distance to a neighboring DSA is 700 NM.

The radii for each zip-code cluster in the zip-code version (which is too large to include as a table) can be obtained by contacting the authors.