

Genomic data compression

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Outline

- FASTQ compression – SPRING
 - Introduction and motivation
 - FASTQ format and compression results
 - Algorithms - SPRING and others
 - SPRING as a practical tool
 - Next steps: preliminary work on noisy long read compression
- Lossy compression for nanopore raw signal data
 - Background
 - Evaluation pipeline
 - Results

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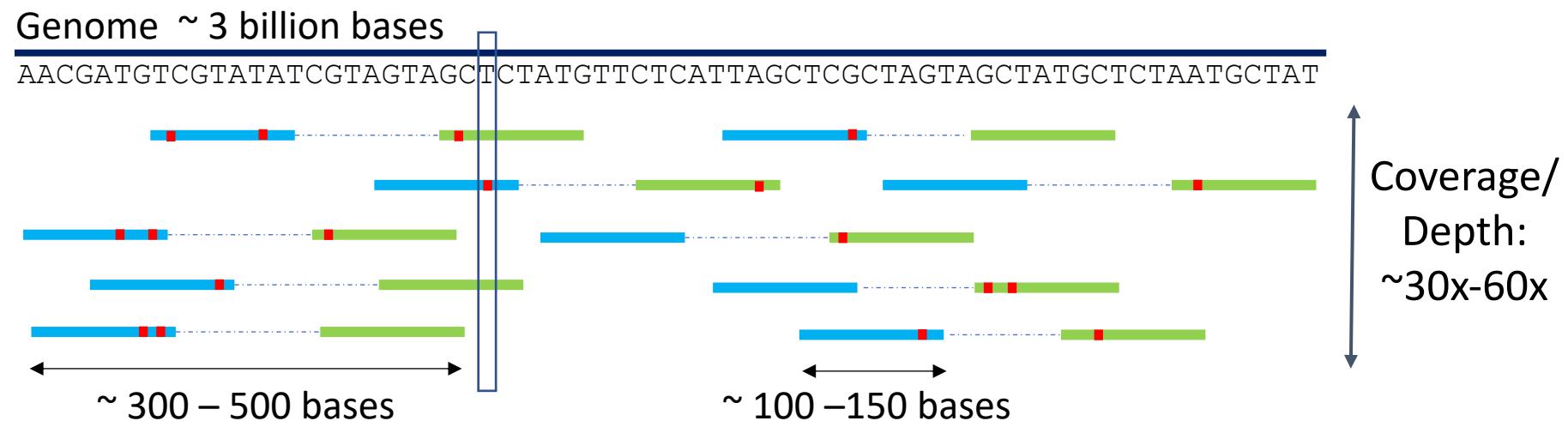
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Joint work with

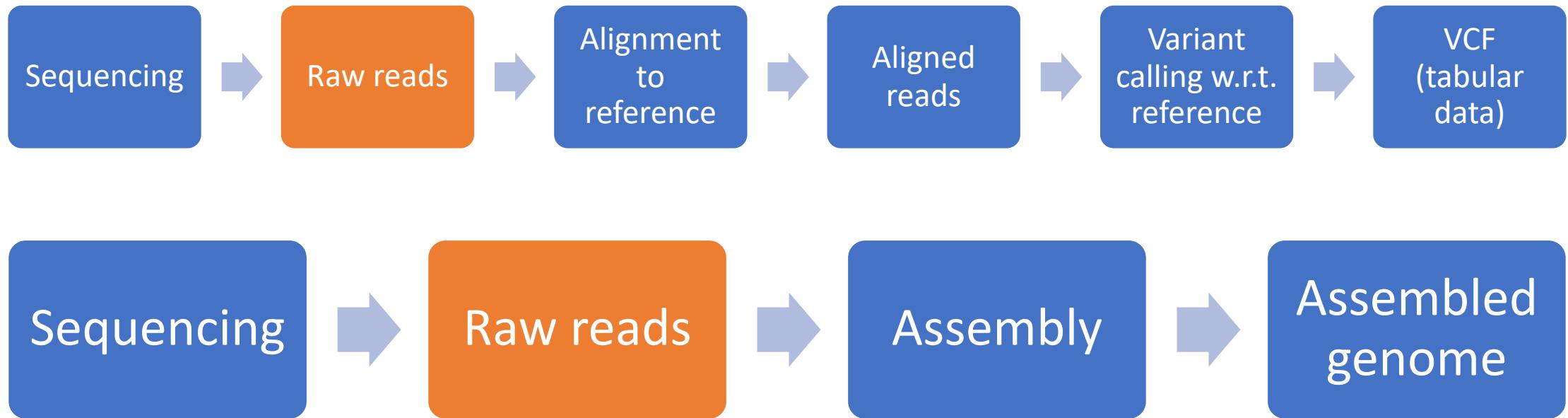
- Kedar Tatwawadi, Stanford University
- Idoia Ochoa, UIUC
- Mikel Hernaez, UIUC
- Tsachy Weissman, Stanford University

Genome sequencing

- Genome: long string of bases {A, C, G, T}
- Sequenced as noisy paired substrings (*reads*):



Typical workflows



Why store raw reads?

- Pipelines improve with time - need raw data for reanalysis
- For temporary storage - alignment and assembly time-consuming
- Can't perform alignment when reference genome not available – e.g., de novo assembly or metagenomics
- Can get better compression than aligned data compression if significant variation from reference (more on this later)!

FASTQ format

File 1

@ERR174324.1 HSQ1009_86:1:1101:1192:2116/1
ATTCNGTCACTTCTCACCGAGGCCCTCATTCAACACTGGGAATTAAAATTCGAC...

Read

3

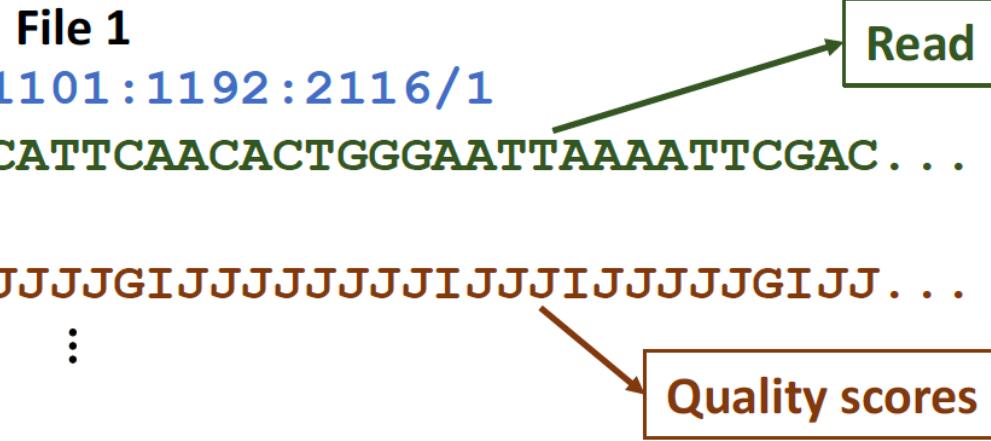
File 2

@ERR174324.2 HSQ1009 86:1:1101:1192:2116/2

CAGANAGAGACTCTGTCTAAAAAAACAAACAAACAAACAAAAGTCTTA . . .

+

•



We'll mostly focus on **reads** in this talk.

Read compression

Read compression

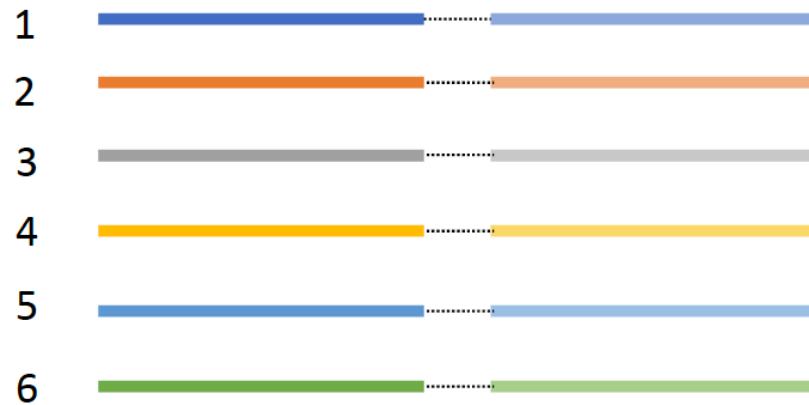
- For a typical 25x human dataset:
 - Uncompressed: 79 GB (1 byte/base)

Read compression

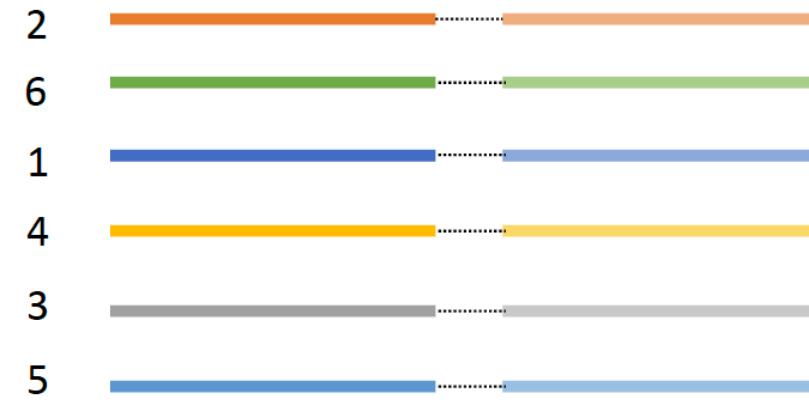
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Read compression

- For a typical 25x human dataset:
 - Uncompressed: 79 GB (1 byte/base)
 - Gzip: ~20 GB (2 bits/base) – still far from optimal
- Order of read pairs in FASTQ irrelevant – can this help?



Original order in FASTQ



New order (preserves read pairing
but pairs ordered arbitrarily)

Read compression results

Compressor	25x human
Uncompressed	79 GB
Gzip	~20 GB

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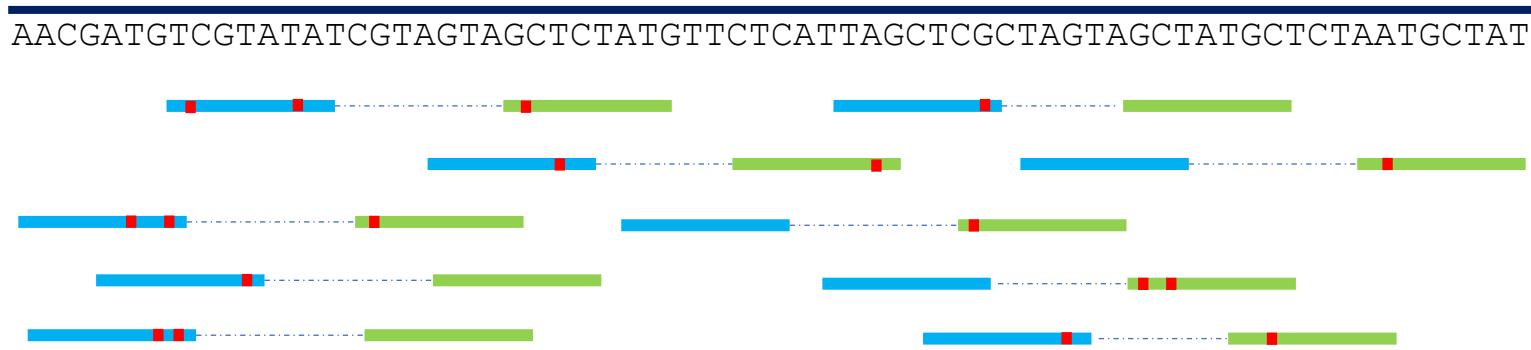
Read compression results

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SPRING (no reordering)	3 GB
SPRING (allow reordering)	2 GB

Read compression results

Compressor	25x human	100x human
Uncompressed	79 GB	319 GB
Gzip	~20 GB	~80 GB
FaStore (allow reordering)	6 GB	13.7 GB
SPRING (no reordering)	3 GB	10 GB
SPRING (allow reordering)	2 GB	5.7 GB

Key idea



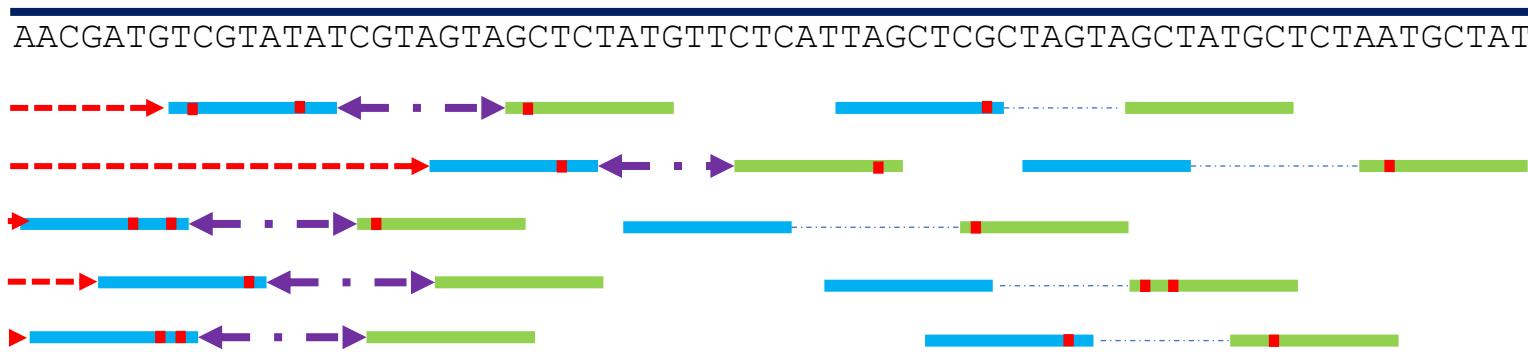
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 - Store genome

Key idea



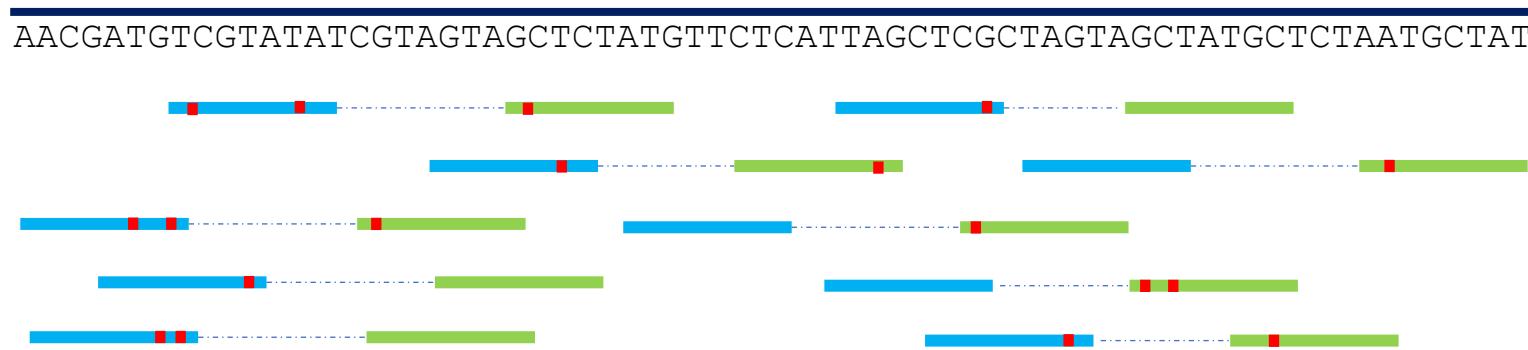
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 - **Store noise in reads**

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- Storing reads equivalent to
 - Store genome
 - Store read positions in genome (+ gap between paired reads)
 - Store noise in reads
- Entropy calculations show this outperforms previous compressors

Key idea

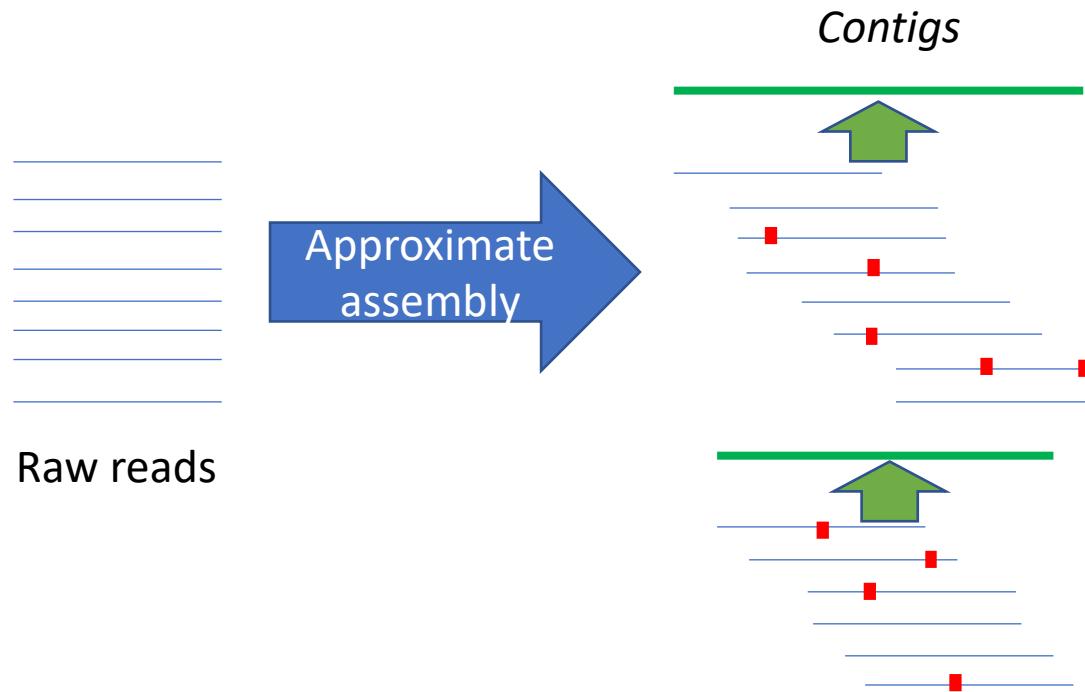
- But... How to get the genome from the reads?
- Genome assembly too expensive - big challenges:
 - resolve repeats
 - get very long pieces of genome from shorter assemblies
- Solution: Don't need perfect assembly for compression!

SPRING workflow

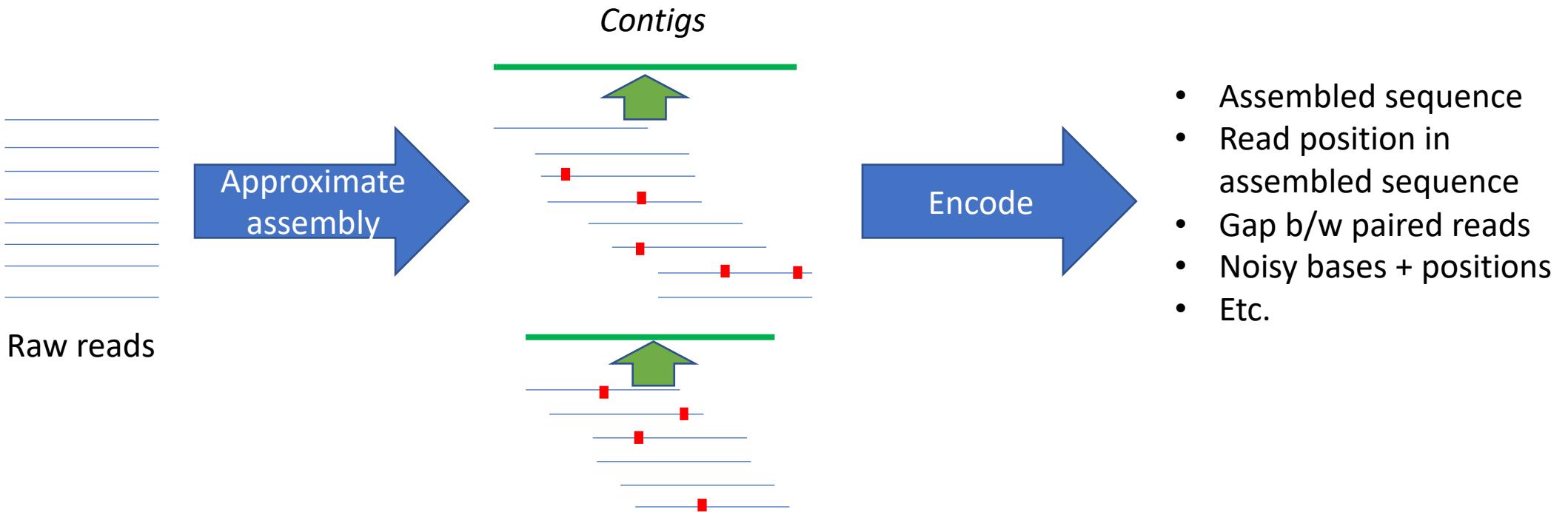


Raw reads

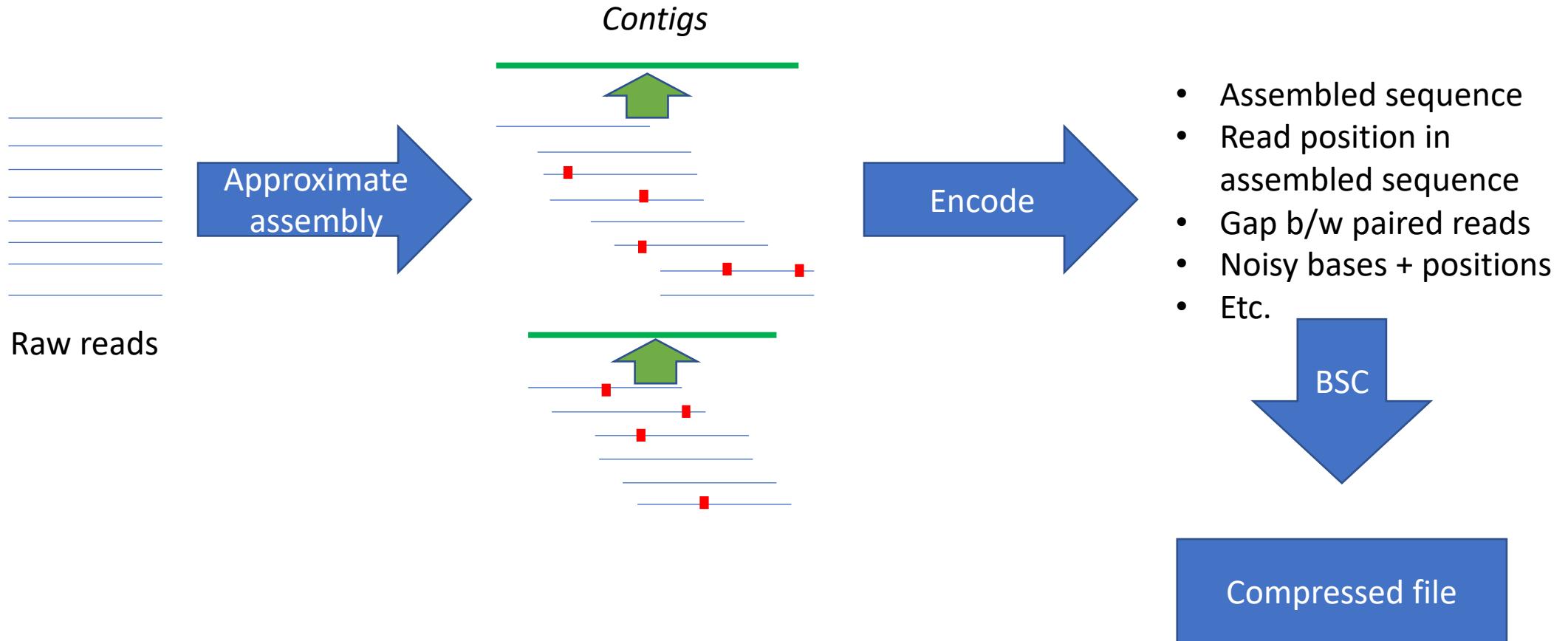
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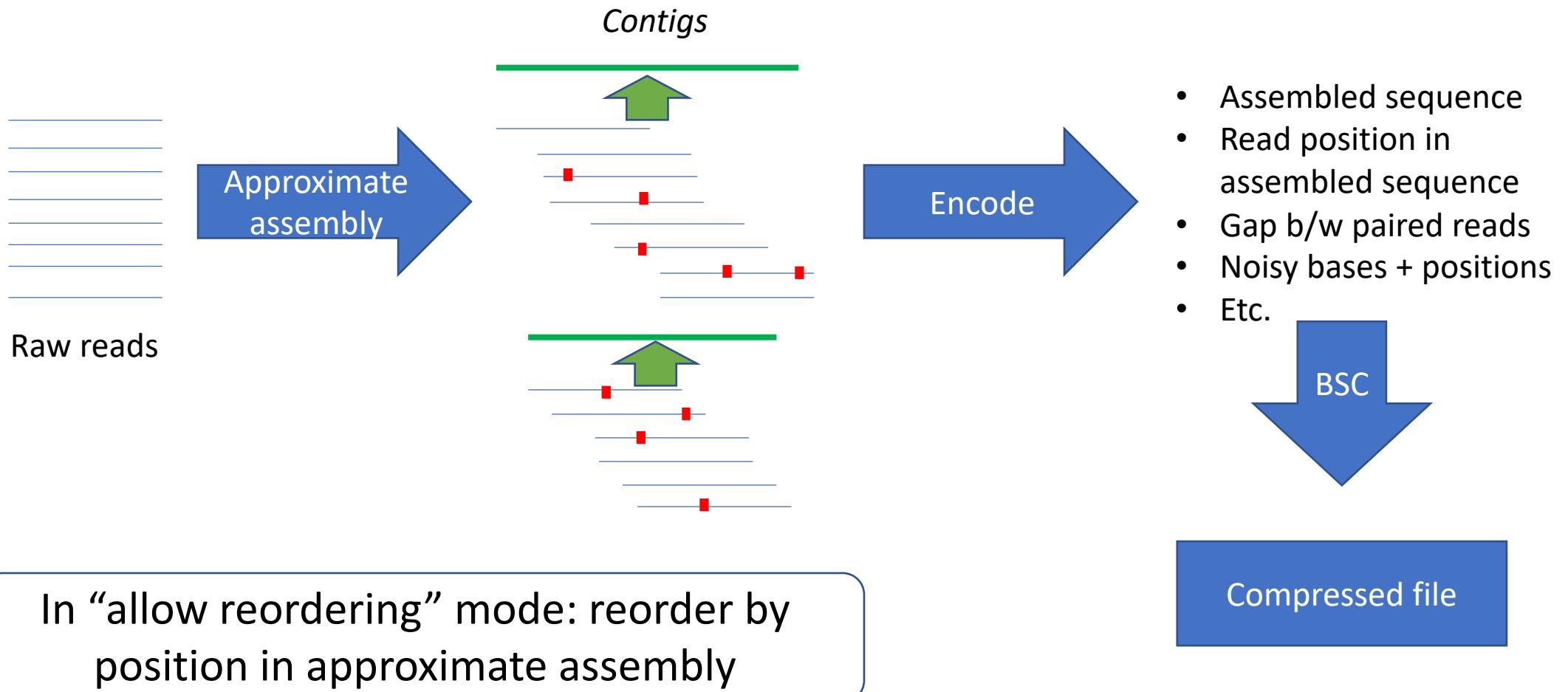


SPRING workflow



<https://github.com/IlyaGrebnev/libbsc>

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 - Index match found but Hamming distance too large → shift search substring by one

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 - Next read found!
- Repeat process with the new read.
- If no match found at any shift, pick arbitrary remaining read & start new *contig*

Consensus + encoding stage (simplified)

	<i>pos</i>	<i>noise</i>	<i>noisepos</i>		
ACTGCT G GCTGCTGC T AGC	0	GT	7, 16		7, 9
CT C CTAGCTGCTGCCAGCC	1	C	3	→	3
GCTAGCT A CTGCCAGCCTA	3	A	8	Delta encoding	8
GCT C GCT A CTG T C GCCTA	3	CATC	4, 8, 12, 14		4, 4, 4, 2
Majority					
ACTGCTAGCTGCTGC C AGCCTA		→	<i>seq</i>		
(Reference Sequence)					

Some technical details

- Hash 2 substrings per read to improve recall rate
- Handle reverse complement reads by searching both orientations
- Specialized hash table structure (BBHash) to reduce memory usage
 - Utilize fact that all keys are known in advance
- Parallelized – each thread works on a different contig
- For reads that are left out in assembly step – try to realign with less strict threshold after consensus
- Several other heuristics to increase speed without sacrificing compression

Quality and read identifier compression

Quality and read identifier compression

- Quality – use general purpose compressor BSC (optionally apply quantization)
- Read identifier – split into tokens and use arithmetic coding [1]

1. Bonfield, James K., and Matthew V. Mahoney. "Compression of FASTQ and SAM format sequencing data." *PLoS one* 8.3 (2013): e59190.

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Dataset	Reads	Quality	Read identifier
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All human datasets. Sizes in GB.

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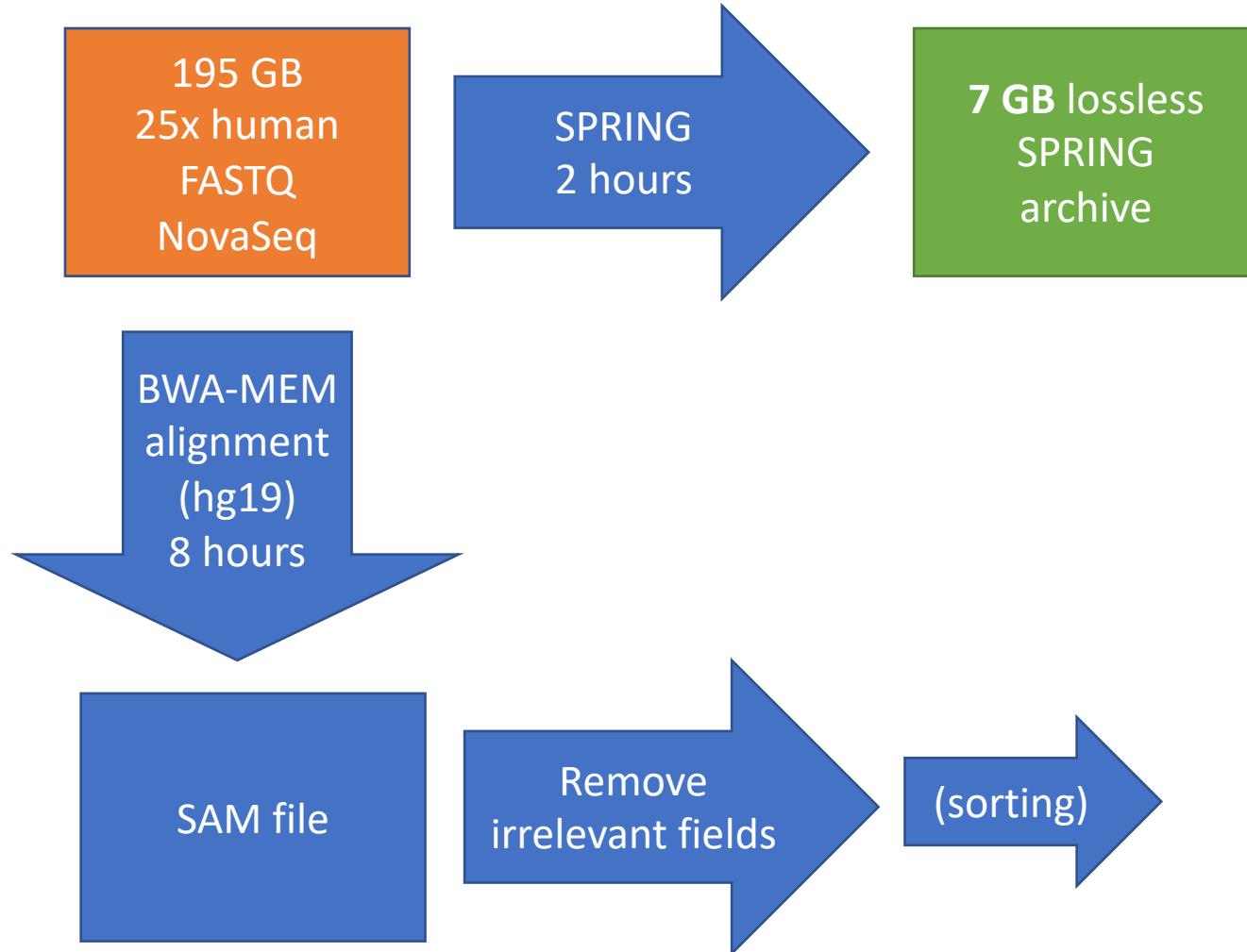
SPRING vs. reference-based compression

195 GB
25x human
FASTQ
NovaSeq

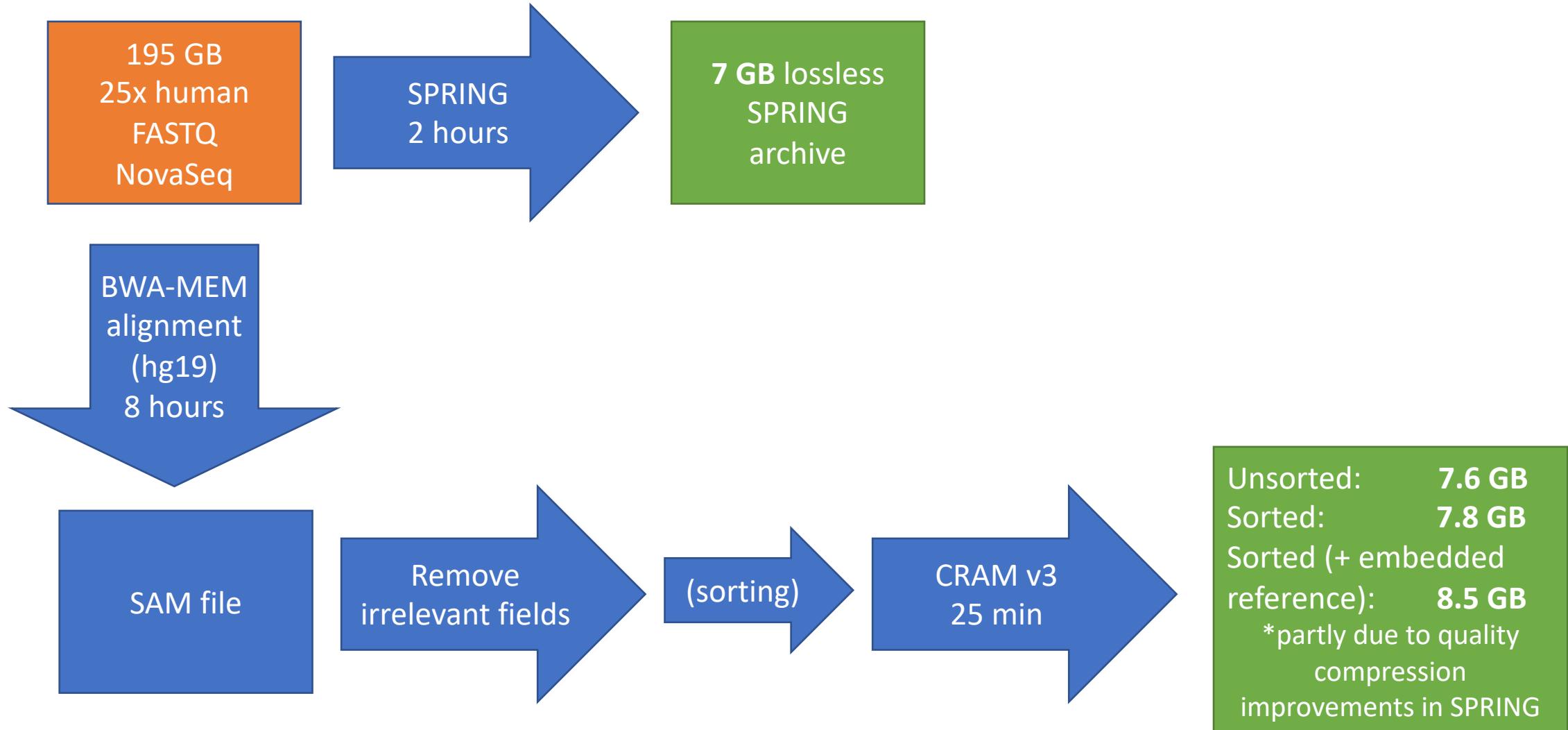
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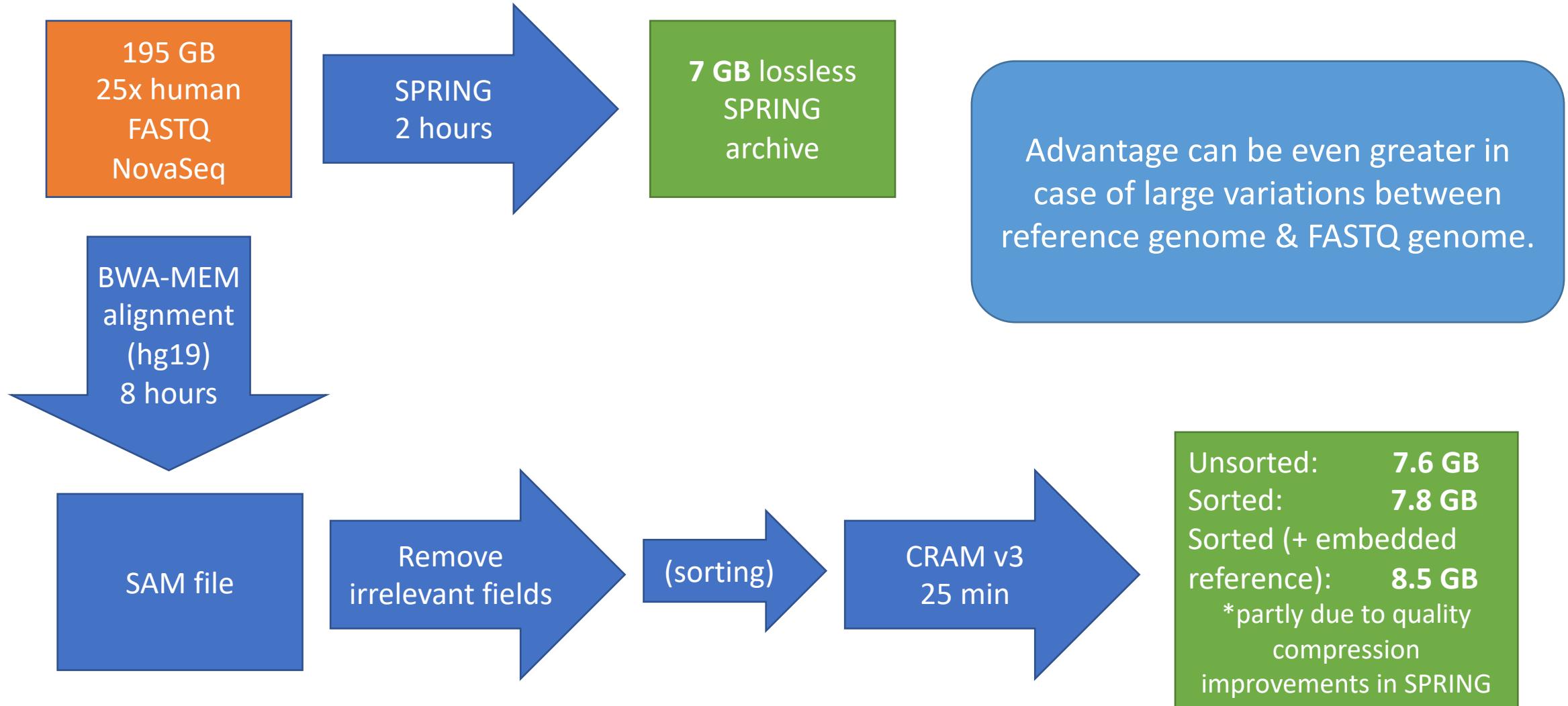
SPRING vs. reference-based compression



SPRING vs. reference-based compression



SPRING vs. reference-based compression



Other approaches for FASTQ compression

- gzip/bzip2
- Context-based arithmetic coding: DSRC 2, Fqzcomp, Quip
- Assembly based: Leon, Quip, Assembletrie
- Reordering based:
 - Reordering based on substrings/minimizers: Orcom, Mince, FaStore, SCALCE
 - BWT-based reordering: BEETL

Numanagić, Ibrahim, et al. "Comparison of high-throughput sequencing data compression tools." *Nature Methods* 13.12 (2016): 1005.

Hernaez, Mikel, et al. "Genomic Data Compression." *Annual Review of Biomedical Data Science* 2 (2019).

Recent FASTQ compressors: FQSqueezer

- FQSqueezer [2]: Adapt general-purpose compressors such as prediction by partial matching (PPM) and dynamic Markov coding (DMC) to read compression
 - 10-30% improvement over SPRING for bacterial datasets
- But requires significantly more time and memory than SPRING
 - Not tested on moderate to high coverage human datasets

1. Deorowicz, Sebastian. "FQSqueezer: k-mer-based compression of sequencing data." *bioRxiv* (2019): 559807.

Recent FASTQ compressors: PgRC

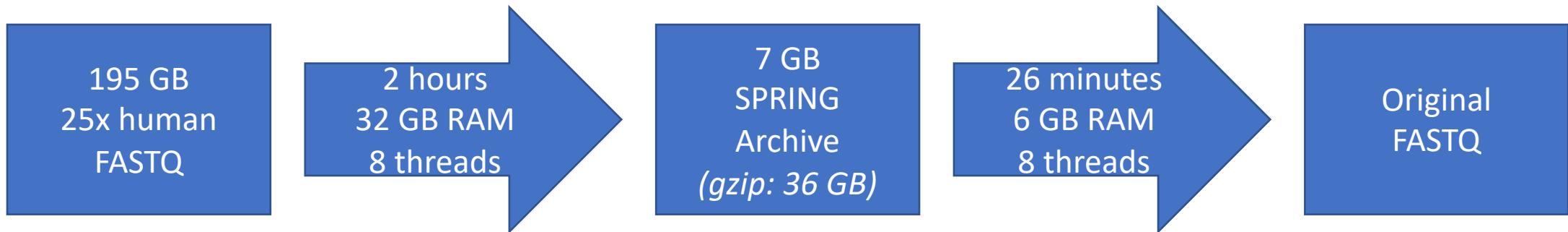
- Pseudogenome-based Read Compressor
- Similar framework as SPRING, but different “assembly” algorithm
- ~10-15% better compression than SPRING
- ~40% slower than SPRING
- Currently only supports read sequences

Kowalski, Tomasz, and Szymon Piotr Grabowski. "Engineering the Compression of Sequencing Reads." bioRxiv (2020).

Recent FASTQ compressors: alignment-based

- Setting: reference of same/related species available
- Approach:
 - Perform quick, inaccurate alignment
 - Much faster than bwa mem or minimap
 - Perform local assembly (optional)
 - Perform reference-based encoding
- Results:
 - Much better computational performance than SPRING
 - Compression generally a bit worse (even worse when reference is included in size)
- References:
 - Jammula, Nagakishore, and Srinivas Aluru. "ParRefCom: Parallel Reference-based Compression of Paired-end Genomics Read Datasets." *Proceedings of the 10th ACM International Conference on Bioinformatics, Computational Biology and Health Informatics*. 2019.
 - Enancio (acquired by Illumina)

SPRING as a practical tool



- Easy to use with support for:
 - Lossless and lossy modes
 - Variable length reads, long reads, etc.
 - Compressed in blocks to allow partial/streaming decompression
 - Scalable to large datasets
 - Gzipped I/O
- Github: <https://github.com/shubhamchandak94/SPRING/>

References

- Shubham Chandak, Kedar Tatwawadi, Tsachy Weissman; Compression of genomic sequencing reads via hash-based reordering: algorithm and analysis, *Bioinformatics*, Volume 34, Issue 4, 15 February 2018, Pages 558–567
- Shubham Chandak, Kedar Tatwawadi, Idoia Ochoa, Mikel Hernaez, Tsachy Weissman; SPRING: a next-generation compressor for FASTQ data, *Bioinformatics*, bty1015
- SPRING code: <https://github.com/shubhamchandak94/Spring>
- genie (open-source MPEG-G codec – *under development*): <https://github.com/mitogen/genie>



Preliminary work: Noisy long read compression

- Joint work with Yifan Zhu
- Building a compressor for noisy long reads (e.g., ONT, PacBio)
- Very similar approach as SPRING
 - Much more challenging due to higher error rates (5-10%), including insertion and deletion errors
- Borrow ideas from assemblers but use approximations/heuristics to achieve >100x speedup
- Multi-stage filtering of reads: kmer-based search -> proper alignment
- Preliminary results encouraging, but need to scale up

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Impact of lossy compression of nanopore raw signal data on basecalling and consensus accuracy

Shubham Chandak *, Kedar Tatwawadi, Srivatsan Sridhar and Tsachy Weissman *

Background

- (Oxford) nanopore sequencing gaining popularity
 - Long reads -> better assembly , structural variant discovery
 - Sequence native DNA and detect modifications
 - Real-time & portable
- Sequencer generates raw current signal that is decoded to base sequence
 - Often need to retain raw intermediate data for (re)analysis
 - Noisy – lossless compression difficult
 - Typical human whole genome exp: terabytes of raw data – 10x more than base sequence

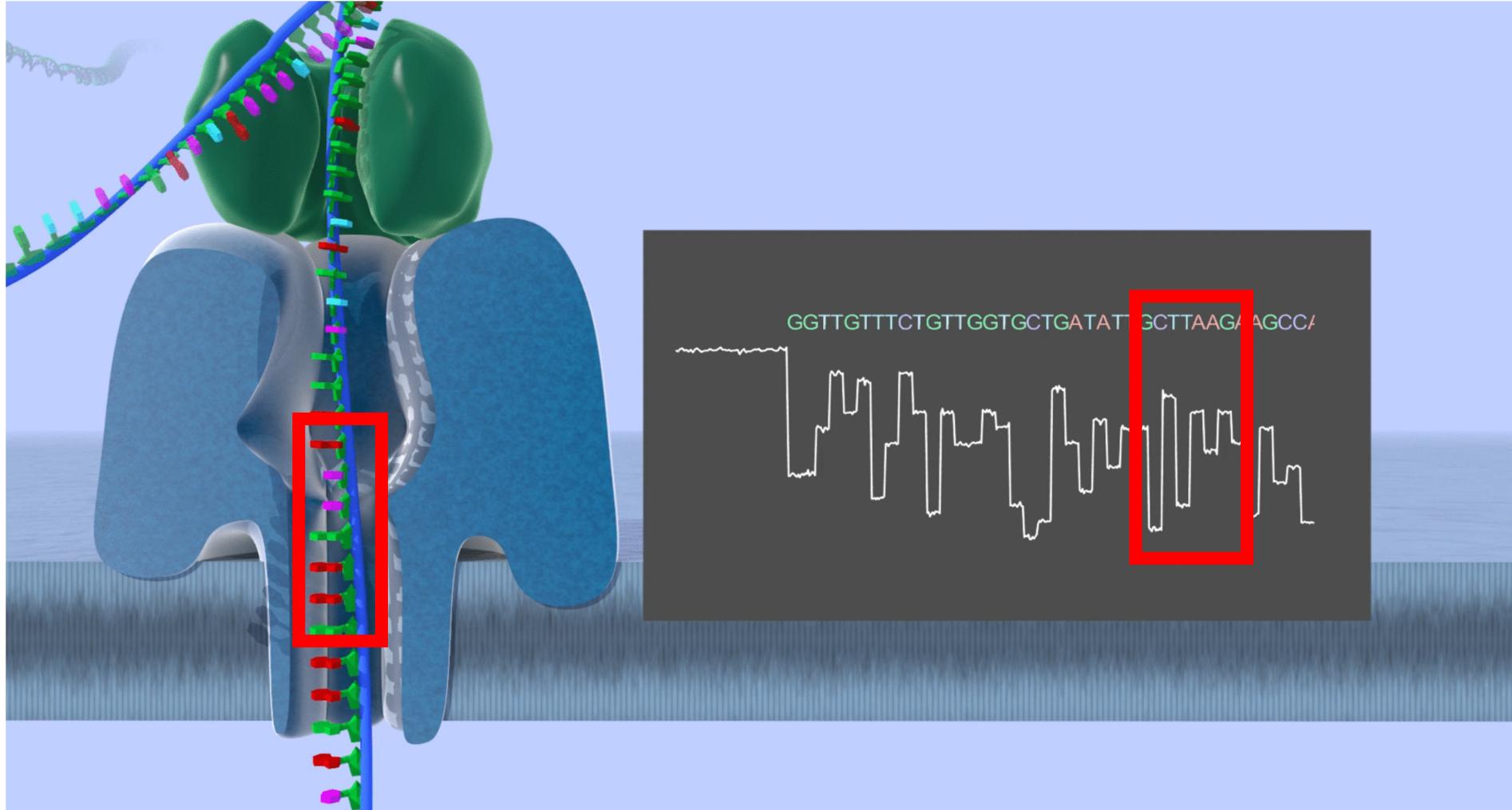
Oxford Nanopore Sequencing

- Nanopore sequencing: portable, real time



<https://directorsblog.nih.gov/2018/02/06/sequencing-human-genome-with-pocket-sized-nanopore-device/>

Nanopore Sequencing Process

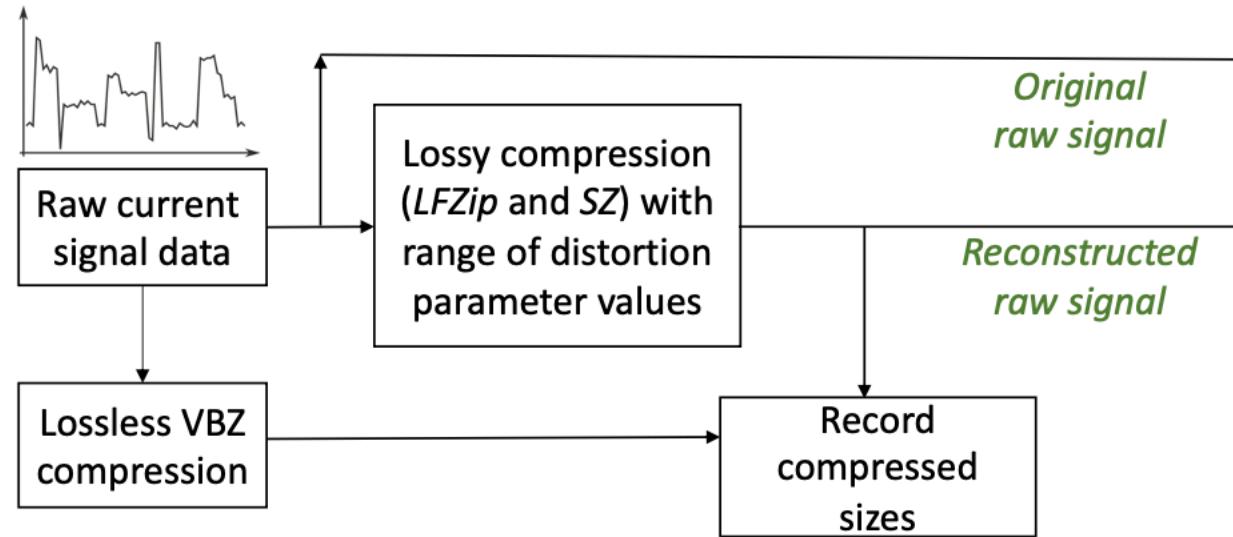


Source: <https://youtu.be/E9-Rm5AoZGw>

Raw data format

- HDF5 file (“.fast5”) with signal stored as series of 16-bit integers
- 5-15 current samples per base -> ~18 bytes/base (uncompressed)
- VBZ: state-of-the-art lossless compressor
 - Variable byte integer encoding followed by zstd
 - 60% size reduction over uncompressed representation
 - Still require 1 TB for 30x human whole genome data

Evaluation pipeline: part 1

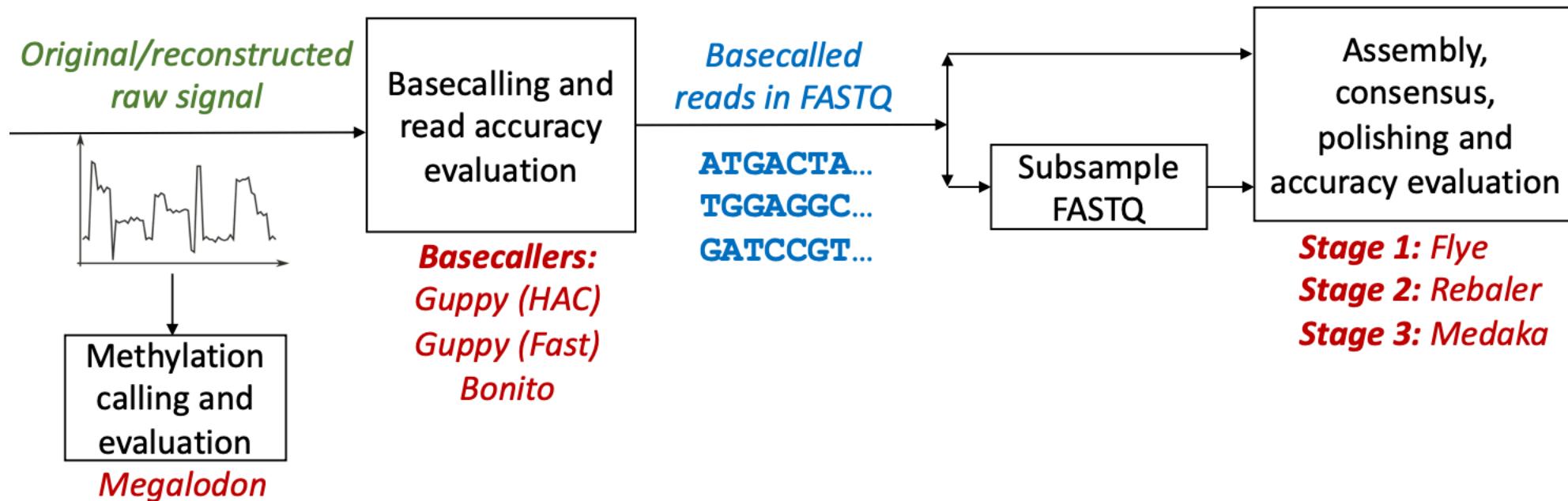


(a) Lossless and lossy compression of raw signal data

Note on lossy time-series compressors LFZip and SZ:

- Guarantee reconstruction at each time step is within ϵ of true value (ϵ user defined parameter)
- Rely on simple prediction/quantization followed by entropy coding (gzip/bzip2...)
- LFZip simply performs uniform scalar quantization (“rounding”) followed by entropy coding

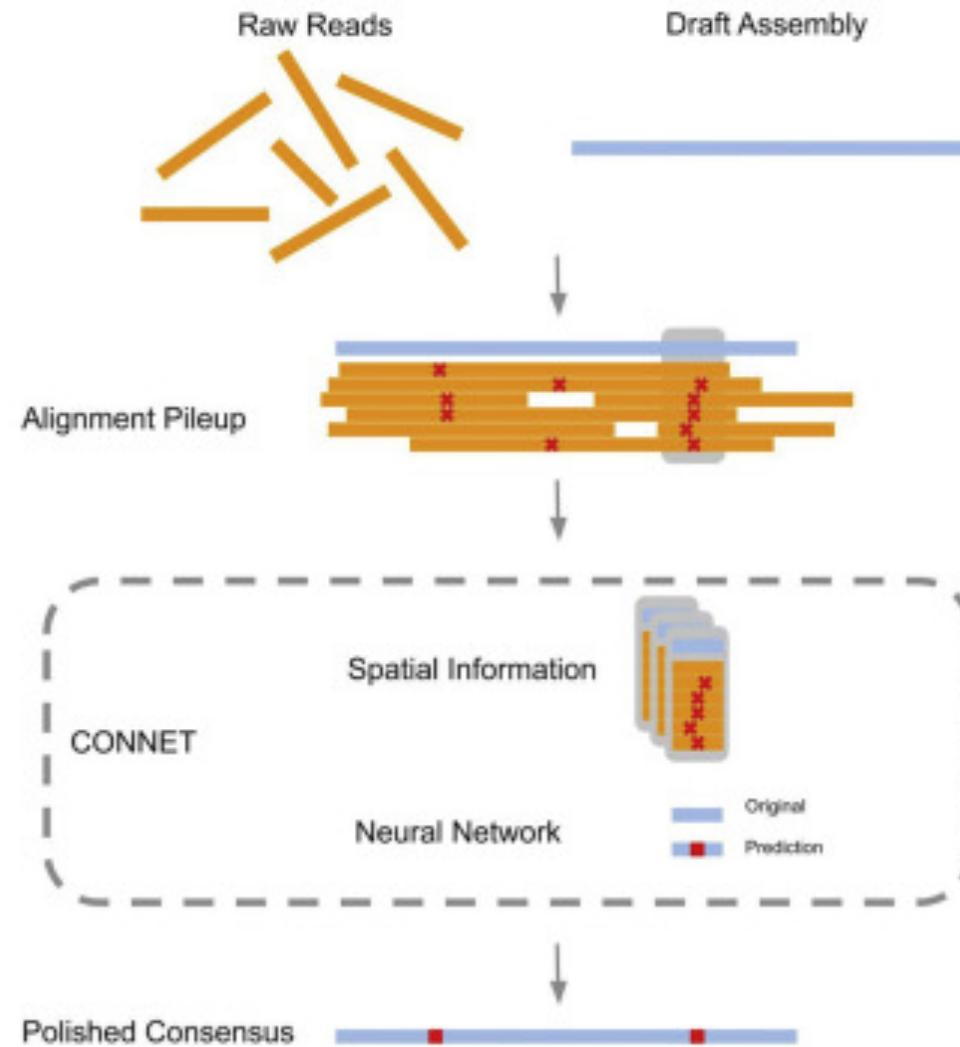
Evaluation pipeline: part 2



(b) Basecalling, consensus and methylation calling accuracy analysis

Note: Attempt to “future-proof” by testing various tools/use cases

Why consensus accuracy might differ from basecall (read) level accuracy:
systematic vs. random errors



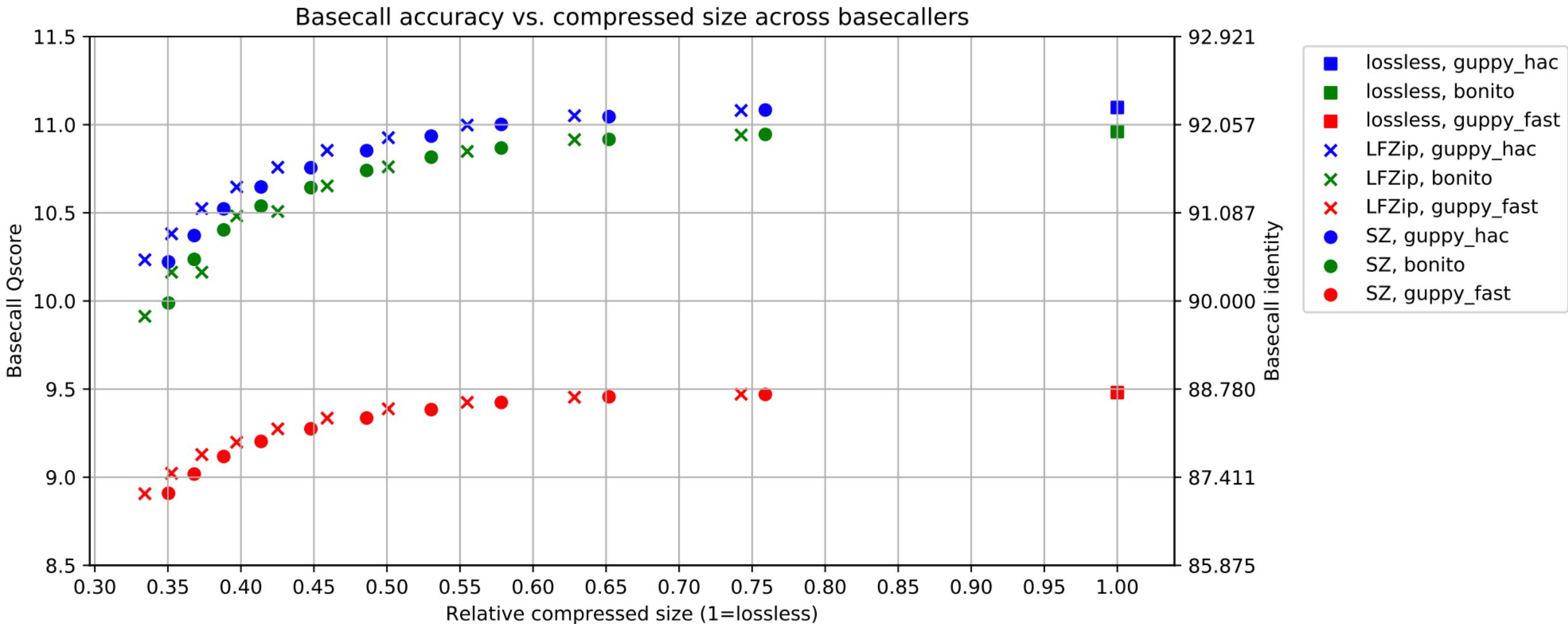
Source: <https://www.sciencedirect.com/science/article/pii/S2589004220303138>

Datasets

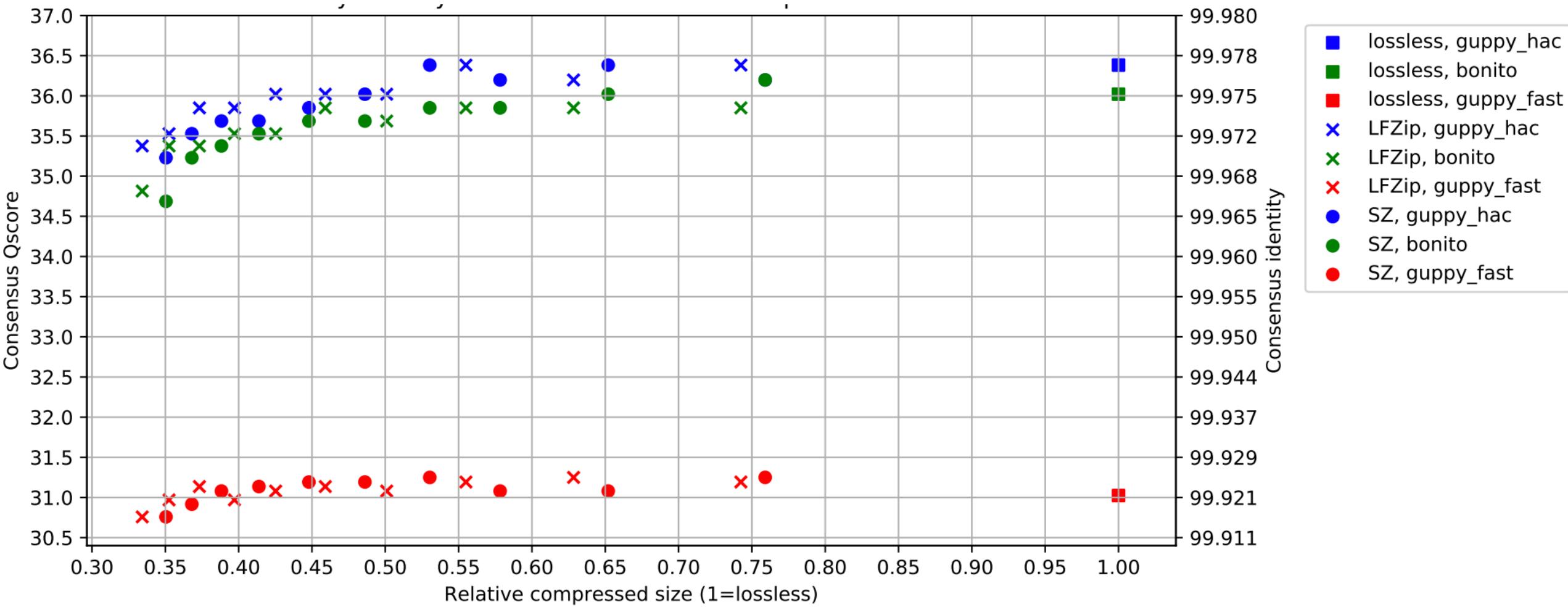
- Human and bacterial datasets for basecall accuracy
- Bacterial datasets for consensus accuracy
- Human dataset with bisulfite benchmark for methylation accuracy

Species	Sample	Genome size (bp)	GC-content	Flowcell type	Read count	Read length N50 (bp)	Approx. depth
<i>Staphylococcus aureus</i>	CAS38_02	2.9×10^6	32.8%	R9.4.1	11,047	24,666	83x
<i>Klebsiella pneumoniae</i>	INF032	5.1×10^6	57.6%	R9.4	15,154	37,181	108x
<i>Escherichia coli</i>	K-12 MG1655	4.6×10^6	50.8%	R10.3	92,000	7,431	128x
<i>Homo sapiens</i>	NA12878	3.1×10^9	40.9%	R9.4	128,314	11,404	0.29x

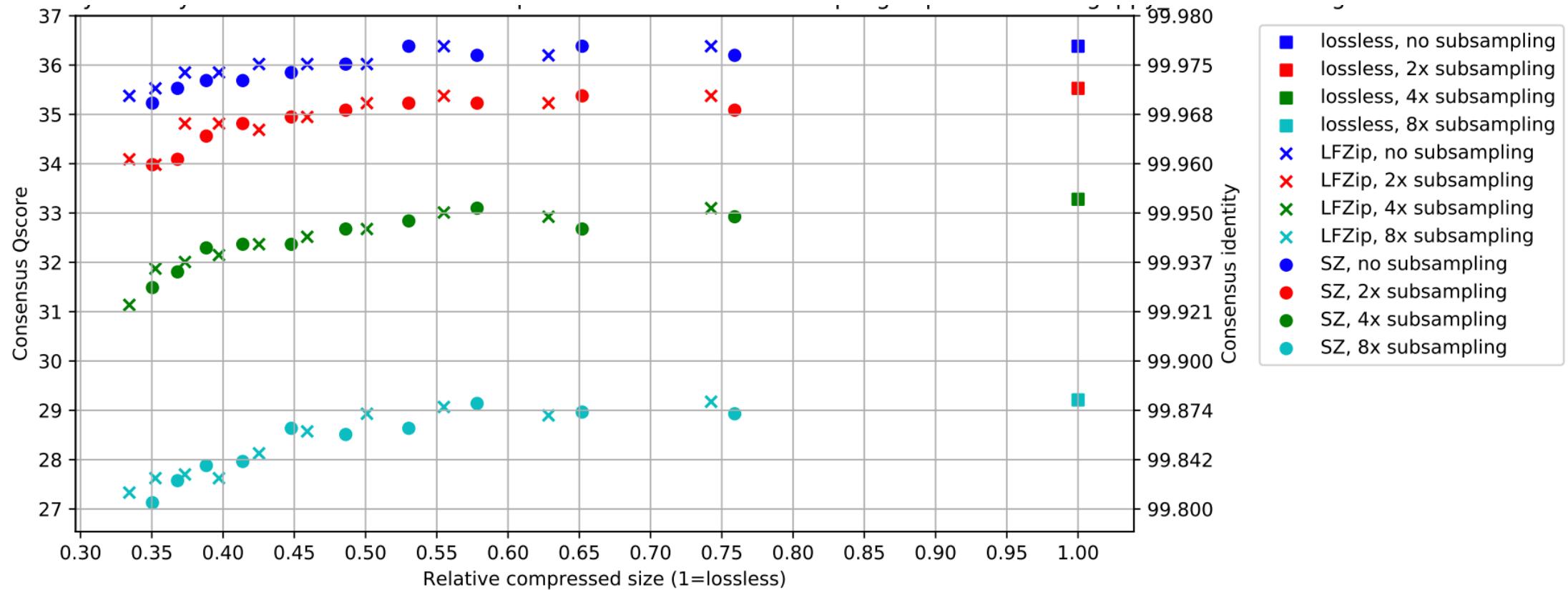
Basecall accuracy



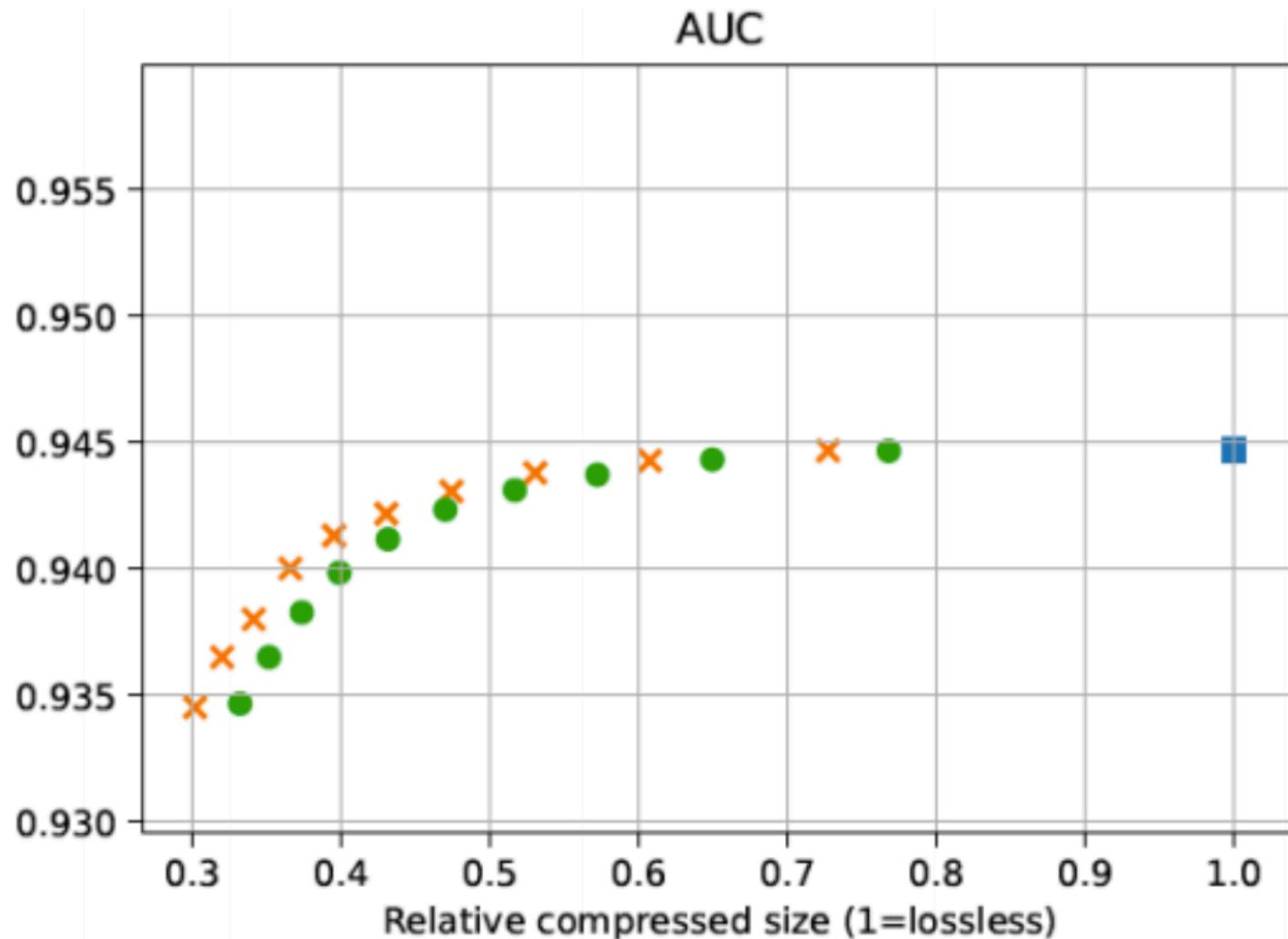
Consensus accuracy



Subsampling experiments



Per-read methylation calling accuracy



Summary

- Lossy compression achieves 35-50% reduction over current best lossless compression:
 - <0.2% reduction in basecall (read) accuracy
 - <0.002% reduction in consensus accuracy (even better for high coverage)
- Highly practical – LFZip simply reduces the data resolution!
- Can be adopted at the nanopore sequencer device itself
 - Similar to Illumina reducing quality score resolution from 40 to 4.
- Future work:
 - Specialized lossy compressors for this data
 - Further evaluation on human data with improved benchmark datasets

Availability

- Biorxiv:
<https://www.biorxiv.org/content/10.1101/2020.04.19.049262v3>
- Evaluation scripts, data, plots:
https://github.com/shubhamchandak94/lossy_compression_evaluation

Thank you!