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Subject: CNS Lab Assignment 03 Writeup and performance screenshots

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Subject: Computer Networks & Security
Assignment : 03

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Title: Error detection

Problem Statement:

Write a program for error detection & correction for 7/8 bits ASCII codes using Hamming codes & CRC

Requirements

IDE: Eclipse (version: 2020-release)

Compiler: gcc (version: 6.3.0)

Hardware: Intel i5-8265U CPU @ 1.60 GHz
8GB RAM & 512GB SSD.

Description:

Data Link Layer:

- In OSI model, data link layer is the 4th layer from top & 2nd from bottom.
- Data link layer gets the raw data from physical layer, try to make physical link reliable.
- It attempts to detect & recover the data that it gets from physical layer.
- Data link layer protocol defines the format of the packet exchanged across nodes, error detection, retransmission, flow control & random access.

Services provided by data link layer:

↳ Framing & Link Access:

↳ data link layer takes packets from network layer & encapsulate them as frames

1) It sends each frame bit-by-bit on the hardware.

1) At receiver's end, it assembles bit into frames.

2) Addressing

→ provides layer-2 hardware addressing mechanism

1) Hardware address is unique & provided at the time of manufacturing

3) Synchronization: machines should be synchronized while sending message on link.

4) Error Control: Errors are detected & attempt is made to recover actual bits. provides error reporting mechanism to sender.

5) Flow control: machines on same link may have different speed/capacity. Data link layer ensures flow control.

6) Multi-Access: To avoid/reduce collisions data link layer provides mechanism such as CSMA/CD to equip capability of accessing a shared media among multiple systems.

Types of Errors:

There are three types of errors.

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Error Type	Diagram	Description
Single bit Error		<p>Only one bit of given data-unit is corrupted.</p> <p>Least-likely in serial data transmission.</p> <p>happen in parallel data transmission</p>
Burst Error		<p>Two or more bits in data unit are corrupted</p> <p>length is measured from first corrupted bit to last.</p> <p>generally happen in serial data transmission.</p>

Methods of error detection & correction:

error detection: involves checking wheather any error has occured or not.

error correction: involves ascertaining the exact number of bits that has been corrupted & location of corrupted bits.

Error detection techniques:

Name	Description
1) parity check.	<p>→ done by adding extra bit called parity bit.</p> <p>ii) there are two types even parity & odd parity.</p> <p>iii) Sender counts no. of ones & make parity either even/odd by adding extra 1.</p> <p>iv) On receiver end, is parity is same as sender then message is accepted.</p>
2) checksum	<p>i) Data is divided into fixed sized frames or segments.</p> <p>ii) Sender: adds the segment using 1's complement arithmetic to get sum. Then it complements sum to get checksum & sends it along data frames.</p> <p>iii) Receiver: adds incoming segment along with the checksum using 1's complement & then complements it.</p> <p>iv) IF result is zero, data is accepted.</p>
3) Cyclic Redundancy Check (CRC)	<p>i) It involves binary division of data bits by predetermined divisor which is generated using polynomials.</p> <p>ii) At sender side after division, remainder is appended at the end of data segment.</p> <p>iii) At receiver end, incoming data unit is divided by divisor. IF there is 0 remainder data segment gets accepted.</p>

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Error Correction Techniques

Technique	Description
1) Backward Error Correction	<p>1) Receiver request the sender to retransmit frame if he detects error.</p> <p>2) It is simple technique which is efficient when retransmission is not expensive.</p>
2) Forward Error Detection Correction	<p>1) Receiver executes error correcting code if he detects error.</p> <p>2) This saves bandwidth requirement for retransmission.</p> <p>3) If there are too many errors retransmission is needed.</p>

Error correction codes:

- 1) Hamming Codes
- 2) Binary convolution Code
- 3) Reed-Solomon code
- 4) Low-density parity-check code.

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here,

r = redundant bits

m = message length

3) Parity Bits.

Parity bits are appended to make parity of 1's either even or odd. For our analysis let's consider even parity from now.

Note that:

even parity: no. of 1's should be even. For odd 1's extra 1 should be added.

odd parity: no. of 1's should be odd. For even 1's extra 1 should be added.

Algorithm:

1) Calculate the value of

r = redundant bits by formula $2^r \geq d+1$

& d = length of data word.

2) Effective length of codeword will be

$$n = r + d$$

3) where Let parity bits be $p_1, p_2, p_3, \dots, p_r$.

calculate p_1, p_2 as follows:

p_1 = consider all the bits positions whose binary representation includes 1 in least pos. significant. eg. (1, 3, 5, 7, 9, ...)

p_2 = consider all the bits positions whose binary representation includes 1 in second position from LSB.

$p_3 =$

4) As we are considering for even parity, set a parity bit to 1 if the total number of 1's

in the positions it checks is odd.

→ Place these redundant bits at the positions of powers of 2 i.e. 1, 2, 4, ... etc. other positions will be occupied by data bits.

At receivers end:

1) Calculate the values of P_1, P_2, \dots, P_z .

if they are of opposite parity errors exist.

2) Calculating the position of corrupted bit.
convert the value of $(P_z P_{z-1} \dots P_1)$ to decimal.

This gives the ^{position} ~~value~~ of corrupted bit.

example:

Even parity

consider the message 1101011

here $d=7$

using formula $2^z \geq d+1$

we get $z=4$.

hence our message length will be $z+d=11$

		1	1	0	1	0	1	0		
--	--	---	---	---	---	---	---	---	--	--

$P_1 \ P_2 \ d_3 \ P_4 \ d_5 \ d_6 \ d_7 \ P_8 \ d_9 \ d_{10} \ d_{11}$

consider:

$$P_1 = d_3 d_5 d_7 d_9 d_{11} = \underline{10100} = 1$$

$$P_2 = d_2 d_6 d_7 d_{10} d_{11} = \underline{01100} \text{ odd 1's} = 1$$

$$P_3 = d_1 d_4 d_6 d_9 d_{11} = \underline{11010} \text{ odd 1's} = 1$$

$$P_4 = d_5 d_6 d_7 = \underline{101} = 0$$

even 1's

$$P_8 = d_9 d_{10} d_{11} = 010 = 1$$

our codeword to send is:

1 1 1 1 0 1 0 1 1 0 1 0

On receiver's end:

let we received: $\begin{matrix} p_1 & p_2 & d_3 & d_4 & d_5 & d_6 & d_7 & d_8 & p_9 & p_{10} & d_{11} \\ 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \end{matrix}$

here $p_1 = \underbrace{111100}_{\text{even 1's}} = 0$

$p_2 = 110110 = 0$

$p_9 = 0101 = 0$

$p_8 = 1010 = 0$

As all parity bits satisfy parity condition, the received message is accepted.

let we received $\begin{matrix} p_1 & p_2 & d_3 & d_4 & d_5 & d_6 & d_7 & d_8 & p_9 & p_{10} & d_{11} \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \end{matrix}$

here $p_1 = 110100 = 1$

$p_2 = 110110 = 0$

$p_9 = 0001 = 1$

$p_8 = 1010 = 0$

As bits p_1 & p_9 violates parity condition, there must be an error in received message

Error position:

consider decimal representation of $p_8 p_4 p_1$:

i.e. 0101 which is 5.

Hence bit-5 is corrupted.

Correct msg. should be 11101011010. ✓

Conclusion

In this assignment, we have learned about various methods of error detection & correction in data link layer along with their types. Also implemented program for Hamming code error correction & detection with the help of our lab guide.


```
**Welcome to Hamming Code Error Detector and Corrector created by Shubham.**
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Enter the message: pict
Message Transfer is started...
```

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----->Tranferring: p

====This is senders end
Choose Parity:
    1. for Odd Parity.
    2. for Even Parity.
: 1
Encoding Your Message with Odd Parity...
Number of redundant bits used are: 4
Encoded Message is:
1  1  0  1  0  0  0  0  1  1  1
p1 p2 d3 p4 d5 d6 d7 p8 d9 d10 d11
Encoded Successfully.

Your Message is Being Transferred to Receiver. Plz Wait..
====This is Medium====
Is any bit corrupted (if yes enter pos else enter -1)
: -1

This is receiver End====
Calculated value of parity bits:
0  0  0  0
p1 p2 p4 p8
No Errors detected. Suceessfully Accepted.
Received Character: p
```

Screenshot-1: Odd Parity with no error

```

----->Tranferring: i

====This is senders end
Choose Parity:
    1. for Odd Parity.
    2. for Even Parity.
: 1
Encoding Your Message with Odd Parity...
Number of redundant bits used are: 4
Encoded Message is:
0  1  1  0  0  0  1  1  0  1  1
p1 p2 d3 p4 d5 d6 d7 p8 d9 d10 d11
Encoded Successfully.

Your Message is Being Transferred to Receiver. Plz Wait..
====This is Medium====
Is any bit corrupted (if yes enter pos else enter -1)
: 7

This is receiver End====
Calculated value of parity bits:
1  1  1  0
p1 p2 p4 p8
Corrupted bit found at pos: 7
Correcting the corrupted bit...
Received Character: i

```

Screenshot-2: Odd Parity with error


```

----->Tranferring: t

====This is senders end
Choose Parity:
    1. for Odd Parity.
    2. for Even Parity.
: 2
Encoding Your Message with Even Parity...
Number of redundant bits used are: 4
Encoded Message is:
0  1  0  1  0  1  0  1  1  1  1
p1 p2 d3 p4 d5 d6 d7 p8 d9 d10 d11
Encoded Successfully.

Your Message is Being Transferred to Receiver. Plz Wait..
====This is Medium====
Is any bit corrupted (if yes enter pos else enter -1)
: -1

This is receiver End====
Calculated value of parity bits:
0  0  0  0
p1 p2 p4 p8
No Errors detected. Suceessfully Accepted.
Received Character: t

```

Screenshot-3: Even Parity with no error

```

----->Tranferring: c

====This is senders end
Choose Parity:
    1. for Odd Parity.
    2. for Even Parity.
: 2
Encoding Your Message with Even Parity...
Number of redundant bits used are: 4
Encoded Message is:
1  1  1  1  1  0  0  0  0  1  1
p1 p2 d3 p4 d5 d6 d7 p8 d9 d10 d11
Encoded Successfully.

Your Message is Being Transferred to Receiver. Plz Wait..
====This is Medium====
Is any bit corrupted (if yes enter pos else enter -1)
: 10

This is receiver End====
Calculated value of parity bits:
0  1  0  1
p1 p2 p4 p8
Corrupted bit found at pos: 10
Correcting the corrupted bit...
Received Character: c

```

Screenshot-4: Even Parity with error

```

****Your Message is Suceessfully Decoded****
Decoded Message is: pict

Error Detection and Correction Process is completed.

Do you want to check again?
: 0
Thank You!

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Thank You!