

CS6124D: Topics in Programming Languages

Sums 20.03.2020

A *Sum* type denoted $T_1 + T_2$, describes a set of values, where each value is either of type T_1 or of type T_2 . Such a type is useful in situations where a collection of values of different types is required, but the collection allows only values of same type (as in the case of a *list*). A value of the *Sum* type $Nat + Bool$ can be either of type Nat or of type $Bool$. But now a question arises: is the value *true* to be typed as $Bool$ or $Nat + Bool$? To resolve this, tags are attached to terms of type *Sum*. We use two tags - *inl* attached to terms of type T_1 and *inr* to terms of type T_2 . The term *inl true* is of sum typ $Bool + T_2$, for some type T_2 .

New syntactic forms

$t ::= \dots$ *terms*
 inl t
 inr t
 case t of *inl* x \Rightarrow t | *inr* x \Rightarrow t

$v ::= \dots$ *values*
 inl v
 inr v

$T ::= \dots$ *types*
 T+T

Evaluation Rules

$$\frac{t_1 \rightarrow t'_1}{inl\ t_1 \rightarrow inl\ t'_1} \quad \text{E-INL}$$
$$\frac{t_1 \rightarrow t'_1}{inr\ t_1 \rightarrow inr\ t'_1} \quad \text{E-INR}$$
$$\frac{t_0 \rightarrow t'_0}{case\ t_0\ of\ inl\ x_1 \Rightarrow t_1 \mid inr\ x_2 \Rightarrow t_2 \rightarrow case\ t'_0\ of\ inl\ x_1 \Rightarrow t_1 \mid inr\ x_2 \Rightarrow t_2} \quad \text{E-CASE}$$
$$case\ (inl\ v_0)\ of\ inl\ x_1 \Rightarrow t_1 \mid inr\ x_2 \Rightarrow t_2 \rightarrow [x_1 \mapsto v_0]t_1 \quad \text{E-CASEINL}$$
$$case\ (inr\ v_0)\ of\ inl\ x_1 \Rightarrow t_1 \mid inr\ x_2 \Rightarrow t_2 \rightarrow [x_2 \mapsto v_0]t_2 \quad \text{E-CASEINR}$$

Typing Rules

$$\frac{\Gamma \vdash t_1 : T_1}{\Gamma \vdash inl\ t_1 : T_1 + T_2} \quad \text{T-INL}$$
$$\frac{\Gamma \vdash t_1 : T_2}{\Gamma \vdash inr\ t_1 : T_1 + T_2} \quad \text{T-INR}$$
$$\frac{\Gamma \vdash t_0 : T_1 + T_2 \quad \Gamma, x_1 : T_1 \vdash t_1 : T \quad \Gamma, x_2 : T_2 \vdash t_2 : T}{\Gamma \vdash case\ t_0\ of\ inl\ x_1 \Rightarrow t_1 \mid inr\ x_2 \Rightarrow t_2 : T} \quad \text{T-CASE}$$

Exercise: A list is a collection of items of the *same* type. Suppose we need to keep a list of courses where each element is a record type with certain set of attributes. But suppose we

defined different types of records for *Core* and *Elective* as *Elective* = { *CourseName:string*, *Credits:Nat*, *Cot:Bool* } and *Core* = { *CourseName:string*, *Credits:Nat* }. Now in order to maintain a list of courses, where each course can be either a *Core* or *Elective*, the Sum type *Core+Elective* will be helpful. Write concrete instances of values of type *Core+Elective*. Define an abstraction *isCotNeeded* that when applied on to a *Course* type term, returns *true* if *cot*(consent of teacher) is needed and *false* otherwise.