

Rain Water Trapped

02 December 2023 12:08 AM

Problem Description

Given a vector **A** of non-negative integers representing an elevation map where the width of each bar is 1, compute how much water it is able to trap after raining.

From < <https://www.scaler.com/academy/mentee-dashboard/class/90012/assignment/problems/47>>

Problem Constraints

$1 \leq |A| \leq 100000$

Input Format

First and only argument is the vector **A**

Output Format

Return one integer, the answer to the question

Input 1:

A = [0, 1, 0, 2]

Input 2:

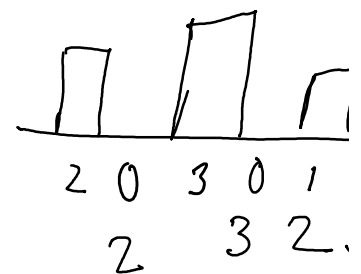
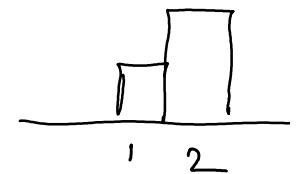
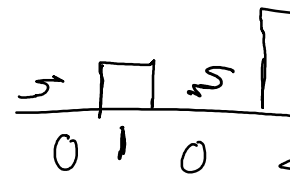
A = [1, 2]

Output 1:

1

Output 2:

0



Brute Force :-

iterate over each index of an vector .

Find out the left highest & right highest element

take the min from both & Subtract it from Current index & Calculate the sum of all values .

T.C. $O(n^2)$

S.C. $O(1)$

_____ x _____ x _____ x _____ x _____ x _____

Optimization :- Prefix Sum

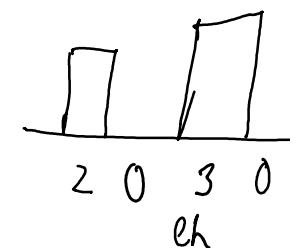
eg:- [2, 0, 3, 0, 1, 0, 1, 2, 4] = 13

lh = [4, 4, 4, 4, 4, 4, 4, 4, 4]

$\min(lh, rh) - A[i]$

If $(i > rh)$

lh = i



0 2 0 3 :

eg:- [1, 2, 0, 1, 0, 3, 1, 0, 1, 0, 4, 1] = 15

lh = 1 2 2 2 2 2 2 2 2 2 4 4



0 0 2 1 2 0 2 3



0 0 2 1 2 0 2 3

Solution : - Calculate the right highest element using pre
 & iterate over the input array calculate the value
 of element if it is greater than the current
 Replace it with current element taking the min
 & subtracting the current element.
 & summing all the values in the end.

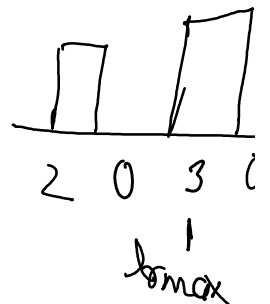
T.C. $O(n)$ - creating prefix right highest
 $O(n)$ - to find the answer.
 $= O(2n) = O(n)$

S.C. $O(n)$ - for prefix array.

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Optimization using pointers : -

eg:- $[2, 0, 3, 0, 1, 0, 1, 2, 4] = 13$



$l = 0$

$r = n - 1$

$lmax = 0$

$rmax = 0$

total = 0

```

while (l <= r) {
    if (lmax < rmax) {
        int sum = lmax - A.get(l);
        if (sum > 0) {
            total += sum;
        }
        lmax = A.get(l);
    }
    l++;
    else {
        int sum = rmax - A.get(r);
        if (sum > 0) {
            total += sum;
        }
        rmax = A.get(r);
    }
    r--;
}
return total;

```

l = 0
r = 8

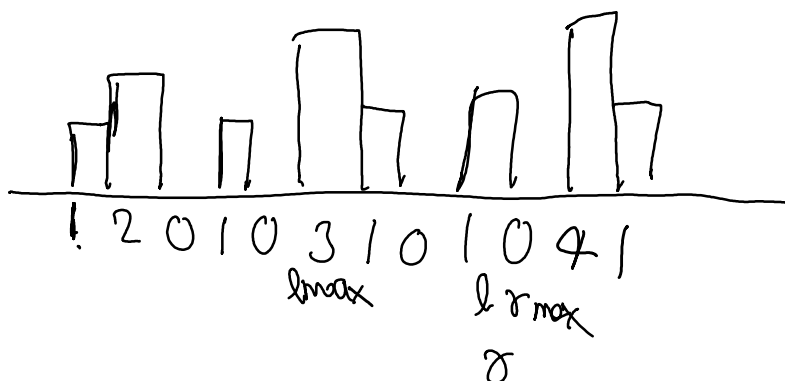
lmax = 3

rmax = 4

total = 2 + 3 + 2 +

return total;

2)



l = 0

r = 11

lmax = 3

rmax = 4

total = 2

1

>

Solution: The first part of the solution is to find the maximum value in the array.

algorithm :- Use two pointer approach use two variables
initialise l to be 0 & r to last index of vector
be 0 & iterate till $r \leq l$ If $l_{max} \leq r_{max}$ do
Subtracting current element from l_{max} If it is
add it in the output total, else assign the current
 l_{max} . & Increment the l

Whereas If $r_{max} > l_{max}$ Calculate the sum by
current element from l_{max} if it is greater
add it in the output total, else assign the
to the r_{max} & decrement r

$$T.C. = O(n)$$

$$S.C = O(1)$$