

Rain Water Trapped

02 December 2023 12:08 AM

Problem Description

Given a vector **A** of non-negative integers representing an elevation map where the width of each bar is 1, compute how much water it is able to trap after raining.

From < <https://www.scaler.com/academy/mentee-dashboard/class/90012/assignment/problems/47> >

Problem Constraints

$1 \leq |A| \leq 100000$

Input Format

First and only argument is the vector **A**

Output Format

Return one integer, the answer to the question

Input 1:

A = [0, 1, 0, 2]

Input 2:

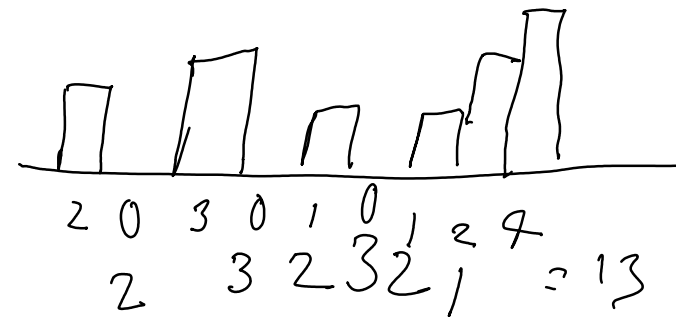
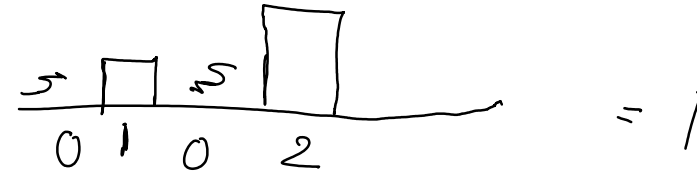
A = [1, 2]

Output 1:

1

Output 2:

0



Brute force :-

iterate over each index of an vector .

Find out the left highest & right highest element

take the min from both & Subtract it from current index element .

& Calculate the sum of all values .

T.C. $O(n^2)$

S.C. $O(1)$

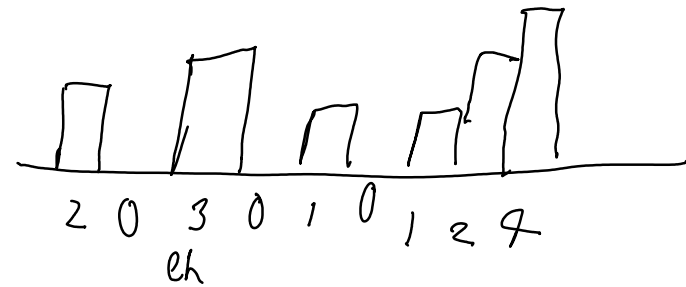
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Optimization :- Prefix Sum

eg:- $[2, 0, 3, 0, 1, 0, 1, 2, 4] = 13$

$sh = [4, 4, 4, 4, 4, 4, 4, 4, 4]$

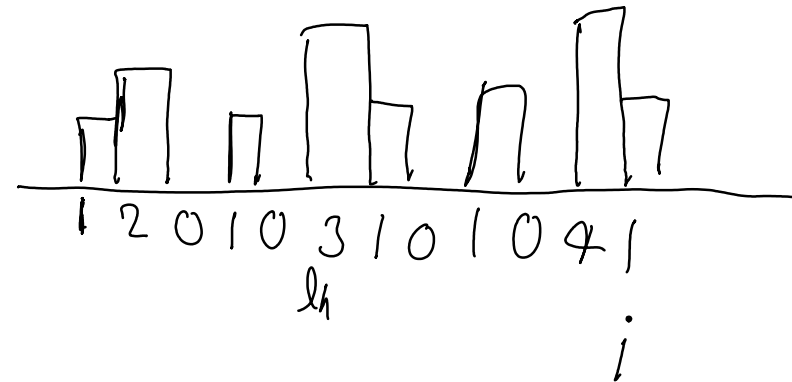
$\min(sh, sh) - A[i]$
 $\text{if } (i > sh)$
 $sh = i$



$0 2 0 3 2 3 2 1 0 = 13$

eg:- $[1, 2, 0, 1, 0, 3, 1, 0, 1, 0, 4, 1] = 15$

$sh = [4, 4, 4, 4, 4, 4, 4, 4, 4, 4, 4, 1]$



$0 0 2 1 2 0 2 3 2 3 0 0 = 15$

Solution :- Calculate the right highest element using prefix sum array.
 Iterate over the input array calculate the value by checking.
 sh element if it is greater than the current element
 Replace it with current element taking the minimum of $n-1$ and $n-2$

& Subtracting the current element,
 & summing all the values in the end,

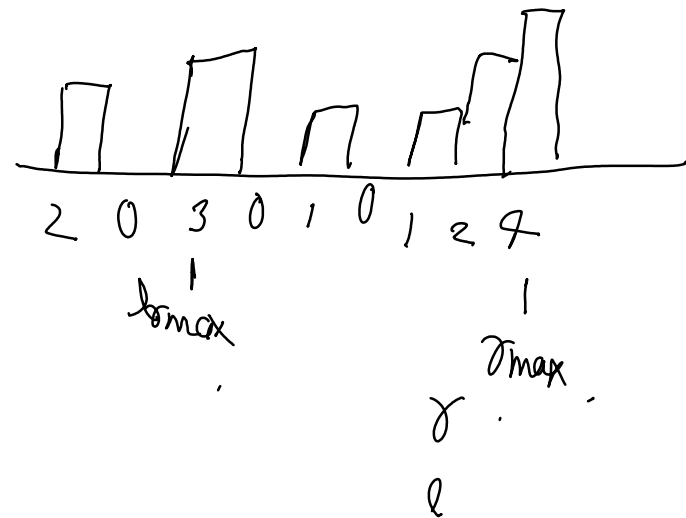
T.C. $O(n)$ — creating prefix & suffix array.
 $O(n)$ — to find the answer.
 $= O(2n) = O(n)$

S.C. $O(n)$ — for prefix array.

— x — x — x — x — x — x — x — x — x —

Optimization using pointers :-

eg:- $[2, 0, 3, 0, 1, 0, 1, 2, 4] = 13$



$l = 0$

$r = n - 1$

$lmax = 0$

$rmax = 0$

total = 0

$l = 0$
 $r = 8$

... ..

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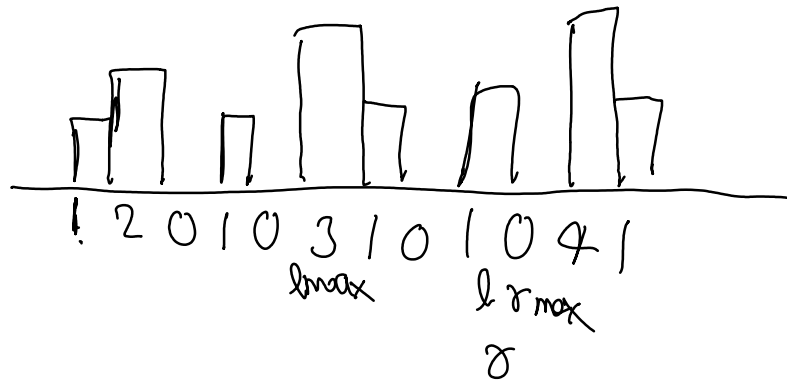
while (l <= r) {
    if (lmax < rmax) {
        int sum = lmax - A.get(l);
        if (sum > 0) {
            total += sum;
        } else {
            lmax = A.get(l);
        }
        l++;
    } else {
        int sum = rmax - A.get(r);
        if (sum > 0) {
            total += sum;
        } else {
            rmax = A.get(r);
        }
        r--;
    }
}
return total;

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lmax = 3

rmax = 4

total = 2 + 3 + 2 + 3 + 2 + 1 = 13



$$l = 0$$

$$r = 11$$

$$l_{max} = 3$$

$$r_{max} = 4$$

$$\text{total} = 2 + 1 + 2 + 2 + 3 + 2 + 3$$

$$= 15$$

Solution :- Use two pointer approach use four variables l , r , l_{max} & r_{max} initialize l to be 0 & r to last index of vector & l_{max} & r_{max} to be 0 & iterate till $r \leq l$ If $l_{max} \leq r_{max}$ calculate the sum by subtracting current element from l_{max} If it is greater than zero add it in the output total, else assign the current element to the l_{max} & increment the l

Whereas If $r_{max} > l_{max}$ calculate the sum by subtracting current element from r_{max} If it is greater than zero add it in the output total, else assign the current element to the r_{max} & decrement r

$$T.P. = O(n)$$

$$S.C = O(1)$$