Sunbeam_Day4_OS_MCQ's

1	Physical memory is broken into fixed-sized blocks called a) frames b) pages
	c) backing store
	d) none of the mentioned
	Ans:A
2	Logical memory is broken into blocks of the same size called
	
	a) frames
	b) pages
	c) backing store
	d) none of the mentioned
	Ans:B
3	Every address generated by the CPU is divided into two parts:
J	a) frame bit & page number
	b) page number & page offset
	c) page offset & frame bit
	d) frame offset & page offset
	Ans:B
4	The is used as an index into the page table.
	a) frame bit
	b) page number
	c) page offset
	d) frame offset

	Ans:B
5	The table contains the base address of each page in physical memory. a) process b) memory c) page d) frame
	Ans:C
6	The size of a page is typically: a) varied b) power of 2 c) power of 4 d) none of the mentioned
	Ans:B
7	With paging there is no fragmentation. a) internal b) external c) either type of d) none of the mentioned Ans:B
8	Virtual memory allows :
	a) execution of a process that may not be completely in memory

	b) a program to be smaller than the physical memory
	c) a program to be larger than the secondary storage
	d) execution of a process without being in physical memory
	Ans:A
9	Virtual memory is normally implemented by
	a) demand paging
	b) buses
	c) virtualization
	d) all of the mentioned
	Ans:A
10	The valid – invalid bit, in this case, when valid indicates :
	a) the page is not legal
	b) the page is illegal
	c) the page is in memory
	d) the page is not in memory
	Ans:C
11	A page fault occurs when :
	a) a page gives inconsistent data
	b) a page cannot be accessed due to its absence from memory
	c) a page is invisible
	d) all of the mentioned
	Ans:B
12	In segmentation, each address is specified by :
	a) a segment number & offset

	b) an offset & value c) a value & segment number d) a key & value Ans:A
13	In paging the user provides only which is partitioned by the hardware into and and a) one address, page number, offset b) one offset, page number, address c) page number, offset, address d) none of the mentioned Ans:A
14	Each entry in a segment table has a : a) segment base b) segment peak c) segment value d) none of the mentioned Ans:A
15	The segment base contains the : a) starting logical address of the process b) starting physical address of the segment in memory c) segment length d) none of the mentioned

	Ans:B
16	The segment limit contains the : a) starting logical address of the process b) starting physical address of the segment in memory c) segment length d) none of the mentioned
	Ans:C
17	is the concept in which a process is copied into main memory from the secondary memory according to the requirement. a) Paging b) Demand paging c) Segmentation d) Swapping
1.0	Ans:B
18	A process is thrashing if a) it is spending more time paging than executing b) it is spending less time paging than executing c) page fault occurs d) swapping can not take place

	Ans:A
19	The three major methods of allocating disk space that are in wide
	use are :
	a) contiguous
	b) linked
	c) indexed
	d) all of the mentioned
	Ans:D
20	In contiguous allocation :
	a) each file must occupy a set of contiguous blocks on the disk
	b) each file is a linked list of disk blocks
	c) all the pointers to scattered blocks are placed together in one
	location
	d) none of the mentioned
0.1	Ans:A
21	In linked allocation:
	a) each file must occupy a set of contiguous blocks on the disk
	b) each file is a linked list of disk blocks
	c) all the pointers to scattered blocks are placed together in one
	location
	d) none of the mentioned

	Ans:B
22	In indexed allocation: a) each file must occupy a set of contiguous blocks on the disk b) each file is a linked list of disk blocks c) all the pointers to scattered blocks are placed together in one location d) none of the mentioned
	Ans:C
23	Contiguous allocation of a file is defined by : a) disk address of the first block & length b) length & size of the block c) size of the block d) total size of the file
	Ans:A
24	and are the most common strategies used to select a free hole from the set of available holes. a) First fit, Best fit b) Worst fit, First fit c) Best fit, Worst fit d) None of the mentioned

	Ans:A
25	For each file there exists a that contains information about the file, including ownership, permissions and location of the file contents. a) metadata b) file control block c) process control block d) all of the mentioned
0.0	Ans:B
26	Metadata includes: a) all of the file system structure b) contents of files c) both file system structure and contents of files d) none of the mentioned Ans:C
27	In the linked allocation, the directory contains a pointer to the:
	I. first block
	II. last block
	a) I only
	b) II only
	c) Both I and II

	d) Neither I nor II
	Ans:C
28	There is no with linked allocation. a) internal fragmentation b) external fragmentation c) starvation d) all of the mentioned
	Ans:B
29	The major disadvantage with linked allocation is that: a) internal fragmentation b) external fragmentation c) there is no sequential access d) there is only sequential access
20	Ans:D
30	Contiguous allocation has two problems and that linked allocation solves. a) external – fragmentation & size – declaration b) internal – fragmentation & external – fragmentation c) size – declaration & internal – fragmentation d) memory – allocation & size – declaration

	Ans:A
31	Each has its own index block.
	a) partition
	b) address
	c) file
	d) all of the mentioned
	d) all of the mentioned
	Ans:C
32	Indexed allocation direct access.
	a) supports
	b) does not support
	c) is not related to
	d) none of the mentioned
	Ans:A
33	A memory page containing a heavily used variable that was
	initialized very early and is in constant use is removed, then the
	page replacement algorithm used is :
	a) LRU
	b) LFU
	c) FIFO
	d) None of the mentioned
	a) None of the mentioned

	Ans:C
34	The aim of creating page replacement algorithms is to :
	a) replace pages faster
	b) increase the page fault rate
	c) decrease the page fault rate
	d) to allocate multiple pages to processes
	a, to amount manufer pages to proceed
	Ans:C
35	A FIFO replacement algorithm associates with each page the
33	ATITO replacement algorithm associates with each page the
	a) time it was brought into momery
	a) time it was brought into memory
	b) size of the page in memory
	c) page after and before it
	d) all of the mentioned
	Ans:A
36	Optimal page – replacement algorithm is :
	a) Replace the page that has not been used for a long time
	b) Replace the page that has been used for a long time
	c) Replace the page that will not be used for a long time
	d) None of the mentioned
	Ans:C
37	Optimal page – replacement algorithm is difficult to implement,

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	because: a) it requires a lot of information b) it requires future knowledge of the reference string c) it is too complex d) it is extremely expensive
	Ans:B
38	In the algorithm, the disk arm starts at one end of the disk and moves toward the other end, servicing requests till the other end of the disk. At the other end, the direction is reversed and servicing continues. a) LOOK b) SCAN c) C-SCAN d) C-LOOK
	Ans:B
39	In the algorithm, the disk head moves from one end to the other , servicing requests along the way. When the head reaches the other end, it immediately returns to the beginning of the disk without servicing any requests on the return trip. a) LOOK b) SCAN c) C-SCAN d) C-LOOK

	Ans:C
40	In the algorithm, the disk arm goes as far as the final
	request in each direction, then reverses direction immediately
	without going to the end of the disk.
	a) LOOK
	b) SCAN
	c) C-SCAN
	d) C-LOOK
	Ans:A