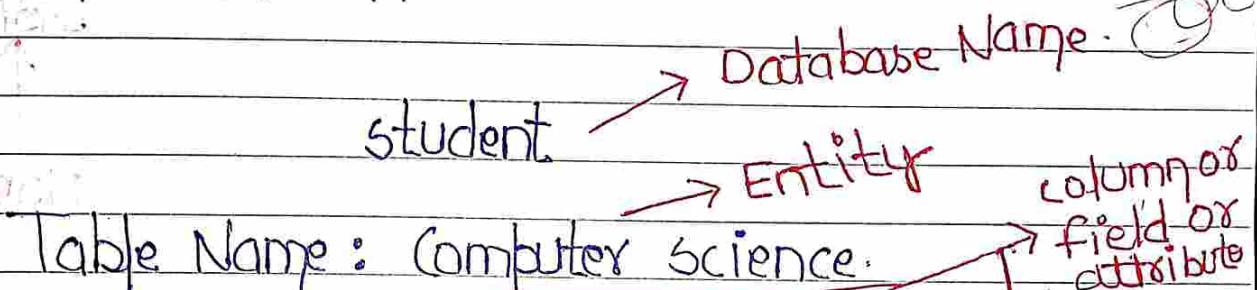


Unit → 1Database Management SystemUnit → 1→ Introduction to Database

→ Database is integrated collection of related information along with the details so that it is available to the several user for the diff' application.



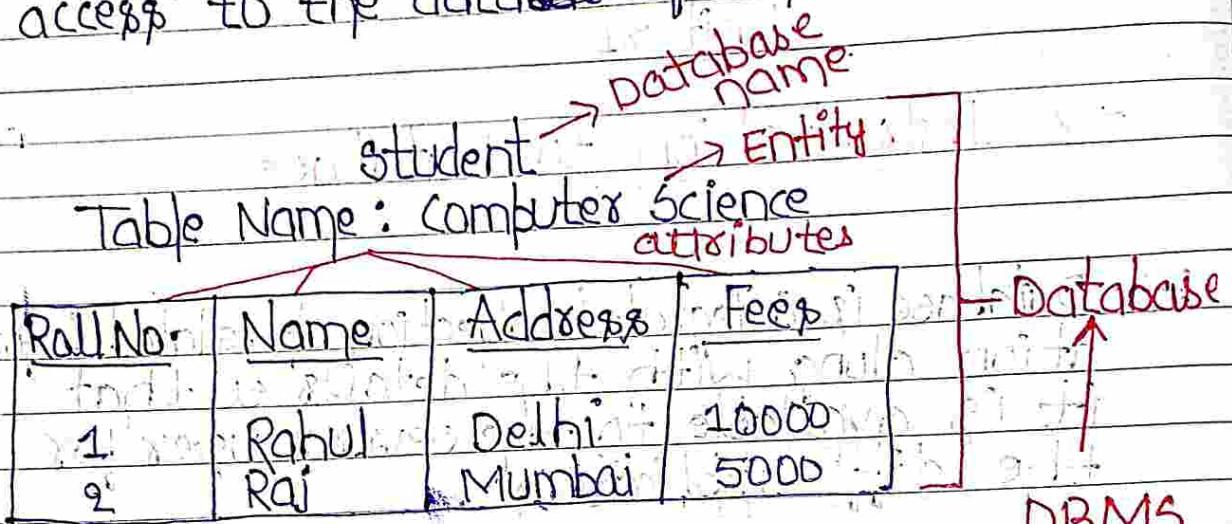
Roll number	Name	Address	Fees
1	Rahul	Delhi	10000
2	Raj	Mumbai	5000
3	Rohit	Kolkata	15000

Note: → Database is a collection of data.

- Database is in the form of table.
- One database may have one table or more than one table.
- One database may have one or more than one entity (Table Name)
- Attributes → Features of entity.
- describe the

DBMS (Database Management System)

- Database is the software system that allows the access to the data of the database.



Rahul Raj Rohit Scholarship

Roll No.	Name	Address	Fees	Scholarship
1	Rahul	Delhi	10000	5000
2	Raj	Mumbai	5000	3000
3	Rohit	Kolkata	15000	8000

File based System

→ In file based system, we have to open individual file for Rahul, Raj, Rohit and scholarship. Then, we can check their details.

Roll No.	Name	Address	Fees	Scholarship
1	Rahul	Delhi	10000	5000
2	Raj	Mumbai	5000	3000
3	Rohit	Kolkata	15000	8000

Database Management System

DBMS

:- Database Management System

- Database requires **data and information**.
- Data :-
 - It can be anything like name, place or number etc.
 - It refers to raw data or unprocessed data.
- Information :-
 - It is organized or classified data so that it gives some meaningful values.
 - It is a collection of data.

Data

- Data is raw facts and figures.

Information

- Information is a processed form of data.

- Data does not help in decision making.

- Information helps in decision making.

Eg :- University Database

:- Data about students, faculty, courses, research laboratories, registration etc.

It helps to keep an accurate track of the academic activities of the university.

Atomicity
either it is complete
done, either not
done, done at all

Page No.

Vikram

Date

Note :-

Before DBMS, organizations usually stored information in file based system, but file based system has many disadvantages:-

- Difficulty in accessing data



We have to access them individually and open the file for each of them
(Raj, Rahul, Rohit)

If I choose file related to Rahul then it can only access ~~use~~ details of Rahul only.

→ But if we use DBMS then we can access all the details of Raj, Rahul and Rohit at same time.

- Security problems → Hard to provide user access to some, but not all data
- Data inconsistent

→ If the file of a student is modified by two different users then there may be inconsistency of their data.

- Anyone update only 10 lines and others may update 20 lines

→ Each student has separate file. We keep all the details of students in separate file.

Page No. KRISHNA
Date / /

anything will either be done completely or even not at all.

• Atomicity problem

→ The information which is entered to any system may fail at any time.

Failure occurs like ATM withdrawal.

• Difficulty in concurrency control

→ Simultaneous execution of many at a time.

• Integrity problem → Hard to add new constraint

• Data Redundancy or Duplicate of data

→ the same information may be duplicated in different files.

• Multiple file formats.

• Difficulty in accessing data

• Need to write a new program to carry out each new task.

Database Management System

→ It is the software system that allows users to define, create and maintain a database and provides controlled access to the data.

• It is a collection of programs that enables user to store, modify and extract information from a database.

Features

• To add new information.

• To view or retrieve the stored information.

• To modify or edit the existing data.

• To remove or delete the information.

Syntax:-

CREATE TABLE table-name

(column-name1 data-type(size),

column-name2 data-type(size),

column-name3 data-type(size)

);

Eg :-

CREATE TABLE my-tab

name varchar(30),

roll number(4),

address varchar(100)

);

Output:

Name	Roll	Address

components of Database

5 major components :-



:- Hardware → It is the actual computer system used for keeping and accessing the database. It consists of secondary storage devices as hard disks.

:- Software → It is actual DBMS.

:- Data → It is the bridge b/w machine components and user components.

:- Users → There are number of users who can access or retrieve data on demand using the applications and the interfaces provided by DBMS.

:- Procedures → It refers to the instructions and the rules that govern the design and the use of database.

Applications of DBMS

- Airlines or railway :- Reservations, schedules, search flight/train between two particular station.

- Universities :- Registration, Grades, Library books etc.

- Sales :- Customers, products, purchases, online tracking.

- Human resources :- employee record, salaries, tax.

Data model :- A set of concepts to describe the structure of a database.

- data
- data relationships
- data constraints

Scheme :- The overall description of the database.

Instance :- Data in the database at the particular instant of time.

Physical Data Independence

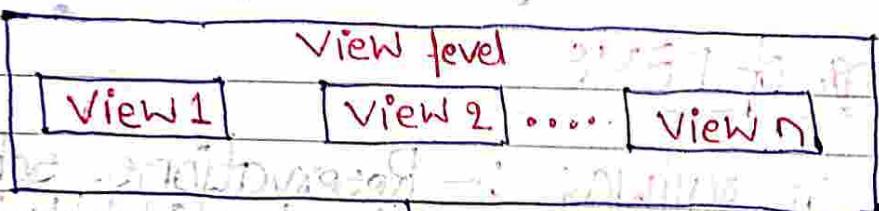
The ability to modify the physical schema without changing the logical schema.

* Level of Abstraction

• Physical level :- It describes how a record is stored.

• Logical level :- It describes what data stored in database and relationship among data.

• View level :- It describes only part of database. Application programs hide details of data types. View can also hide information for security purposes.



logical level

~~Physical level~~

✓ Instances → At a particular time, what I do → the actual content of the database
✓ schemas → plan to create or work on a database.
logical schema
structure of logical level

Database Languages

Page No. 10
Physical schema
structure of physical level

→ Database languages are used to create and maintain database on computer.
(DDL)

* **Data Definition Language** :- It is a language that allows user to define data and their relationship to other types of data.

- CREATE
- ALTER
- DROP
- TRUNCATE
- RENAME

physical data

independence

the ability to modify the physical schema without changing the logical schema

✓ Specification notation for defining database schema
(DML)

* **Data Manipulation Language** :- It provides a set of operations to support the basic data manipulation operations.
It allows user to insert, update, delete and retrieve data from the database.

- DELETE
- INSERT
- SELECT
- UPDATE

• It also known query language
(DCL)

* **Data Control Language** :- DCL statements control access to data and the database.

- GRANT
- REVOKE
- COMMENT

• **Transaction Control Language (TCL)**
— Commit

X ER MODEL

→ The ER model defines the conceptual view of a database.

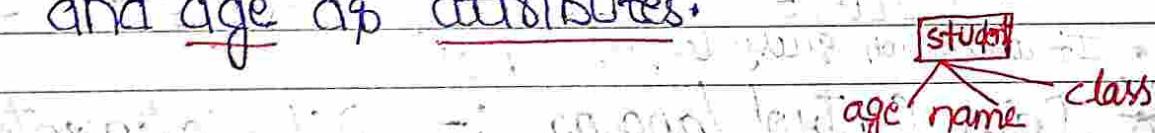
→ Entity : It can be a real-world object, either animate or inanimate, that can be identifiable. Or it has independent existence.

e.g. In school database; students, teachers, classes and courses offered can be considered as entities. All these entities have some attributes or properties that give their identity.

→ Attributes : Entities are represented by means of their properties.

or It is a property of an entity, relationship.

e.g. In school database; students entity may have name, class and age as attributes.



→ Entity type : It is a collection of entity which has common attributes.

student → entity type OR schema

ID	Name	Age	Attributes
e1	Ram	14	
e2	Shyam	12	
en	John	13	

entity

Basic Terminology of ER MODEL or Element

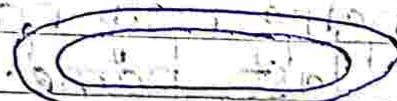
Rectangle



An object is in real world Entity, or strong Entity.

Weak Entity.

oval

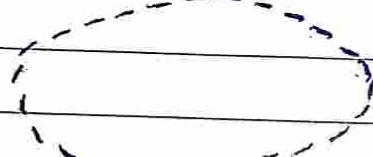


Attribute
(which describe entity)

Multivalued Attribute.



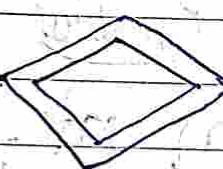
Composite Attribute.



Derived Attribute.



Relationship
(Association among the entities)



Weak Relationship.



Key attribute (Primary key)



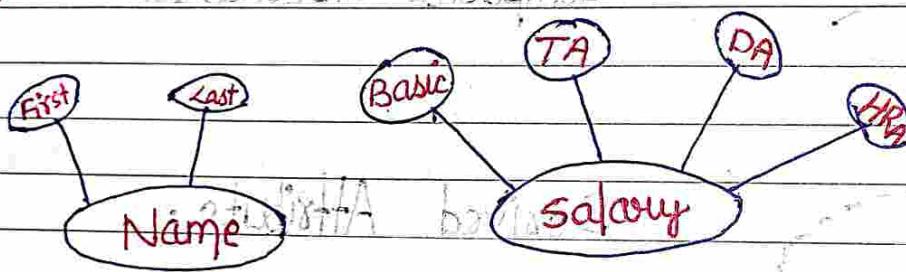
Partial Key.



Aggregation box

* Types of Attributes

- **Simple attribute** :- It cannot be divided further.
eg :- student's phone number is an atomic value of 10 digits. Mobile No
- **Composite attribute** :- It consists of more than one simple attribute.
eg :- student's complete name may have first_name and last_name.
eg :- salary includes basic, TA, DA, HRA.



- **Derived attribute** :- The value of an attribute can be derived from other attributes present in database.

Note :- Age can be derived from date-of-birth
Area can be derived from radius.

- **Single-value attribute** :- It takes one value per entity.

eg :- Gender

Gender is an entity which can hold either male or female.

circle

ox

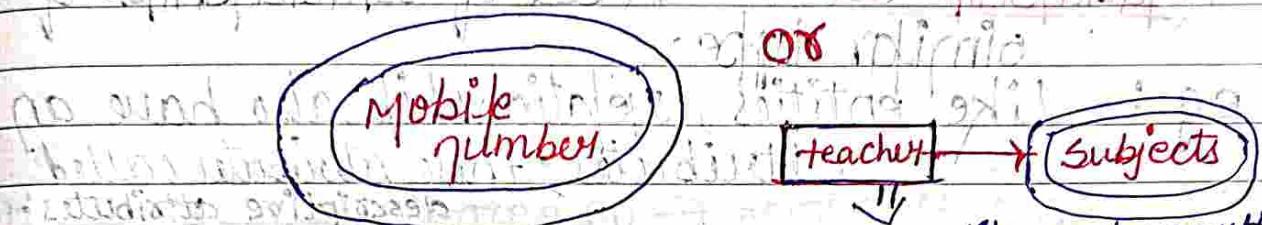
Radius

area

Age

- **Multivalued attribute :-** It may takes more than one values.

eg:- A person may have more than one phone number or email address.



Teacher entity can have multiple subject values.

- **StoredProcedure :-** It does not take any updation in database.

eg:- Date_of_birth doesn't take any updation in future.

(dob)

- **Key :-** It is an attribute or collection of attributes that uniquely identify an entity among entity set.

or which uniquely identify in the entity set.

eg:- roll_number of student makes him/her identifiable among students.

(rollno.)

- **Superkey :-** A set of attributes (one or more) that collectively identifies an entity.

*(Required notes for an entity) writing
about its own behaviour*

It describes how entities interact.

Page No.

Date

KRISHNA

- Relationship :- Association among the entities.

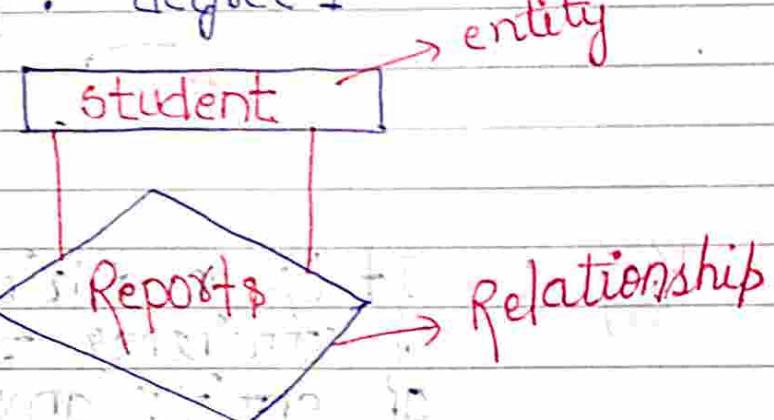
e.g. :- student enrolls in a course, where Enrolls is called relationship.

- Relationship Set :- A set of relationships of similar type.

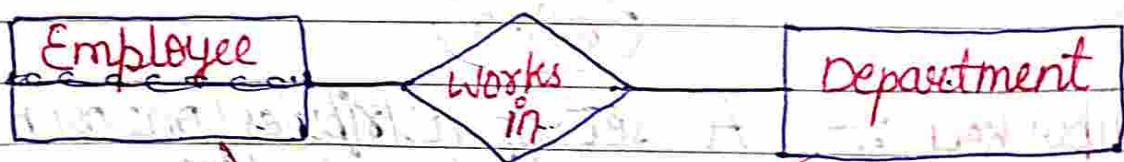
e.g. :- Like entities, relationship also have attributes; these attributes called descriptive attributes.

- Degree of Relationship :- specifies the no. of entity sets participates in relationship set.

- i) Unary :- degree 1

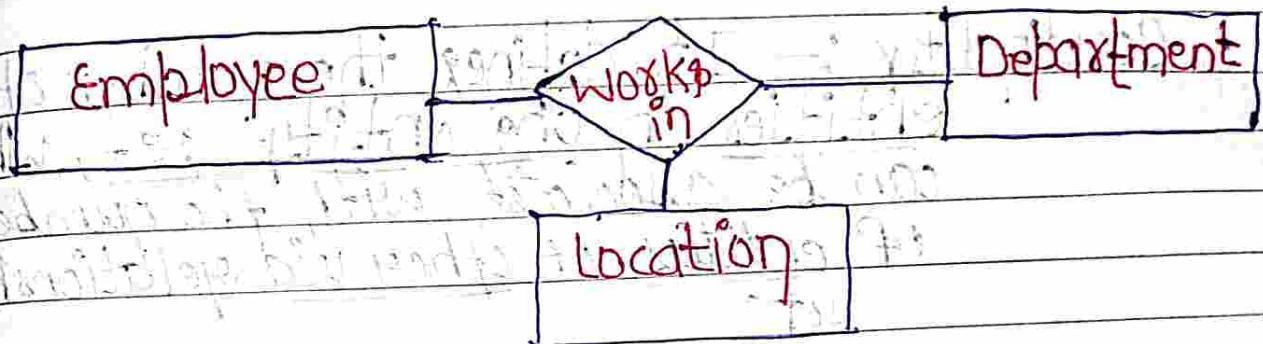


- ii) Binary :- degree 2 → Relationship between two entity sets -

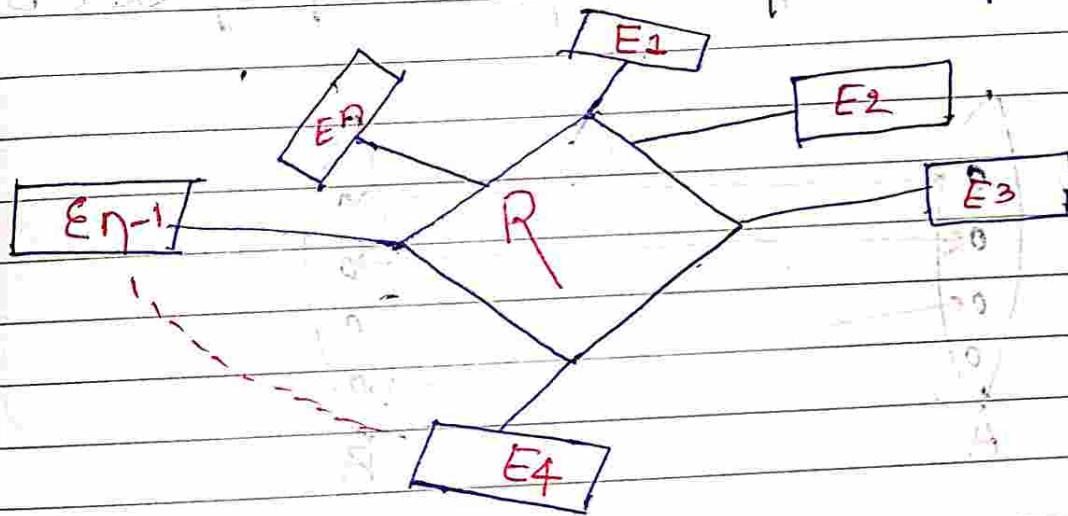


entities (there is relation between employee and department)

iii) Tertiary :- degree 3 \rightarrow Relationship among three entity sets.



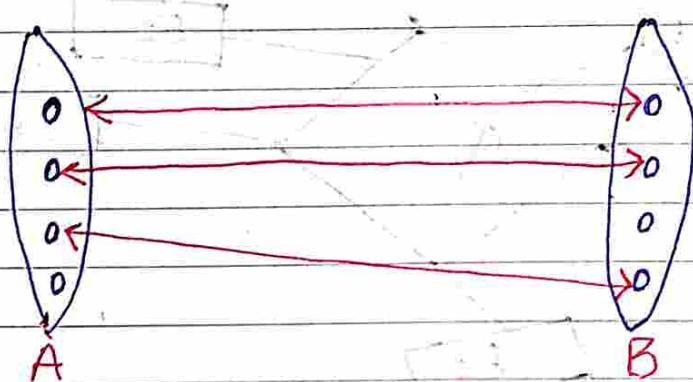
iv) n -ary \rightarrow degree $n \rightarrow$ Relationship among n entity sets.



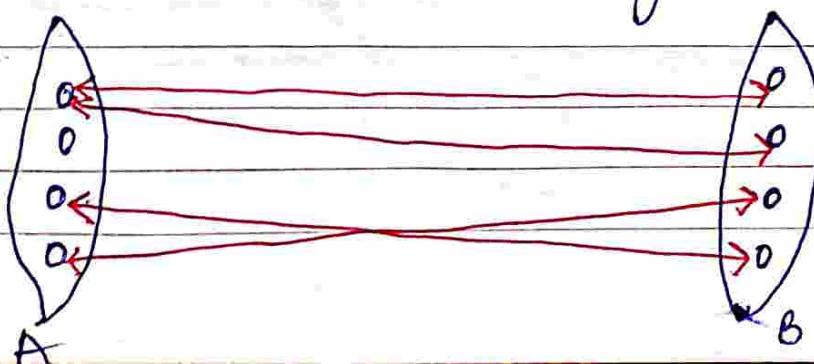
* Mapping

- Cardinality :- It defines the number of entities in one entity set, which can be associated with the number of entities of other via relationship set.

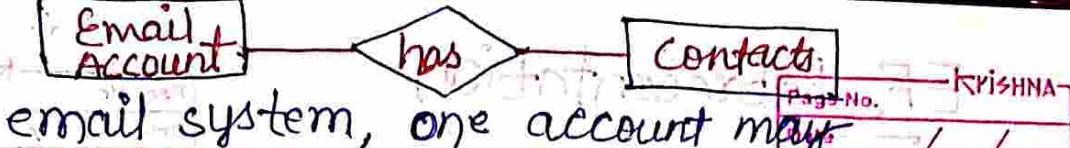
- One-to-One :- One entity from entity set A can be associated with at most one entity of set B & vice-versa.



- One-to-many :- One entity from entity set A can be associated with more than one entities of entity set B however an entity from set B, can be associated with at most one entity.



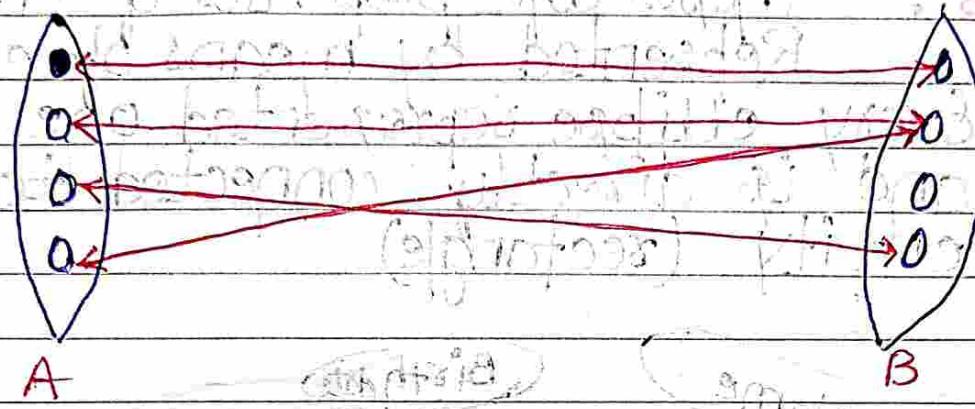
Eg:-



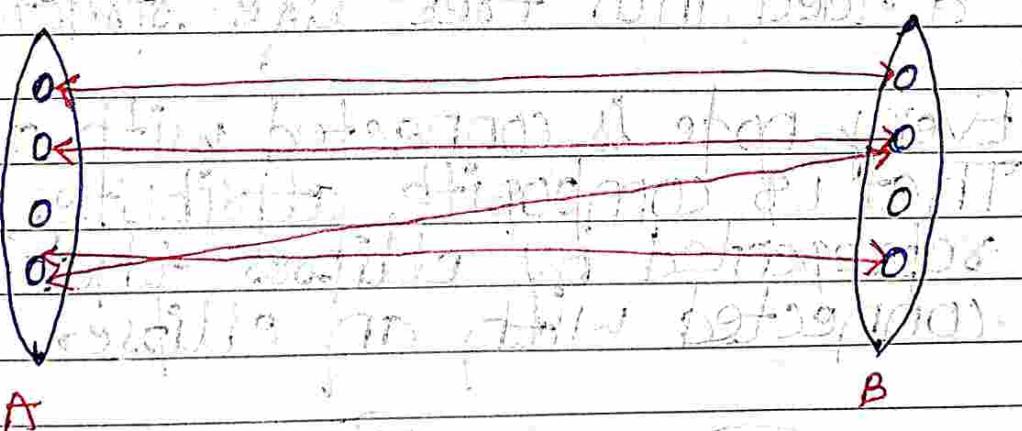
In email system, one account may

have multiple contacts.

- Many-to-one :- More than one entity from set A can be associated with at most one entity of entity set B.



- Many-to-Many :- One entity from A can be associated with one or more than one entity from B and vice-versa.



ER Representation

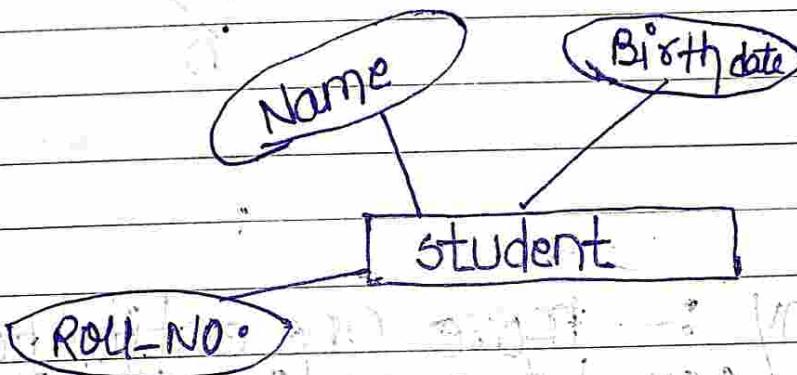
Page No.
Date

KRISHNA

- Entity :- Represented by means of rectangles.

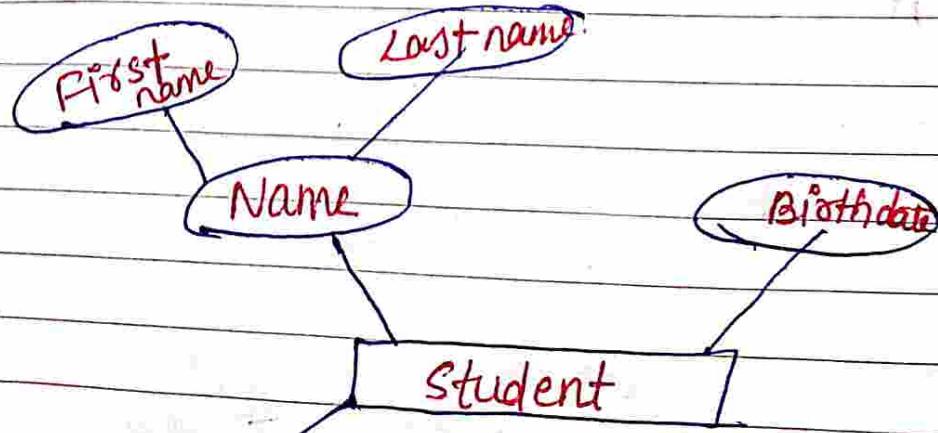


- Attributes :- These are properties of entity.
Represented by means of ellipses
⇒ Every ellipse represents one attribute
and is directly connected to its entity (rectangle)

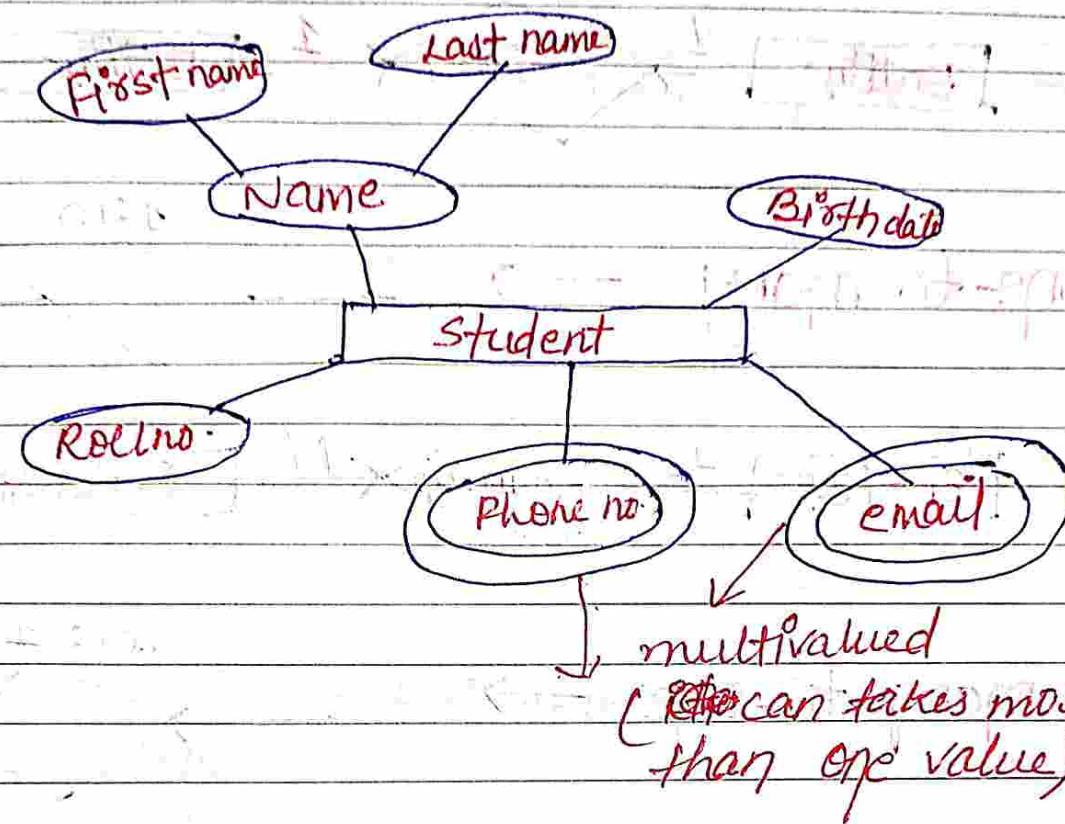


- Composite attributes :- They are further divided into tree like structure.

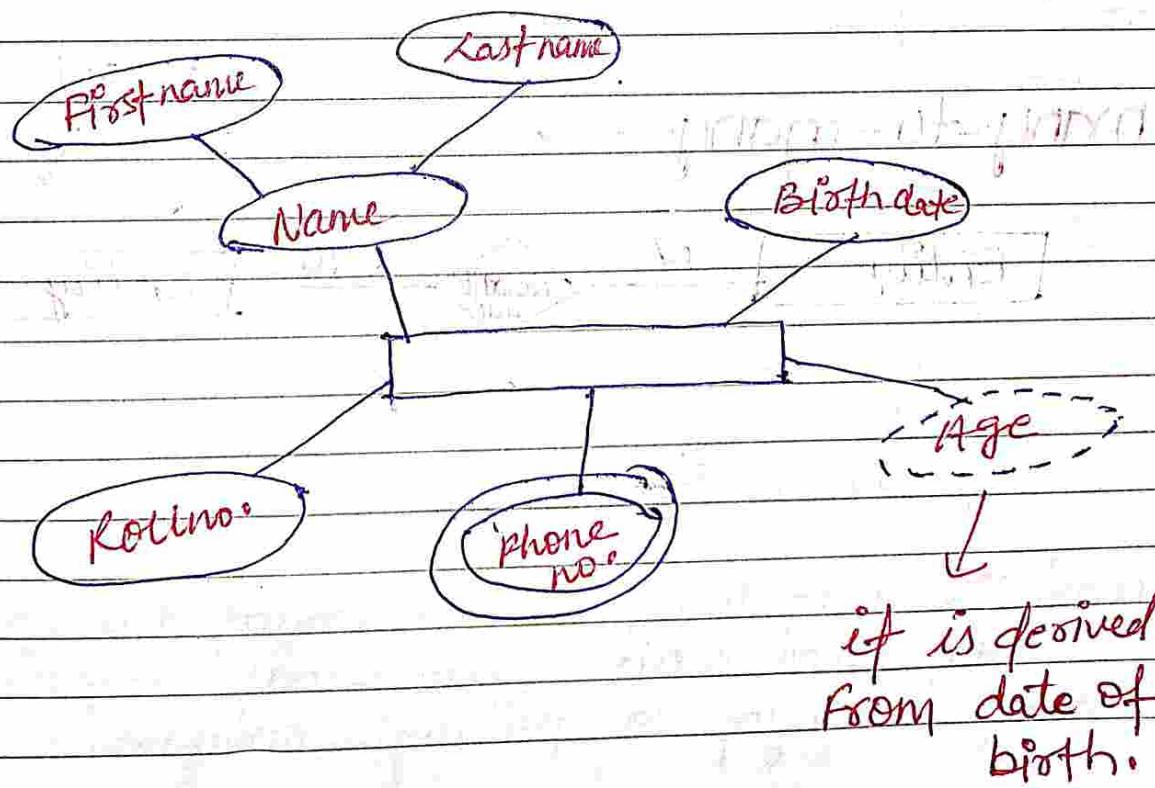
⇒ Every node is connected with attribute. That is composite attributes, are represented by ellipses that are connected with an ellipse.



- Multivalued :— These attributes are depicted by double ellipse.



- Derived :— These can be depicted by dashed ellipse.



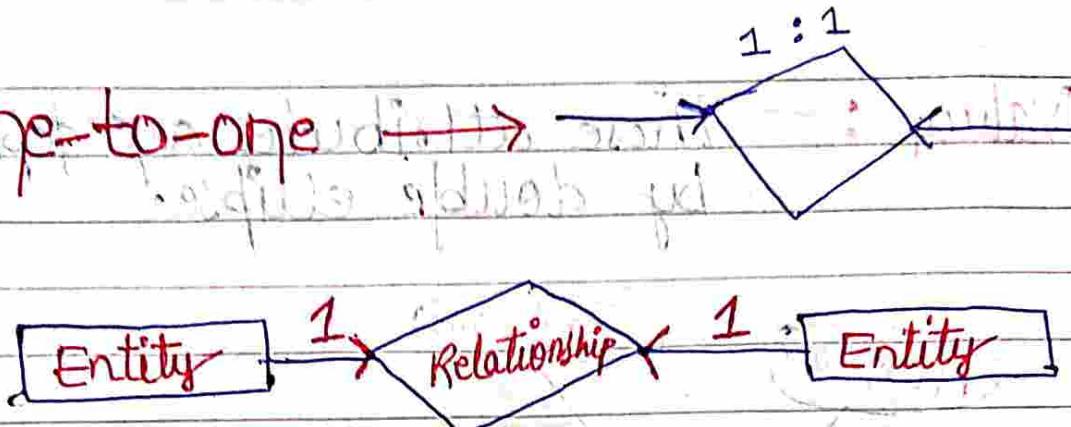
ER Representation of Relationship

Page No.

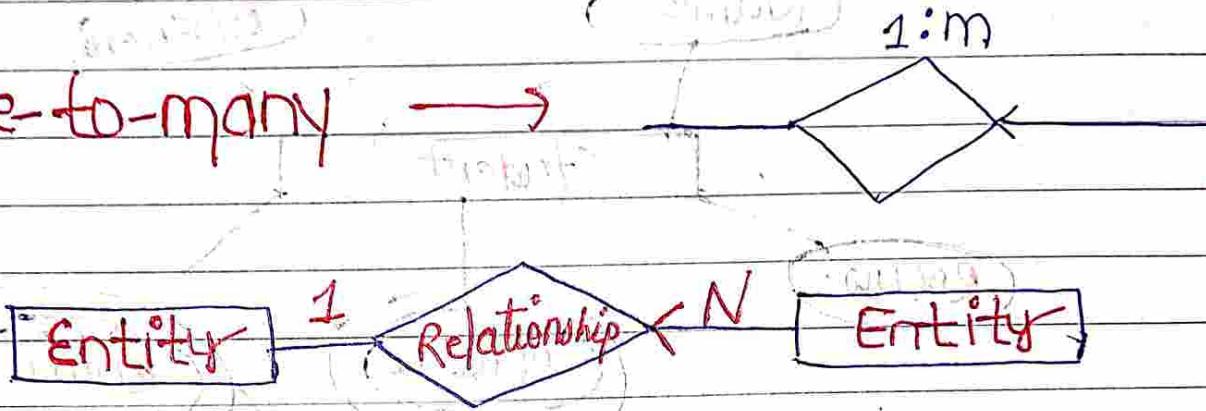
KRISHNA

Date / /

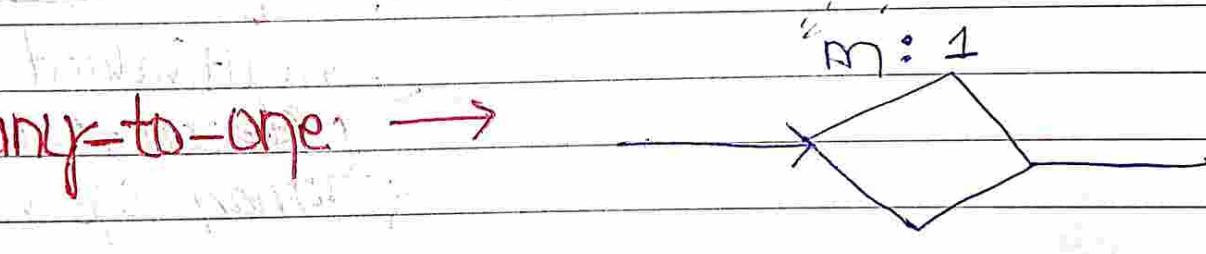
- One-to-one →



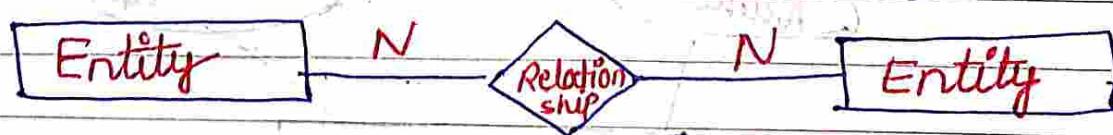
- One-to-many →



- many-to-one →



- many-to-many →



Weak and Strong Entity

- Weak entity is an entity that depends on the existence of other entity.

Or, we can say that it is an entity which cannot be identified by its own attributes.

Order

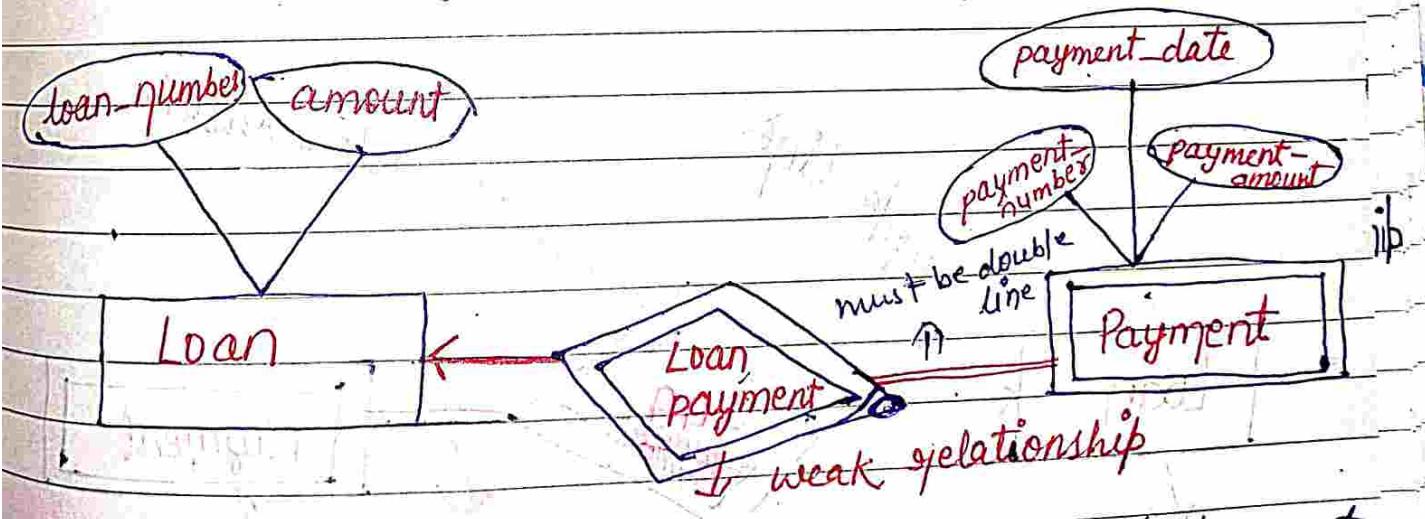
Order item

Weak entity

Here, the order item will be meaningless without an order so it depends on the existence of Order = weak entity set which does not have a key attribute

- An entity set has a primary key is called as Strong entity set.

→ Consider an entity set payment Payment which has three attributes :- payment-number, payment-date and payment amount



Although each payment entity is distinct but payment for different loans may have same payment-date, payment-number or payment-amount.

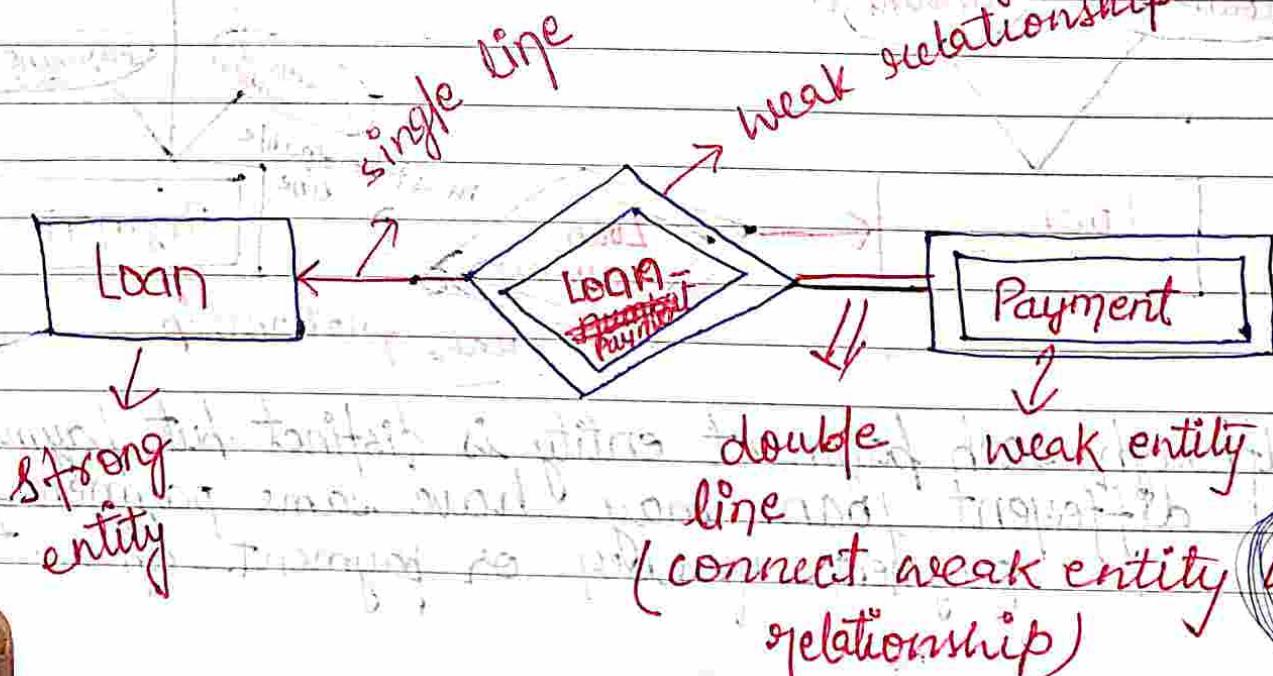
- A member of strong entity set is called **dominant entity** and member of weak entity set is called **subordinate entity**.

Weak entity set

- It is represented by a rectangle.
- Member called dominant.
- Relationship b/w two strong entity is represented by diamond symbol.
- Line connecting strong entity with relationship is single.

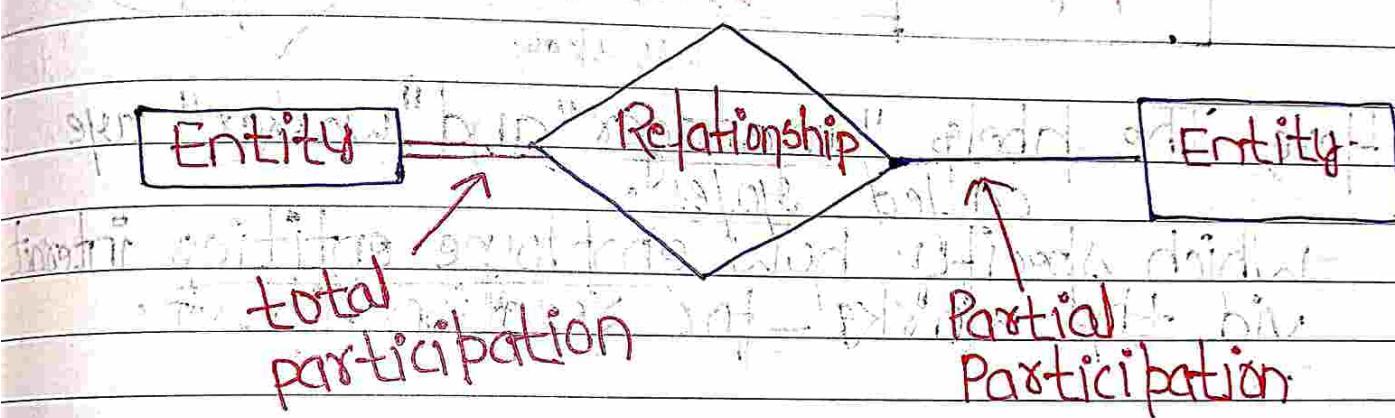
Weak entity set

- It is represented by a double rectangle.
- Member is called subordinate.
- Relationship b/w one strong and one weak set is represented by double diamond symbol.
- Line connecting weak entity set with relationship is double.



Total and Partial Participation

- **Total Participation :-** Each entity is involved in the relationship.
Total participation is represented by double line.
- **Partial Participation :-** Not all entities are involved in the relationship.
Partial participation is represented by single line.

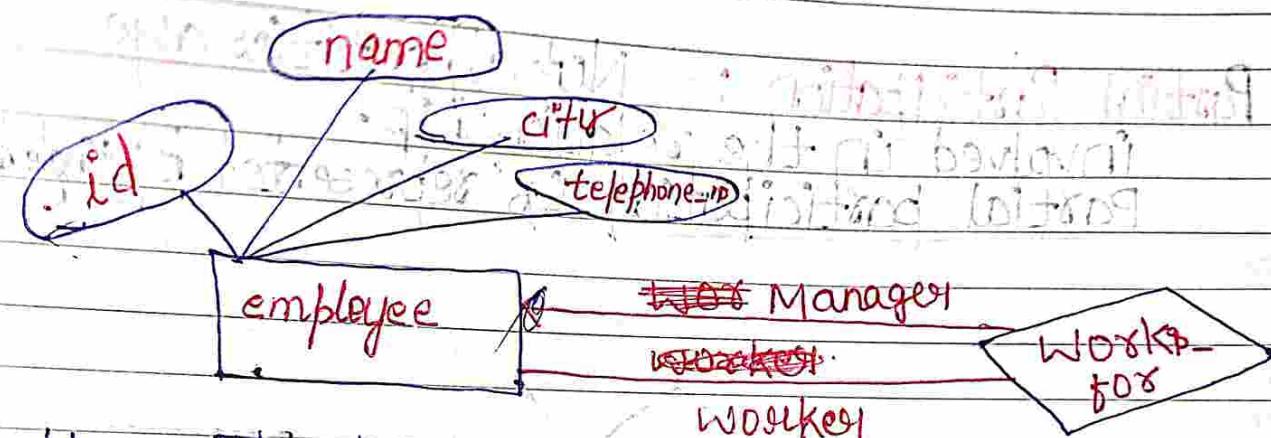


eg :- Employee head of department

- Not all employees become a head, but department will always be headed by one employee.

so, employee participated partially in relationship

Relationships are ~~roles~~ type which specify how employee entities interact in relationship set.

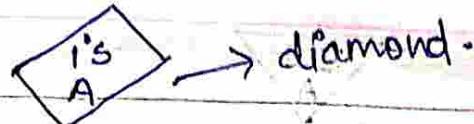


Here, The labels "manager" and "worker" are called roles.

which specify how employee entities interact via the works-for relationship set.

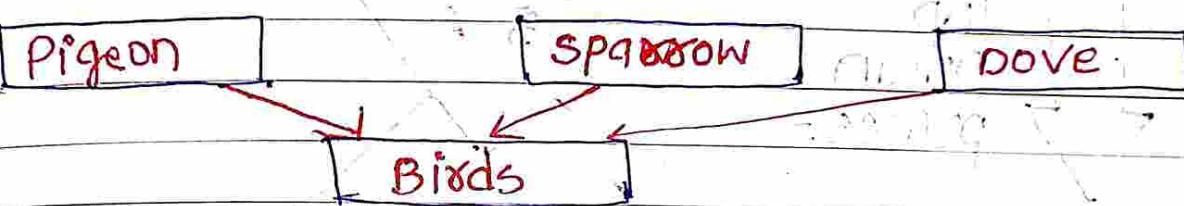
- Roles are indicated in E-R diagrams by labelling the lines that connect diamonds to rectangles.

Generalization



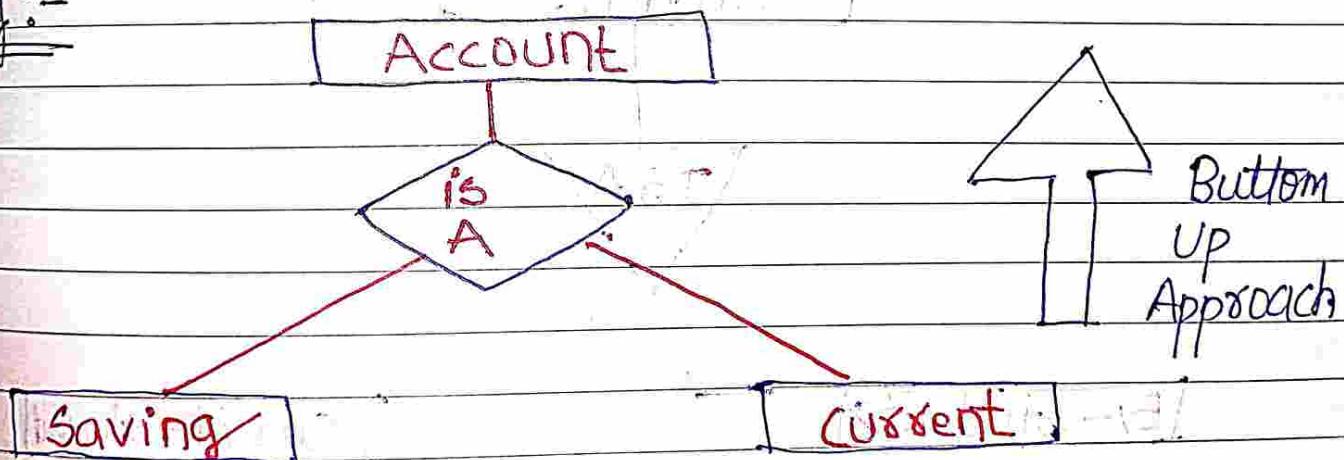
→ The process of generalizing entities, where the generalized entities contain the properties of all the generalized entities.

- In this, a number of entities are brought together into one generalized entity based on similar characteristics.

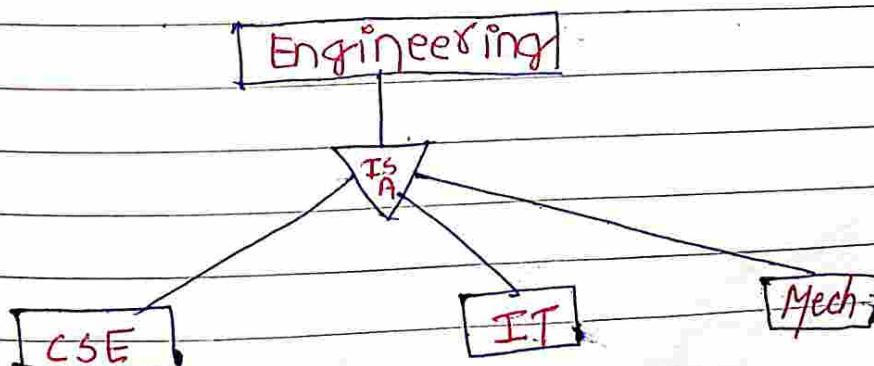


Here, Pigeon, Sparrow, Dove can all be generalized as birds (all have almost same features).

Eg :-



Eg:-



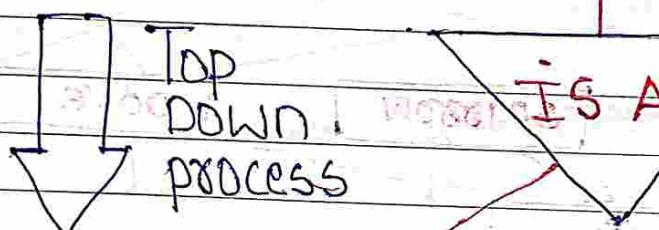
Specialization



02/02/1973

- The group of entities is divided into sub-group based on their special characteristics.
- Eg:- In school, person may be divided as teacher or student.

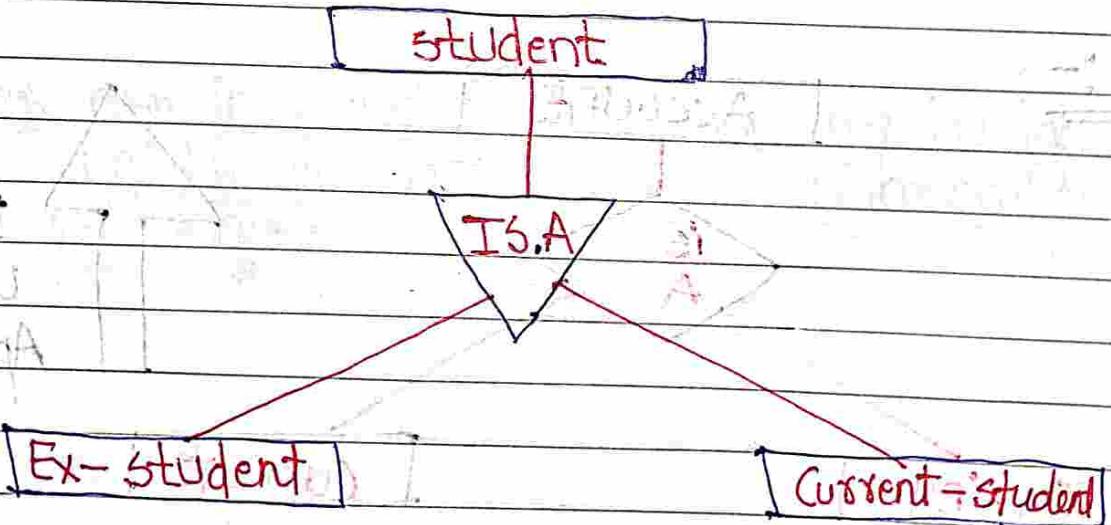
Person



Student

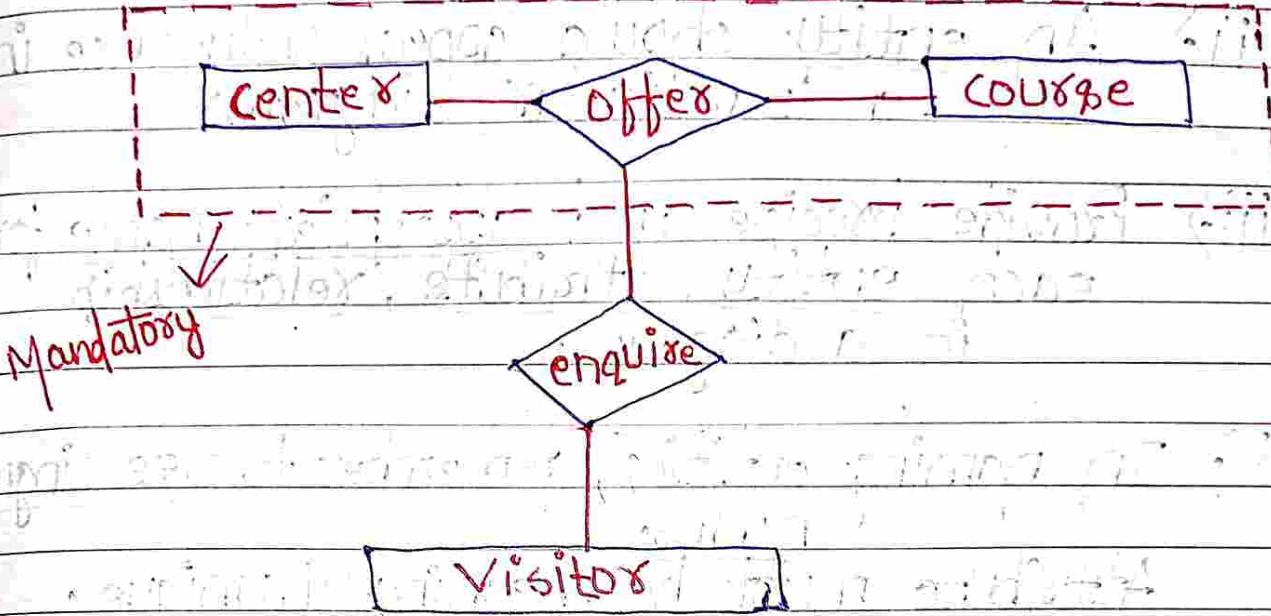
Teacher

Eg:-

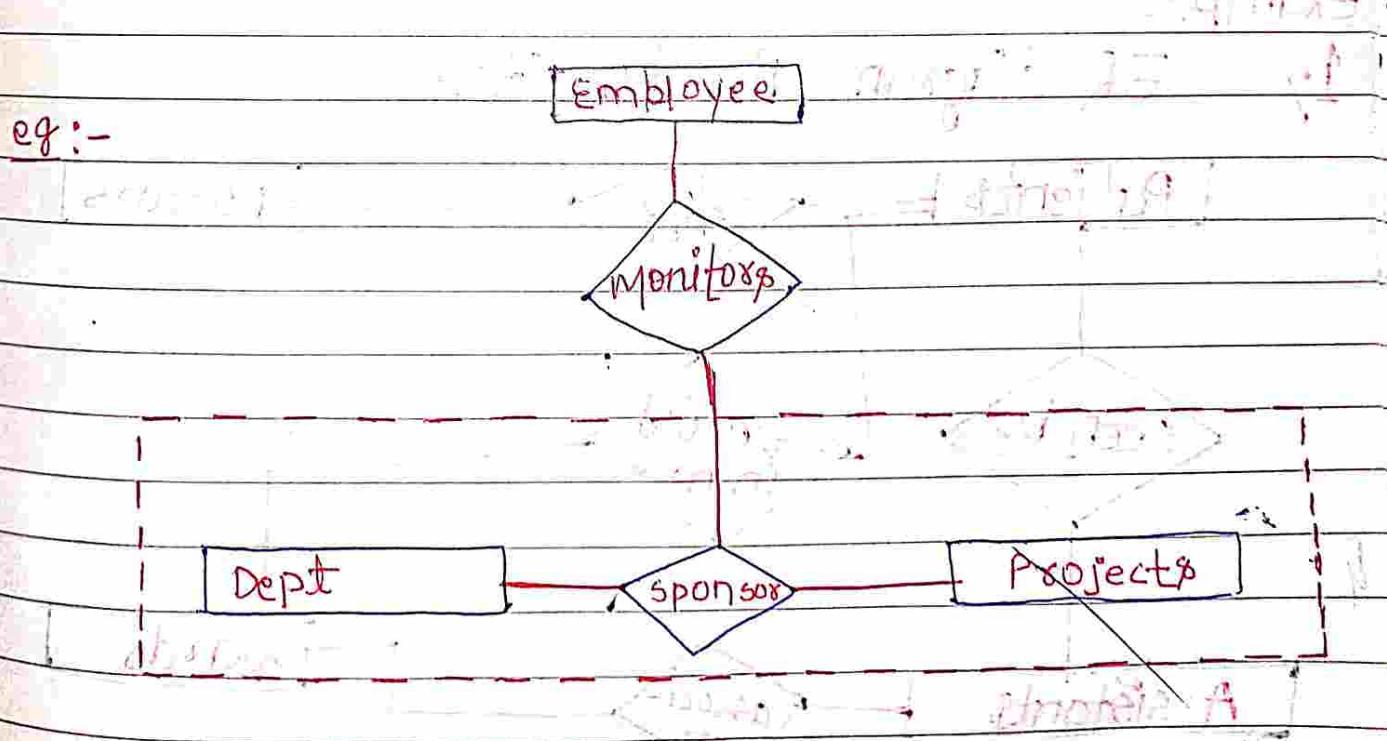


Aggregation

- A process when relation between two entity is treated as a single entity.



Here, the relation b/w center and course, is acting as an entity in relation with visitor.

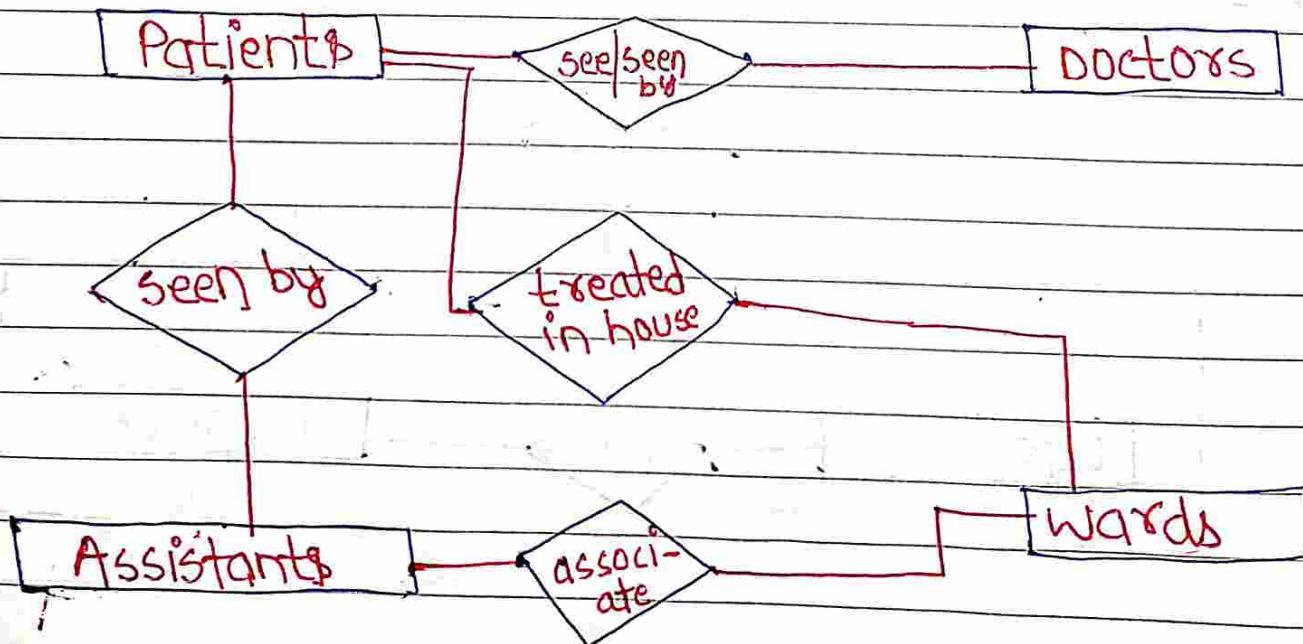


* Rules to draw ER diagram

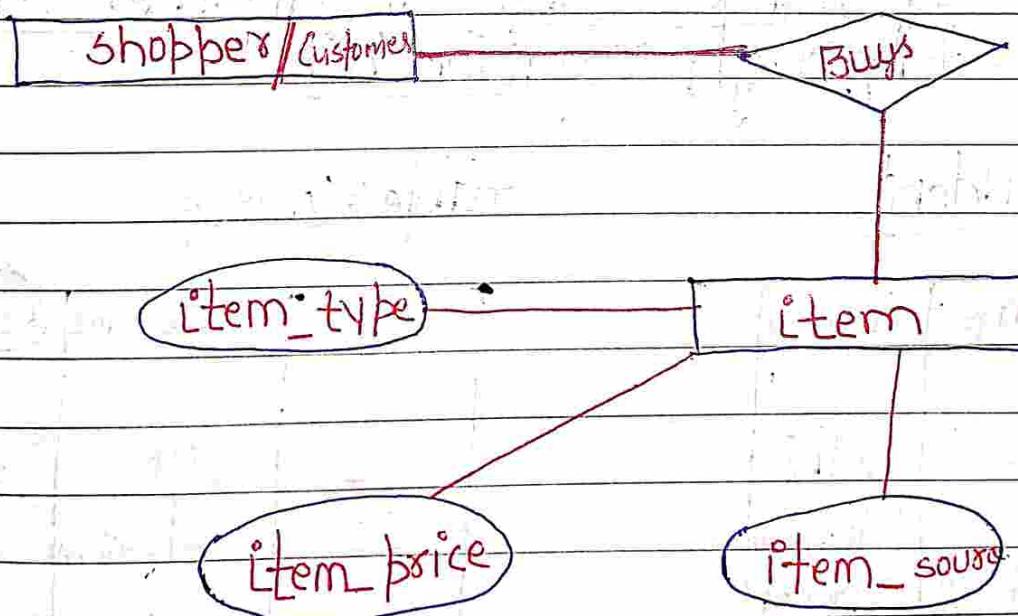
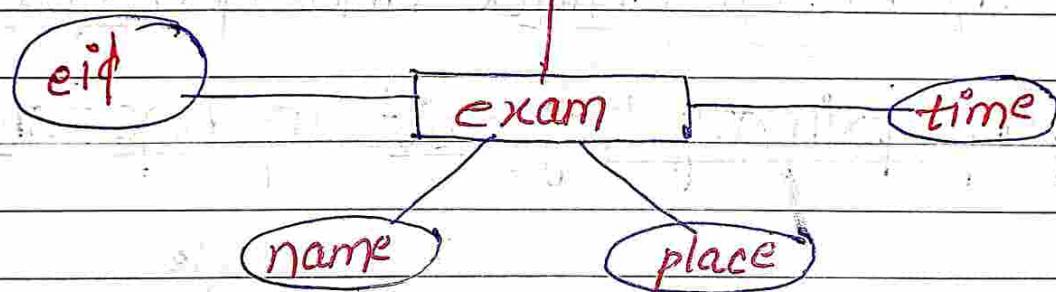
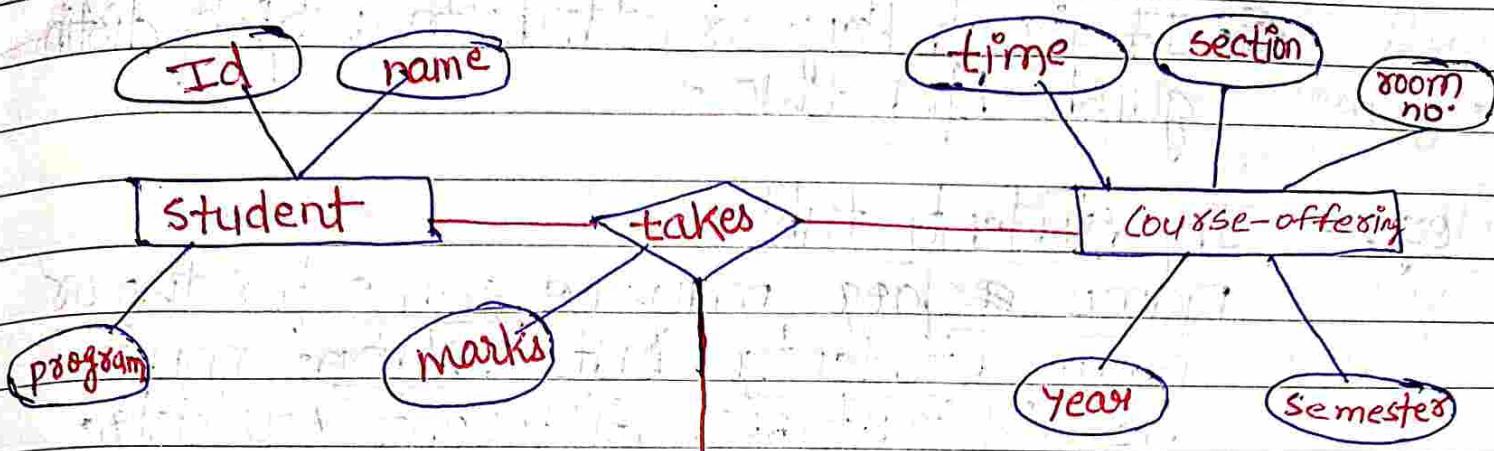
- i) Identify all the relevant entities in a given system.
- ii) An entity should appear only once in a particular diagram.
- iii) Provide precise and appropriate name for each entity, attribute, relationship in a diagram.
- iv) In naming entities, remember to use singular nouns.
Attribute must be meaningful, unique.
- v) Remove unnecessary relationships b/w entities.

Examples

1) ER diagram for hospital



2. ER diagram for student database



- Primary key : \rightarrow It is unique as well as not null value.
- It is a key or attribute which distinguishes between two.

e.g.: - In Students table,

name, class may be same for two or more students but Rollno. must be different ^{or unique} for different students. Hence, Rollno. is a primary key as it is unique and non-null value.

- Only one primary key exist in one table.

• Foreign key : \rightarrow It is the attribute or set of attribute that references to primary key of same table or another table through relationship.

• table \rightarrow student

Rollno	Name	Address
1	A	Delhi
2	B	Mumbai
3	A	Chennai

table \rightarrow Course

Course Id	Course Name	Rollno
C1	DBMS	1
C2	Networks	2

Base table or
referenced table

(Table which has primary key)

Referencing table

Table which has primary key

Referential Integrity

→ Let us suppose, If we want to buy mobile phone then we either check it online or manually. We observed that the same mobile phones have different prices in online and manually. so, we can say that there is no integrity here. (It is because amazon, flipkart have diff database)

But; In University database, If we like to search a particular student with their name and registration number then student will be same either we search them in library, placement, course and other department.

so, There is referential integrity mandatory in database.

eg:-	Table: Student			Table: Course		
	Rollno.	Name	Address	CourseID	CourseName	Rollno.
1	A	Delhi		C1	DBMS	1
2	B	Mumbai		C2	Networks	2
3	A	Chennai		C3	C++	7
4	D	Goa				

Base table
or referenced table

• Referenced table

- Insertion (No violence)
- Deletion (may be violence)
- Updation (may be violence)

Referencing table

- Insertion → May cause Violence
- Delete → Will not cause Violence
- Updation → May cause Violence

on delete cascade
then it also delete

Candidate Key :- It is a minimal set of attributes which uniquely identifies a tuple in a relation.

Super key :- A super key is a combination of all possible attributes which can uniquely identify two tuples in a table.

• Super set of any candidate key is a super key.

eg:-

Candidate key		Attributes	
	(RollNo)	Name	Age

• Candidate key
→ RollNo.

• Super key will be collection of tuples of one candidate key and any other attributes.

eg. :- (RollNo.) → It is candidate key as well as super key.
(RollNo; name)
(RollNo; age)
(RollNo; name, age)

{
Super key
(It contains Candidate key and other attributes)

(name, age) → It is not a super key because it does not contain candidate key.

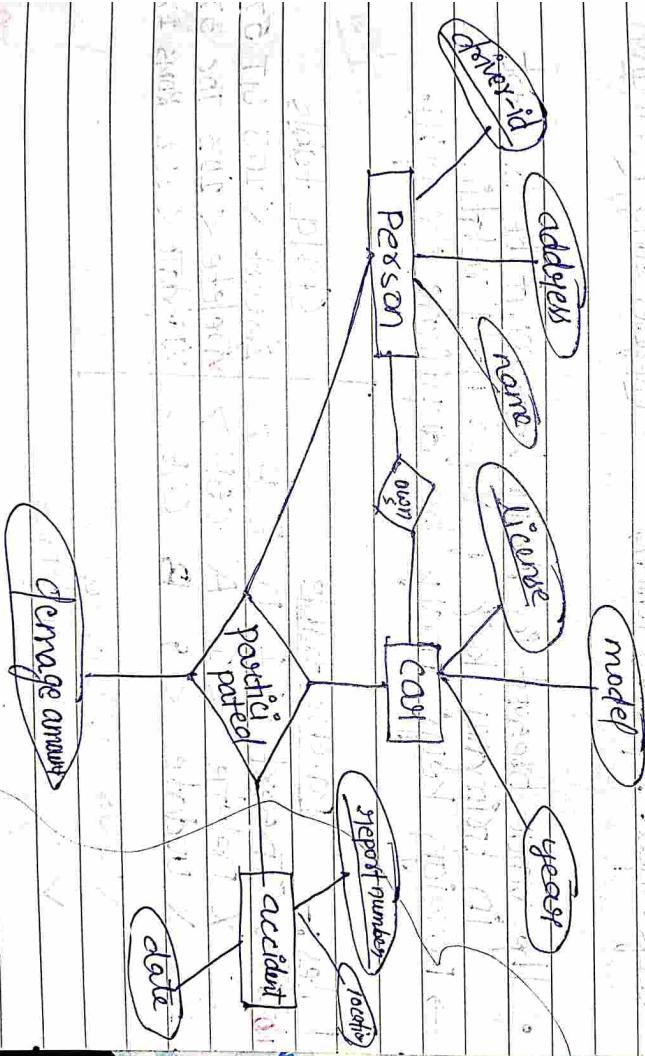
Note :- Every candidate key is called super key but every super key need not be a candidate key.

Ques:- R(A₁, A₂, A₃, A₄, ..., A_n)
then how many super keys are possible?

- ① If A₁ is candidate key.
 ② If A₁, A₂ are candidate keys.

Sol:-

at 168505:



~~on foreign key~~

→ It is based on Referential Integrity constraint.

table :- Student

Primary key

R. No.	Name	Branch
1	A	CSE
2	B	IT
3	C	CSE

table :- Registration

C. NO	C. Name	R. NO.
101	DBMS	1
102	CD	1
103	TOC	3
104	PL	null

solution

- In the

eg:-

Parent key

(Referenced relation) (Parent)

(Referencing) (child)

Let us

- The value present in foreign key must be present in primary key of referenced relation.

→ Foreign keys may be duplicate or null.

Rules :-

Parent table

- ✓ Insert < 4 D ECE >
- ✗ Delete < 1 A CSE >
- ✗ Update < 2 B CSE >

✗ → Violence

✓ → No violence

Child table

- ✗ Insert < 105 GT 5 >
- ✓ Delete < 103 TOC 3 >
- ✗ Update < 102 RDMs 1 >

Note

Generalisation

Specialisation

Here, Deletion and updation from referenced relation and insertion and updation from child table referencing relation may cause violence due to foreign key constraint.

SCD

Solution :- Foreign key with OnDelete Cascade

- In this process, if data from parent key is deleted or updated then the related data from child table also be deleted (^{as updated}) cascadingly simultaneously.

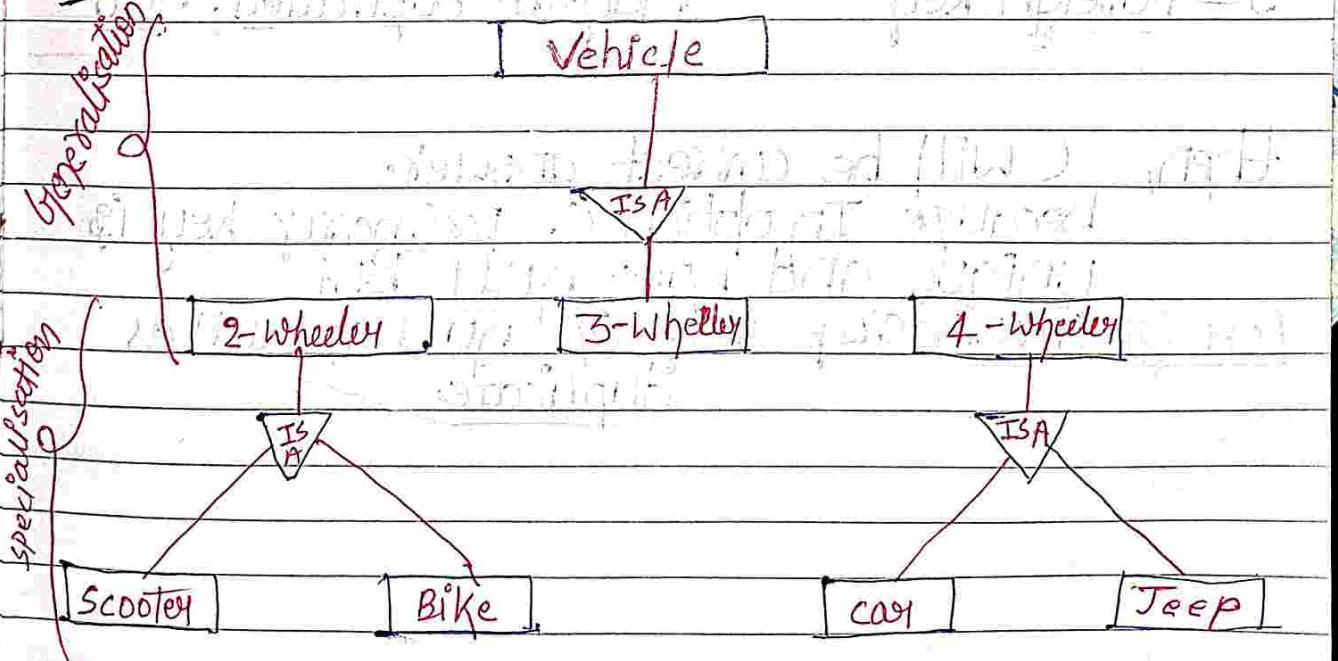
e.g. :-

Parent table (A, B) is child table ; Foreign key

1 of dupl	RNO	Name	Age	class	Course	RNO
1 - cat	1	A	17	101	CSE	1
3	2	B	15	102	CD	2
new	3	A	18	103	TOC	3

Let us suppose, if we want to delete record of student whose rollno. is 2 ~~then~~ from parent table then there need be to delete the related data from child table simultaneously.

Note



Ques :- Consider a relation R(A,B) where A is primary key and B is foreign key referencing the same relation then which of the following row sequence can be inserted successfully into R.

- a) $\begin{matrix} A & B \\ \text{(1, null)} & \end{matrix}$ $\begin{matrix} A & B \\ (2, 1) & \end{matrix}$ $\begin{matrix} A & B \\ (2, 2) & \end{matrix}$ $\begin{matrix} A & B \\ (3, 2) & \end{matrix}$
- b) $\begin{matrix} A & B \\ (\text{null}, 1) & \end{matrix}$ $\begin{matrix} A & B \\ (1, 2) & \end{matrix}$ $\begin{matrix} A & B \\ (2, 3) & \end{matrix}$ $\begin{matrix} A & B \\ (3, 4) & \end{matrix}$
- c) $\begin{matrix} A & B \\ (1, \text{null}) & \end{matrix}$ $\begin{matrix} A & B \\ (2, 1) & \end{matrix}$ $\begin{matrix} A & B \\ (3, 2) & \end{matrix}$ $\begin{matrix} A & B \\ (4, 2) & \end{matrix}$
- d) all.

→ On keeping the concept of primary key and foreign key as :-

A = Primary key → must be unique, non-null, non-duplicate.

B = Foreign key → May be duplicated, null.

then, C will be correct answer.

because In option C, primary key is unique and non-null but foreign secondary key is null as well as duplicate.

Ques Consider a relation $R(A, B)$ where A is a primary key and B is a foreign key referencing A with on delete cascade. Consider the following relation instance on R:

$R(A, B)$

- (2 4)
- (3 4)
- (4 3)
- (5 2)
- (7 2)
- (9, 5)
- (6 4)

What are the tuples that must be additionally deleted to preserve referential integrity constraint when the tuple (2, 4) deleted?

- (3, 4) (4, 3)
- (3, 4) (4, 3) (6, 4)
- (5, 2) (7, 2)
- (5, 2) (7, 2) (9, 5)

First which tuple is deleted



* ER model to Relational model

i) Entity set:

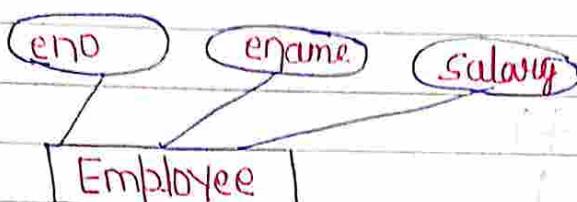


Table: Employee

eno	ename	salary
-----	-------	--------

- An entity set is mapped with a relation.
- The key attribute of an entity becomes primary key of the relation.

ii) Entity set with composite attributes

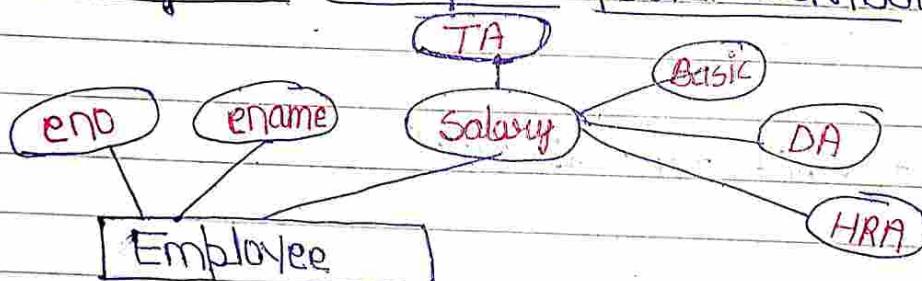
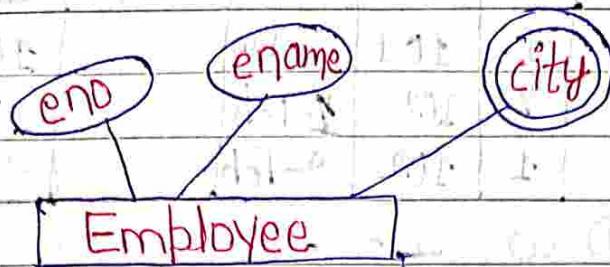


Table :- Employee

eno	ename	Basic	TA	DA	HRA
1	A	5000	3000	5000	1200

- The salary attribute of a relation includes the simple attributes Basic, TA, DA, HRA etc.

iii) Entity set with multivalued attribute

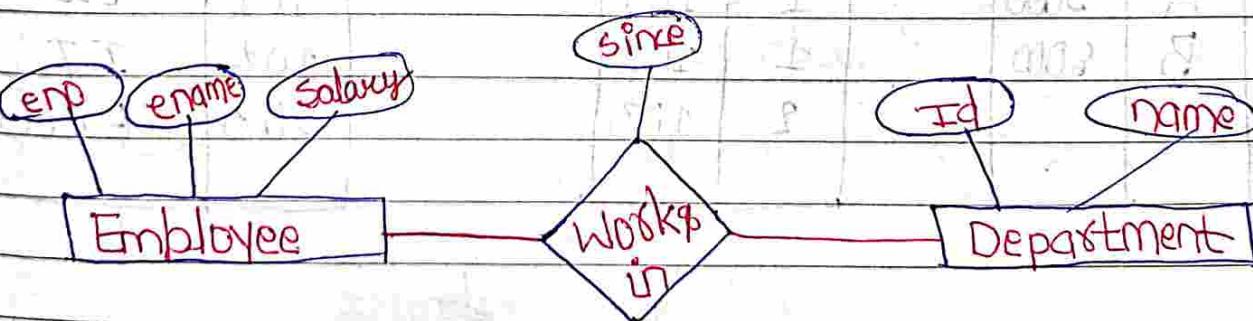


Employee ~~Employee - address~~ Employee-address

e.no	ename	e.no	city
1	A	1	Hyd
2	B	1	Bang
		2	Bang
		2	Pune

- If an entity contains multivalued attributes.
- It is represented with two relations one with all simple attributes and the other with key and multivalued attributes.

iv) Translating relationship set into a table



EmployeeForeign key Works inForeign keyDepartment

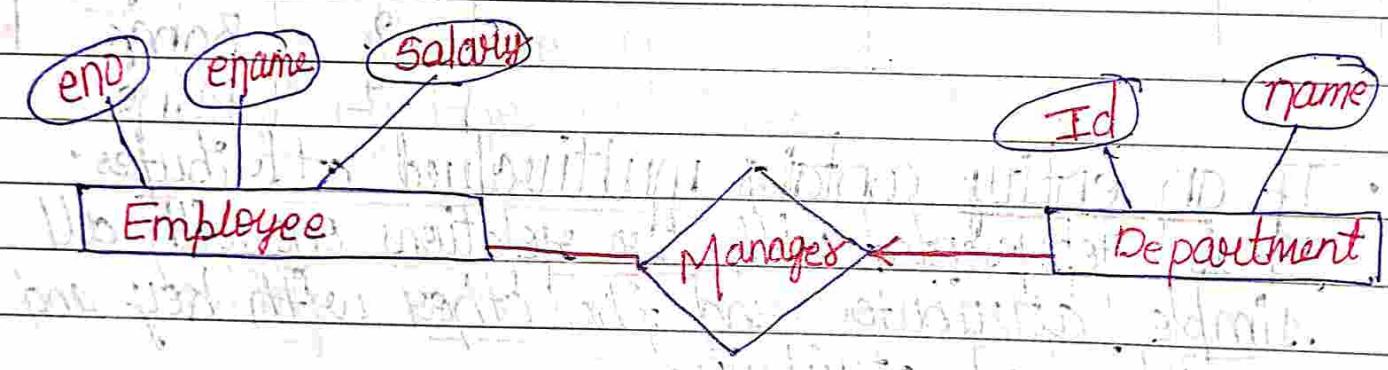
eno	ename	Salary
1	A	5000
2	B	9000

eno	Did	Since
1	101	1-Jan
2	102	1-Feb
	102	2-Feb

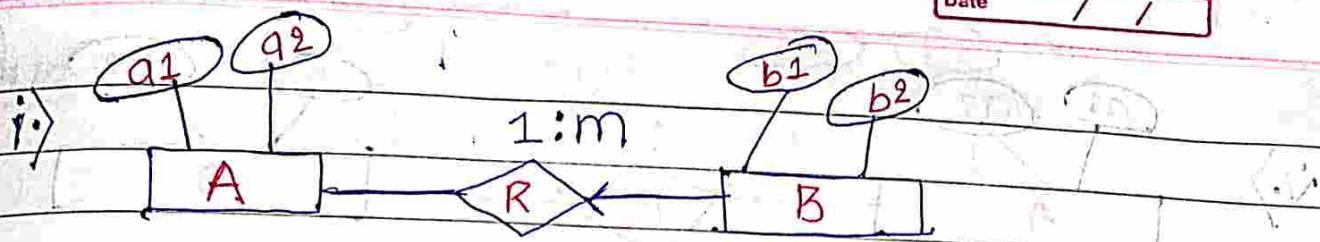
Id	Name
101	CSE
102	IT

v) Relationship set with key constraint

→ Each department is required to have at most one employee as a manager.

Employee Manager Department

eno	ename	Salary	eno	Id	name	Id	name
1	A	5000	1	101		101	CSE
2	B	8000	1	102		102	IT
			2	103		103	ECE

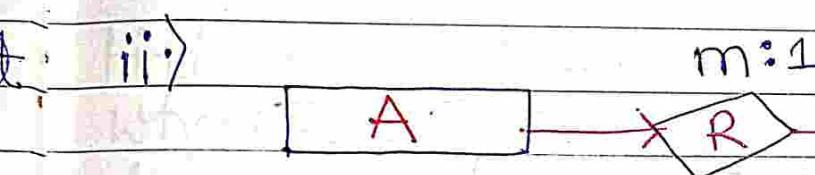


2 tables
⇒ A, BR

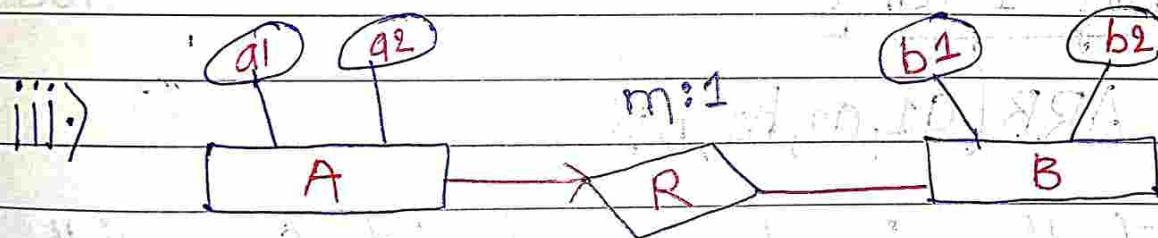
A(a₁, a₂), BR(b₁, b₂, q₁)

Foreign key

It is a foreign key
of B because
a₁ belongs to A.

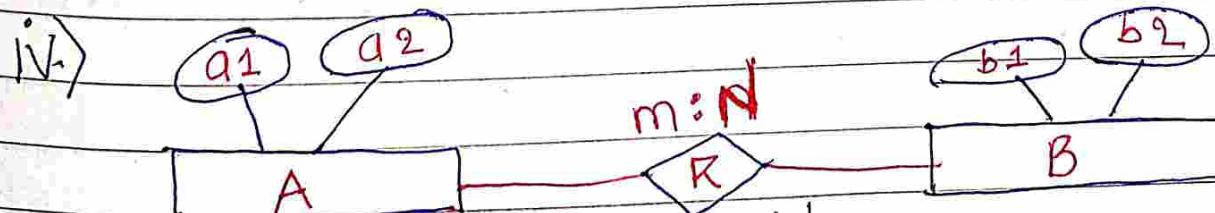


2 tables
⇒ (AR, B).

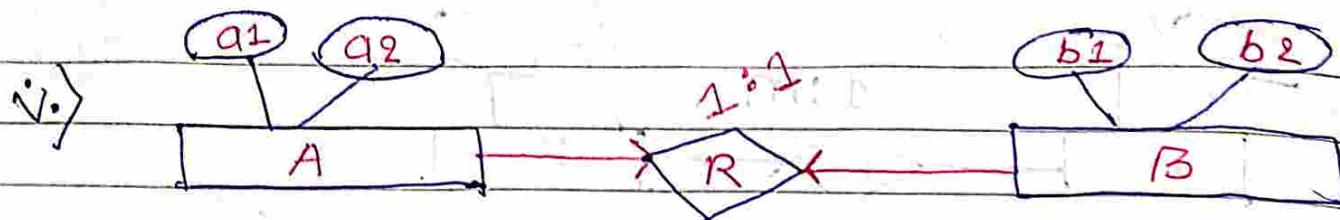


2 tables
⇒ AR(a₁, a₂, b₁), B(b₁, b₂)

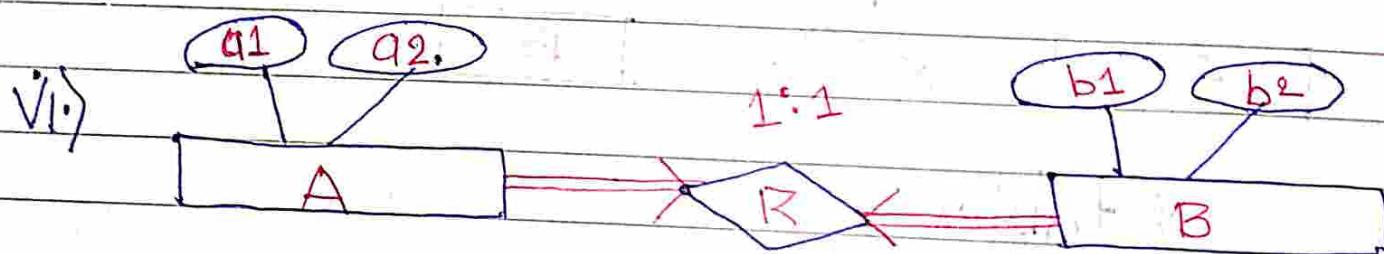
It is a foreign key for A
because b₁ belongs to B.



~~2 tables~~ 3 tables.
i.e.; A(a₁, a₂) R(a₁, b₁) B(b₁, b₂)



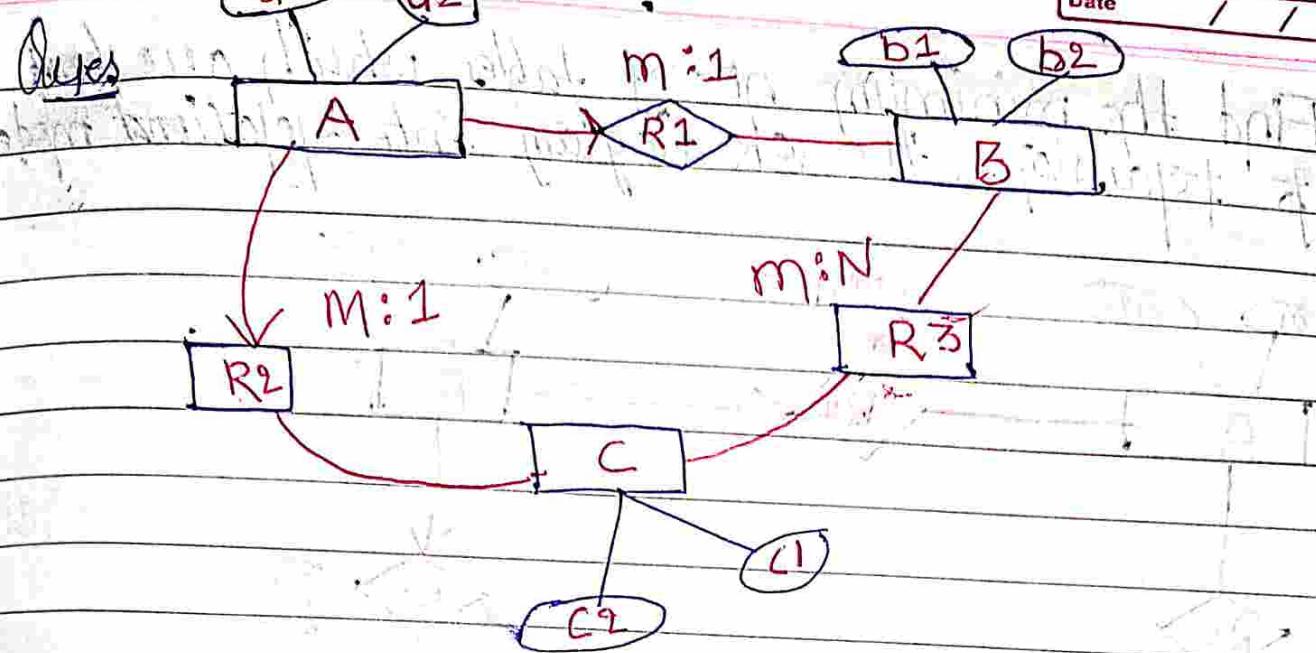
2 tables
i.e; $AR(a_1, a_2, b_1)$, $B(b_1, b_2)$ or
 $A(a_1, a_2)$; $BR(b_1, b_2, a_1)$



Only 1 table

i.e; $ARB(a_1, a_2, b_1, b_2)$

Note :- If there is a key constraint from both the sides of an entity set with total participation then we represent that binary relationship using single table.



Find the minimum no. of tables that one possible when you translate the above E-R diagram into relational model.

Sol: 4 tables

Foreign key because

b_1 belongs to B &
 c_1 belongs to C .

i) $AR_1R_2(a_1, a_2, b_1, c_1)$

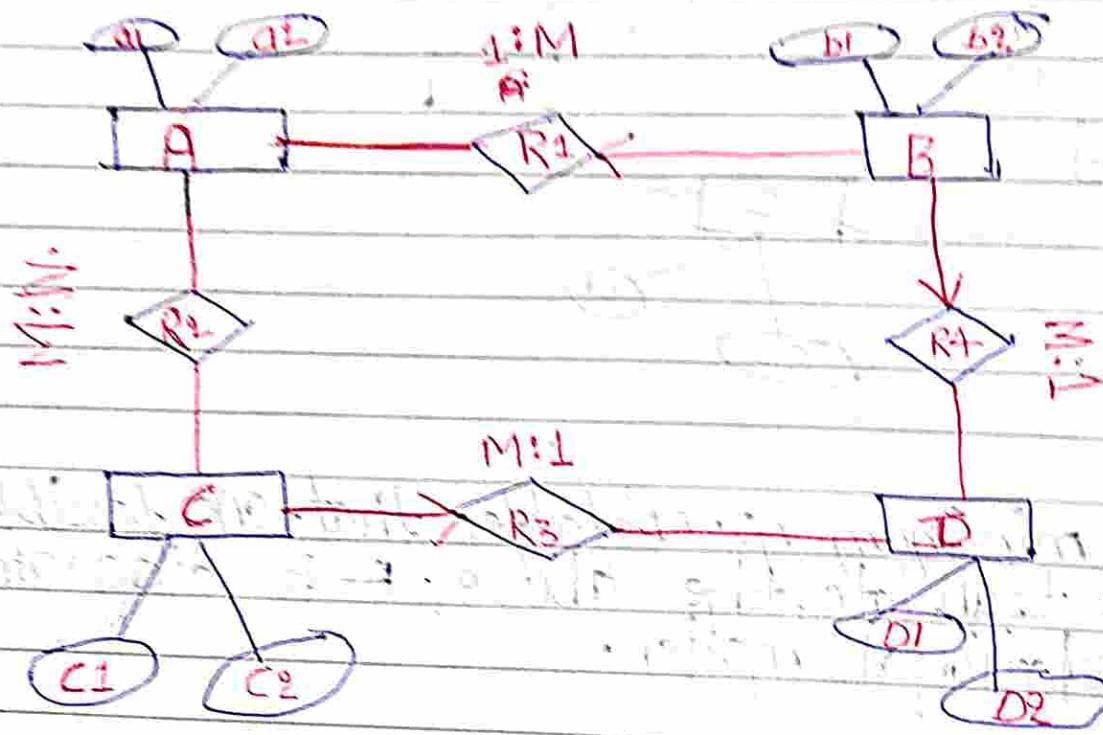
ii) $B(b_1, b_2)$

Foreign key

iii) $R_3(b_1, c_1)$

iv) $C(c_1, c_2)$

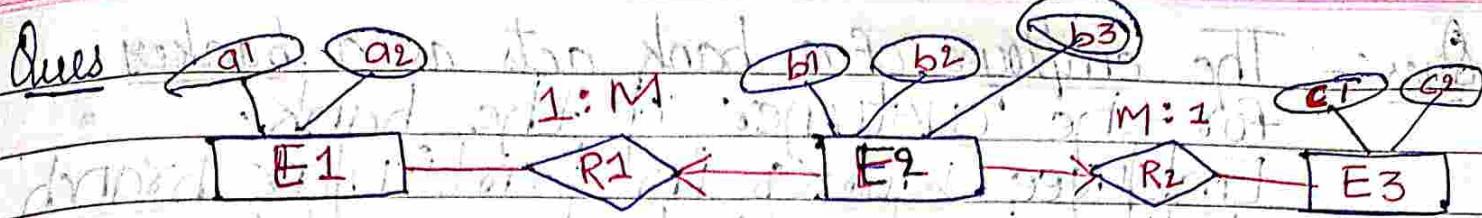
Ques Find the minimum no. of tables which are required to translate the ER diagram into relational model.



Ans:- 5 tables

foreign key since a_1 belongs to A
 d_1 belongs to D
 c_1 belongs to C

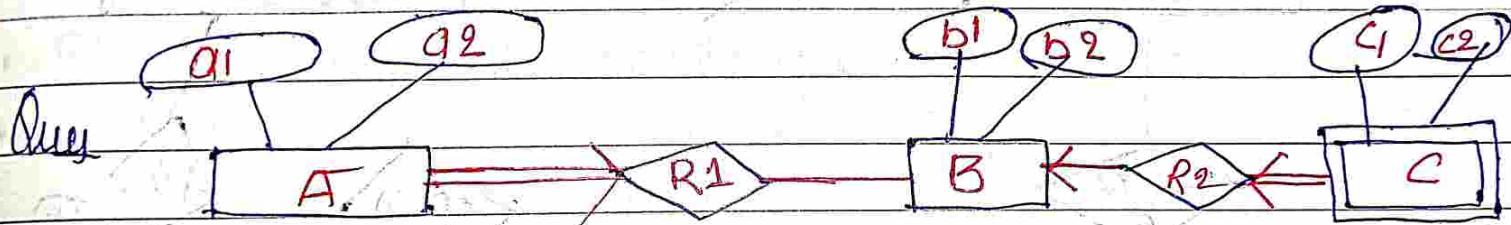
- i) $B \sqcup R_1 \sqcup R_4 \sqcup R_3 (b_1, b_2, a_1, c_1, d_1)$
- ii) $A(a_1, a_2)$
- iii) $R_2(a_1, c_1)$ Foreign key
- iv) $D(d_1, d_2)$
- v) $C R_3 (c_1, c_2, d_1)$ Foreign key.



Find total no. of table (minimum).

Sol:- E1(a₁, a₂) E2(b₁, b₂, b₃, a₁, c₁) E3(c₁, c₂)

so, min^m no. of tables = 3.

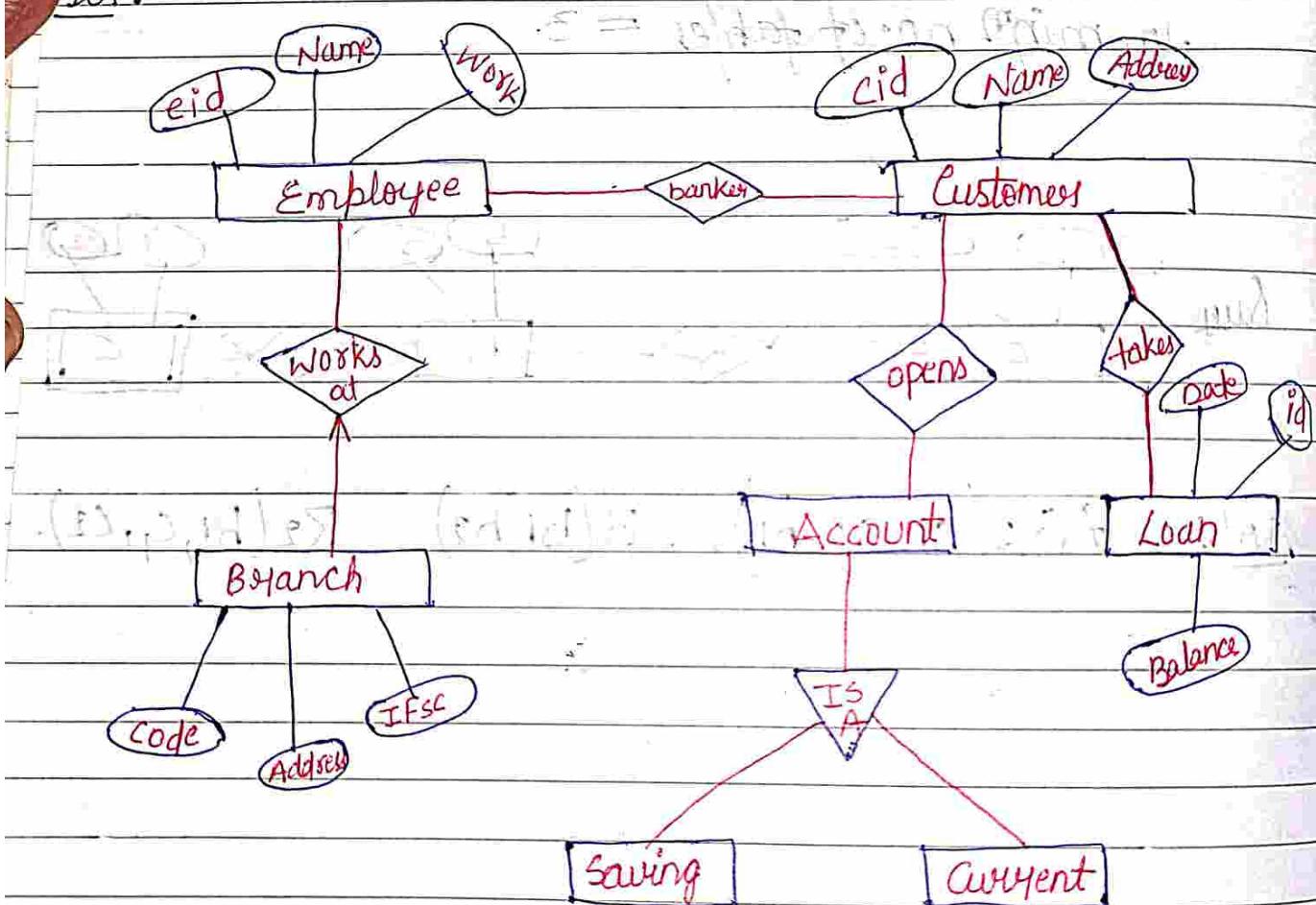


Sol:- AR1(a₁, a₂, b₁) , B(b₁, b₂) R2(b₁, c₁, c₂).

Ques :- The employee of a bank acts as a banker for the customer of the bank.

Employee works in particular branch of the bank. Customer can take loan from the bank. At the same time, a customer can open either saving or current account.

Sol :-



Here, Entity

Employee
Customer
Branch
Account
Loan

Attributes

eid, Name, work.
cid, Name, Address.
code, Address, IFSC
saving, current as type
note id, balance.

CREATE DATABASE

→ It is used to create database.

Syntax :- CREATE DATABASE database_name;

Eg :- CREATE DATABASE my_db;

Guidelines for database name

- Name should be start with an alphabet
- Blank spaces and single quotes, are not allowed
- Reserve words of that RDBMS| DBMS cannot be used as database name.

Steps :-

- Open mysql command line client - Unicode.

MySQL Command Line Client

mysql > CREATE DATABASE my_db;
Query OK, 1 row affected (0.00 sec)

It signify that database is created.

This is used to tell your RDBMS/DBMS that you want to use this database.

Syntax :- USE database_name;
Ex:- USE my_db;

Steps :- Open mysql command line client

```
mysql > CREATE DATABASE my_db;
Query OK, 1 row affected (0.00 sec)
```

```
mysql > USE my_db;
Database changed
```

✓ Feedback
It is saved in this database.

CREATE TABLE

Syntax :- CREATE TABLE table_name
 (column_name1 data_type(size),
 column_name2 data_type(size),
 column_name3 data_type(size)
);

(column_name1 data_type(size),
 column_name2 data_type(size),
 column_name3 data_type(size)
);

Q8

(CREATE TABLE table_name

 column_name1 data_type(size)[constraints],
 column_name2 data_type(size)[constraints]
);

NOT NULL
PRIMARY
KEY

Keyword to create
table

eg:- CREATE TABLE (my_tab)

field of name Varchar(30),
 column roll number (4),
 name address Varchar(100);
 datatypes (Varchar, number).
 Semicolon indicates that the above statement is ended.

Guidelines for Create Table

- Table name should start with an alphabet.
- Blank spaces and single quotes are not allowed.

DESC

- ⇒ This is used to describe your table. DESC only describes structure of table, not the information (rows) inside table.
- DESC is short form of Describe.

Syntax : DESC tablename;

eg:- DESC my_tab;

Notes

~~SHOW DATABASES | TABLES~~

SHOW DATABASES | TABLES

- SHOW DATABASES — This command is used to view all the databases name;

Syntax :- SHOW DATABASES;

- SHOW TABLES — This command is used to view all the tables of current database

Syntax :- SHOW TABLES;

It displays all the database created

mysql> SHOW DATABASES;

| Database |

| db |

| information_schema |

| my_db |

| performance_schema |

9 rows in set (0.00 sec)

will open my database

It named my db

mysql> USE my_db;

mysql> SHOW TABLES; → If shows all the tables present in my current database :-

| my_tab |

USE my_db

Data Types

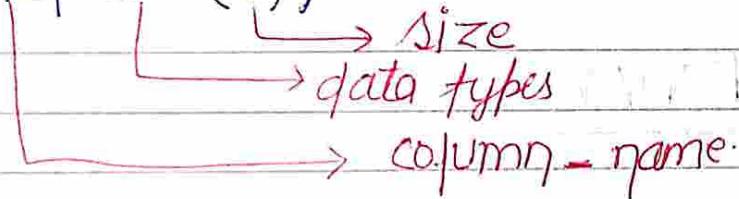
INT OR INTEGER

:- It holds whole number between -32768 and 32767 either it is negative or positive.

- It cannot hold any decimal value.
- The maximum number of digits may be specified in parenthesis. by size.

syntax :- column-name INT(size);

e.g. :- you INT(5);



DEC OR DECIMAL

→ column-name DECIMAL(T,D)

where, T = Total digits

Range = 1-65

- It holds fixed point numbers.
- Size is the total number of digits and p is the number of digit after decimal point.
- ✓ The decimal pt and -ve sign are not counted in size.

syntax :- column-name DECIMAL(size,p)

e.g. :- price DECIMAL(4,2)

→ size ↗
of digit
No. after decimal

56.89 or -12.34
→ -ve sign and
decimal pt is not
counted

NO. of Digit	NO. of Bytes
0	0
1-2	1
3-4	2
5-6	3
7-9	4

Counting of no. of digit starts from left of Dec. Imp.

$$\text{Eg: } \frac{1699.81}{2+1} = 3 \text{ bytes}$$

no. of digit 4
no. of bytes 2

KRISHNA

Page No. 1

Date

Eg: - 14564345.66999.81

No. of digit 9
No. of bytes 4

4
2
1

$$\text{Total bytes} = 4+9+1=7$$

• CHAR OR CHARACTER

:- It holds a fixed length string (can contain letters, numbers, special character etc).

- The fixed size is specified with in parenthesis.

Syntax :- column_name CHAR(20);

eg :- name CHAR(20);

size

datatype

column_name

• VARCHAR

:- It holds a variable length, string (can contain letters, numbers, and special character).

- The maximum size is written in parenthesis.
- If can store upto 255 characters.

Note :- If we put a greater value than 255 then it will be converted to a TEXT type.

Syntax :- column_name VARCHAR(size);

eg :- name VARCHAR(50);

size

datatype

column_name

• TEXT

:- It holds a string with a maximum length of 65,535 characters.

Syntax :- column-name TEXT;

Eg :- address TEXT;

→ column-name → data types.

• DATE

:- It displays Date values in yyyy-mm-dd format.

Syntax :- column-name DATE;

Eg :- age DATE;

For DOB (Date of birth).

• TIMESTAMP

:- It displays date and time.

Syntax :- column-name TIMESTAMP;

Eg :- login_dt TIMESTAMP;

* CREATE TABLE named as saket

mysql > USE my_db;
Database changed

mysql > SHOW TABLES;

| my-tab { table exists }

mysql > CREATE TABLE saket

→ student_id INT(5),
 → name VARCHAR(50),
 → address TEXT,
 → dob DATE,
 → fee DEC(10,2) *NO. of digit after decimal*
 →); *size*

Query OK, 0 rows affected.

mysql > DESC saket;

Field	Type	Null	Key	Default
student_id	int(5)	Yes		NULL
name	VARCHAR(50)	YES		NULL
address	text	YES		NULL
dob	date	YES		NULL

Database Manipulation Languages

These commands are applied on data stored in the table.

- Different DML commands are :-

Statement	Description
SELECT	Retrieves data from the database, enter new rows, changes existing rows, and removes unwanted rows from tables in the database resp.
INSERT	
UPDATE	
DELETE	
MERGE	

SELECT

- It is also considered as a Data Query Language (DQL).

Select data or

- To retrieve data from database we use SELECT command.
- Table name is used with FROM clause.
- Define condition with WHERE clause.
- To select all columns.

→ SELECT * FROM Persons WHERE P_ID <= 5;

To select specific columns

→ Select P_ID, Pname FROM Persons WHERE P_ID <= 5;

1. To Select all columns from the table

Syntax :- $\text{SELECT } *$ FROM table-name;

↙
all

2. To Select particular columns from the table

Syntax :- $\text{SELECT column}_1, \text{column}_2, \dots$
FROM table-name;

eg:- $\text{SELECT name, city, pin}$ ↗
FROM bholi; ↘
particular
table

3. To SELECT with LIMIT

→ It is used to specify some number of records to display.

Eg :-

Syntax :- $\text{SELECT column-name}_1, \text{column-name}_2$
or $\text{SELECT } *$
FROM table-name
LIMIT record-number.

eg:- → $\text{SELECT } *$ FROM bholi
→ LIMIT 5; ↗
It displays data or records

of table upto 5.

Syntax :- $\text{SELECT } *$ FROM bholi
→ LIMIT 5; ↗
no. of records
It displays data or
records of table from
index no. 5 to 6.

Syntax :- $\text{SELECT } *$ FROM bholi
→ LIMIT 3, 4; ↗
no. of records
Index no. 3

Note

• INSERT INTO

a) Insert with specifying column name

:- The INSERT INTO statement is used to insert new records/rows in table.

Syntax :- `INSERT INTO table-name (col-name1, col-2, ...)`
`VALUES (value1, 'value2', 'value3', ...);`

If insert value into table datatype (data type (so, value is within ''))

table-name → These are all column present in table

Eg :- `INSERT INTO my-tab (stu_id, name, address)`
`VALUES (05, 'Anu', 'Delhi');`

Rules :- column and value datatype should be same.

- Any value that goes into a VARCHAR, CHAR, DATE or TEXT column has single quotes around it.
- No need of quotes for numeric values for INT, DEC.

Note • To see all the information inserted into table by using :

`SELECT * FROM table-name;`

→ Insert into table without specifying column name

Syntax :-

INSERT INTO table-name
VALUES ('value1', 'value2', 'value3', value4)

⇒ Sequence is very important in this case.

Note:- To see all the information inserted into table by using :-

SELECT * FROM table-name;

c) changing the order of column

stu-id	name	city	pin
1	lakshmi	hyd	500001

Syntax :-

INSERT INTO table-name (column2, column1, column3, ...)
VALUES ('value2', value1, value3);

see above operation

SELECT * FROM table-name;

d) Insert into table with specifying columns

Syntax :-

```
INSERT INTO table-name(column1, column2,
column3)
VALUES ('value 1', 'value 2', 'value 3');
```

Note :- If there are 4 columns which are present in table but we want to insert data in only 3 columns.

- table → bholi

Eg :-

stu_id	name	address	pin
1	saket		

→ INSERT INTO bholi (stu_id, name, pin)
 → VALUES (1, 'saket', '211019');

Query OK

mysql > SELECT * FROM bholi;

stu_id	name	address	pin
1	saket	NULL	211019

We do not pass any values for column address of the table bholi.

c) Insert multiple record at one time

syntax :-

```
INSERT INTO table-name (column1, column2, column3, col4)
VALUES (value1, 'value2', 'value3',
        value4),
       (value1, 'value2', 'value3', value4),
       (value1, 'value2', 'value3', 'value4);
```

If is used to add multiple record within the column of table at one time.

to see ~~the~~ insertion / data

```
.SELECT * FROM table-name.
```

Single quotes Problem

Eg :- `INSERT INTO table-name (column 1, column 2, column 3)
VALUES (142, 'K.K's company', 'Delhi');`

Due to this single quote, it shows error because system does not fetch detail due to multiple quotes.

There are two ways to overcome this problem,

i) USE backslash

`:- 'K.K\'s company'`

→ backslash

ii) Use two times single quotes

`:- 'K.K''s Company'`

→ single quotes → table-name

Eg :- `INSERT INTO bholi (id, c-name, city)
VALUES(142, 'K.K's company', 'Delhi');`

Or,

`INSERT INTO bholi (id, c-name, city)
VALUES (142, 'K.K''s company', 'Delhi');`

~~KRISHNA~~

WHERE Clause :- It is used to filter specific data.

→ WHERE is used to search for a specific data.

- It extracts only those records which satisfy specified condition

Syntax :- a) Specific data from all column

SELECT * FROM table-name
WHERE column-name [operator] value;

eg:- SELECT * FROM bholi → table-name
WHERE stu_id = 9; value operator

stu_id	name	city	pin
9	Saket	Sasaram	821115

b) Specific data from specific column

(multiple
columns
may
possible)

Syntax :- SELECT ~~rows~~ FROM table-name,
WHERE column-name = 'value'. operator

eg:- SELECT name FROM bholi → table-name
WHERE name = 'Anu' value

name
Anu

→ column-name

Note :- Value can be text or numeric. If it is text then put it into single quote.

operator

operator	Description
=	Equal
<> or !=	Not equal
>	Greater than
<	Less than
>=	Greater than or equal
<=	Less than or equal
BETWEEN	Between an inclusive range
LIKE	Search for a pattern
IN	To specify mul. values for column

Equal

SELECT * FROM bholi
 WHERE stu_id = 9;

table name

value

column name

operator

Not Equal

SELECT * FROM bholi
 WHERE stu_id != 9;

Greater than

SELECT * FROM bholi
 WHERE stu_id > 9;

operator

or WHERE name > 'R'

Less than

SELECT * FROM bholi
 WHERE stu_id < 9;

operator

- `SELECT * FROM bholi WHERE name > 'J'` → It displays all the names which starting with letter which are greater than J.
 - BETWEEN

\rightarrow It selects values within a range.

SELECT * FROM bholi

WHERE column-name BETWEEN value1 AND value2

a) Between text

Syntax: SELECT * FROM bholi;

**WHERE name BETWEEN 'J'
AND 'S'**

IN operator

→ It specifies multiple values in a WHERE clause

Syntax :-

```
SELECT * FROM table_name  
WHERE column_name IN ('Value1', 'Value2', ...)
```

eg:- `SELECT * FROM bholi WHERE name IN ('Anu', 'Sonu');`

stu_id	name	city	pin
2	Anu	Delhi	11102
3	Sonu	Mumbai	546879
4	Anu	Kolkata	568974

It displays all the records in which

- LIKE :- It is used to search for a specified in a column.

Syntax :-

SELECT * FROM table-name
WHERE column-name
LIKE 'pattern';

eq :- ELIRE pallash;

~~58~~ • SELECT * FROM bholi
WHERE name LIKE '%nu';

Wildcards :- It is used to search data with incomplete information.

% → zero or more character

→ One single character

`[]` → Ranges of character to
`[^]` → Matches all like a

• $\%o$ → zero or more char
 • 'Greek%' → All start with Greek
 • '% shows' → Ending with shows
 • '%.sh%' → contain with sh

- One single character
 • 'Show' → starting with show
 • 'eek' → any character
 then eek

Page No. KRISHNA
 Date / / / /

zero

NULL Value :- It is diff from value.

→ NULL values represent missing unknown data

• A field with NULL value is one that has been left blank during creation.

a) IS NULL

→ It is used to test for empty values.

→ This is used to select only the records with NULL Values in the column.

Syntax :-

SELECT * FROM table_name
 WHERE column_name IS NULL;

It displays records with NULL Value.

eg:-

stu_id	name	city	pin
1	Saket	Patna	211019
2	Ankit	NULL	821115
3	Sumit	Sasaram	821115

mysql > SELECT * FROM bholi
 → WHERE city = 'NULL'; X It shows an error

mysql > SELECT * FROM bholi
 → WHERE city IS NULL; ✓ correct way to display

stu_id	name	city	pin	display
2	Ankit	NULL	821115	NULL Val

match. character NOT enclosed within the bracket

b) IS NOT NULL

→ This is used to select only the records with no NULL values in the column.

Syntax :- `SELECT column_name FROM table-name
WHERE column_name IS NOT NULL;`

Eg :- mysql> `SELECT * FROM bholi
WHERE city IS NOT NULL;`

table :- bholi

stu_id	name	city	pin
1	Saket	Prayagraj	211019
2	Sumit	Sasaram	821115

c) IS NULL Operator

:- It is used to test for empty values.

Eg :- `SELECT CustomerName, ContactName,
Address
From customers
WHERE Address IS NULL;`

✓

It lists all customers with a NULL value in the "Address" field.

AND Operator

→ AND operator displays a record if both the first condition AND the second condition are true.

Syntax :- `SELECT * FROM table-name WHERE column-name = 'value'
AND column = 'value';`

If displays those columns which satisfy both conditions.

table :- bholi

eg :-

stu_id	name	city
1	Saket	Prayag
2	Shubham	Sasaram
3	Saket	Patna

mysql > `SELECT * FROM bholi WHERE name = 'Saket' AND stu_id = 3;`

It shows only those records which satisfy both conditions.

stu_id	name	city
3	Saket	Patna

OR operator

→ OR operator displays a record if either the first condition OR the second condition is true.

Syntax :-

```

    SELECT * FROM bholi
    WHERE column_name = 'Value'
    OR column_name = 'Value';
  
```

1st

2nd condition

It displays if either 1st OR 2nd condition is true. (any condition is true)

Combinations of AND & OR

Syntax :-

```

    SELECT * FROM table_name
    WHERE column_name = 'Value'
    AND(column_name = 'Value' OR
        column_name1 = 'Value');
  
```

eg:- `SELECT * FROM bholi
WHERE name = 'Anu' OR
AND (stud_id = 5 OR address = 'Kolkata');`

- name must be Anu but stud_id may be 5

address may be Kolkata.] at least one of them is true.

table :- bholi

	stu_id	name	city	pin
either	2	Anu	Delhi	111025
second	3	Sonu	Mumbai	546879
iii	4	Anu	Kolkata	568974

eg:-

mysql > SELECT * FROM bholi
WHERE name = 'Anu'
AND (city = 'Delhi' OR stu_id = 4);

	stu_id	name	city	pin
either	2	Anu	Delhi	111025
condition	4	Anu	Kolkata	568974

Here, [name must be Anu] but city and stu_id are connected with OR operator
{ so, city may be Delhi or not
stu_id may be 4 or not

NOT LIKE

syntax :-

SELECT * FROM table_name
WHERE column_name NOT LIKE 'pattern';

eg :- SELECT * FROM bholi
WHERE name NOT LIKE '%nu';

It displays all the name
not ending with nu.

~~ORDER BY clause~~

NOT Operator

:- NOT operator displays a record if the condition(s) is NOT TRUE.

Syntax :-

SELECT column1, column2, ...

FROM table_name

WHERE Condition NOT Condition

Eg:-

SELECT * From customers
Where NOT country = "Germany"

ORDER BY

- ASC → It sorts in ascending order.
- DESC → It sorts in descending order.

Syntax :-

SELECT * FROM table_name
ORDER BY column_name ASC or DESC

• ORDER BY descending order

- Using keyword :- DESC

It sorts in descending order

Syntax :- [SELECT * FROM table-name
ORDER BY column-name DESC ;]

Eg :- SELECT * FROM bholi
ORDER BY stu_id DESC;

table :- bholi

stu_id	name	city
3	S	Ptna
4	K	Dlhi
2	A	Kol

After

DESC →

stu_id	name	city
4	K	Delhi
3	S	Patna
2	A	Kolkata

• ORDER BY Ascending order (By default)

- Using keyword :- ASC

It sorts in Ascending Order

Syntax :-

SELECT * FROM table-name
ORDER BY stu_id ASC ;

Eg :-

SELECT * FROM bholi
ORDER BY stu_id ; (ASC is optional)

Eg :-

SELECT * FROM bholi
ORDER BY stu_id DESC, name ASC;

Constraint

→ NOT NULL

→ Unique Key

→ Primary Key

Page No.

KRISHNA

1. NOT NULL

→ NOT NULL constraint enforces a field to always contain a value.

This means that you cannot insert a new record, or update a record without adding a value to this field.

Eg: → CREATE TABLE Student

Fields : Name varchar(30),
Roll integer(5),
Mobile_no integer(10) NOT NULL

These may be empty.

It signifies that mobile no. can't be empty.

mysql > DESC student;

Field	Type	NULL	Key	Default	Extra
Roll	int	YES			
Name	Varchar(30)	YES			
Roll	int(5)	YES			
Mobile_no	int(10)	NO			

• Those must be only value.

• It is compulsory to fill data or value in mobile-no.

2. Unique key (UNIQUE KEY)

- It is a type of constraint; it uniquely identifies each record in a database table.
- There can be many ~~unique~~ UNIQUE constraint per table.
- A UNIQUE key column can contain NULL values.

Eg: → CREATE TABLE student

```
Name varchar(30) UNIQUE KEY;
Roll integer(5);
Mobile_no integer(10) UNIQUE KEY;
);
```

- If uniquely identify.
- May contain null value
- More than one UNIQUE KEY can be possible

mysql> DESC student;

Field	Type	NULL	Key	Default	Extra
Name	varchar(30)	YES	UNI	NULL	
Roll	int(5)	YES		NULL	
Mobile_no	int(10)	YES	UNI		

mysql> SELECT * FROM student

→ Values ('A', 101, 9876543212),
 ('B', 102, 9654387921);

It shows duplicate entry 'A' for Name field

3. PRIMARY KEY

→ It also uniquely identifies each record in a database value.

- PRIMARY KEY must contain a ~~UNIQUE~~ values.
- A primary key column cannot contain NULL.
- Each table can have ONE Primary Key.

Syntax:-

```
CREATE TABLE student  
(  
    Name varchar(30),  
    Roll integer(5) NOT NULL,  
    Mobile_no integer(10),  
    PRIMARY KEY (Roll))
```

↑ used to
write primary
key

Q8

```
CREATE TABLE student  
(  
    Name varchar(30),  
    Roll integer(5) NOT NULL PRIMARY KEY,  
    Mobile_no integer(10))
```

mysql > DESC student;

Field	Type	NULL	Key	Default
Name	VARCHAR(30)	YES		
Roll	int(5)	NO	PRIMARY	
Mobile_no	int(10)	YES		

~~(auto increment)~~ * AUTO INCREMENT

⇒ AUTO INCREMENT

* AUTO INCREMENT

a) General

⇒ AUTO INCREMENT is used to generate a unique number, when a new record is inserted into a table.

Syntax :-

CREATE TABLE .table_name

column_name int NOT NULL AUTO_INCREMENT,
 column_name1 varchar(50) NOT NULL,
 column_name2 varchar(50),
 PRIMARY KEY (column_name)

It is the
 column_name corresponding
 to auto increment

Eg :- CREATE TABLE t-tab

→ (→ table name

→ Roll int NOT NULL AUTO_INCREMENT,
 → Name Varchar(30),
 → city Varchar(50),
 → PRIMARY KEY (Roll) → It is that
 →); column-name which are related to Auto Increment

Field	Type	NULL	Key	Default	Extra
Roll	int	No	PRI	NULL	auto increment
Name	Varchar(30)	YES		NULL	
city	Varchar(50)	YES		NULL	

eg:- mysql > INSERT INTO t-tab(Roll, Name, City)
 → VALUES (1, 'SAKET', 'PRAYAGRAJ'),
 IN AUTO_INCREMENT,
 If we pass NULL
 then it generates next values instead of NULL)

mysql > SELECT * FROM t-tab;

Roll	Name	City
1	Saket	PRAYAGRAJ
2	Aman	Sasaram
3	Bholi	Patna

We pass NULL
 value but
 it generates
 2 and 3 in
 the place of NULL

Note :- If I pass NULL value for AUTO_INCREMENT column then it ignores NULL and generates unique numbers for them.

Case-I If both values are NULL

mysql > INSERT INTO table_name
 VALUES (NULL, 'SHUBHAM', 'Sasaram'),
 (NULL, 'SAKET', 'Prayag');

Here, It ignores NULL value and generates unique number in an order for them

Roll	Name	City
1	SHUBHAM	Sasaram
2	SAKET	Prayag

Case :- II If 1st value is NULL and 2nd value is 1.

mysql > INSERT INTO table_name

VALUES

(NULL, 'Shubham', 'Sasaram'),

(1, 'Saket', 'Patna');

Here, It ignores NULL value and generate 1 but the second which I pass here is already 1. So, (1, 'Saket', 'Patna') is automatically erased from database.

Roll	Name	City
1	Shubham	Sasaram

Case :- III If 1st value is 1 and 2nd value is NULL

mysql > INSERT INTO table_name

VALUES

(1, 'Shubham', 'Patna'),

(NULL, 'Saket', 'Delhi');

Here, The first auto-increment column value is 1 and the second value is NULL then auto-increment ignores NULL value and generate 2.

Roll	Name	City
1	Shubham	Patna
2	Saket	Delhi

AUTO_INCREMENT with particular number

→ By default, AUTO_INCREMENT starts from 1.

But, if we want to start number from the given particular number then we use syntax :-

ALTER TABLE table_name AUTO_INCREMENT = 10;
Keyword

Syntax :-

```
CREATE TABLE table_name
(
    id int NOT NULL AUTO_INCREMENT PRIMARY KEY,
    name varchar(50)
);
```

mysql> DESC table_name;

Field	Type	NULL	KEY	DEFAULT	EXTRA
id	int	NO	PRI	NULL	
name	varchar(50)	YES		NULL	auto_increment

mysql> ALTER TABLE table_name AUTO_INCREMENT=10;

mysql> INSERT INTO table_name (name)
VALUES ('RAM'),
('Ali');

id	name
10	RAM
11	Ali

c) AUTO_INCREMENT WITHOUT specifying column name

Syntax :-

```
• INSERT INTO table_name  
VALUES (NULL, 'Rohan');
```

mysql> SELECT * FROM table_name.

<u>id</u>	<u>name</u>
10	Ram
11	Ali
12	Rohan

Aliases (AS)

→ Aliases are used to temporary rename column name or a table name.

For Table

Syntax :- SELECT column_name
FROM table_name AS alias_name;

Ex :- SELECT name FROM student AS vidyarthi;

Here, This statement change the name of table from student to vidyarthi

For column

Syntax :- SELECT column_name
AS alias_name FROM table_name;

Eg :- SELECT name AS student_name FROM student.

Here, It changes the ~~name~~ column name of name from student table to student_name.

Note :-

student_name
or
"student name" } all are same

Student name

~~Arith~~

Arithmetic Operators (*, /, +, -)

:- We can operate all arithmetic operators in mysql.

Syntax :- → `SELECT column_name, column_name Operator value
FROM table_name;`

→ `SELECT column_name, column_name Operator value
FROM table_name;`

Eg :- `SELECT name, cost, cost + 100 FROM bholi`

operator
value

↓
column-name

↓
table name

Or

`SELECT name, cost, cost + 100 AS "New Cost",
FROM itemstab.`

name	cost	id	Newcost
laptop	5000	11	6000
mobile	8000	12	9000
tv	12000	13	13000
refrigerator	15000	14	16000

Note :- For Operator value
value of

→ We can add two or more columns in any columns.

$$\text{Newcost} = \text{cost} + 100$$

$$\text{or } \text{Newcost} = \text{cost} + \text{id}$$

SELECT DISTINCT \rightarrow If it is not used in primary & Unique key

\rightarrow SELECT DISTINCT statement is used to display only distinct (different) values
 \rightarrow It does not show duplicate value

Syntax :- SELECT DISTINCT column-name
 FROM table-name

Eg : \rightarrow SELECT DISTINCT name FROM bholi;

column-name table-name

table :- emp

id	name	dept	salary	city
1	Ram	IT	30000	Delhi
2	Amit	HR	35000	Mumbai
3	Sunil	Trainer	40000	Kolkata
4	Ram	Executive	10000	Bhopal
5	Dev	IT	25000	Patna

mysql > SELECT DISTINCT.name FROM emp

\Rightarrow Ram
 Amit

Sunil
 Dev

Here; In emp table, Ram exist two times
 but it displays only once.

\Rightarrow It ignores duplicate entry.

mysql> SELECT DISTINCT dept FROM emp.

→ IT HR

Trainer

Executive

table name

Here; In dept column, IT exists two times but it ignores duplicate entry of IT and show only once.

ALTER TABLE

Data Definition Language (DDL)

:- CREATE DATABASE

ALTER DATABASE

DROP DATABASE

CREATE TABLE

ALTER TABLE

DROP TABLE

CONSTRAINT (Keys)

CREATE INDEX

DROP INDEX

CREATE VIEW

ALTER VIEW

DROP VIEW

ALTER TABLE

→ This command is used to Add | change | Modify | DROP existing structure of Table.

1. ADD Column
2. Enable | Disable constraints.
3. Change Column
4. Modify Column.
5. Drop Column

1. ADD Column

(a) Add One column

→ When a new column is to be added to the existing table structure without constraints

Syntax :-

```
ALTER TABLE table_name
ADD column_name datatype (size);
```

	Name	Roll
(10)		

Eg:-

```
ALTER TABLE bhopali
ADD stu_id INT(5);
```

{ It adds new column in existing table.

COLUMN

NEW ROW

	Name	Roll	stu_id

(b) Add more than one column

Syntax :-

```
ALTER TABLE table_name
ADD COLUMN column-name datatype(size),
ADD COLUMN column-name datatype(size);
```

↙

It adds two columns in existing table

Eg:-

```
ALTER TABLE bholi
ADD COLUMN city varchar(20),
ADD COLUMN email varchar(30);
```

	Name	Roll	stu_id	city	email
1					
2					
3					
4					
5					
6					
7					
8					
9					
10					

(c) ADD COLUMN by Position

LAST

- ✓ — ~~last~~ (By default)
- ✓ — FIRST
- ✓ — AFTER
- FIRST

Syntax :-

```
ALTER TABLE table_name
ADD COLUMN column-name datatype(size) FIRST;
```

↙
It adds new column at the FIRST Field.

AFTER

:->

Syntax :- ALTER TABLE table_name
ADD COLUMN column_name datatype(size)
AFTER column_name;

If adds new column after the ~~first~~ any other column.

Eg:- ALTER TABLE table_name
ADD COLUMN state varchar(30)
AFTER city;

Here, It adds new column named as state after the column city.

S.No	Id	Name	City	State	Pin

(d) Add column with constraints

:- When a new column is to be added to the table with constraints.

Syntax :- ALTER TABLE table_name
ADD COLUMN column_name datatype(size)
constraint_name;

Eg:- ALTER TABLE items
ADD COLUMN item_no int(5) NOT NULL;

new column
which we want to add

Q. How to CHANGE COLUMN

A: Change column name and its data-type with constraint

a) One column and its datatype

→ This is used to change name and its data-types of an existing column.

Syntax :-

```
ALTER TABLE table_name
CHANGE COLUMN oldcolumn_name
newcolumn_name new_data_type (size);
```

b) more than one column name and its datatype

Syntax = ALTER TABLE table_name

```
CHANGE COLUMN old-name newcolumn-name
new_data_type (size),
CHANGE COLUMN old-column-name new-column-name
new_data_type (size);
```



It creates two new column with their datatypes in place of old column.

B: change column name and its datatype with constraint.

Syntax:- ALTER TABLE table-name
CHANGE COLUMN old-column new-column
datatype(size) constraint-name,
ADD constraint-name (column-name),
newdatatype

eg:- ALTER TABLE my_tab
CHANGE COLUMN sall id int(5) NOTNULL,
ADD PRIMARY KEY (new column-name);

Here, my_tab is table-name.

sall is old-column-name.

id is new-column-name

NOT NULL and PRIMARY KEY ARE constraints.

C. MODIFY COLUMN

a) modify column datatype and its size

:- This is used to modify size of the datatype or the datatype itself of an existing column. (but column-name must be empty)

Syntax :-

[ALTER TABLE table-name
MODIFY COLUMN column-name datatype(size);

Eg :- ALTER TABLE my-tab → table-name
MODIFY column-name char(20); → existing column-name ↓ size new datatype

D. DROP COLUMN

:- It is used to delete any column in a table.

a) Delete column in a table

Syntax :-

[ALTER TABLE table-name
DROP COLUMN column-name ;

Eg :- ALTER TABLE my-tab → table-name
DROP COLUMN roll; → column-name

b) Removing constraint from a column

Syntax :-

[**ALTER TABLE** table-name
 DROP constraint-name column-name;]

Eg:- [**ALTER TABLE** my-tab
 DROP UNIQUE KEY (you);]

↓ → column-name
new constraint-name

E. DROP TABLE

Syntax :-

DROP TABLE table name;
(It deletes ~~the~~ table with their data)

F. TRUNCATE TABLE

:- It is used to, when we only want to delete the data inside the table, and not the table itself.

Syntax :-

TRUNCATE TABLE table name;

{ DESC table-name :→ It shows table is not deleted ;
SELECT * FROM table name → It deletes data from table }

(DROP TABLE) Drop table
It deletes a table with data.

TRUNCATE TABLE (TRUNCATE TABLE)
• It deletes ~~only~~ ^{Data} table.
• Table exists after Truncate.

RENAME TABLE

:- It is used to rename one or more table.

Syntax :-

Id	name	city
1	Akash	Priyay
2	Sumit	Sassyam

RENAME TABLE old-table-name to new-table-name

Eg :-

RENAME TABLE bhali to saket.

old table-name

new

table-name

ALTER DATABASE

:- It is used to change the overall characteristics of the database.

These characteristics are stored in db.opt file in the database directory.

Syntax :- ALTER DATABASE database-name;

DROP DATABASE

:- It is used to delete a database.

Syntax :- `DROP DATABASE database_name;`

Eg :- `DROP DATABASE my-db;`

database name.

Here, Here, my-db database is deleted from database list.

SHOW COLUMNS

:- It is just like DESC.

• It shows all the columns of table, their data types and specific details.

Syntax :- `SHOW COLUMNS FROM table_name;`

SHOW CREATE DATABASE

→ It is used to show the command which you have written while creating your database.

Syntax :- `SHOW CREATE DATABASE database-name;`

Eg :- `SHOW CREATE DATABASE my-db;`

Output :-	Database	CREATE Database
	my_db	CREATE DATABASE my_db.

SHOW CREATE TABLE

→ It shows command which you have written while creating your table

Syntax :- [SHOW CREATE TABLE table name]

Ex :- SHOW CREATE TABLE emp; ↳ table name

UPDATE Record

→ It is used to update existing records in a table.

Syntax :-

UPDATE table name
SET column1 = value1, column2 = value2, ...
WHERE any column = any value.

table :- emp	emp_id	name	city	salary
	101	Saket	Delhi	80,000
	102	Shubham	Chennai	60,000
	103	Amit	Mumbai	55,000

Eg :-

UPDATE emp.

SET name = 'Saket', salary = 86000
WHERE emp_id = 101

emp_id	name	city	salary
101	Saket	Delhi	86000
102	Shubham	Chennai	60000

any column

which must

be unique and reference to Saket

Here, First of all, it searches emp_id = 101 then it updates their name and salary. If we don't use WHERE clause then all salary records are replaced by 86000 and name by Saket.

Note :- WHERE is necessary otherwise all records will be replaced by given name and salary.

UPDATE with CASE

Syntax :-

```
UPDATE table_name
SET new_column (which we want to update)
CASE
```

```
WHEN column_name1 = some_value1
```

```
THEN new_value1
```

```
WHEN column_name2 = some_value2
```

```
ELSE new_value3.
```

```
END;
```

Eg:- UPDATE student \rightarrow table-name
 SET Result =
 CASE \rightarrow column-name (which we want to update)
 WHEN mark ≥ 300 THEN 'First'
 WHEN mark < 300 AND mark ≥ 250 THEN 'Second'
 WHEN mark < 250 AND mark ≥ 150 THEN 'Third'
 ELSE 'Fail'
 END;

~~DELETE~~

DELETE Record

: \rightarrow It is used to delete records in a table.

1. DELETE one record

Syntax :-

[DELETE FROM table-name
 WHERE column-name = value1 AND
 column-name = value2
 must be unique]

Eg:- DELETE FROM emp
 WHERE emp_name = 'Rohan' AND
 id = 5;

Here, at least ^{one} column must be unique.

2. DELETE all records

Syntax :-

DELETE FROM table_name;

Or

DELETE * FROM table_name;

✓ It deletes ~~permanently~~ permanently from table.

(Copy Old table to New table)

1. Within same database

Syntax :-

CREATE TABLE newtable_name LIKE old-table;

INSERT new_table SELECT * FROM old_table;

Eg:- CREATE TABLE teacher LIKE student;

INSERT teacher SELECT * FROM student;

✓ Here, teacher is new-table-name
student is old-table-name.

So, All table from student is copied to new_table teacher.

2. In different database

Syntax :-

CREATE TABLE new_table LIKE old_database.oldtable
name;

INSERT new_table SELECT * FROM olddatabase.
oldtablename;

Eg:- CREATE TABLE teacher, LIKE my-db.student;
new-tablename old-database old-table
name

INSERT teacher SELECT * FROM my-db.student;

Eg:-

my sq

Note :-

FUNCTIONS

- MIN(column_name) → Smallest value of selected column.
- MAX(column_name) → Largest value of selected column.
- SUM(column_name) → The total sum of a numerical value of column.
- AVG(column_name) → The average value of a numerical column.
- SQRT(column_name) → The square root of a numerical value of this column.
- ROUND(column_name, decimal) → It is used to round a numerical value.

Syntax

:-

Eg:-

If find
minimum
maximum

MIN(column_name)

Syntax :-

[SELECT MIN(column_name) FROM table;]

table :- emp

Name	Registration	city	salary
Saket	11903813	Prayag	90000
Amit	11904652	Sasaram	46000
Sulekh	11903020	Patna	65100

↓

It finds minimum value from salary column.

Eg:-

mysql → [SELECT MIN(salary) FROM emp;]

MIN(salary)

46000

Note :- If we want to rename MIN(salary) then we have to use Alias (AS).

Syntax :- [SELECT MIN(~~salary~~) AS newsalary
FROM emp;]

MAX(column_name)

Syntax :- SELECT MAX(column_name) FROM table_name;

Eg:- SELECT MAX(salary) FROM emp;

✓

MAX(salary)

90000

If finds maximum

maximum value from salary column.

SUM (column_name)

Syntax :- SELECT SUM(column_name) FROM table_name;

Eg :- SELECT SUM(salary) FROM emp;

If adds all
the numerical
value from
Salary column

SUM(salary)

201100

Avg (column_name)

It displays
average value
of salary
column

Syntax :- SELECT AVG(column_name), FROM table_name;

Eg :- SELECT AVG(salary) FROM emp;

Avg(salary)

67033.33

Or

SELECT AVG(salary) AS av_salary FROM emp;

It renames
AVG(salary)
as
av_salary.

av_salary

67033.33

SQRT (column-name) → SELECT column,
SQRT(column) FROM
table-name;

Syntax :- [SELECT *, SQRT(column) FROM tablename;
Eg :- SELECT *, SQRT(salary) AS sroot FROM emp;

Name	Registration	Salary	sroot
Saket	11903813	90000	300.00
Amit	11904652	46000	214.47
Sulekh	11903020	65100	255.15

ROUND (column-name, decimal)

:- Syntax :-

[SELECT *, ROUND(column, decimal)]
FROM table-name;

table :- products

Id	Price	P-name
1	6999.55	Mobile
2	16000.56	Computer
3	10000.93	Laptop

Eg:- SELECT *, ROUND(price, 1) AS approx
FROM products.

Id	Price	P-name	approx
1	6999.55	Mobile	6999.6
2	16000.56	Computer	16000.6
3	10000.93	Laptop	10000.9

COUNT Function

- The COUNT(column-name) function returns the no. of values (NULL values will not be counted) of the specified column;

Syntax :- [SELECT COUNT(column-name)
FROM table-name ;]

- COUNT(*) :- It returns all the number of records in a table.

Syntax :- [SELECT COUNT(*)
FROM table-name ;]

- The COUNT(DISTINCT column-name) :- It returns the no. of distinct value of the specified column. (How many distinct values are present in a specific value)

Syntax :- [SELECT COUNT(DISTINCT column-name)
FROM table-name ;]

table :- orders

Example

	Order Id	Cust Id	Seller Id
	1	101	201
	2	102	202
	3	103	203
	4	101	205
	5	101	204
	6	102	301
	7	103	301

Eg:- `SELECT COUNT(CUSTID)`
`FROM orders;`

→ It tells how many customer ID are present in orders table.

COUNT(CUSTID)
7

Eg:- `SELECT COUNT(DISTINCT CUSTID)`
`FROM orders;`

COUNT(DISTINCT CUSTID)
3

→ It tells how many distinct custId are present in orders table

Here; In CustId, there are 3 distinct value
 $101, 102, 103$.

Multiple value is not counted.

- To calculate how many (no.) of orders by each specific custId

Syntax :- `SELECT COUNT(CUSTID) FROM ORDER`
`WHERE CUSTID = 101;`

COUNT(CUSTID)
3

UPPER Function

:- It converts the value of a field to Uppercase.

- It can written as UCASE(column-name)

syntax :-

SELECT UPPER(column_name)
FROM table_name.

08

SELECT UCASE (column_name)
 FROM tablename.

Eg :- SELECT UPPER(emp_name) AS Name, city
FROM emp.

column name which we want to convert in uppercase

table-name

LOWER Function

:- If converts the value of a field to lowercase.

It can be written as LCASE(column_name)

Syntax :- **SELECT LOWER (column_name)**
FROM table_name;

Def

SELECT LCASE (column_name)
 FROM table_name;

MID Function

Id	city
101	Sasaram
102	Delhi
103	Kolkata

:- It is used to extract characters from a text field.

Syntax :- `SELECT MID(column_name, start, length)`
`FROM table_name;`

Ex, `SELECT MID(city, 1, 3)` → end with
`FROM emp;` ↓ starting from
 column-name
 It displays

	Id	city
	101	sas
	102	Del
	103	Ko

only 3 letters
 of city-name
 from City column

Ex, `SELECT SUBSTRING(city, 1, 3)`
`AS shortcity FROM emp;`

LENGTH Function

It also
 calculate
 space chara
 cter

→ It displays the length of the value in
 a text field.

Syntax :- `SELECT column_name, LENGTH(column)`
`FROM emp;`

Eg:- `SELECT city, LENGTH(city)`
`FROM emp;`

Id	City	LENGTH(city)
101	Sasaram	7
102	Delhi	5
103	Kolkata	7

It displays
 total length of
 value present
 in specific
 column

Eg :- [SELECT emp_name,
 CONCAT(city, ', ', salary)
 FROM emp ;]

table :- emp

emp_name	CONCAT(city, ', ', salary)
Dev	Mumbai, 30000
Jai	Patna, 15000
Vinay	Chennai, 20000
Saket	Delhi, 25000
Bholi	Ranchi, 35000

Eg:- SELECT emp_name,
 CONCAT(salary, SELECT MID(city, 1, 3)))
 As address FROM emp ;

table :- emp

It extracts
 only 3 letters
 starting from 1 to
 3 from
 city column

emp_name	address
Dev	30000 MUM
Jai	15000 PAT
Vinay	20000 CHE
Saket	25000 DEL
Bholi	35000 RAN

REVERSE Function

:- It ~~order~~ reverses the order of letter in String of any column.

Syntax :- `SELECT REVERSE(column-name) FROM table :- emp`

	<u>Id</u>	<u>emp-name</u>	<u>city</u>	<u>column-name</u>
	108	Dev	Delhi	
	109	Jai	Mumbai	
	110	Vimal	Patna	
	111	Saket	Chennai	
	112	Bholi	Ranchi	

Eg:- `SELECT REVERSE(city) FROM emp;`

<u>REVERSE(city)</u>
ihleD
iabmuM
antap
iannehC
ihcnAR

NOW Function

→ It displays time and date of each record from table.

:- It returns the current system data and time.

Syntax :- `SELECT column-name,
 (column-name1,`

`NOW() AS DateTime`

`FROM table-name;`

Eg:- `SELECT emp-name,`

`city, NOW() AS DateTime`

`FROM emp;`

If there are two or more than ~~two~~ data having different values then it displays ~~values~~ ~~on the basis of function used with GROUP BY~~ values on the basis of function used with GROUP BY

GROUP BY → Always uses with Function

→ Syntax :-

SELECT column-name,
MIN(column-name)
FROM table-name
GROUP BY column-name;

// Need a function

table :- emp

	ID	Name	Dept	salary
	185	Dev	IT	45000
	189	Raj	HR	32000
	201	Dev	Accountant	30000
	203	Ram	IT	35480
	206	Sunil	HR	62000

Eg :-

SELECT Name, → column-name.
MIN(salary) ← Function → table-name.
FROM emp → column-name;
GROUP BY Name; → column-name.

ID	Name	Dept	salary
201	Dev	Accountant	30000
189	Raj	HR	32000
203	Ram	IT	35480
206	Sunil	HR	62000
203	Ram	IT	35480

→ IT groups
each dept
together

Here; In emp table, There are 2 employees whose name is Dev.
But by using, GROUP BY with MIN fn then it displays Dev having minimum salary.

If we use MAX(column_name)

Eg:- SELECT Name, MAX(salary), both will be same column names.
 FROM emp
 GROUP BY Name;

<u>Id</u>	<u>Name</u>	<u>Dept</u>	<u>Salary</u>
SC1	Dev	Account	30000
189	Raj	HR	32000
203	Ram	IT	35480
206	Sunil	HR	62000
285	Dev	IT	45000

It displays Dev having maximum salary.

<u>Id</u>	<u>Name</u>	<u>Dept</u>	<u>Salary</u>
189	Raj	HR	32000
206	Sunil	HR	62000
203	Ram	IT	35480
285	Dev	IT	45000

It groups employee from ~~each department~~ together.

But, if the ~~same~~ department has same employee name and we want to GROUP BY Name then it displays Name of employee ~~by~~ on the basis of function.

HAVING Clause

:- It always uses in function.

Syntax :-

```
SELECT column_name1,  
Function_name(column_name2)  
FROM table-name  
GROUP BY column_name1  
HAVING Function_name(column_name2) > value
```

same details

Eg:-

```
SELECT Name,  
MIN(salary),  
FROM emp  
GROUP BY Name  
HAVING MIN(salary) > 40000;
```

It displays ~~one~~ Name
which having ~~is~~
 $\text{MIN}(\text{salary})$ greater
than 40000.

Table - Orders

Order Id	Cust Id	Seller Id
1	101	201
2	102	202
3	103	203
4	101	205
5	101	204
6	102	301
7	103	301

Eg :- SELECT cust_id,
 count(*)

FROM orders
GROUP BY cust_id

~~having count(*) > 2~~

} It groups
all the distinct
value of cust_id

→

CUSTID	COUNT(*)
101	3
102	2
103	3

Eg :-

If groups
all the distinct
value of cust_id

It displays which
satisfy fn.

SELECT cust_id,
 count(*)

FROM orders

GROUP BY cust_id

~~HAVING COUNT(*) > 2;~~

} HAVING
comes with
GROUP BY

function

→

CUSTID	COUNT(*)
101	3
103	3

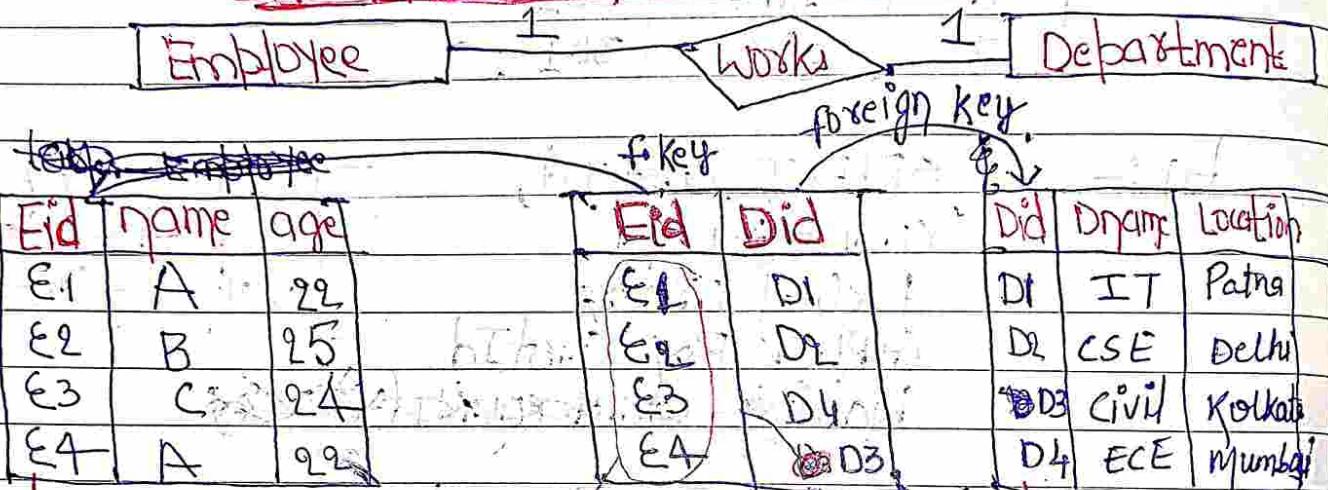
* Primary key → Uniquely +

* Relationship

→ It is the association between two or more entities.

One to One

- ER ~~relationship~~ to relationship model



Primary Key

Each Id of employee related uniquely to Did.

Primary Key

Primary key may be either Eid or Did.

For practical

Step:-1

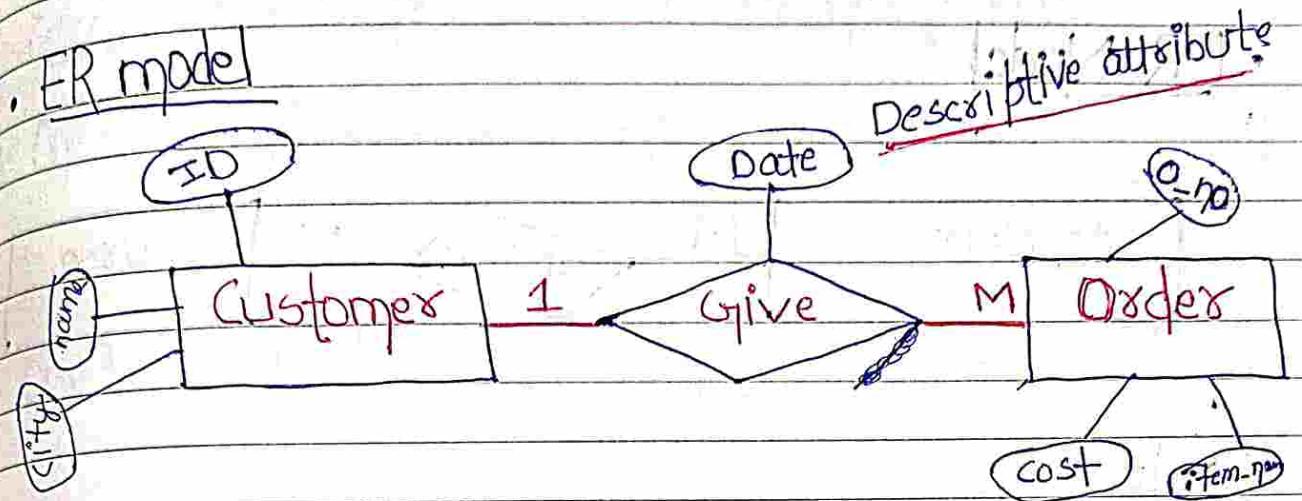
Open any database

Step:-2

Create 2 tables, one for employee and other for Department.

Step:-3

Again, create a table for works and create foreign key with reference of Eid and Did.

One to ManyER modeRelationship table

Note

Relationship table \leftrightarrow primary key, def \leftrightarrow many side & primary key

Customer			Give			Order		
ID	Name	City	ID	O-70	Date	O-70	Item-No	Cost
C1	A	Jalandhar	C1	O1	26-2	O1	Bucket	1,000
C2	B	Patna	C1	O2	27-2	O2	Shoes	2000
C3	C	Delhi	C2	O3	25-2	O3	Shirt	1500
C4	D	Pune	C2	O4	21-2	O4	Jeans	2000

Note :- primary key of relationship table will associate with primary key of many entity

Ex:- One to Many

- One customer can give more than one order so, C1 can give order having O-70 01 and 02.

Similarly, C2 can give order having O-70 03 and 04.

But, One order O-70 associate with only one customer.

Many to Many (M-N)

• ER Model



Table

student			study		course		
Rollno	Name	age	Rollno	c-id	c-id	name	credit
1	A	16	1	C1	C1	Maths	4
2	B	17	2	C2	C2	Phy	4
3	A	16	1	C2	C2	Chem	4
4	C	17	2	C1	C1	Hindi	4
5	D	15	3	C3	C3		

Referencing
or relational
table

- Primary key in referencing or relationship table

:- Primary key will be :-

composite key of Rollno & c-id.

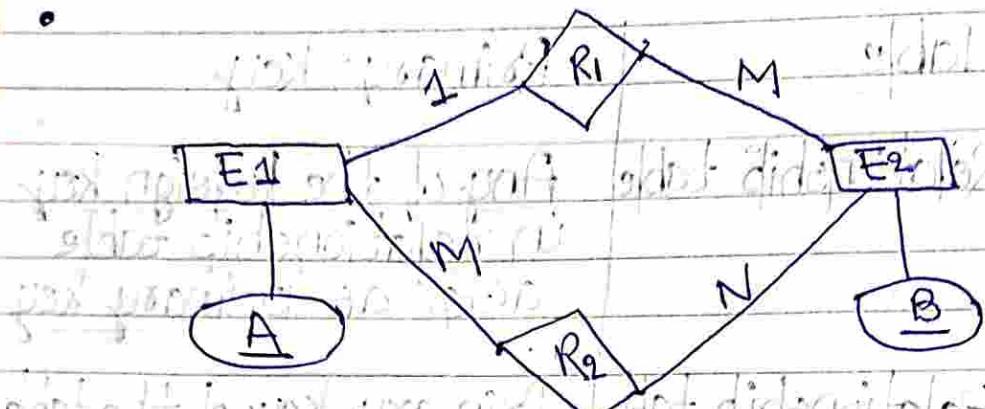
Here, One ~~can~~ study by many students
and each ~~one~~ student ~~can~~ study many courses.

Note :-

<u>Cardinality</u>	<u>Table</u>	<u>Primary key</u>
1. One to one	Relationship table	Any of the foreign key in relationship table acts as primary key.
2. One to many	Relationship table	Primary key of the table in many side acts as primary key of relationship or referencing table.
3. Many to many	Relationship table	→ composite key of both foreign key of Relationship table acts as primary key.
4. Many to one	Relationship table	Primary key of the table in many side acts as primary key of relationship table.

GATE

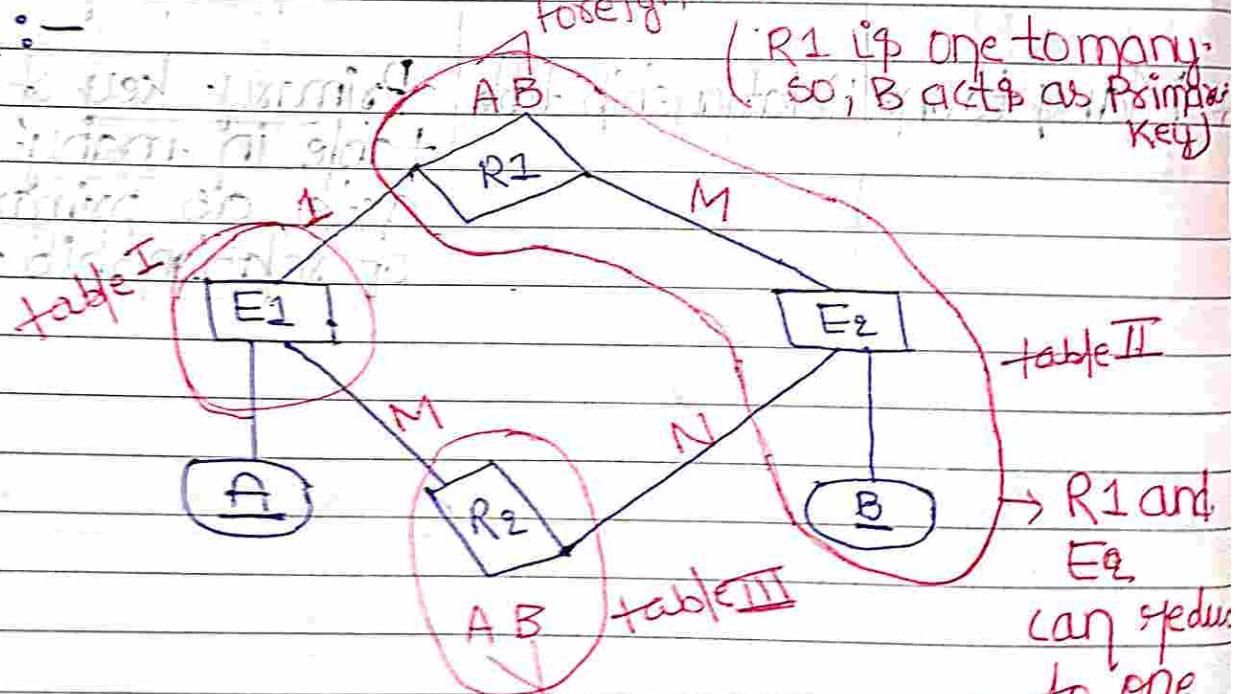
KRISHNA
Page No. / /
Date / /



What is the minimum no. of tables required to represent this ER Model into Relational model?

- a) 2 b) 3 c) 4 d) 5

Sol:-



Foreign key
R2 is many to many relation.
so, primary key will be composite key of A & B.
so, R2 cannot reduce in table.

Ques. Find Nth highest salary using SQL.

Ans. table I = E1
table II = E2
table III = R2.

SELECT COUNT(DISTINCT column)
FROM table_name
AS alias_name

Syntax :-
Select column_name1, column_name2
From table_name AS alias_name
Where N-1 = (SELECT COUNT(DISTINCT column)

From table_name
AS alias_name
Nth
(eg:- For 4th,
where
SAL
alias_name2.column2
alias_name1.column1)

EMP e1

ID	Salary
1	10000
2	20000
3	20000
4	30000
5	40000
6	50000

EMP e2

ID	Salary
4	10000
2	20000
3	20000
4	30000
5	40000
6	50000

For 4th highest salary i.e., N=4
Eq:- SELECT ID, salary

FROM emp e1
WHERE 3 = (SELECT COUNT(DISTINCT salary)
FROM emp e2
WHERE e2.salary > e1.salary)

⇒ (2, 20000).

~~Primary key~~

FOREIGN KEY

→ A Foreign key in one table points to a Primary Key in another table (Parent table)

- The primary key used by a foreign key is also known as parent key.
The table where the primary key is from is known as parent table.

- The foreign key can be used to make sure that the row in one table have corresponding row in another table.

- Foreign key value can be null, even though primary key value can't.

↑
primary key

↑
Foreign key

Emp_Id	Name	Address
1	Rani	Delhi
2	Jai	Mumbai
3	Veeru	Kolkata

Dept_Id	DName	Emp_ID
1	IT	1
2	HR	1
3	Admin	2

CREATE FOREIGN KEY

Page No.
Date

Krishna

new-father
old-child

CREATE TABLE department
D_id int NOT NULL AUTO_INCREMENT PRIMARY
D_Name varchar(40),
E_id int; columns
values
PK
CONSTRAINT employee_Eid_FK
FOREIGN KEY (E_id) REFERENCES employee(E_id); old-father
PK
It is optional;

Here we can see that we have given constraint name as employee_Eid_FK over E_id column of child table.

The CONSTRAINT clause allows to define constraint name for the foreign key constraint. It is optional.

The REFERENCES clause specifies the parent table and its columns to which the columns in the child table refer.

CREATE PARENT TABLE & CHILD TABLE

Parent table → table-name

CREATE TABLE employee → table-name

```
( eid INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
    ename VARCHAR(40),
    address VARCHAR(40),
);
```

child table

child table-name

CREATE TABLE department → table-name

```
( did INT NOT NULL AUTO_INCREMENT PRIMARY KEY,
```

```
    dname VARCHAR(40),
```

```
    empid INT NOT NULL,
```

```
    CONSTRAINT employee_eid_fk
```

```
    FOREIGN KEY (empid) REFERENCES
```

```
    employee (eid)
```

parent

table

column

primary key

of parent table

>>> DESC department;

Field	Type	Null	Key	Default	Extra
did	int(11)	No	PRI	NULL	auto-increment
dname	varchar(40)	YES		NULL	
empid	int(11)	No		NULL	

DROP FOREIGN KEY

syntax :- ALTER TABLE table-name
DROP FOREIGN KEY [parent-table-column-fk]
constraint name;

E :- ALTER TABLE department
DROP FOREIGN KEY employee-eid-fk;
constraint name

I
O

Add Foreign key constraint in a table

:- CREATE TABLE department;

(
Create table | did int not null auto increment
primary key,
dname varchar(40),
empid int not null
) ;

Syntax :- ALTER TABLE department
ADD CONSTRAINT employee-eid-fk
FOREIGN KEY (empid) REFERENCES
employee(eid);

mysql> SELECT * from INFORMATION_SCHEMA

mysql> SELECT * from INFORMATION_SCHEMA.TABLE_CONSTRAINTS
WHERE TABLE_NAME = 'department';

CONSTRAINT NAME	TABLE SCHEME	TABLE NAME	CONSTRAINT TYPE
PRIMARY	geek	department	PRIMARY KEY
PRIMARY	geekyshows	department	PRIMARY KEY
employee-eid-fk	geekyshows	department	FOREIGN KEY

Note :- We can't delete Parent table.

If we want to delete records then we can delete records of child table first then we can delete from parent table.

mysql > USE my_db;

mysql > SHOW TABLES.

Output :-

Tables in my_db	
department	child table
employee	Parent table

mysql > SELECT * FROM employee;

Output :-

Primary key	eid	employee		Parent table, we cannot delete any records from it
		ename	address	
	1	Rani	Delhi	
	2	Sumit	Kolkata	
	3	Rahul	Mumbai	
	4	Raj	Delhi	
	5	Jay	Hyd	
	6	Veeru	Patna	

mysql > SELECT * FROM department;

Output :-

Primary Key	did	department		Foreign key
		dname	empid	
	1	IT	1	child table
	2	HR	1	
	3	Admin	1	
	4	IT	2	
	5	Exe	2	
	6	IT	3	
	7	Exe	3	
	8	IT	4	
	9	HR	5	
	10	Admin	6	

mysql> DELETE FROM employee
WHERE eid = 1;

error, cannot delete or update a parent row.

Note :- Here, eid references to empid.

so, we have to delete records from child table first then we can delete records from the parent table.

- Delete emp = 1 from department table.
so, we use did = 1, 2 and 3.

mysql> SELECT * FROM department
WHERE did = 1;

mysql> DELETE FROM department
WHERE did = 2;

mysql> DELETE FROM department
WHERE did = 3;

These
syntax
delete
emp=1

mysql> SELECT * FROM department;

department → child table

Primary key	did	dname	empid	Foreign key
	4	IT	2	
	5	Exe	2	
	6	IT	3	
	7	Exe	3	
	8	IT	4	
	9	HR	5	
	10	Admin	6	

→ empid = 1,
all records
related to empid = 1
are deleted from
child table or
Referencing table

Now, we can delete the same records from parent table

mysql> SELECT * FROM employee
WHERE eid=1;

Query OK, 0 rows affected.

employee → parent table

eid	ename	address
2	sumit	Kol
3	Rahul	Mum
4	Raj	Delhi
5	Jay	Hyd
6	Veeru	Patna

Here,
eid = 1 is
deleted not
from emp
table.

Note :- If we want to delete records from
parent table then we have to
delete ^{some} records from child table first.

On DELETE clause

- ON DELETE CASCADE
- ON DELETE SET NULL
- ON DELETE NO ACTION
- ON DELETE RESTRICT

ON DELETE CASCADE

→ As we know that, if we want to delete records from Parent table then we have to delete records from child table first.

But, with the help of ON DELETE CASCADE we can delete the records of Parent table directly and also delete the record of child table automatically.

Syntax :- ~~DELETE FROM parent_table_name~~ WHERE column_name = 'Value';

Eg:- ~~DELETE FROM employee~~
WHERE eid = 1;

Note :- We have to add ON DELETE CASCADE during the creation of tables.

compulsory

Eg:- mysql> SHOW TABLES

Tables
employee

mysql> SELECT * FROM employee;

eid	ename	address
1	Rani	Delhi
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

mysql> CREATE TABLE department

did int not null auto_increament primary key,
 dname varchar(40),
 empid int not null,
 CONSTRAINT employee_eid_fk FOREIGN KEY (empid) REFERENCES
 employee(eid)
 ON DELETE CASCADE

mysql> INSERT INTO department(did, dname, empid)
 → VALUE (1, 'IT', 1),
 (2, 'HR', 2),
 (3, 'Accountant', 3),
 (4, 'IT', 4);
 (5, 'HR', 5);

mysql> SELECT * FROM department;

did	dname	empid
1	IT	1
2	HR	2
3	Accountant	3
4	IT	4
5	HR	5

mysql> DELETE FROM employee WHERE eid = 1; → Parent table # If deletes records of parent table as well as child table.

eid	ename	address
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

mysql> DELETE FROM department;
~~WHERE empid=1~~

did	dname	empid
2	HR	2
3	Accountant	3
4	IT	4
5	HR	5

Here, record correspondence of ~~empid=1~~ deleted from child table also.

ON DELETE SET NULL

→ With the help of ON DELETE CASCADE, we can delete the records of parent table and child table automatically.

• ON DELETE CASCADE, it deletes all the (whole row) data related to Foreign key.

I But, if we want to delete the ~~the~~ particular details of a record from parent table only then the ~~whole row~~ same details are removed from child table only.

- It does not delete whole row.
- It sets NULL VALUE at the place from where data is deleted from child table. → Parent table

Eg:- mysql> SELECT * FROM employee; → employee

eid	ename	address
1	Rani	Delhi
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

mysql> CREATE TABLE department → child table

did int not null auto_increment primary key,
 dname varchar(40),
 empid int,
 constraint employee_eid_fk
 foreign key (empid) references employee(eid)

* ON DELETE SET NULL

mysql> INSERT INTO department (dname, empid)
VALUES ('IT', 1),
('HR', 1),
('Admin', 1),
('IT', 2),
('Exe', 2),
('IT', 3),
('Exe', 3),
('IT', 4),
('HR', 5);

of did auto-increment

department		
did	dname	empid
1	IT	1
2	HR	1
3	Admin	1
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

mysql> ~~DELETE FROM employee~~

mysql > DELETE FROM employee
where eid=1;

eid	ename	address
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

mysql>

SELECT * FROM department;

Records do

Not deleted
from child
table, only
add NULL
value.

did	dname	empid
1	IT	NULL
2	HR	NULL
3	Admin	NULL
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

When we delete records corresponding to empid=1 then it deletes data from parent table and add NULL in child table

ON DELETE NO ACTION / ON DELETE NO ACTION

:-

~~now~~

mysql> SELECT * FROM employee;

eid	ename	address
1	Rani	Del
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

Create child table with ON DELETE NO ACTION

mysql> CREATE TABLE department

did int not null auto_increment primary key,
 dname varchar(40),
 empid int,
 constraint employee_eid_fk
 foreign key(empid) references employee(eid)
 * ON DELETE RESTRICT
);

mysql> INSERT INTO department(dname, empid)
 VALUES ('IT', 1),
 ('HR', 1),
 ('Admin', 1),
 ('IT', 2),
 ('Exe', 2),
 ('IT', 3),
 ('Exe', 3),
 ('IT', 4),
 ('HR', 5);

mysql> SELECT * FROM department;

did	dname	empid
1	IT	1
2	HR	1
3	Admin	1
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

mysql> DELETE FROM employee
Where eid = 1;

Error : cannot delete or update a parent row

so: At first, we have to delete child table first
and then delete the respective parent table

mysql> DELETE FROM department
Where did = 1;

mysql> DELETE FROM department
Where did = 2;

mysql> DELETE FROM department
Where did = 3;

mysql> SELECT * FROM department;

did	dname	empid
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

Now, delete parent table

mysql> DELETE FROM employee
where eid=1;

did	dname	address
2	Sumit	Kol
3	Sonam	Hyd
4	Priti	Mum
5	Kunal	Kol

ON UPDATE Clause

ON UPDATE CASCADE

UPDATE IN PARENT TABLE FIRST

ON UPDATE SETNULL

UPDATE IN CHILD TABLE FIRST

ON UPDATE NO ACTION

ON UPDATE RESTRICT

ON UPDATE CASCADE

→ It is used to update the records of parent as well as child table.

- If we update any records in parent table then the same records get updated from child table automatically.

Syntax :- Add ON UPDATE CASCADE (add at the end of creation of table).

mysql> SELECT * FROM employee; → parent table

eid	ename	Address
1	Rani	Del
2	Sumit	Kol
3	Sonam	Mum
4	Priti	Hyd
5	Kunal	Kel

mysql> CREATE TABLE department
→ child table

```

CREATE TABLE department
(
    did int not null auto_increment primary key,
    dname varchar(40),
    empid int,
    CONSTRAINT employee_eid_fk,
    FOREIGN KEY (empid) REFERENCES employee(eid)
    * ON UPDATE CASCADE
);

```

mysql > INSERT INTO department (dname, empid)
VALUES ('IT', 1),

('HR', 1),
(‘Admin’, 1),
(‘HR’, 5);

Page No.

KRISHNA

Date

mysql> SELECT * FROM department;

→ child table

did	dname	empid
1	IT	1
2	HR	1
3	Admin	1
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

mysql> UPDATE employee
SET eid = 10
WHERE eid = 1; { It updates the
value of eid as 10 }

eid	ename	Address
2	Sumit	Kol
3	Sonam	Mum
4	Priti	Hyd
5	Kunal	Kol
10	Rani	Del

mysql> SELECT * FROM department;

did	dname	empid
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5
10	IT	2

mysql> SELECT * FROM department;

did	dname	empid
1	IT	10
2	HR	10
3	Admin	10
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

values updated

ON UPDATE SET NULL

- If update updates the records in parent table and puts SET NULL value in updated records in child table.

mysql> SELECT * FROM employee;
→ payent table

eid	ename	Address
1	Rani	Del
2	Sumit	Kol
3	Sonam	Mum
4	Priti	Hyd
5	Kunal	Kol

mysql > CREATE TABLE department
 did int not null auto_increment primary key,
 dname varchar(40),
 empid int,
 CONSTRAINT employee_eid_fk,
 FOREIGN KEY (empid) REFERENCES
 employee eid
 * ON UPDATE SET NULL
);

mysql > INSERT INTO department(dname, empid)
 VALUES ('IT', 1),
 ('HR', 1),
 ('Admin', 1),
 ('IT', 2),
 ('Exe', 2),
 ('IT', 3),
 ('Exe', 3),
 ('IT', 4),
 ('HR', 5);

mysql > ~~SELECT~~ SELECT * FROM department;

did	dname	empid
1	IT	1
2	HR	1
3	Admin	1
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

mysql > UPDATE employee
 SET eid = 10
 WHERE eid = 1;

<u>eid</u>	<u>ename</u>	<u>Address</u>
2	Sumit	Kol
3	Sonam	Mum
4	Priti	Hyd
5	Kunal	Kol
10	Rani	Del

mysql > SELECT * FROM department;

<u>did</u>	<u>dname</u>	<u>empid</u>
1	IT	NULL
2	HR	NULL
3	Admin	NULL
4	IT	2
5	Exe	2
6	IT	3
7	Exe	3
8	IT	4
9	HR	5

It is used when two or more columns have same records.

Composite Key • The combination of two or more columns make primary key.

→ It is a type of primary key which is created by the composition of two or more columns.

Syntax :- PRIMARY KEY (column_1, column_2);

Eg:- CREATE TABLE course → table-name

composite key [coursecode varchar(40),
Date date,
Cname varchar(40),
Seat int,
Remain int,
Room int,
Rcapacity int,
* PRIMARY KEY (coursecode, Date)

mysql > DESC course;

Field	Type	NULL	Key	Default.
coursecode	varchar(40)	NO	PRI	NULL
Date	date	NO	PRI	NULL
Cname	varchar(40)	YES		NULL
Seat	int(11)	YES		NULL
Remain	int(11)	YES		NULL
Room	int(11)	YES		NULL
Rcapacity	int(11)	YES		NULL

If two tables have same column name then we use shorthand notation to identify from which table the ~~the~~^{some} columns belong to.

Shorthand notations

student

Rollno	Name	Branch
1	Jay	Computer science
2	Suhani	Electronic and Comm.
3	Kriti	Electronic and Comm.

Exam

Rollno	s_code	Mark	p_code
1	CS11	50	CS
1	CS 12	60	CS
1	CS 13	55	CS
1	CS 14	70	CS
2	EC 101	66	EC
2	EC 102	70	EC
2	EC 103	50	EC
3	EC 101	45	EC
3	EC 102	50	EC

Here both tables student and Exam have same column Rollno.

~~so, we write~~ ~~SELECT Rollno, Name~~

If there is same column then RDBMS unable to detect from which table, the same column belongs to.

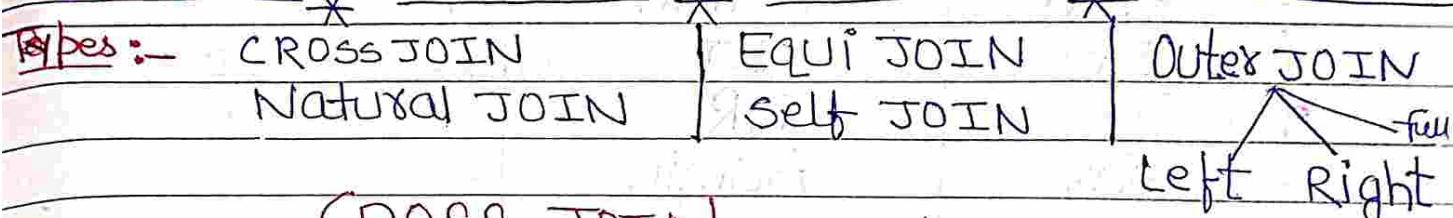
so, we use shorthand notation for fully qualified notation of any column.

Syntax :- [table-name] [column-name]

Eg:- student.Rollno // RDMs easily detect this Rollno belongs to student table

Eg:- Exam.Rollno // This Rollno belongs to Exam table.

JOINS



CROSS JOIN

:- The Cross join returns every row from one table crossed with every row from the second.

Employee			Department	
Embid	Name	City	did	DeptName
1	Rahul	Delhi	101	IT
2	Krish	Kol	102	HR
3	Jay	Mum	103	Admin

Syntax :- [SELECT * FROM Employee] [CROSS JOIN] Department;

Or,

[SELECT * FROM Employee, Department];

Or,

[SELECT Name, Deptname FROM Employee] [CROSS JOIN] Department;

krishna

Q:- SELECT name, Depname from department, employee;

	Name	Depname
1	Rahul	IT
	Rahul	HR
2	Rahul	Admin
	Krish	IT
3	Krish	HR
	Krish	Admin
4	Jay	IT
	Jay	HR
5	Jay	Admin

Eg:- SELECT * FROM Department, Employee;

	Employee			Department	
Emplid	Name	City		did	DepName
1	Rahul	Delhi		101	IT
2	Krish	Kol		102	HR
3	Jay	Mum		103	Admin

Eg:- SELECT * FROM Employee, Department;

↓
start from department

Emplid	Name	City	did	DepName
1	Rahul	Delhi	101	IT
2	Krish	Kol	102	HR
3	Jay	Mum	103	Admin

mysql> SELECT * FROM Department, Employee;

	did	DepName	Empid	Name	city
	101	IT	1	Rahul	Delhi
	102	HR	1	Rahul	Delhi
	103	Admin	1	Rahul	Delhi
	101	IT	2	Krish	Kol
	102	HR	2	Krish	Kol
	103	Admin	2	Krish	Kol
	101	IT	3	Jay	Mum
	102	HR	3	Jay	Mum
	103	Admin	3	Jay	MUM

Self JOIN

⇒ A table is joined itself.

Eg:- Find student id who is enrolled in at least two courses.

study

	s_id	c_id	since
	S1	C1	2016
	S2	C2	2017
	S1	C2	2017

JOIN

= Cross product

⊗ + some condition

⑥ Apply, concept of self Join

T1			T2		
S_id	C_id	since	S_id	C_id	since
S1	C1	2016	S1	C1	2016
S2	C2	2017	S2	C2	2017
S1	C2	2017	S1	C2	2017

~~Self Join~~
Self Join :-

T1		T2	
S1	C1	S1	C1
S1	C1	S2	C2
S1	C1	S1	C2
S2	C2	S1	C1
S2	C2	S2	C2
S2	C2	S1	C2
S1	C2	S1	C1
S1	C2	S2	C2
S1	C2	S1	C2

CROSS PRODUCT

mysql > SELECT T1.S_id FROM
 STUDY AS T1,
 STUDY AS T2] CROSS PRODUCT
 WHERE
 T1.S_id = T2.S_id] condition
 AND
 T1.C_id <> T2.C_id]
 NOT equal

Acc. to condition

T1

T2

~~S1 C1~~~~S1 C1~~~~S1 C1~~~~S2 C2~~~~S1 C1~~~~S1 C2~~~~S2 C2~~~~S1 C1~~~~S2 C2~~~~S2 C2~~~~S1 C2~~~~S1 C2~~~~S1 C2~~~~S2 C2~~~~S1 C2~~~~S1 C2~~~~S1 C2~~~~S2 C2~~~~S1 C2~~~~S1 C2~~

satisfy condition

satisfy condition

so, Output :- S1

Equi JOIN

Q → Find the Emb_name who worked in a department having location same as Address ?

primary key Emp

E-no	E-name	Address	Dept	Location	E-no
1	Ram	Delhi	D1	Delhi	1
2	Vasun	Chd	D2	Pune	2
3	Ravi	Chd	D3	Patna	4
4	Amit	Delhi			

Sol:- JOIN = CROSS PRODUCT + condition.

SELECT E-name
FROM Emp, Dept

Emp · E-no = Dept · E-no
AND
Emp · Address =
Dept · Location

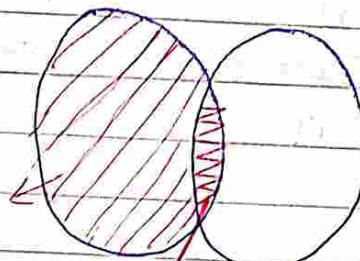
→ CROSS

Left outer join

→ It gives the matching rows and the rows which are in left table but not in right table.

emp

emp-no	e-name	Dept-no
E1	Varun	D1
E2	Amrit	D2
E3	Ravi	D1
E4	Nitin	-



Dept-no	d-name	Loc
D1	IT	Delhi
D2	HR	Hyd.
D3	Finance	Pune

mysql > SELECT emp-no, e-name, d-name, LOC
FROM emp → left table-name
left outer join dept
ON (emp.dept-no = dept.dept-no);
↓
condition

so/7

Emp-no	e-name	d-name	LOC
E1	Varun	IT	Delhi
E2	Amrit	HR	Hyd.
E3	Ravi	IT	Delhi
E4	Nitin	—	—

matching rows

Natural JOIN

(CROSS PRODUCT + condition)

emp			Dept		
e-no	e-name	Address	dept-no	Name	e-no
1	Ram	Delhi	D1	HR	1
2	Vasun	Chd	D2	IT	2
3	Ravi	Chd	D3	MRKT	3
4	Amrit	Delhi			

Ques Find the employee names who are working in a department.

Sol:- $\text{SELECT e-name emp, Dept} \rightarrow \text{cross product}$
~~WHERE $e.no = dept.e.no$~~ \square condition

Or,

SELECT e-name
 FROM emp
 Natural Join Dept.

① Aadhaar card, Pan card, Voter ID, mobile
 Candidate key \rightarrow collection of all such keys which
 can distinguish b/w two entity.
 Primary key extracts from candidate key.

Ques Query the Name of students in table STUDENTS who scored higher than 75 marks. Order your output by the last character of each name. If two or more students have same names ending in the last 3 characters (Robby, Bobby) secondary sort them in descending order. Ascending ID.

Sol:

<u>column</u>	<u>Type</u>
ID	integer
Name	string
Marks	Integer

```

SELECT Name
FROM STUDENTS
where Marks > 75
ORDER BY SUBSTR(name, -3, 3), ID;
  
```

Ques Query the list of name from Employee table in alphabetical order.

Sol:-

```

SELECT name
FROM Employee
ORDER BY name ASC;
  
```

Ques Query the name of employee from Employee table having salary greater than 2000 and have been employee for less than 10 months. Sort your result by ascending employee-id.

→ `SELECT name
FROM Employee
WHERE salary > 2000 AND months < 10
ORDER BY employee_id.`

Ques Write a query to find the maximum total earnings to all employees as well as the total number of employees who have max[↑] total earnings. Then print these values as 2 space-separated integers.

→ `Select * From
(Select months * salary, count(*)
From employee
group by months * salary
order by 1 desc)
Where rownum = 1 ;`

Ques Query the list of CITY names starting with vowels (A,e,i,o,u) from STATION.
Your result cannot contain duplicates.



```
SELECT DISTINCT CITY
FROM STATION
WHERE CITY LIKE 'A%' OR
CITY LIKE 'E%' OR CITY LIKE 'I%' OR
OR CITY LIKE 'O%' OR
CITY LIKE 'U%';
```

Ques Query the list of CITY names starting with vowels and ending with vowels.

→ SELECT distinct city from station
 where (city LIKE 'A%' or city LIKE 'E%',
 OR city LIKE 'I%' OR city LIKE 'O%',
 OR city LIKE 'U%') and
 (city LIKE '%a' OR city LIKE '%e'
 OR city LIKE '%i%' OR city LIKE
 '%o%' OR city LIKE '%u'),

Ques Query the list of city names from station that do not start with vowels.

→ SELECT DISTINCT CITY
FROM STATION

WHERE CITY NOT LIKE 'A%' AND
CITY NOT LIKE 'E%' AND
CITY NOT LIKE 'I%' AND
CITY NOT LIKE 'O%' AND CITY
NOT LIKE 'U%';

Query the list of city names from station that do not start with vowel and do not end with vowel.

→ SELECT ~~REFRESH~~ DISTINCT CITY

FROM STATION WHERE

(city NOT LIKE 'A%' AND city not like 'E%'
AND city not like 'O%' AND city
not like 'I%' AND city not like
'U%') OR

(city NOT like '%A' AND city not like '%E'
AND city not like '%I' AND city
not like '%O' AND city not like
'%U');

Database Constraints

→ A constraint is a rule that the database manager enforces.

Categories of Integrity

i) Domain integrity rules

It is concerned with maintaining the correctness of attribute values within relations.

- It ensures the data values inside a database follow defined rules for values, range and format.

ii) Entity integrity

It ensures that each row in a table is a uniquely identifiable entity.

We can apply entity integrity to a table by specifying PRIMARY KEY constraints.

- It can be enforced through indexes, UNIQUE constraints and PRIMARY KEY constraints.

The primary key for a row is unique, it does not match the primary key of any other row in the table.

The primary key is not NULL.

The uniqueness property ensures that the primary key of each row uniquely identifies there are no duplicate.

iii) Referential Integrity

- It is concerned with maintaining the correctness of relationships between relations.
- It is the state of the database in which all values of all foreign keys are valid.
- We can apply referential integrity using a Foreign KEY constraint.

Note :-

A table in which a referential ~~table~~ constraint and a foreign key are defined is called referencing table while a table is referenced from a ~~referencing table~~ with a foreign key is called a referenced table.

Rules :-

1. We can't delete a record from a primary table if matching records exist in a related table.
2. We can't change a primary key value in the primary table, if that record has related records.