

Part 1:

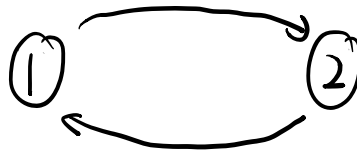
T ₁	T ₂
	R(x)
	R(y)
	W(y)
R(x)	
R(y)	
	R(x)
	R(y)
W(y)	
	W(x)
	R(z)
	W(z)

X: R₂ R₁ R₂ W₂

Y: R₂ W₂ R₁ R₂ W₁

Z: R₂ W₂

thus we have :



which is a cycle and thus
not conflict serializable

S: R₂(x) ; R₂(y) ; W₂(y) ; R₁(x) ; R₁(y) ; R₂(x) ; R₂(y) ; W₁(y) ; W₂(x) ; R₂(z) ; W₂(z)

(B)

If we have multiple transaction and one of the transaction only contains a single query, we may encounter the situation where this single-query transaction is working on a table that is also used by other transactions. If we do not wrap the single query statement in a transaction and utilize lock to insure the conflict serializability, the whole schedule may run into conflicts and further become non conflict serializable. Thus we do need a transaction for a single query statement.