Part 2:

A: R_1 W_1 R_3 W_3 R_2

B: R_1 R_3 W_1 W_3 R_2

thus we have the precedence graph:
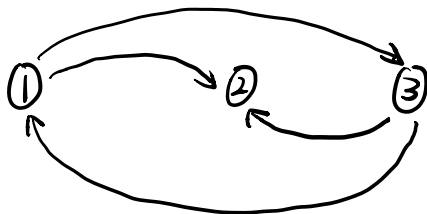
And since we have ① → ③ → ① to form a cycle, the schedule is non-conflict serializable