# Universidad Nacional Autónoma de México

FACULTAD DE CIENCIAS

# Competitive Programming Template

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# 1. Initial Template

```
#include < bits / stdc++.h>
using namespace std;
//using\ namespace\ \_gnu\_pbds;
using bint = _{-}int128;
using ll = long long; // \% ld
using ld = long double; // \%0.5 lf
using vi = vector<ll>;
using pi = pair<11, 11 >; // mp(a,b); \{a,b\};
using ps = pair<string , string >;
using ti = tuple<ll, ll, ll >; // mt(a,b,c); \{a,b,c\};
using ts = tuple<string , string >;
ll gcd(ll a, ll b) {return b ? gcd(b, a %b) : a;}
11 \operatorname{lcm}(11 \ a, 11 \ b) \ \{ \mathbf{return} \ (a \ / \gcd(a, b)) * b; \}
#define endl '\n'
#define npos string::npos
#define LSOne(a) ((a) \& -(a))
#define sq(a) (a)*(a)
\#define sz(a) ((int)(a).size())
#define fst first
#define snd second
#define pb push_back
#define rep(i,a,b) for(int i = a; i < b; i++)
#define all(a) (a).begin(), (a).end()
#define rall(a) (a).rbegin(), (a).rend()
#define fo(a) find_by_order(a)
#define ok(a) order_of_key(a)
#define lb(v,a) lower_bound(v.begin(),v.end(),a)-v.begin()
\#define ub(v,a) upper_bound(v.begin(),v.end(),a)-v.begin()-1
#define bitcount(a) __builtin_popcountll(a) // 1LL \ll 62-1
#define tzcount(a) __builtin_ctzll(a) // 1LL << 62
#define median (a,b,c) max (min(a,b), min(max(a,b),c))
const 11 \text{ MOD} = 1e9+7;
const 11 INF = 1e9; // 10^9 = 1B is < 2^31-1
const 11 LLINF = 4e18; // 4*10^18 is < 2^63-1
const ld EPS = 1e-9; // 10^-9
const 1d PI = 3.1415926535897932384626433832795028841971;
/** == Notes ==
 * # of bits of n :: (int) log 2(n);
 * char to int :: c - '0';
 * Check if bit is on (i--) :: 1 & (n >> i);
 * When nothing left, cursor next line :: cin.iqnore();
 */
```

```
// < ascending, > descending; sort(all(v), cmp);
bool cmpPi(pi a, pi b) {return a.f \Longrightarrow b.f ? a.s \gt b.s : a.f \gt b.f;}
bool cmpStr(str a, str b) {
        return a. size() = b. size() ? a > b : a. size() > <math>b. size();
bool cmpTi(ti a, ti b) {
        return get < 0 > (a) = get < 0 > (b)?
          (get < 1 > (a) = get < 1 > (b) ? get < 2 > (a) < get < 2 > (b)
         = get < 1 > (a) < get < 1 > (b) = get < 0 > (a) < get < 0 > (b);
// priority_queue > ascending, < descending;
// priority_queue < ll, vi, greater < ll >> pq;
struct cmpq {bool operator() (int a, int b) {return a > b;}};
struct cmpqPi {
        bool operator() (pi a, pi b) {
                 return a.f = b.f ? a.s > b.s : a.f > b.f; }};
struct cmpqTi {
        bool operator() (ti a, ti b) {
                 return get <0>(a) = get <0>(b) ? (get <1>(a) = get <1>(b)
                  ? get < 2 > (a) > get < 2 > (b) : get < 1 > (a) > get < 1 > (b)
                  : get < 0 > (a) > get < 0 > (b); \} ;
int digits (ll n) {return int (\log 10(n)+1);}
bool isnumber (string a) {rep(i,0,a.length()) {
         if(!isdigit(a[i])) return false;}
        return true;}
string tobinary(ll n) {string s = "";
        for(int i = (int) log2(n); i >= 0; i--) {
                  (1\&(n>>i))? s+="1" : s+="0";}
        return s;}
int main() {
    ios::sync_with_stdio(0); cin.tie(0);
    // freopen("input.txt", "r", stdin);
    // freopen ("output.txt", "w", stdout);
    int t; cin >> t;
    \mathbf{while}(t--) \text{ solve}();
    return 0;
}
```

# 2. Syntax

```
// == Char to string ==
string s(1,c);
// == stringstream delimiter ==
rep(i,0,s.length()) { if(s[i]=='/') s[i] = '_';}
stringstream ss(s); string t;
\mathbf{while} (ss >> t) \{ \dots \} (1)
while (\operatorname{getline}(\operatorname{ss}, \operatorname{t}, ', ', ')) {...} (2)
// == Count num of this char in string s==
count (all (s), 'a');
// == erase() certain character ==
s.erase(remove(all(s), c), s.end());
// == Check \ if \ string \ contains \ a \ char ==
s.find(c) != string::npos ? found : !found;
// == Print for loop ==
#define prtfor(a) rep(i,0,a.size()) cout \ll a[i] \ll "\\n"[i\ima.size()-1];
// ==Print for each loop==
#define prtfore(a) for(auto x : a) {cout \ll x \ll "_\n"[x==*(-a.end())];}
// ==How to switch on and off==
bool io = 0;
io \hat{}=1;
// == Sort \ k \ characters \ starting \ from \ i-th==
for (int i = 0; i \le n-k; i++) {
         sort(v.begin()+i,v.begin()+i+k);
}
// ==How to not print a newline in the last case==
while(cin >> a) {
         (flag)? flag=0: printf("\n");
         solve(a);}
// == Regex for ip address ==
string inputString = "172.16.254.1";
regex a("(([0-9]|[1-9][0-9]|1[0-9][0-9]|2[0-4][0-9]|25[0-5])\setminus.)
\{3\}([0-9]|[1-9][0-9]|1[0-9][0-9]|2[0-4][0-9]|25[0-5])");
return regex_match (inputString, a);
```

### 3. Bitmask

```
// ==Binary to decimal==
stoi(s, 0, 2);
// ==Decimal to binary==
for(int i = (int) log 2(n); i >= 0; i---) {
         (1&(n>>i))? s+="1": s+="0";
// ==ASCII Char to binary==
bitset < 7 > (c).to_string();
// == Left and Right shift operator==
n << 1 - n*2^1
n << 2 - n*2^2
n >> 1 - n/2^1
n >> 2 - n/2^2
// i-th bit is counting from right to left
// == Turn \ on \ i-th \ bit \ and \ add ==
n = (1 << i);
// == Check \ if \ i-th \ bit \ is \ set \ on ==
n \& (1 << i);
// == Clear/turn \ off \ i-th \ bit==
n \& = (1 << i);
// ==Flip i-th bit with XOR==
n = (1 << i);
// == Value \ of \ LSB \ that \ is \ on ==
((n) \& -(n)); // LSOne(n);
// == Turn \ on \ all \ n \ bits ==
(1 << n) - 1;
// ==Enumerate all proper subsets
for (int i = n; i; i = (n \& (i-1))) {cout << i << '\n';}
// ==Count how many bits are on==
_builtin_popcount(n);
// == Count \ how \ many \ trailing \ zeroes ==
_{-}builtin_ctz(n);
```

#### 4. Order Statistics Tree

```
//==Sorts first k+1 elems in ascending or descending order==
nth_{element}(v.begin(), v.begin()+k, v.end());
nth_element(v.begin(), v.begin()+k, v.end(), greater < int > ());
//==Policy-Based\ Data\ Structures\ (pbds)==
#include <bits/extc++.h>
                                                    // pbds
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
              tree_order_statistics_node_update > ost;
vi v = \{2, 4, 7, 10, 15, 23, 50, 65, 71\};
ost tree;
rep(i,0,n) tree.insert(v[i]);
*tree.find_by_order(0)
                         // smallest = 2
*tree.find_by_order(n-1) // largest = 71
*tree.find_by_order(4)
                          // 5th smallest = 15
tree.order_of_key(2)
                          // index 0 (rank 1)
                          // index 8 (rank 9)
tree.order_of_key(71)
                          // index 4 (rank 5)
tree.order_of_key(15)
//==Kadane's Algorithm==
tuple < ll, int, int > maxSubarray (vi &v) {
    int n = v.size(), l=0, r=0;
    11 ans = v[0], sum = 0, mp = -1;
    rep(i,0,n) {
        sum += v[i];
        if(sum > ans) \{ans = sum; l = mp+1; r = i;\}
         if(sum \le 0) \{sum = 0; mp = i;\}
    return {ans, l, r};
}
tuple < ll, int, int > minSubarray (vi &v) {
    int n = v. size(), l=0, r=0;
    11 ans = v[0], sum = 0, mp = -1;
    rep(i, 0, n)  {
        sum += v[i];
        if(sum < ans) \{ans = sum; l = mp+1; r = i;\}
         if(sum >= 0) \{sum = 0; mp = i;\}
    return {ans, l, r};
}
```

## 5. Strings

```
// == Longest Prefix Suffix ==
vi LPS(string s) {
     vi v(s.length(),0);
     rep(i,1,s.length()) {
          int j = v[i-1];
          while (j \&\& s[i]! = s[j]) \quad j = v[j-1];
          \mathbf{i}\mathbf{f}(\mathbf{s}[\mathbf{i}]==\mathbf{s}[\mathbf{j}]) \mathbf{j}++;
          v\,[\,\,i\,\,]\ =\ j\,\,;
     } return v;
}
// ==Knuth-Morris-Pratt==
vi KMP(string s, string t) {
     int n = s.length(), m = t.length();
     vi lps = LPS(t), kmp;
     int i=0, j=0;
     \mathbf{while} (\mathbf{n-i} >= \mathbf{m-j})  {
          if(s[i] = t[j]) \{i++; j++;\}
           \textbf{if} \, (\, j \underline{\hspace{-1mm}} m) \  \, \{ kmp.\, pb \, (\, i-j \,) \, ; \  \, j \, = \, lp \, s \, [\, j-1]; \} 
          else if(i < n && s[i] != t[j]) {
                j ? j = lps[j-1] : i++;
     return kmp;
}
// == Z ==
vi Z(string s, string t) {
     string a = t+"\$"+s; int n = a.length(), m = t.length();
     vi p(n), z;
     for (int i=1, l=0, r=0; i < n; ++i) {
          if(i \le r) p[i] = min(r-i+1, p[i-l]);
          while (i+p[i] < n \&\& a[p[i]] = a[i+p[i]] + p[i];
          if(i+p[i]-1 > r) l=i, r=i+p[i]-1;
     rep(i,0,n) \{ if(p[i] = m) z.pb(i-m-1); \}
     return z;
}
```

```
// == Longest \ Palindromic \ Subsequence ==
string LPS(string &s) {
    string t = s;
    reverse (all(t));
    return LCS(s,t);
}
// == Longest \ Palindromic \ Subsequence ==
// a, b = s \cdot length(); vvi dp(MAXN, vi(MAXN, -1));
int lps(string &s, int a, int b) {
    if (a==0 | b==0) return 0;
    \mathbf{if}(dp[a][b] \stackrel{!}{=} -1) \mathbf{return} dp[a][b];
    if(s[a-1] = s[n-b])
         return dp[a][b] = 1 + lps(s, a-1, b-1);
    } else {
        return dp [a] [b] = max(lps(s,a-1,b),lps(s,a,b-1));}
// ==Longest Substring without Repeating Chars==
int lengthOfLongestSubstring(string s) {
         set < char > alphabet;
         int ans = 0, l = 0;
         for(int i = 0; i < s.length(); i++) {
             while (alphabet.count(s[i])) {
                  alphabet.erase(s[1]); l++;
             alphabet.insert(s[i]);
             ans = \max(\text{ans}, i-l+1);
         return ans;}
// == Longest Substring without repeating chars==
string LSWRC(string &s) {
    int l=0, r=0, i=0, j=0, m=0;
    unordered_set < char > a;
    \mathbf{while}(\mathbf{r} < \mathbf{s.size}()) {
         if(a.find(s[r]) = a.end()) {
             a.insert(s[r]);
             if(r-l+1 > m) \{m = r-l+1; i = l; j = r;\}
             r++:
         else \{a.erase(s[1]); 1++;\}
    return s.substr(i,m);}
```

```
// ==Longest Substring with K unique chars==
string LKS(string &s, int k) {
    int i=0, j=0, ans=-1;
    unordered_map<char, int> m;
    while (j < s.size())
        m[s[j]]++;
        while (m. size() > k) {
            m[s[i]] - -;
             if (!m[s[i]]) m. erase(s[i]);
             i++;
        if(m. size()==k) ans=max(ans, j-i+1);
        j++;
    return s.substr(i,ans);}
// == Minimum \ substring \ of \ s \ that \ contains \ all \ chars \ of \ t==
string minWindow(string s, string t) {
    vi m(256,0);
    int a=0, b=INT\_MAX, c=0, i=0, j=0;
    rep(i,0,t.length()) \{if(!m[t[i]]) c++; m[t[i]]++;\}
    \mathbf{while}(j < s.length()) {
        m[s[j]]--;
        if (!m[s[j]]) c--;
        while (!c) {
             if(b > j-i+1) \{a = i; b = min(b, j-i+1);\}
            m[s[i]]++;
             if(m[s[i]] > 0) c++;
             i++;
        j++;
    return (b != INT_MAX) ? s.substr(a,b) : "-1";}
```

# 6. Range Queries

```
//==Sparse Table==
vvi sparse_table(vi &v) {
    int n = v.size(), k = (int)log2(n)+1; vvi st(n, vi(k));
    rep(i,0,n) st[i][0] = v[i];
    rep(j,1,k+1) {
        for (int i=0; i+(1<< j)<=n; i++) {
            // RMQ :: Minimum <, Maximum >
            st[i][j-1] < st[i+(1<<(j-1))][j-1] ? st[i][j]=st[i][j-1]
                                     : st[i][j]=st[i+(1<<(j-1))][j-1];
            // RSQ
// st[i][j] = st[i][j-1] + st[i+(1<<(j-1))][j-1];
    } return st;}
// Minimum <=, Maximum >=
int RMQ(int l, int r, vvi &st) {
    int m, j = (int) log2(r-l+1);
    st[1][j] \le st[r-(1<< j)+1][j]? m=st[1][j]: m=st[r-(1<< j)+1][j];
    return m;}
ll RSQ(int l, int r, vvi &st) {
    11 ans = 0; int k = (int) \log 2(st.size()) + 1;
    for (int j=k; j>=0; j--) {
        if(1+(1<< j)-1 <= r) {ans} += st[1][j]; 1 += 1<< j;}
    } return ans;}
```

```
//==Fenwick Tree==
class Fenwick {
private:
     vll ft;
public:
     Fenwick (int m) \{ft.assign(m+1,0);\}
     void build (const vll &f) {
          int m = (int) f.size()-1; ft.assign(m+1,0);
          for(int i=1; i \le m; ++i) {
                    ft[i] += f[i];
                   if (i+LSOne(i) <= m) ft [i+LSOne(i)] += ft [i]; } }
     Fenwick(const vll &f) {build(f);}
     Fenwick (int m, const vi &s) {
          v11 f(m+1,0);
          for (int i=0; i < (int) s. size(); ++i) ++f[s[i]];
          build (f);}
     11 \operatorname{rsq}(\mathbf{int} \ \mathbf{j}) \{11 \operatorname{sum} = 0;
         for(;j;j=LSOne(j)) sum += ft[j]; return sum;}
     \operatorname{ll} \operatorname{rsq}(\operatorname{int} i, \operatorname{int} j) \{\operatorname{return} \operatorname{rsq}(j) - \operatorname{rsq}(i-1);\}
     void update(int i, ll v) {
         for (; i < (int) ft. size(); i+=LSOne(i)) ft [i] += v;
     int select(ll k) {
         int l = 1, h = ft.size()-1;
         for (int i=0; i < 30; ++i) {
                   int m = (1+h)/2; (rsq(1,m) < k)? l=m : h=m; return h;
};
class RUPQ {
private:
    Fenwick ft;
public:
    RUPQ(int m) : ft(Fenwick(m)) \{ \}
     void range_update(int ui, int uj, int v) {ft.update(ui, v);
                                                         ft.update(uj+1, -v);}
     11 point_query(int i) {return ft.rsq(i);}
};
class RURQ {
private:
    RUPQ rupq; Fenwick purq;
public:
    RURQ(int m) : rupq(RUPQ(m)), purq(Fenwick(m))  {}
    void range_update(int ui, int uj, int v) {
         rupq.range_update(ui, uj, v);
         purq.update(ui, v*(ui-1));
         purq.update(uj+1, -v*uj);}
     11 \operatorname{rsq}(\mathbf{int} \ j) \ \{\mathbf{return} \ \operatorname{rupq.point\_query}(j) * j - \operatorname{purq.rsq}(j); \}
     Il rsq(int i, int j) \{return rsq(j) - rsq(i-1);\}
};
```

```
//==Segment Tree RMQ==
class Segtree {
private:
    int n; vll A, st, lazy;
    int l(int p) {return p \ll 1;} int r(int p) {return(p \ll 1)+1;}
    ll conquer(ll a, ll b) {
        if(a==-1) return b; if(b==-1) return a;
        return min(a,b);
    void build (int p, int L, int R) {
        \mathbf{if}(L = R) \operatorname{st}[p] = A[L];
        else {
             int m = (L+R)/2;
             build (l(p), L, m); build (r(p), m+1,R);
             st[p] = conquer(st[l(p)], st[r(p)]); \}
    void propagate (int p, int L, int R) {
        \mathbf{if}(\text{lazy}[p]!=0) {
             // st[p] = lazy[p];
             // if(L!=R) lazy[l(p)] = lazy[r(p)] = lazy[p];
             st[p] += lazy[p];
             if (L!=R) { lazy [l(p)]+=lazy [p]; lazy [r(p)]+=lazy [p]; }
             lazy | p | = 0;
    11 RMQ(int p, int L, int R, int i, int j) {
        propagate(p, L, R);
        if(i > j) return -1;
        if((L >= i) \&\& (R <= j)) return st[p];
        int m = (L+R)/2;
        return conquer (RMQ(l(p), L, m, i, min(m, j)),
                        RMQ(r(p), m+1, R, max(i, m+1), j);
    void update(int p, int L, int R, int i, int j, ll val) {
        propagate(p, L, R);
        if(i > j) return;
        if((L >= i) \&\& (R <= j)) 
             lazy[p] = val;
             propagate(p, L, R);
        } else {
             int m = (L+R)/2;
             update(l(p), L, m, i, min(m, j), val);
             update(r(p), m+1, R, max(i, m+1), j, val);
             int lsubtree = (lazy[l(p)] != 0) ? lazy[l(p)] : st[l(p)];
             int rsubtree = (lazy[r(p)] != 0) ? lazy[r(p)] : st[r(p)];
             st[p] = (lsubtree \ll rsubtree) ? st[l(p)] : st[r(p)];
             // \ll min, \gg max
        }
    }
```

```
public:
    Segtree (int sz) : n(sz), st(4*n), lazy(4*n, 0) {}
    Segtree (const vll &initialA) : Segtree ((int)initialA.size()) {
         A = initialA;
         build (1, 0, n-1);
    void update(int i, int j, ll val) {update(1,0,n-1,i,j,val);}
    11 RMQ(int i, int j) {return RMQ(1,0,n-1,i,j);}
};
//==Segment Tree RSQ==
const int MAXN = 2e5+5;
struct node { ll val, lzAdd, lzSet; node() { }; } st[MAXN << 2];
class Segtree {
private:
    int l(int p) \{return p \ll 1;\}
    int r(int p) \{return(p \ll 1)+1;\}
    void conquer(int p) \{ st[p]. val = st[l(p)]. val + st[r(p)]. val; return; \}
    void propagate (int p, int L, int m, int R) {
         \mathbf{if}(\operatorname{st}[p].\operatorname{lzSet} != 0) {
              st[l(p)].lzSet = st[r(p)].lzSet = st[p].lzSet;
              st[l(p)].val = (m-L+1)*st[p].lzSet;
              st[r(p)].val = (R-m)*st[p].lzSet;
              st[l(p)].lzAdd = st[r(p)].lzAdd = 0;
              st[p].lzSet = 0;
         else\ if(st[p].lzAdd != 0) 
              if(st[l(p)].lzSet == 0) st[l(p)].lzAdd += st[p].lzAdd;
                   st[l(p)].lzSet += st[p].lzAdd;
                   st[l(p)].lzAdd = 0;
              \mathbf{if}(\operatorname{st}[r(p)] \cdot \operatorname{lzSet} = 0) \operatorname{st}[r(p)] \cdot \operatorname{lzAdd} + \operatorname{st}[p] \cdot \operatorname{lzAdd};
              else {
                   st[r(p)].lzSet += st[p].lzAdd;
                   st[r(p)].lzAdd = 0;
              st[l(p)].val += (m-L+1)*st[p].lzAdd;
              st[r(p)].val += (R-m)*st[p].lzAdd;
              st[p].lzAdd = 0;
         } return;
    }
```

```
void build (int p, int L, int R) {
        st[p].lzAdd = st[p].lzSet = 0;
        if(L = R) \{ st[p]. val = A[L]; return; \}
        int m = (L+R) >> 1;
        build (l(p), L, m); build (r(p), m+1,R);
        conquer(p); return;}
    void add(int p, int L, int R, int i, int j, ll val) {
        if (i > R \mid j < L) return;
        if (i <= L && R <= j) {
            st[p].val += (R-L+1)*val;
             if(st[p].lzSet == 0) st[p].lzAdd += val;
            else st[p].lzSet += val; return;}
        int m = (L+R) \gg 1;
        propagate(p, L, m, R);
        add(l(p), L, m, i, j, val); add(r(p), m+1, R, i, j, val);
        conquer(p); return;}
    void set(int p, int L, int R, int i, int j, ll val) {
        if (i > R \mid | j < L) return;
        if (i <= L && R <= j) {
            st[p].val = (R-L+1) * val;
            st[p].lzAdd = 0; st[p].lzSet = val; return;
        int m = (L+R) \gg 1;
        propagate(p, L, m, R);
        set(l(p), L, m, i, j, val); set(r(p), m+1, R, i, j, val);
        conquer(p); return;}
    ll query(int p, int L, int R, int i, int j) {
        if(i > R \mid | j < L) return 0;
        if(i \le L \&\& R \le j) return st[p].val;
        int m = (L+R) >> 1;
        propagate(p, L, m, R);
        return query (l(p), L, m, i, j) + query (r(p), m+1, R, i, j); 
public:
    int n;
    11 A[MAXN];
    Segtree (int sz) : n(sz) \{ \}
    void build() \{build(1,1,n);\}
    void add(int i, int j, ll val) \{add(1,1,n,i,j,val);\}
    void set(int i, int j, ll val) \{ set(1,1,n,i,j,val); \}
    If query (int i, int j) {return query (1,1,n,i,j);}
};
```

### 7. Math

```
const ld pi = 3.1415926535897932384626433832795028841971;
//==Spiral\ Square\ Perimeter==
2 + n*(n+1) + n; // n is length of side of square
//==Dice\ Moves\ to\ get\ n==
if(n = 1) \{ cout << -1 << endl; return; \}
int ans = (n/11)*2;
ans += ((n\%11)/5);
if (((n\%11)\%5) != 0) ans++;
if(n = 7) ans++;
//==a^b < c^d==
(b*log(a) \le d*log(c))
//==Gauss's Circle Problem==
// # of lattice points within circle of radius r.
ll lattice_points(ll r) { // 1e9 lattice points with r = 18000
    11 \text{ points} = 0;
    rep(i,1,r+1) points += floor(sqrt(sq(r)-sq(i)));
    return 111+411*(r+points);
}
//==Circumscribed Circle==
// Polygon of v vertices, each side length s, return area.
ld circumscribed_circle(ld v, ld s) {
    1d area = s/(2*sin(pi/v));
    area*=(area*pi);
    return area;
//==XOR from 1 to n==
int computeXOR(int n) {
    if (n %4==0) return n;
    if (n %4==1) return 1;
    if (n\%4==2) return n+1;
    return 0;}
//==Two nums that are sum s and xor x==
pi sumxor(11 s, 11 x) 
    11 z = (s-x)/2, a = 0, b = 0;
    for (int i = 0; i < (int) log 2 (s); i++) {
        11 xi = (x & (1 << i));
        11 \ zi = (z \& (1 << i));
        if (xi && zi) {return \{-1,-1\};}
        if (! xi && zi) {
            a = ((1 << i) | a);
            b = ((1 << i) | b); 
        if(xi \&\& !zi) \{a = ((1 << i) | a);\}\}
    return \{a,b\};\}
```

# 8. Graph

```
//==Graph (int)==
vi graph[n];
vi visited (n, 0);
vi dist(n, 0);
//==Graph (string)==
map<string, vector<string>> graph;
map<string, bool> visited;
map<string, int> distance;
//==Graph (matrix)==
vector < pi> graph [n][m];
vvi visited (n, vi(m, 0));
//==Generating Edges (matrix graph)
 graph[i][j].pb(mp(i,j));
 // left
 if(j-1 >= 0) \{ graph[i][j].pb(\{i,j-1\}); graph[i][j-1].pb(\{i,j\}); \}
 // right
 if(j+1 \le m-1) \{graph[i][j].pb(\{i,j+1\}); graph[i][j+1].pb(\{i,j\});\}
 //up
 if(i-1 >= 0) \{ graph[i][j].pb(\{i-1,j\}); graph[i-1][j].pb(\{i,j\}); \}
 //up left
 if(i-1) = 0 \& j-1 > = 0) \{graph[i][j].pb(\{i-1,j-1\});
                              graph[i-1][j-1].pb(\{i,j\});
 // up right
 if(i-1) = 0 \& j+1 \le m-1  {graph[i][j].pb({i-1,j+1});
                                 graph[i-1][j+1].pb(\{i,j\});
void bfs(int s) {
    visited[s] = 1;
    queue < int > q; q.push(s);
    \mathbf{while}(!q.empty()) {
         int u = q. front(); q. pop();
         for(auto v : graph[u])  {
             if (! visited [v]) {
                  visited[v] = 1;
                  \operatorname{dist}[v] = \operatorname{dist}[u] + 1;
                  q.push(v);}
         }
    }
}
```

```
void dfs(int u) {
     visited[u] = 1;
     for (int v : graph[u]) {
          \mathbf{if} (! visited | v | ) {
               dfs(v); dist[v] = dist[u]+1;
          }
     }
}
void dfs(int a) {
     visited[a] = 1;
     stack < int > s; s.push(a);
     while (!s.empty()) {
          int u = s.top(); s.pop();
          visited[u] = 1;
          for(auto v : graph[u])  {
               \mathbf{if}(! \text{ visited}[v]) {
                    \operatorname{dist}[v] = \operatorname{dist}[u] + 1;
                   s.push(v);
              }
         }
    }
}
//==Graph (dijkstra)==
vector < pi> graph [n];
vi dist(n, INF);
//==Edges (w :: weight)==
graph[u].pb(\{v,w\});
void dijkstra(int s) {
     priority_queue<pi, vector<pi>, greater<pi>> pq;
     pq.push(\{0,s\}); dist[s] = 0;
     \mathbf{while}(! pq.empty())  {
         \mathbf{auto} \ \mathbf{p} = \mathbf{pq.top()}; \ \mathbf{pq.pop()};
          int u = p.second;
          if(p.first != dist[u]) continue;
          for (auto z : graph [u])
              int v = z. first;
              int weight = z.second;
               if(dist[v] > dist[u] + weight) {
                    dist[v] = dist[u] + weight;
                   pq.push({ dist [v], v});
              }
         }
    }
}
```

```
//==Disjoint Set Union==
class DSU {
private:
    vi p, rank, sz; int numSets;
public:
   DSU(int n) {
        p.assign(n,0);
        rep(i, 0, n) p[i] = i;
        rank.assign(n,0);
        sz.assign(n,1);
        numSets = n;
    int find (int i) {return (p[i]==i) ? i : (p[i] = find(p[i]));}
    bool same(int i, int j) {return find(i) == find(j);}
    int numDisjoint() {return numSets;}
    ll size(int i) {return sz[find(i)];}
    void unite(int i, int j) {
        if(same(i,j)) return;
        int x = find(i), y = find(j);
        if(rank[x] > rank[y]) swap(x,y);
        p|x| = y;
        if(rank[x] = rank[y]) + rank[y];
        sz[y] += sz[x];
        --numSets; \};
//==Strongly Connected Components==
class SCC {
private:
    vector < vi > AL, AL_T; vi A, dfs_num; 11 MAXN, ans = 0;
public:
   SCC(int n) {
        AL. assign(n, \{\}); AL_T. assign(n, \{\});
        dfs_num.assign(n,0);
        MAXN = n;
    void Kosaraju (int u, int pass) { //==pass== 1 original, 2 transpose
        dfs_num[u] = 1;
        vi &neighbor = (pass==1) ? AL[u] : AL_T[u];
        for (auto &v : neighbor) {
            if (!dfs_num[v]) Kosaraju(v, pass);
        } A.pb(u);}
    void addEdge(int u, int v) {AL[u].pb(v); AL_T[v].pb(u);}
    11 components() {
        rep(u,0,MAXN) \{if(!dfs_num[u]) Kosaraju(u,1);\}
        dfs_num . assign (MAXN, 0);
        for (int i = MAXN-1; i >= 0; i---) {
            if(!dfs_num[A[i]]) Kosaraju(A[i],2), ans++;
        return ans; } };
```

```
//==Minimum Spanning Tree==
struct Edge { ll u, v, w; };
// < Minimum Spanning Tree, > Maximum Spanning Tree
bool cmp(Edge a, Edge b) {return a.w > b.w;}
class MST {
private:
    vi p, rank, sz;
    11 \text{ sets}, \text{ cost} = 0;
public:
    vector < Edge > edges, mst;
    MST(int n) {
        p. assign (n, 0);
        rep(i, 0, n) p[i] = i;
        rank.assign(n,0);
        sz.assign(n,1);
        sets = n;
    int find(int i) {return (p[i]==i) ? i : (p[i] = find(p[i]));}
    bool same(int i, int j) {return find(i) == find(j);}
    11 size(int i) {return sz[find(i)];}
    void unite(int i, int j) {
        int a = find(i), b = find(j);
        if (a=b) return;
        if(rank[a] < rank[b]) swap(a,b);
        p[b] = a;
        if(rank[a] = rank[b]) ++rank[a];
        sz[a] = sz[b];
        --sets;
    11 total_cost() {return cost;}
    void addEdge(int u, int v, int w) {Edge e; e.u = u; e.v = v;
                                                   e.w = w; edges.pb(e);
    vector < Edge > Kruskal() {
        sort (all (edges), cmp);
        for (Edge e : edges) {
            —e.u; —e.v;
             \mathbf{if}(!same(e.u,e.v)) {
                 cost += e.w; cost \% = MOD; if(cost < 0) cost += MOD;
                 mst.pb(e);
                 unite (e.u, e.v);
        return mst;
    }
};
```

```
//==Articulation Points and Bridges==
class Bridges {
private:
    vector < vi > graph; ll dfsNumberCounter, dfsRoot, rootChildren,
                                 MAXN, UNVISITED = -1;
    vi dfs_num, dfs_low, dfs_parent, articulation_vertex,
                                       articulation_points;
    vector<pi> bridges;
public:
    Bridges(int n) {
        dfsNumberCounter = 0, MAXN = n;
        graph.assign(n, \{\});
        dfs_num.assign(n,UNVISITED);
        dfs_low.assign(n,0);
        dfs_parent.assign(n,-1);
        articulation_vertex.assign(n,0);
    \mathbf{void} addEdge(\mathbf{int} u, \mathbf{int} v) {graph[u].pb(v); graph[v].pb(u);}
    void articulationPointAndBridge(int u) {
        dfs_num[u] = dfsNumberCounter++;
        dfs_low[u] = dfs_num[u];
        for (auto &v : graph [u]) {
             if(dfs_num[v] = UNVISITED) {
                 dfs_parent[v] = u;
                 if (u == dfsRoot) ++rootChildren;
                 articulationPointAndBridge(v);
                 if(dfs_low[v] >= dfs_num[u]) articulation_vertex[u] = 1;
                 if(dfs_low[v] > dfs_num[u]) \{bridges.pb(\{u,v\});\}
                 dfs_low[u] = min(dfs_low[u], dfs_low[v]);
            else\ if(v != dfs_parent[u])\ dfs_low[u] = min(dfs_low[u],
                                                              dfs_num[v];
        }}
    void artibridge() {
        rep(u, 0, MAXN) {
             if(dfs_num[u] = UNVISITED)  {
                 dfsRoot = u; rootChildren = 0;
                 articulationPointAndBridge(u);
                 articulation_vertex[dfsRoot] = (rootChildren > 1);
        rep(u, 0, MAXN) {
             if(articulation_vertex[u]) articulation_points.pb(u);}
    vector<pi> getBridges() {return bridges;}
    vi getArticulation() {return articulation_points;}
};
```

```
//==Topological Sort==
class Toposort {
private:
    vector < vi> graph; int MAXN;
    vi ts, tl, vis, indegree;
    priority_queue<ll , vi , greater<ll> > pq;
public:
    Toposort(int n) {
        MAXN = n;
        tl.assign(MAXN, 0);
        graph.assign(n, \{\});
        vis. assign (n, 0);
        indegree.assign(n,0);
    void addEdge(int u, int v) {graph[u].pb(v); indegree[v]++;}
    void addEdgeTL(int v, int u) {graph[u].pb(v); indegree[v]++;}
    void topo(int a) {
        vis[a] = 1;
        for(auto \&u : graph[a]) \{if(!vis[u]) topo(u);\}
        ts.pb(a);}
    vi toposort() {
        for (int i = MAXN-1; i >= 0; i---) { if (! vis [i]) topo(i);}
        //rep(i, 0, MAXN) \{ if(!vis[i]) topo(i); \}
        reverse (all (ts)); return ts;
    }
    vi khan() {
        rep(i,0,MAXN) { if(!indegree[i]) pq.push(i);}
        while (!pq.empty()) {
             int u = pq.top(); pq.pop(); ts.pb(u);
             for (auto &v : graph [u]) {
                 -indegree [v];
                 if (!indegree[v]) pq.push(v);}
        } return ts;
    vi topolabel() { // order of node appearance lexicographically
        set < ll, greater > s; int c = MAXN-1;
        rep(i,0,MAXN) { if(!indegree[i]) s.insert(i);}
        while (! s . empty ()) {
             ll\ u = *s.begin();\ s.erase(u);\ tl[u] = c;
             for (auto &x : graph [u]) {
                 if(--indegree[x]==0) s.insert(x);
             } c--;
        return tl;
    }
};
```

```
//==Lowest Common Ancestor==
class LCA {
private:
     vector < vector < pi>> graph; vector < vi> up;
     vi depth, dist; ll l, MAXN;
public:
     LCA(ll n) {
          graph.\,assign\left(n\,,\{\,\}\,\right);\;\;up\,.\,assign\left(n\,,vi\left(\,2\,0\,\right)\,\right);
           depth.assign(n,0); dist.assign(n,0);
          FOR(i, 0, n) \{ up[i][0] = i; \}
          l = ceil(log2(n)); MAXN = n;
     void addEdge(ll u, ll v, ll w) \{graph[u].pb(\{v,w\}); graph[v].pb(\{u,w\})\}
     void init() {
           dfs(0);
          FOR(i, 1, 1) \{FOR(j, 0, MAXN) | up[j][i] = up[up[j][i-1]][i-1]; \}
     void dfs(ll u) {
          for(auto \&[v,w] : graph[u])  {
                if(v != up[u][0])  {
                     depth[v] = depth[u]+1;
                      \operatorname{dist}[v] = \operatorname{dist}[\operatorname{up}[v][0] = u] + w;
                      dfs(v);
                }}}
     11 \text{ jump}(11 \text{ x}, 11 \text{ d}) 
          FOR(i, 0, 1) \{ if((d>>i)&1) | x = up[x][i]; \}
          return x;
     }
     ll lca(ll u, ll v) 
           if(depth[u] < depth[v]) swap(u,v);
          u = jump(u, depth[u]-depth[v]);
           if (u==v) return u;
          RFOR(i,0,1) {
                if (up [u] [i] != up [v] [i])
                     u = up[u][i], v = up[v][i];
          return up [u][0];
     }
     ll distance (int u, int v) {
           return \operatorname{dist}[\mathbf{u}] + \operatorname{dist}[\mathbf{v}] - 2 * \operatorname{dist}[\operatorname{lca}(\mathbf{u}, \mathbf{v})];
     }};
```

## 9. Combinatorics

```
const 11 MOD = 1e9+7;
const 11 MAXN = 1e6;
vll fac(MAXN), inv(MAXN);
// ==Binary Exponentiation for Modular Multiplicative Inverse==
// bcpow(a, MOD-2) :: Moduler Inverse of a
11 bcpow(11 a, 11 b){
    11 r = 1;
    while(b) {
         if(b\&111) r = r * a %MOD;
         a = a * a \%MOD; b >>= 111;
    } return r;}
// == Binary Exponentiation ==
11 binpow(11 a, 11 b) {
    ll r = 1;
    while(b) {
         if(b \& 111) r = r * a;
         a = a * a; b >>= 111;
    } return r;}
//==Binomial Coefficient==
11 \text{ nCk}(11 \text{ n}, 11 \text{ k}) 
    ld r = 1;
    rep(i,1,k+1) r = r * (n-k+i)/i;
    return (11) (r+0.01);
// == Calculates Factorials and Inverses ==
void facinv() {
    fac[0] = 1; fac[1] = 1;
    rep(i, 2, MAXN) fac[i] = fac[i-1] * i %MOD;
    inv[MAXN-1] = bcpow(fac[MAXN-1], MOD-2);
    for (ll i = MAXN-2; i \ge 0; i--) inv[i] = inv[i+1] * (i+1ll) %MOD;}
// == Binomial Coefficient with Modular Arithmetic ==
ll \ nCk(ll \ n, ll \ k)  {
    return n < k ? 0 : fac[n] * inv[k] %MOD * inv[n-k] %MOD;}
vll inverses (ll n, vll &f) {
    vll \operatorname{inv}(n+1); \operatorname{inv}[n] = \operatorname{bcpow}(f[n], MOD-2);
    for (ll i = n-1; i >= 0; i--) inv[i] = inv[i+1] * (i+1ll) %MOD;
    return inv;}
```

## 10. **Geometry**

```
//==Convex Hull Graham Scan==
struct pt\{ld x, y;\};
bool operator < (const pt &p, const pt &q) {
         return mp(p.y, p.x) < mp(q.y, q.x);
}
// Direction :: collinear 0, clockwise -1, counterclockwise 1;
int orientation(pt p, pt q, pt r) {
    1d val = (r.y-p.y)*(q.x-p.x)-(q.y-p.y)*(r.x-p.x);
    return ! val ? 0 : (val < 0 ? -1 : 1);
}
vector<pt> grahamScan(vector<pt> &pts) {
    pt p0 = *min_element(all(pts), [](pt p, pt q) 
              return mp(p.y, p.x) < mp(q.y, q.x);});
    sort(all(pts), [&p0](const pt &p, const pt &q) {
         int o = orientation(p0, p, q);
         if (!o) return (p0.x-p.x)*(p0.x-p.x)+(p0.y-p.y)*(p0.y-p.y)
                      < (p0.x-q.x)*(p0.x-q.x)+(p0.y-q.y)*(p0.y-q.y);
        return o > 0; // o < 0 for cw
    });
    vector <pt> ch;
    rep(i,0,pts.size()) {
         while (ch.size() > 1 \&\& orientation(ch[ch.size()-2],
               \operatorname{ch.back}(), \operatorname{pts}[i]) \stackrel{!}{=} 1 /* -1 \text{ for } cw */) \operatorname{ch.pop\_back}();
         ch.pb(pts[i]);
    return ch;
}
set < pt > tmp;
vector<pt> pts;
rep(i,0,n) {
    cin \gg x \gg y;
    pt p; p.x = x; p.y = y;
    tmp.insert(p);
for (auto a : tmp) pts.pb(a);
vector<pt> ans = grahamScan(pts);
cout << ans. size() << endl;
for (auto a : ans) {
    cout << a.x << "" << a.y << endl;
}
```

# 11. Number Theory

```
// == Factorials ==
vll factorials (ll n) {
    vll f(n+1,1);
    rep(i,1,n+1) f[i] = f[i-1] * i %MOD;
    return f;}
// == Trailing Zeros for n!==
11 trailingZeros(ll n) {
    11 z = 0; for (11 i = 5; i \le n; i \le 5) z + n/i; return z;
//==Check if n is prime==
bool isPrime(ll n) {
    if(n \le 1) return 0;
    if(n \le 3) return 1;
    if (n %2==0 | | n %3==0) return 0;
    for (11 i = 5; i*i \le n; i+=6) {
        if(n\% ==0 \mid | n\%(i+2)==0) return 0;
    } return 1;
}
//==Check if a^2-b^2 is prime==
bool is DiffPrime (ll a, ll b) {
    return (isPrime(a+b) && a-b==1) ? 1 : 0;
}
//==Sieve of Eratosthenes==
vll primes(ll n) {
        vll p(n+1,1); p[0]=0; p[1]=0;
        for (11 i=2; i*i \le n; i++) {
        if(p[i]) \{for(ll j=i*i; j<=n; j+=i) p[j]=0;\}\}
        rep(i,0,n+1) \{if(p[i]) q.pb(i);\}
        return q;}
//==Divisors of N==
vll divisors(ll n) {
        vll d; d.pb(1); d.pb(n);
        for (11 i=2; i*i \le n; i++) {
                 if(n\% i==0) i*i==n? d.pb(i) : (d.pb(i), d.pb(n/i));
        return d;}
//==Practical Numbers==
bool practicalNum(vll d, ll n) {
        ll r = 0; bool p;
        for (auto a : d) { if (r+1 < a) break; r+=a;}
        (r+1<n*2)? p=0: p=1;
        return p;}
```

```
// == K - divisible Sum==
// Sum of n numbers div by k and max elem of arr is min possible
(n+(k*((n+k-1)/k))-1)/n
//==Divisible pairs in an Array==
int divisible_pairs(vi &A) {
    int n = A. size(), ans = 0;
    unordered_map<ll, ll> frq;
    rep(i, 0, n) frq[A[i]]++;
    rep(i, 0, n)  {
         for (int j = 1; sq(j) <= A[i]; j++) {
              if(A[i]\%j==0) {
                  (A[i]==sq(j))? ans+=frq[j]: ans+=frq[j]+frq[A[i]/j];
         ans--;
    { return ans; }
//==Numbers of sums of distinct powers of 3 til 1e5==
vector < int > v;
rep(i, 0, 9) {
    int a = pow(3,i); v.pb(a);
    rep(j, 0, pow(2, i)) {
         int tmp = a;
         rep(k,0,i)
              if (j & (1 << k)) {tmp += pow (3, k);}
         v.pb(tmp);
\} v.pb(pow(3,9));
//==Number\ of\ pairs\ that\ can\ matrix\ multiply==
map < int, ll > snd; vi v(n);
rep(i,0,n) \{cin >> v[i] >> b; snd[b]++;\}
11 \text{ ans} = 0;
\operatorname{rep}(i,0,n) \left\{ \mathbf{if}(\operatorname{snd}.\operatorname{count}(v[i])) \text{ ans } += \operatorname{snd}[v[i]]; \right\}
//==Number of numers that have only 1 zero in binary==
ll a, b; cin >> a >> b;
11 i = (11) \log 2(a), j = (11) \log 2(b);
ll first = (i*(i-1))/2, last = (j*(j-1))/2;
i++; j++;
11 x = (1LL \ll i) - 1, y = (1LL \ll j) - 1;
for (11 k = i-2; k >= 0; k--)
    if((x \& ~(1LL << k)) < a) first++;
for (11 k = j-2; k >= 0; k--) {
    if((y \& ~(1LL << k)) <= b) last++;
} cout << last-first << endl;</pre>
```

```
//==Maximal AND==
// can turn on k bits in total in any elem in array
int n, k; cin \gg n \gg k;
     vi bits (31,0);
     vector < ll > v(n);
     rep(i,0,n) {
          cin >> v[i];
          for (int j = (int) \log 2(v[i]); j >= 0; j--) {
               if(v[i] & (1 << j)) bits[j]++;
     int bit = 30;
     \mathbf{while}(k > 0) {
          \mathbf{while}(\mathbf{bits}[\mathbf{bit}] = \mathbf{n} \mid \mid \mathbf{n-bits}[\mathbf{bit}] > \mathbf{k}) \quad \mathbf{bit} = \mathbf{k}
          if(bit < 0) break;
          rep(i,0,n) {
               v[i] = (1 << bit);
              k = (n-bits[bit]);
               bits[bit] = n;
     }
     11 ans = pow(2,31)-1;
     rep(i,0,n) ans &= v[i];
     cout << ans << endl;
//==Partition Problem==
bool findPartition(vi &v) {
     int n = v.size(), sum = 0;
     rep(i,0,n) sum += v[i];
     if (sum %2) return false;
     int p = sum/2;
     vector < bool > part(p+1,0);
     rep(i,0,n) {
          for(int j = p; j >= v[i]; j--) 
               \mathbf{if}(\operatorname{part}[j-v[i]]==1 \mid | j==v[i]) \operatorname{part}[j] = 1;
     return part[p];
//==2 max elems==
int a = 0, b = 0;
rep(i, 0, n)  {
          cin >> c;
          if(c > a) \{b=a; a=c;\}
          else if (c > b) {b=c;}
}
```

```
//==Smallest number with length m and sum of digits s==
string getMin(int m, int s) {
    if(m == 1) \{return to_string(s);\}
    string a = ""; bool first = 1;
    while ((m-1)*9 > s \&\& m > 0)
        if(first) {first = 0; a += "1"; s--; m--; continue;}
        a += "0"; m--;
    a += to_string(s-(m-1)*9); m--;
    rep(i, 0, m) = "9";
    return a;
}
//==Largest number with length m and sum of digits s==
string getMax(int m, int s) {
    string a = "";
    while (s >= 9) {a += "9"; m--; s-=9;}
    if(s >= 1) \{a += to_string(s); m--;\}
    rep(i, 0, m) = "0";
    return a;
}
//==2D Prefix Sum==
int c = 0; // til 1e6
rep(i,0,1414) {
    int a = i, b = 0;
    \mathbf{while} (\mathbf{a} >= 0)  {
        11 \operatorname{can} = 0;
        if(a > 0) \{can \neq pyramid[a-1][b];\} // up
        if(b > 0) \{can += pyramid[a][b-1];\} // left
        if(a > 0 \&\& b > 0) \{can = pyramid[a-1][b-1];\} // up-left
        can += binpow(c, 2);
        pyramid[a][b] = can;
        dp[c] = can;
        c++;
        a--; b++;
}
```

# 12. Implementation

```
// == Make \ Palindrome \ with \ given \ string ==
string palindrome(string s) {
        unordered_map<char, int> m;
        rep(i, 0, s. length()) m[s[i]]++;
        int odd=0; char c;
        for (auto a : m) {
                 if(a.second \%2!=0) \{odd++; c=a.first;\}\}
        if (odd>1 || odd&&s.length()%2==0) return "Null";
        string x=",", y="";
        for (auto a : m) {
                 string s(a.second/2,a.first);
                 x=x+s; y=s+y;
        return (odd) ? (x+c+y) : (x+y);
// ==All triplets v[i]+v[j]+v[k] that equal n==
vvi threeSum(vi &v, int n) {
    sort(all(v)); vvi ans;
    rep(i, 0, v. size()-2) {
        if ( i > 0 \& v[i-1] = v[i]) continue;
        int k = -v[i]+n, a = i+1, b = v \cdot size()-1;
        \mathbf{while}(\mathbf{a} < \mathbf{b})
             if(v[a] + v[b] == k) {
                 ans.push\_back(\{v[i], v[a], v[b]\});
                 a++, b--;
                 while (a < b \& v [a] = v [a-1]) a++;
                 while (a < b \& v[b] = v[b+1]) b--;
             if(v[a] + v[b] > k) b--;
             if(v[a] + v[b] < k) a++;
    } return ans;}
//==Number\ of\ cars(b)\ that\ overtook\ other\ cars(a)
int overrun (vi &a, vi &b) {
    rep(i,0,n) c[b[i]] = i; // turn of that index that goes out
    rep(i,0,n) res[i] = c[a[i]]; // order of that car that goes out
    int fined = 0, x=-1;
    rep(i,0,n) {
        (res[i] > x) ? x = res[i] : fined++;
    return fined;
}
```

```
//==Maximize difference with other player==
11 \text{ asum} = 0, \text{ bsum} = 0;
int player = 0;
while (!a.empty() || !b.empty()) {
    int ai = a.top(), bi = b.top();
    if(ai > bi \&\& !player) \{asum += ai; a.pop();\}
    else if(ai < bi && player) \{bsum += bi; b.pop();\}
    else { player ? a.pop() : b.pop(); }
    player = 1;
}
// Number of different strings by removing 2 consecutive characters
int diffStrings(string s) {
    int ans = 1;
    rep(i,0,s.length()-2) ans += s[i]!=s[i+2];
    return ans;
}
// Number of pairs of lower and upper case
// Can swap lower to upper or viceversa max k times
int pairsLowerUpper(int k, string s) {
    map < char, int > m; int ans = 0; int available = 0;
    rep(i, 0, s. length()) m[s[i]]++;
    for (auto a : m) {
        char l = tolower(a.first), u = toupper(a.first);
        if (!m[1] \&\& !m[u]) continue;
        else if (m[1] \&\& m[u]) {
             ans += \min(m[1], m[u]);
             available += abs (m[1]-m[u])/2;
            m[1] = 0; m[u] = 0;
        } else {
             available += \max(m[1], m[u])/2;
            m[1] = 0; m[u] = 0;
    ans += \min(\text{available}, k);
    return ans;}
// Prefix needed from s to form t
map<char, vi> letters;
rep(i,0,n) letters [s[i]]. pb(i+1);
vi alpha (26,0);
rep(i,0,t.length()) alpha[t[i]-'a']++;
int ans = 0;
for (int i = 0; i < 26; i++) {
    if(alpha[i]) \{ans = max(ans, letters[i+'a'][alpha[i]-1]);\}
}
cout << ans << endl;
```

```
// # of operations to make s anti-palindrome
map<char, int> m;
rep(i, 0, n) m[s[i]] + +;
bool cant = 0;
for (auto a : m) {
    if(m[a.first]*2 > s.length()) cant=1;
if(n\%2 \mid | cant) \{cout << -1 << endl; return;\}
vi same (26);
rep(i, 0, n/2) {
    same [s[i]-'a'] += s[i]==s[n-i-1];
    // add to index of the alphabet if same on both sides
// doing just #same+1/2 won't work cuz what if aaaabbcdefghbbaaaa
int ans = 0;
priority_queue < int > pq;
rep(i,0,26) {
    if (same [i]) pq.push (same [i]);
    // if there is repeating alphabet, add to pq
while (pq.size() > 1) {
    int a = pq.top(); pq.pop();
    int b = pq.top(); pq.pop();
    a--; b--; ans++;
    if(a) pq.push(a);
    if(b) pq.push(b);
if(!pq.empty()) ans += pq.top();
// Swap letters of s at distance k or k+1 to make same as t
void canBeSame(string s, string t, int k) {
    bool can = 1; int n = s.length();
    if(k > n) {
         (s=t) ? printf("YES\n") : printf("NO\n"); return;}
    if(k > n/2) {
         if(s.substr(n-k,n-(n-k)*2) != t.substr(n-k,n-(n-k)*2)) can=0;
    \operatorname{map} < \operatorname{char}, \operatorname{int} > \operatorname{ms}, \operatorname{mt}; \operatorname{rep}(i,0,n) \operatorname{ms}[s[i]] + +; \operatorname{rep}(i,0,n) \operatorname{mt}[t[i]] + +;
    for (auto a : ms) \{if(ms[a.first] != mt[a.first]) \{can=0; break;\}\}
    can ? printf("YES\n") : printf("NO\n");
}
```

## 13. **DP**

```
//==Dice\ Combinations==
vi dp(n+1,0); dp[0] = 1;
     FOR(i, 1, n+1) {
           FOR(j, 1, 7) {
                 if(i < j) break;
                 dp[i] += dp[i-j]\%MOD;
                 dp[i] \% MOD;
                 \mathbf{if}(\mathrm{dp}[\mathrm{i}] < 0) \mathrm{dp}[\mathrm{i}] += \mathrm{MOD};
      cout \ll dp[n] \ll endl;
//==Coin Combinations I==
vi coin(n); FOR(i,0,n) cin >> coin[i];
      vi dp(x+1); dp[0] = 1;
     FOR(i,1,x+1) {
           FOR(j,0,n) {
                 if(i-coin[j] < 0) continue;
                 dp[i] += dp[i-coin[j]];
                 dp[i] \ll MOD;
                 \mathbf{if}(\mathrm{dp}[\mathrm{i}] < 0) \mathrm{dp}[\mathrm{i}] += \mathrm{MOD};
      cout \ll dp[x] \ll endl;
//==Coin\ Combinations\ II==
vi coin(n); FOR(i,0,n) cin >> coin[i];
      vi dp(x+1); dp[0] = 1;
     FOR(i,0,n) {
           FOR(j, coin[i], x+1) {
                 dp[j] += dp[j-coin[i]];
                 dp [j] \not = MOD;
                 \quad \textbf{if} \left( \, \mathrm{dp} \left[ \, j \, \right] \, < \, 0 \right) \ \mathrm{dp} \left[ \, j \, \right] \, + = \mathrm{MOD};
      cout \ll dp[x] \ll endl;
```

```
// == Longest Common Subsequence ==
string LCS(string s, string t) {
    int n = s.length(), m = t.length();
    vvi dp(n+1, vi(m+1));
    rep(i, 0, n+1) {
        rep(j, 0, m+1)  {
            !i||!j|? dp[i][j] = 0 : (s[i-1] = t[j-1])
                   ? dp[i][j]=dp[i-1][j-1]+1
                   dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
        }
    int i = n, j = m, k = dp[n][m]; vector < char > c(k+1); c[k] = '\0';
    while(i && j) {
       string lcs(all(c)); lcs.erase(lcs.length()-1,1);
    return lcs;
}
//==Longest\ Increasing\ Subsequence==
vi LIS(vi &v) {
    int n = v.size();
    vi ans, sub, si, path(n, -1);
    rep(i,0,n) {
        if(sub.empty() \mid \mid sub[sub.size()-1] < v[i]) {
            path[i] = sub.empty() ? -1 : si[sub.size()-1];
            sub.pb(v[i]);
            si.pb(i);
        } else {
            int idx = lb(sub, v[i]);
            path[i] = idx = 0 ? -1 : si[idx - 1];
            sub[idx] = v[i];
            si[idx] = i;
        }
    int t = si[si.size()-1];
    while (t != -1)
        ans.pb(v[t]);
        t = path[t];
    reverse (all (ans));
    return ans;
}
```

```
//==Longest Subsequence of A having LCM at most K
int LongestSubsequenceLCM(vi A, int k) {
    int n = A. size(); map<ll, ll>m;
    rep(i, 0, n) m[A[i]] + +;
    vi \ v(k+1,0);
    for (auto a : m) {
         if(a.fst \ll k) {
             for(int i=1;;++i) {
                 if(a.fst*i > k) break;
                 v[a.fst*i] += a.snd;
         } else break;
    11 lcm = 0, ans = 0; //if lcm == 0, no answer
    rep(i,1,n+1) \{ if(v[i] > ans) \{ ans = v[i]; lcm = i; \} \}
    vi u; rep(i,0,n) if (\operatorname{lcm} \Re[i]==0) u.pb(A[i]); // subsequence
    return ans;
}
```

## 14. Search

```
//==Binary Search==
int binarySearch(vi &v, int x) {
    int l = 0, r = v.size()-1;
    \mathbf{while}(l \le r)  {
        int m = 1+(r-1)/2;
        v[m] < x ? l = m+1 : r = m-1;
         if(v[m] = x) return m;
    return -1;
//==Ternary Search==
int ternarySearch(vi &v, int x) {
    int l = 0, r = v.size()-1;
    \mathbf{while}(\mathbf{r} >= \mathbf{l})  {
        int m1 = 1+(r-1)/3, m2 = r-(r-1)/3;
         if(v[m1] > x) r = m1-1;
         else if (v[m2] < x) l = m2+1;
         else \{1 = m1+1; r = m2-1;\}
         if(v[m1] = x) return m1;
         if(v[m2] = x) return m2;
    return -1;
```

#### 15. Prefix

```
v11 v = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\};
int n = v.size();
//==Prefix Sum==
vll pfx (n+1,0);
rep(i,1,n+1) \{pfx[i] = pfx[i-1] + v[i-1];\}
11 \text{ max\_subarray\_sum} = \text{pfx}[1], \text{min\_prefix\_sum} = \text{pfx}[0];
rep(i,1,n+1) {
    max_subarray_sum = max(max_subarray_sum, pfx[i]-min_prefix_sum);
    min_prefix_sum = min(min_prefix_sum, pfx[i]);
}
//==2D Prefix Sum==
vector < vll > v2d = {
     \{0,0,0,0,0,0,0\}
    \{0,1,2,3,4,5\},\
    \{0,1,2,3,4,5\}
     \{0,1,2,3,4,5\}
     \{0,1,2,3,4,5\},
     \{0,1,2,3,4,5\}
\}; n = v2d. size()-1;
\operatorname{vector} < \operatorname{vll} > \operatorname{pfx2d}(n+1, \operatorname{vll}(n+1, 0));
rep(i, 1, n+1) {
    rep(j,1,n+1) {
         pfx2d[i][j] = v2d[i][j]+pfx2d[i-1][j]+
                          pfx2d[i][j-1]-pfx2d[i-1][j-1];
    }
// x = 3^{5}, y = 2^{3}
int x1 = 3, x2 = 5, y1 = 2, y2 = 3;
11 \text{ sum} = \text{pfx2d}[x2][y2] - \text{pfx2d}[x1-1][y2] -
          pfx2d[x2][y1-1] + pfx2d[x1-1][y1-1];
cout << sum << endl;</pre>
//==Prefix Multiplication Except Self==
vi productExceptSelf(vi &v) {
    int n = v.size();
    vi prefix(n), suffix(n), ans;
    int p = 1, s = 1;
    rep(i,0,n) \{ prefix[i] = p *= v[i]; suffix[n-1-i] = s *= v[n-1-i]; \}
    rep(i,0,n) {
         if(!i) \{ans.pb(suffix[i+1]); continue;\}
         if (i=n-1) {ans.pb(prefix[i-1]); continue;}
         ans.pb(prefix [i-1]*suffix [i+1]);
    } return ans;}
```

## 16. Swaps

```
// Number of bubble sort swaps to sort array
//==Inversion Count with Merge Sort==
ll merge (vi &v, vi &A, int l, int m, int r) {
    int i = 1, j = m, k = 1; ll inv = 0;
    while(i <= m-1 && j <= r) {
        if(v[i] \le v[j]) A[k++] = v[i++];
        else \{A[k++] = v[j++]; inv += (m-i);\}
    while ( i \le m-1) A[k++] = v[i++];
    while (j \le r) A[k++] = v[j++];
    rep(i, l, r+1) v[i] = A[i];
    return inv;
}
ll mergeSort(vi &v, vi &A, int l, int r) {
    int m; ll inv = 0;
    if(l < r)
        m = (l+r)/2;
        inv += mergeSort(v,A,l,m);
        inv += mergeSort(v,A,m+1,r);
        inv += merge(v,A,l,m+1,r);
    } return inv;
}
ll inversionCount(vi &v) {
    vi A(v. size());
    return mergeSort (v, A, 0, v. size()-1);
}
//==Inversion Index==
struct comparepq {bool operator() (pi a, pi b) {return a.first > b.first;}
int inversionIndex(vi v) {
    if(v.size() \le 1) return 0;
    /*==Ascending==
    struct\ comparepq\ \{bool\ operator()\ (pi\ a,\ pi\ b)\ \{return\ a.first>b.fir
    priority\_queue < pi, vector < pi >, comparepq > pq; */
    priority_queue <pi> pq;
    rep(i,0,v.size()) \{pq.push(\{v[i],i\});\} v.clear(); int ans = 0;
    \mathbf{while}(!pq.empty()) {
        pi p = pq.top(); pq.pop();
        int y = lb(v, p.second); ans += p.second - y;
        v.insert(lower_bound(all(v),p.second),p.second);
    } return ans;
}
```

```
//==Swaps(any i, j) needed to sort array==
void swap(vi &A, int i, int j) {
    11 \text{ tmp} = A[i];
    A[i] = A[j];
    A[j] = tmp;
}
int min_swaps(vi A, int n) {
    int ans = 0; vi tmp = A; sort(all(tmp));
    map<11, 11> m; // indexes of A
    {\rm rep}\,(\,{\rm i}\,\,,0\,\,,{\rm n}\,)\ m[{\rm A}[\,{\rm i}\,]\,]\ =\ {\rm i}\,\,;\ /\!/\ m[5]\ =\ in\,d\,e\,x\ of\ 5
    rep(i,0,n) {
         \mathbf{if}(A[i] != tmp[i]) 
              ans++;
             int j = A[i];
             // swap elem with idx it should be in
             swap(A, i, m[tmp[i]);
             // update idx
             m[j] = m[tmp[i]];
             m[tmp[i]] = i;
    return ans;
}
//=Swaps in a binary tree to make sorted perfect binary tree=
//==int ans = 0; swapsBT(v, 0, m-1, ans); m = \# elems.
bool mergeBT(vi &v, int l, int m, int r) {
    if(v[m] < v[m+1]) return 0;
    rep (i, l, m+1) swap (v[i], v[i+(r-l+1)/2]);
    return 1;
}
void swapsBT(vi &v, int l, int r, int &ans) {
    if(l < r) {
         int m = 1+(r-1)/2;
         swapsBT(v,l,m, ans);
         swapsBT(v,m+1,r, ans);
         if(mergeBT(v,l,m,r)) ans++;
}
is\_sorted(all(v))? printf("%l\n",ans) : printf("-1\n");
```