Module
2

Embedded Processors and Memory

Lesson
5

Memory-I

Instructional Objectives

After going through this lesson the student would

- o Different kinds of Memory
    - Processor Memory
    - Primary Memory
    - Memory Interfacing

Pre-Requisite

Digital Electronics, Microprocessors

5.1   Introduction

This chapter shall describe about the memory. Most of the modern computer system has been designed on the basis of an architecture called Von-Neumann Architecture[1]

| Input Output Devices | ←→ | Central Processing Unit | ←→ | Memory |
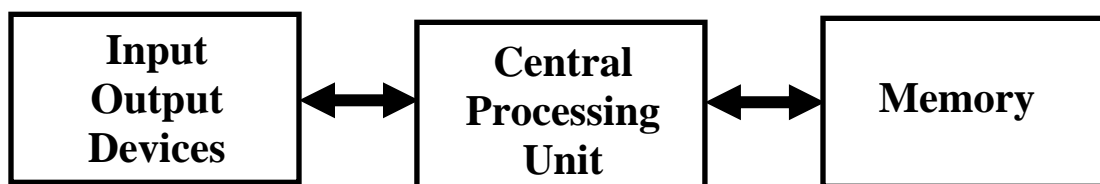
Fig. 5.1 The Von Neumann Architecture

The Memory stores the instructions as well as data. No one can distinguish an instruction and data. The CPU has to be directed to the address of the instruction codes.
The memory is connected to the CPU through the following lines

1. Address

2. Data

3. Control

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[1] http://en.wikipedia.org/wiki/John_von_Neumann.   The so-called von Neumann architecture is a model for a computing machine that uses a single storage structure to hold both the set of instructions on how to perform the computation and the data required or generated by the computation. Such machines are also known as stored-program computers. The separation of storage from the processing unit is implicit in this model.

By treating the instructions in the same way as the data, a stored-program machine can easily change the instructions. In other words the machine is reprogrammable. One important motivation for such a facility was the need for a program to increment or otherwise modify the address portion of instructions. This became less important when index registers and indirect addressing became customary features of machine architecture.
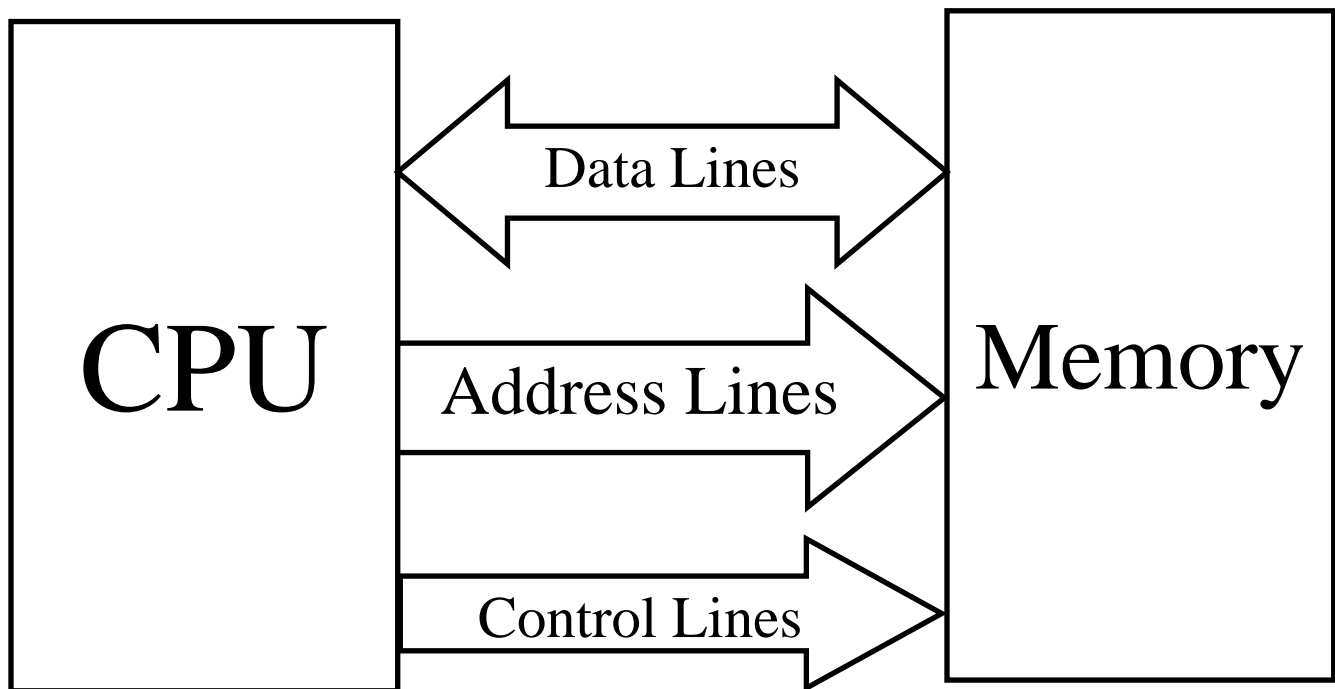
Fig. 5.2 The Memory Interface

In a memory read operation the CPU loads the address onto the address bus. Most cases these lines are fed to a decoder which selects the proper memory location. The CPU then sends a read control signal. The data is stored in that location is transferred to the processor via the data lines.

In the memory write operation after the address is loaded the CPU sends the write control signal followed by the data to the requested memory location.

The memory can be classified in various ways i.e. based on the location, power consumption, way of data storage etc

The memory at the basic level can be classified as

1. Processor Memory (Register Array)
2. Internal on-chip Memory
3. Primary Memory
4. Cache Memory
5. Secondary Memory

Processor Memory (Register Array)

Most processors have some registers associated with the arithmetic logic units. They store the operands and the result of an instruction. The data transfer rates are much faster without needing any additional clock cycles. The number of registers varies from processor to processor. The more is the number the faster is the instruction execution. But the complexity of the architecture puts a limit on the amount of the processor memory.

Internal on-chip Memory

In some processors there may be a block of memory location. They are treated as the same way as the external memory. However it is very fast.

Primary Memory

This is the one which sits just out side the CPU. It can also stay in the same chip as of CPU. These memories can be static or dynamic.

Cache Memory

This is situated in between the processor and the primary memory. This serves as a buffer to the immediate instructions or data which the processor anticipates. There can be more than one levels of cache memory.

Secondary Memory

These are generally treated as Input/Output devices. They are much cheaper mass storage and slower devices connected through some input/output interface circuits. They are generally magnetic or optical memories such as Hard Disk and CDROM devices.
The memory can also be divided into Volatile and Non-volatile memory.

Volatile Memory

The contents are erased when the power is switched off. Semiconductor Random Access Memories fall into this category.

Non-volatile Memory

The contents are intact even of the power is switched off. Magnetic Memories (Hard Disks), Optical Disks (CDROMs), Read Only Memories (ROM) fall under this category.
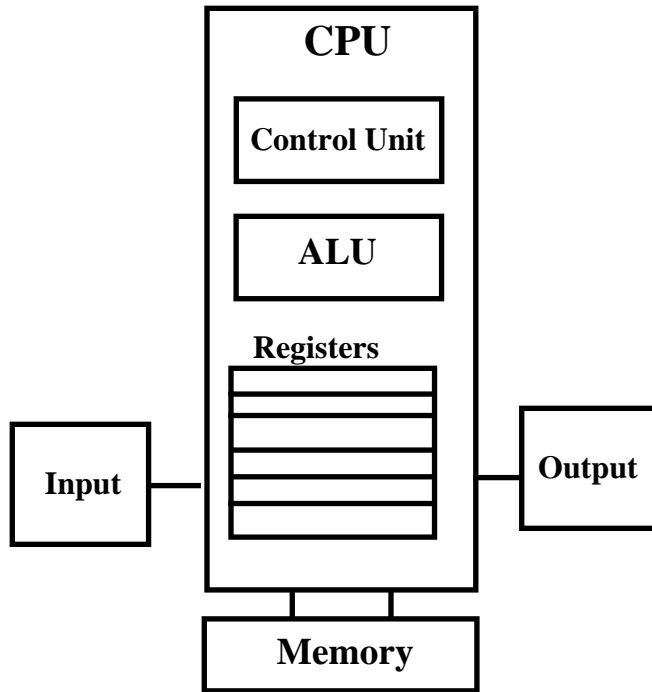
Fig. 5.3 The Internal Registers

5.2 Data Storage

An *m* word memory can store *m* x *n*: *m* words of *n* bits each. One word is located at one address therefore to address *m* words we need.

        k = Log2(*m*) address input signals
        or k number address lines can address $m = 2^k$ words

Example 4,096 x 8 memory:

- 32,768 bits
- 12 address input signals
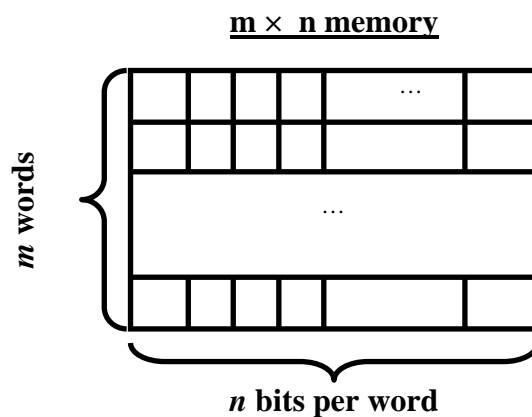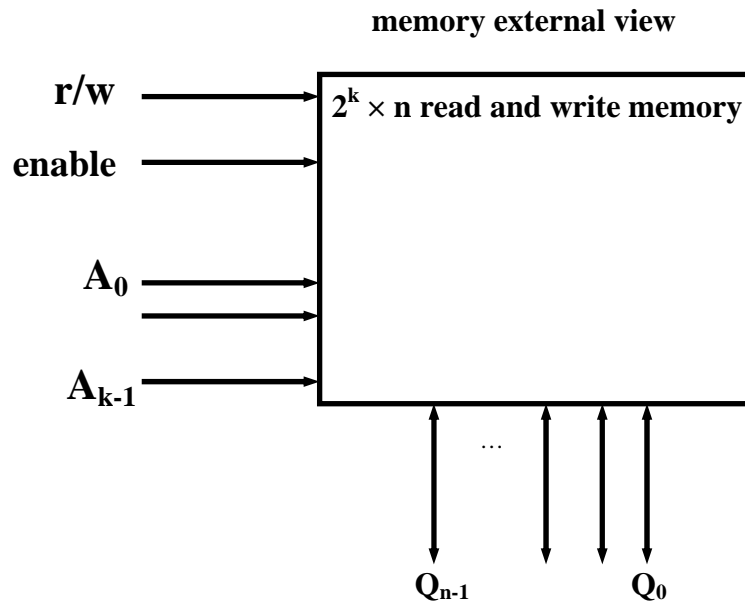- 8 input/output data signals

Fig. 5.4 Data Array

Memory access

The memory location can be accessed by placing the address on the address lines. The control lines read/write selects read or write. Some memory devices are multi-port i.e. multiple accesses to different locations simultaneously

memory external view

Fig. 5.5 Memory Array

Memory Specifications

The specification of a typical memory is as follows
The storage capacity: The number of bits/bytes or words it can store
      The memory access time (read access and write access): How long the memory takes to load the data on to its data lines after it has been addressed or how fast it can store the data upon supplied through its data lines. This reciprocal of the memory access time is known as *Memory*

Bandwidth

      The Power Consumption and Voltage Levels: The power consumption is a major factor in embedded systems. The lesser is the power consumption the more is packing density.
Size: Size is directly related to the power consumption and data storage capacity.

Generation 1

Generation 2

Generation 3

Generation 4

Fig. 5.6 Four Generations of RAM chips

There are two important specifications for the Memory as far as Real Time Embedded Systems are concerned.
  – Write Ability
  – Storage Performance

Write ability

It is the manner and speed that a particular memory can be written

- Ranges of write ability
  – High end
    - processor writes to memory simply and quickly e.g., RAM
  – Middle range
    - processor writes to memory, but slower e.g., FLASH, EEPROM (Electrically Erasable and Programmable Read Only Memory)
  – Lower range
    - special equipment, "programmer", must be used to write to memory   e.g., EPROM, OTP ROM (One Time Programmable Read Only Memory)
  – Low end
    - bits stored only during fabrication e.g., Mask-programmed ROM
- In-system programmable memory
  – Can be written to by a processor in the embedded system using the memory
  – Memories in high end and middle range of write ability

Storage permanence

It is the ability to hold the stored bits.
Range of storage permanence
  – High end
    - essentially never loses bits
    - e.g., mask-programmed ROM

- Middle range
  - holds bits days, months, or years after memory's power source turned off
  - e.g., NVRAM
- Lower range
  - holds bits as long as power supplied to memory
  - e.g., SRAM

- Low end
  - begins to lose bits almost immediately after written
  - e.g., DRAM

Nonvolatile memory
- Holds bits after power is no longer supplied
- High end and middle range of storage permanence

5.3 Common Memory Types

Read Only Memory (ROM)

This is a nonvolatile memory. It can only be read from but not written to, by a processor in an embedded system. Traditionally written to, "programmed", before inserting to embedded system
Uses

- Store software program for general-purpose processor
  - program instructions can be one or more ROM words
- Store constant data needed by system
- Implement combinational circuit

External view
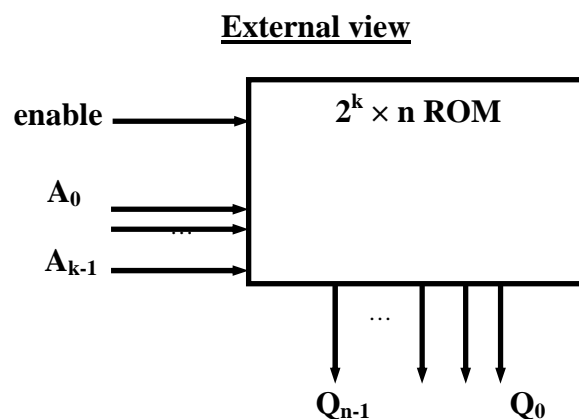
**Fig. 5.7 The ROM Structure**

Example

The figure shows the structure of a ROM. Horizontal lines represents the words. The vertical lines give out data. These lines are connected only at circles. If address input is 010 the decoder sets 2nd word line to 1. The data lines Q3 and Q1 are set to 1 because there is a "programmed"

connection with word 2's line. The word 2 is not connected with data lines Q2 and Q0. Thus the output is 1010

Internal view
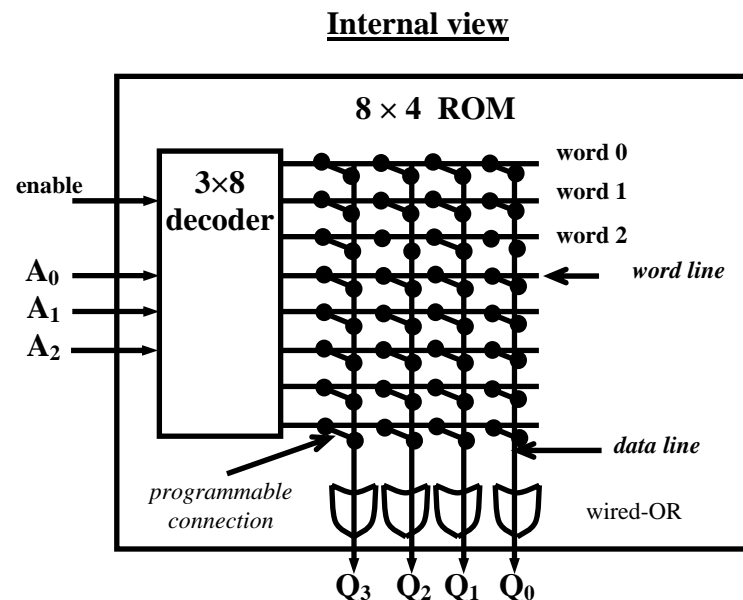
Fig. 5.8 The example of a ROM with decoder and data storage

Implementation of Combinatorial Functions

Any combinational circuit of n functions of same k variables can be done with 2^k x n ROM. The inputs of the combinatorial circuit are the address of the ROM locations. The output is the word stored at that location.
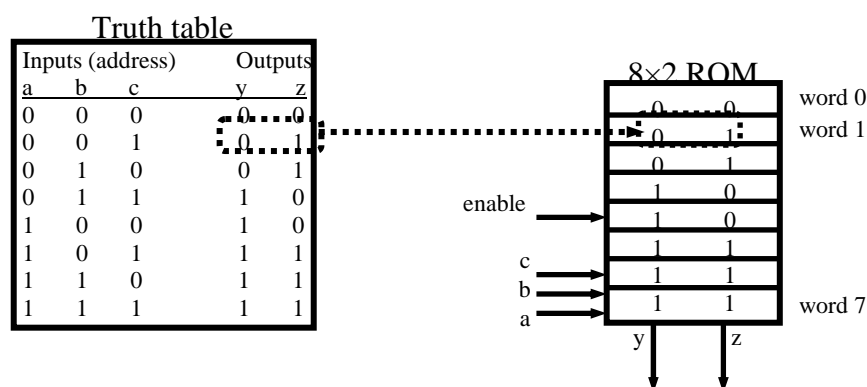
Fig. 5.9 The combinatorial table

Mask-programmed ROM

The connections "programmed" at fabrication. They are a set of masks. It can be written only once (in the factory). But it stores data for ever. Thus it has the highest storage permanence. The bits never change unless damaged. These are typically used for final design of high-volume systems.

OTP ROM: One-time programmable ROM

The Connections "programmed" after manufacture by user. The user provides file of desired contents of ROM. The file input to machine called ROM programmer. Each programmable connection is a fuse. The ROM programmer blows fuses where connections should not exist.

- Very low write ability: typically written only once and requires ROM programmer device
- Very high storage permanence: bits don't change unless reconnected to programmer and more fuses blown
- Commonly used in final products: cheaper, harder to inadvertently modify

EPROM: Erasable programmable ROM

This is known as erasable programmable read only memory. The programmable component is a MOS transistor. This transistor has a "floating" gate surrounded by an insulator. The Negative charges form a channel between source and drain storing a logic 1. The Large positive voltage at gate causes negative charges to move out of channel and get trapped in floating gate storing a logic 0. The (Erase) Shining UV rays on surface of floating-gate causes negative charges to return to channel from floating gate restoring the logic 1. An EPROM package showing quartz window through which UV light can pass. The EPROM has

- Better write ability

  − can be erased and reprogrammed thousands of times
- Reduced storage permanence

  − program lasts about 10 years but is susceptible to radiation and electric noise
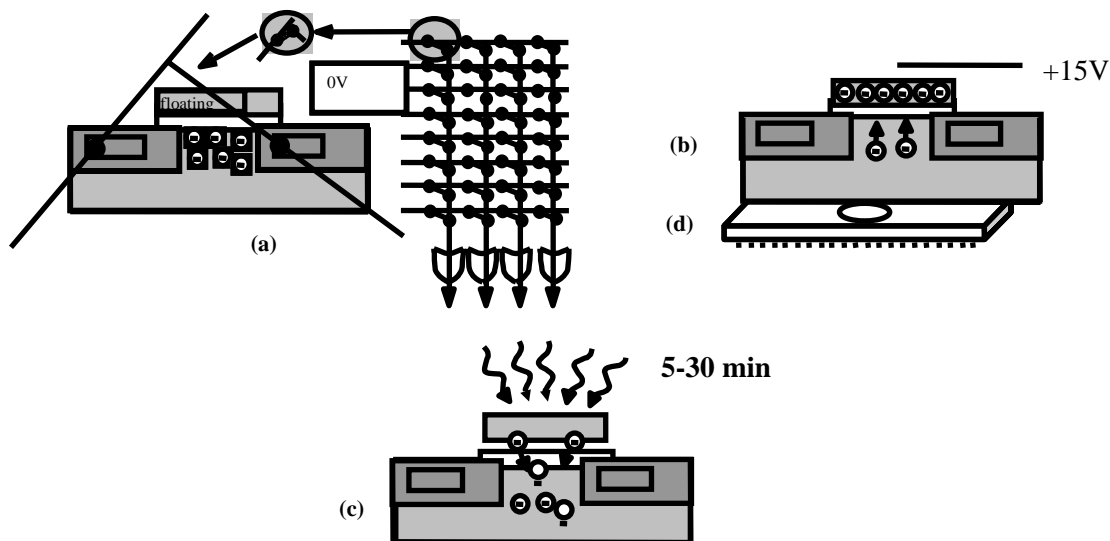- Typically used during design development

Fig. 5.10 The EPROM

EEPROM

EEPROM is otherwise known as Electrically Erasable and Programmable Read Only Memory. It is erased typically by using higher than normal voltage. It can program and erase individual words unlike the EPROMs where exposure to the UV light erases everything. It has

- Better write ability
    - can be in-system programmable with built-in circuit to provide higher than normal voltage
        - built-in memory controller commonly used to hide details from memory user
    - writes very slow due to erasing and programming
        - "busy" pin indicates to processor EEPROM still writing
    - can be erased and programmed tens of thousands of times
- Similar storage permanence to EPROM (about 10 years)
- Far more convenient than EPROMs, but more expensive

Flash Memory

It is an extension of EEPROM. It has the same floating gate principle and same write ability and storage permanence. It can be erased at a faster rate i.e. large blocks of memory erased at once, rather than one word at a time. The blocks are typically several thousand bytes large
- Writes to single words may be slower
    - Entire block must be read, word updated, then entire block written back
- Used with embedded systems storing large data items in nonvolatile memory
    - e.g., digital cameras, TV set-top boxes, cell phones

RAM: "Random-access" memory

- Typically volatile memory
    - bits are not held without power supply
- Read and written to easily by embedded system during execution
- Internal structure more complex than ROM
    - a word consists of several memory cells, each storing 1 bit
    - each input and output data line connects to each cell in its column
    - rd/wr connected to every cell
    - when row is enabled by decoder, each cell has logic that stores input data bit when rd/wr indicates write or outputs stored bit when rd/wr indicates read

external view
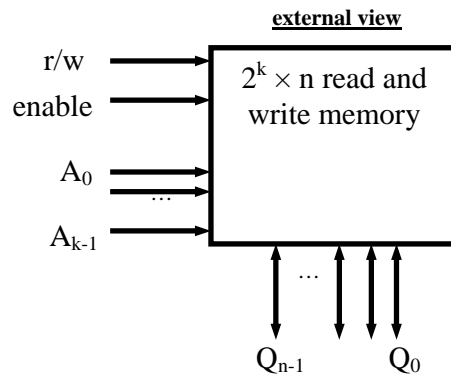
Fig. 5.11 The structure of RAM

internal view
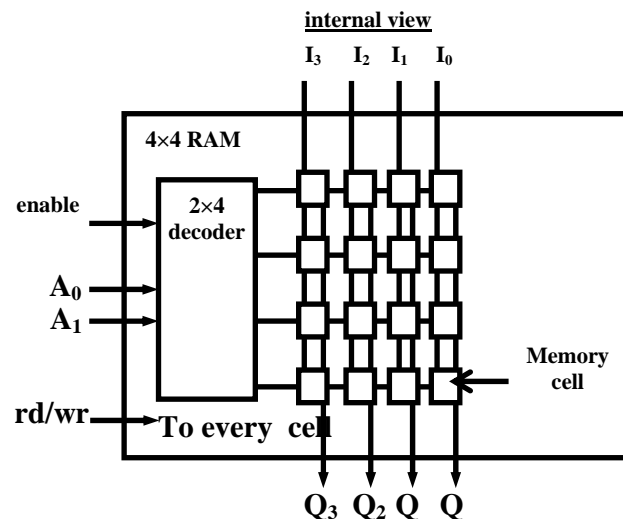
Fig. 5.12 The RAM decoder and access

Basic types of RAM

- SRAM: Static RAM
  - Memory cell uses flip-flop to store bit
  - Requires 6 transistors
  - Holds data as long as power supplied
- DRAM: Dynamic RAM
  - Memory cell uses MOS transistor and capacitor to store bit
  - More compact than SRAM
  - "Refresh" required due to capacitor leak
    - word's cells refreshed when read
  - Typical refresh rate 15.625 microsec.
  - Slower to access than SRAM
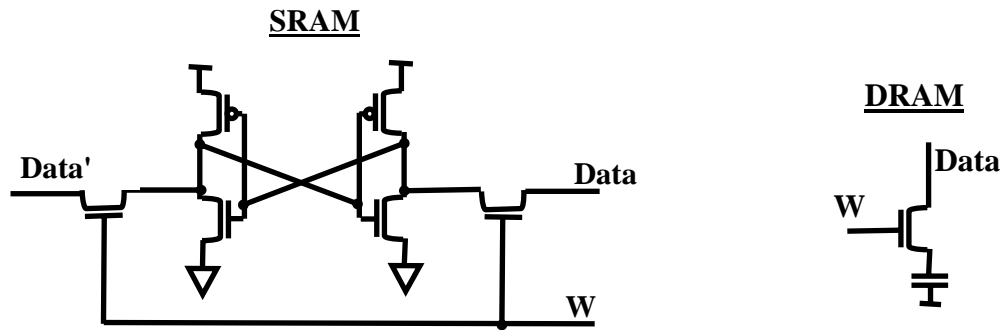
SRAM

DRAM

Ram variations

- PSRAM: Pseudo-static RAM
    - DRAM with built-in memory refresh controller
    - Popular low-cost high-density alternative to SRAM
- NVRAM: Nonvolatile RAM
    - Holds data after external power removed
    - Battery-backed RAM
        - SRAM with own permanently connected battery
        - writes as fast as reads
        - no limit on number of writes unlike nonvolatile ROM-based memory
    - SRAM with EEPROM or flash stores complete RAM contents on EEPROM or flash before power

5.4 Example: HM6264 & 27C256 RAM/ROM devices

- Low-cost low-capacity memory devices
- Commonly used in 8-bit microcontroller-based embedded systems
- First two numeric digits indicate device type
    - RAM: 62
    - ROM: 27
- Subsequent digits indicate capacity in kilobits

11-13, 15-19 → data<7…0>

2,23,21,24, 25, 3-10 → addr<15...0>

22 → /OE

27 → /WE

20 → /CS1

26 → CS2   **HM6264**

11-13, 15-19 → data<7…0>

27,26,2,23,21, 24,25, 3-10 → addr<15...0>
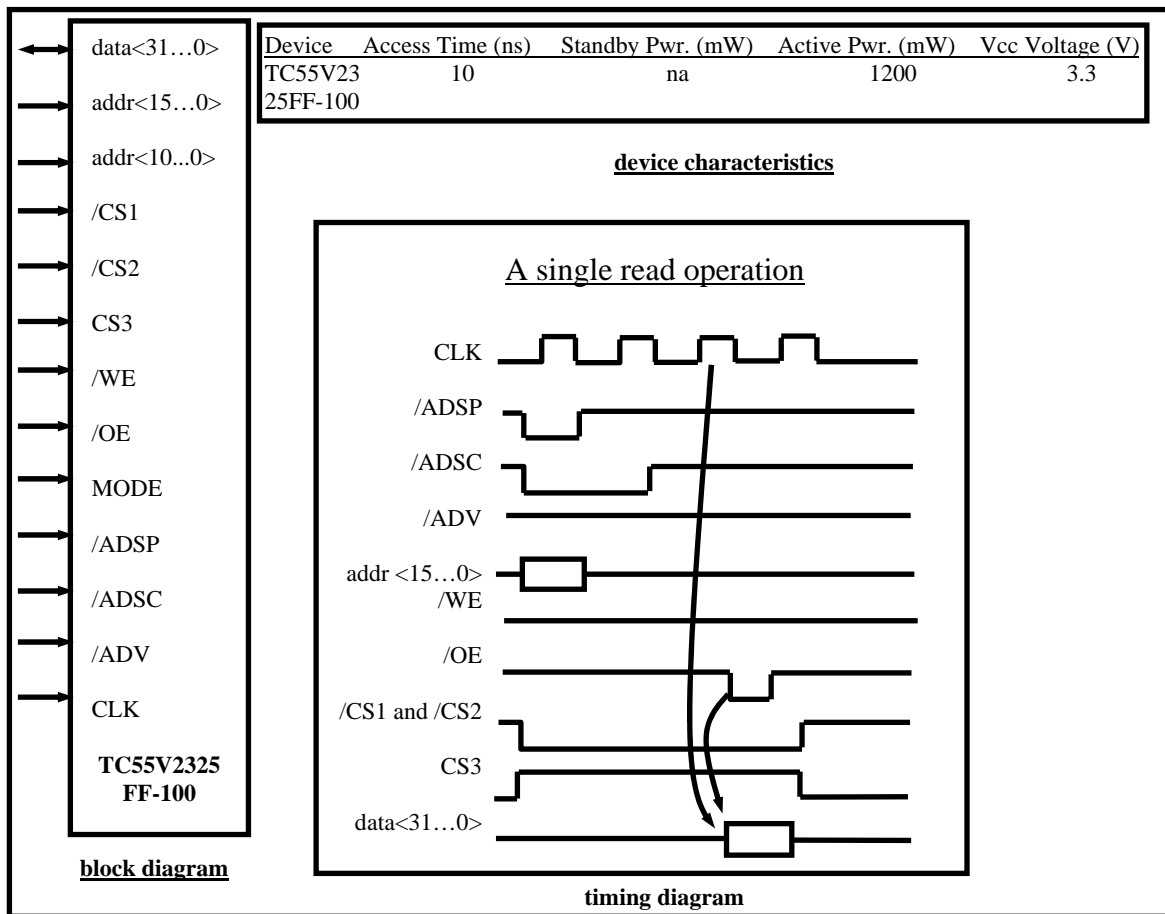
22 → /OE

20 → /CS

27C256

block diagrams

Device	Access Time (ns)	Standby Pwr. (mW)	Active Pwr. (mW)	Vcc Voltage (V)
HM6264	85-100	.01	15	5
27C256	90	.5	100	5

device characteristics

Read operation

data
addr
OE
/CS1
CS2

Write operation

data
addr
WE
/CS1
CS2

timing diagrams
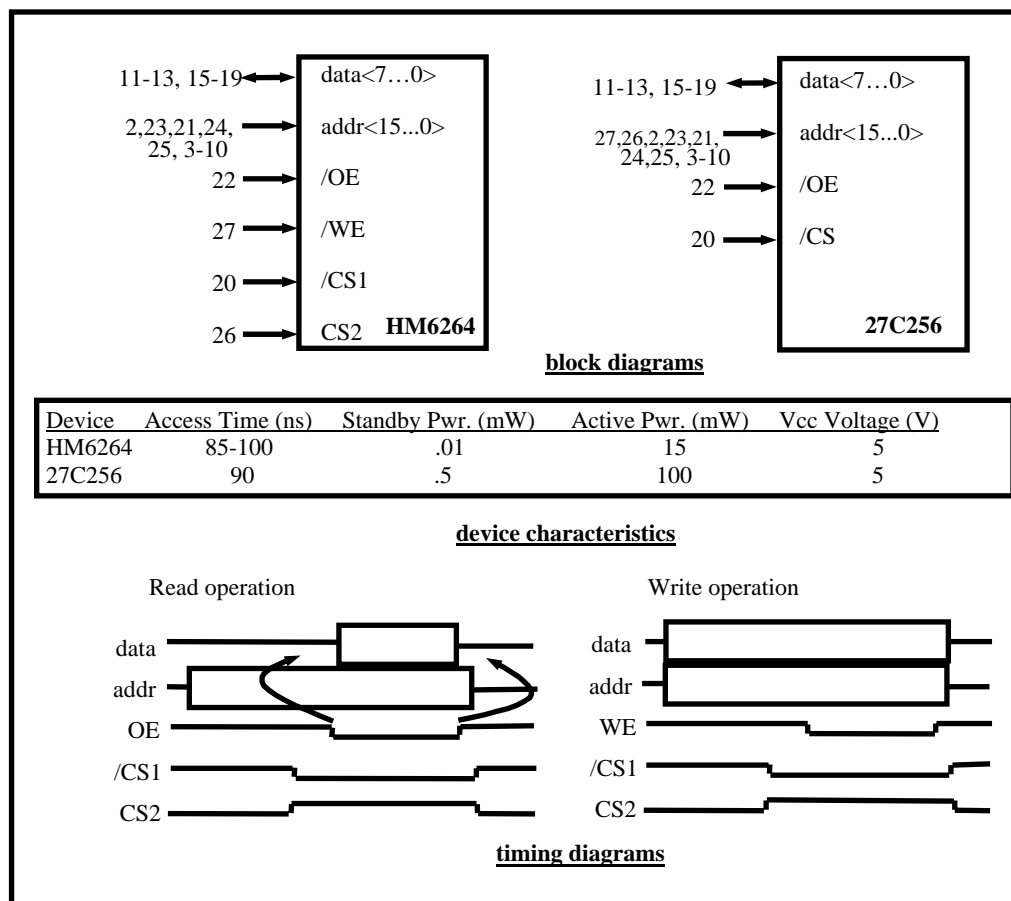
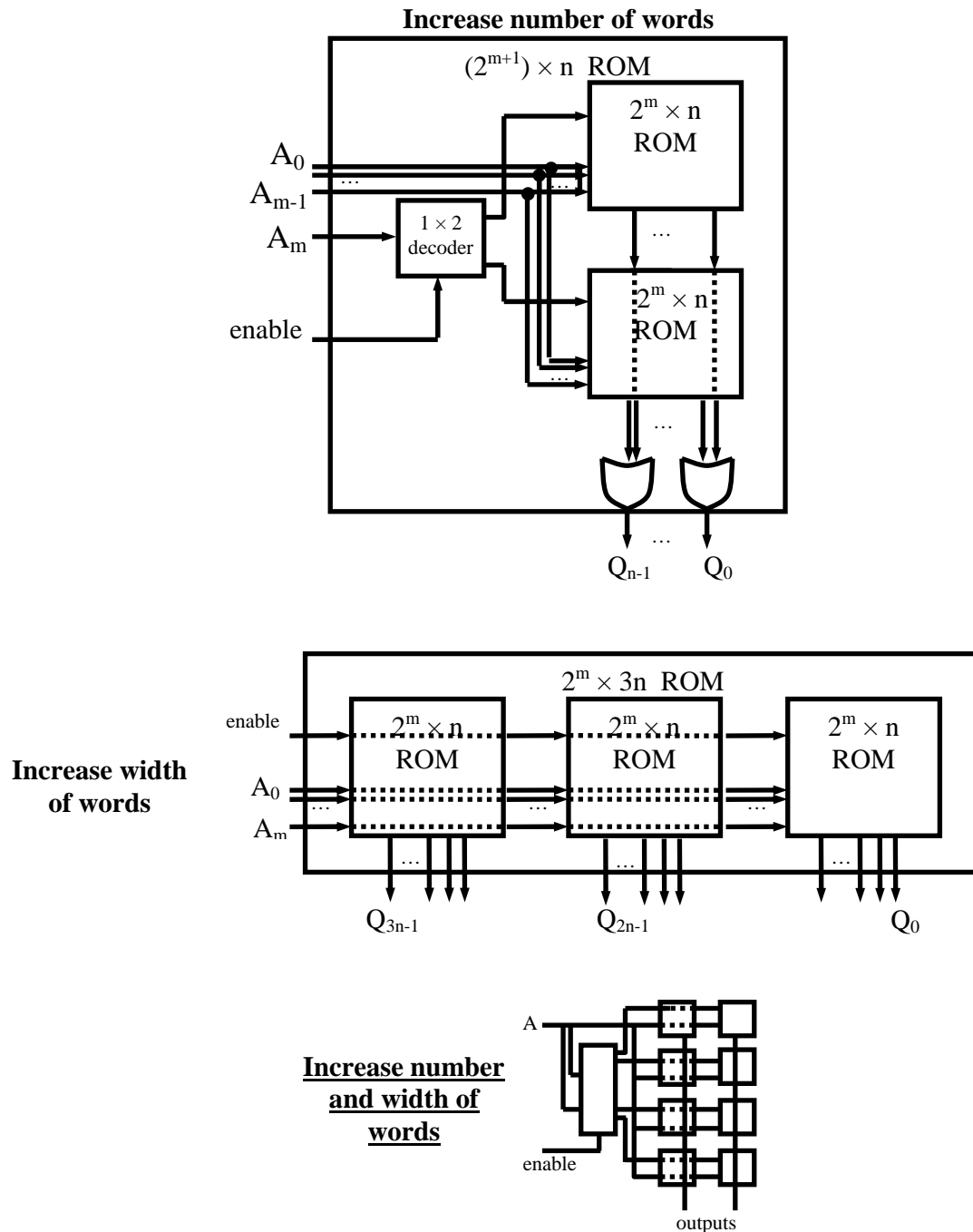# 5.5   Example: TC55V2325FF-100 memory device

- 2-megabit synchronous pipelined burst SRAM memory device
- Designed to be interfaced with 32-bit processors
- Capable of fast sequential reads and writes as well as single byte I/O

data<31…0>

addr<15…0>

addr<10...0>

/CS1

/CS2

CS3

/WE

/OE

MODE

/ADSP

/ADSC

/ADV

CLK

TC55V2325
FF-100

block diagram

Device	Access Time (ns)	Standby Pwr. (mW)	Active Pwr. (mW)	Vcc Voltage (V)
TC55V23 25FF-100	10	na	1200	3.3

device characteristics

A single read operation

CLK

/ADSP

/ADSC

/ADV

addr <15…0>
/WE

/OE

/CS1 and /CS2

CS3

data<31…0>

timing diagram

5.6 Composing memory

- Memory size needed often differs from size of readily available memories
- When available memory is larger, simply ignore unneeded high-order address bits and higher data lines
- When available memory is smaller, compose several smaller memories into one larger memory
  – Connect side-by-side to increase width of words
  – Connect top to bottom to increase number of words
    • added high-order address line selects smaller memory containing desired word using a decoder
  – Combine techniques to increase number and width of words

Increase number of words

$(2^{m+1}) \times n$ ROM

$2^m \times n$ ROM

A_0
A_{m-1}
A_m

...

1×2 decoder

enable

$2^m \times n$ ROM

...

...

Q_{n-1} Q_0

$2^m \times 3n$ ROM

enable

Increase width of words

A_0
A_m

$2^m \times n$ ROM

...

$2^m \times n$ ROM

...

$2^m \times n$ ROM

...

Q_{3n-1} Q_{2n-1} Q_0

A

Increase number and width of words

enable

outputs

Fig. 5.13 Composing Memory

5.7 Conclusion

In this chapter you have learnt about the following

1. Basic Memory types

2. Basic Memory Organization

3. Definitions of RAM, ROM and Cache Memory

4. Difference between Static and Dynamic RAM

5. Various Memory Control Signals

6. Memory Specifications

7. Basics of Memory Interfacing

5.8 Questions

Q1. Discuss the various control signals in a typical RAM device (say HM626)
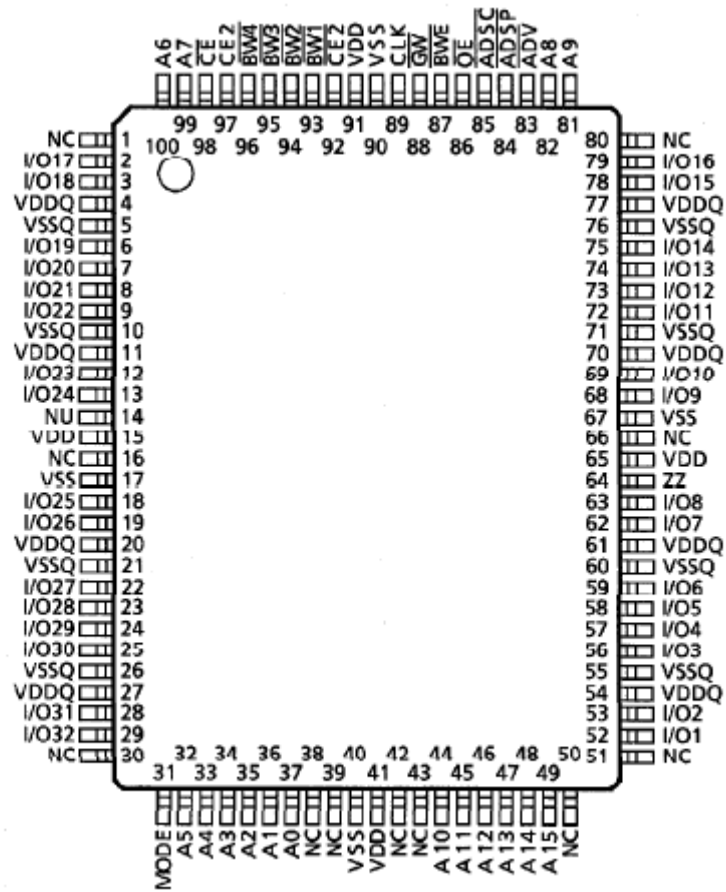
Ans:
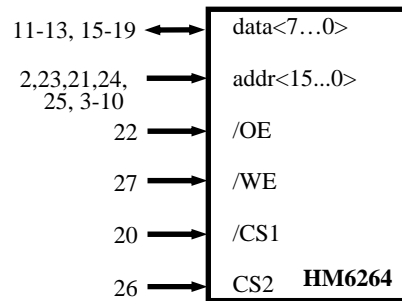
/OE: output enable bar: the output is enables when it is low. It is same as the read bar line
/WE: write enable bar: the line has to made low while writing to this device
CS1: chip select 1 bar: this line has to be made low along with CS2 bar to enable this chip

Q2. Download the datasheet of TC55V2325FF chip and indicate the various signals.

PIN NAMES

A0~A15	Address Inputs
I/O1~I/O32	Data Inputs / Outputs
CLK	Clock Input
$\overline{CE}$, $\overline{CE2}$, CE2	Chip Enable Inputs
$\overline{ADSP}$	Address Status Processor Input
$\overline{ADSC}$	Address Status Controller Input
$\overline{ADV}$	Address Advance Input
$\overline{GW}$	Global Write Enable Input
$\overline{BWE}$	Byte Write Enable Input
$\overline{BW1}$~$\overline{BW4}$	Byte Write Enable Inputs
$\overline{OE}$	Output Enable Input
MODE	Mode Select Input
ZZ	Snooze Input
NU	Not Usable Input
V_{DD}, V_{DDQ}	Power
V_{SS}, V_{SSQ}	Ground
NC	No Connection