

国 SCHEDULING

DCS4103 Operating System



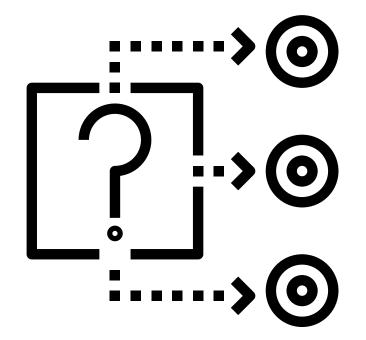
Process Scheduling

execution per unit time

☐ Process of determining which process will own CPU for execution while another process is on hold ☐ Select at least one of the processes available in the ready queue for execution ☐ Selection process carried out by the CPU scheduler Criteria Minimize Maximize ☐ Waiting time ☐ CPU utilization Process needs to wait in the ready CPU remains busy queue. ☐ Throughput ☐ Response time Number of processes that finish their

Request was submitted until the first

response is produced.



Uniprocessor Scheduling

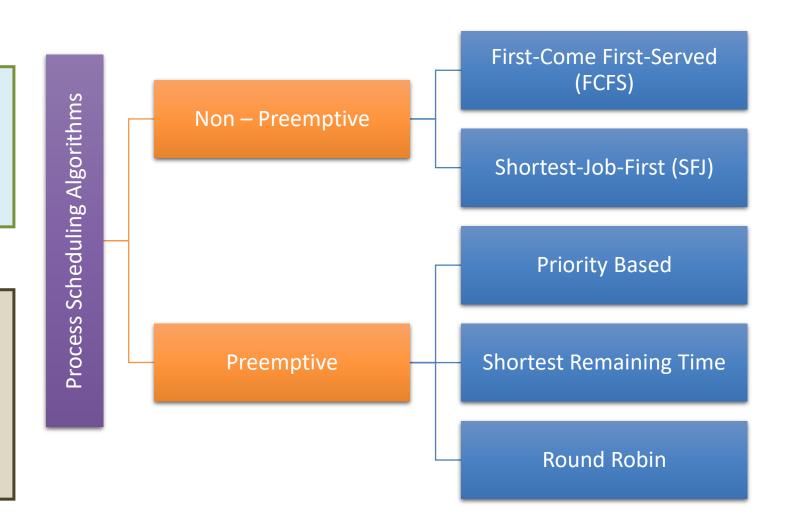


Non - Preemptive

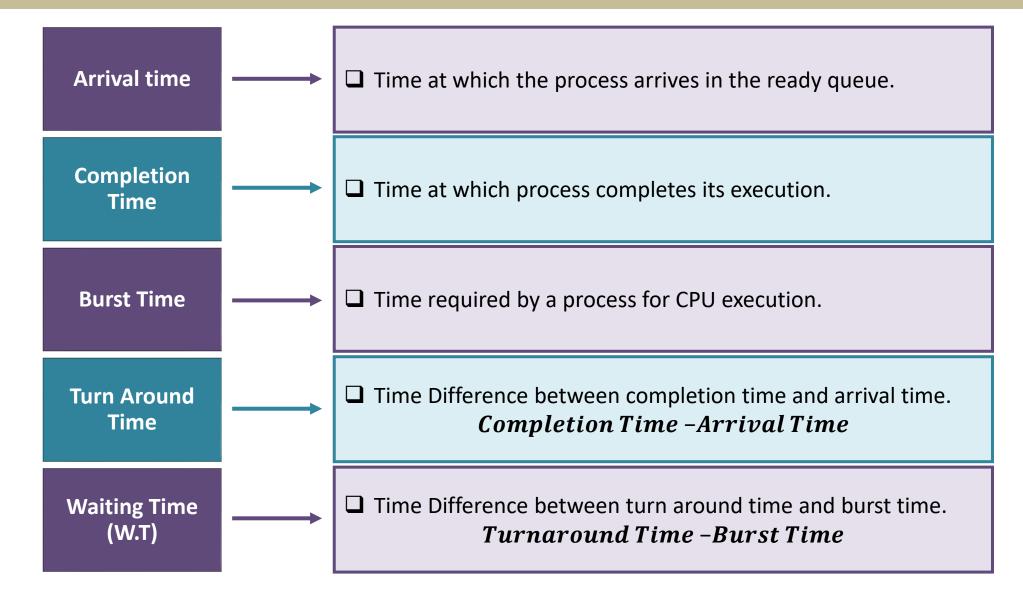
Process cannot be preempted until completes the allotted time

Preemptive

Based on priority where a scheduler may preempt a low priority running process anytime when a high priority process enters a ready state









	First-Come First Served (FSFC)			Disadvantages	
East Pool	plemented non-preemptive ecuted on a first-come, first-serve bases to implement and use. Or in performance sh waiting time ample Buying movie ticket at counter	sis	finish its execut Short processes wait for the long	that are at the back of g process at the front to time-sharing systems	the queue must

Shortest-Job-First (SJF)

Implemented non-preemptive
 The next process or job with the shortest completion time will be executed first.
 Useful for batch-type processing
 Improves job output by offering shorter jobs
 Reduces the average waiting time

Disadvantages

Job completion time must be known earlier
 Long turnaround times or starvation
 Elapsed time should be recorded, that results in more overhead on the processor.
 Hard to know the length of the upcoming CPU request.



Priority Based

- ☐ The scheduler selects the tasks to work as per the priority.
- ☐ Preemptive
 - Run a task with a higher priority
- ☐ Lower the number, higher is the priority.
- ☐ If two jobs having the same priority are READY, it works on a FIRST COME, FIRST SERVED basis.
- ☐ Easy to use
- ☐ Suitable for applications with fluctuating time and resource requirements.

Disadvantages

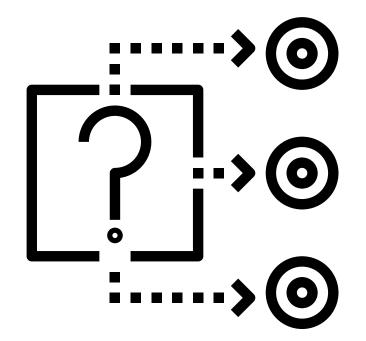
- ☐ If the system eventually crashes, all low priority processes get lost.
- ☐ If high priority processes take lots of CPU time,
 - the lower priority processes may starve and will be postponed for an indefinite time.
- ☐ If a new higher priority process keeps on coming in the ready queue,
 - the process which is in the waiting state may need to wait for a long duration of time.



Shortest Remaining Time				
 □ SJF preemptive scheduling Jobs are put into the ready queue Shortest burst time begins execution If a process with even a shorter burst time arrives, the current process is removed or preempted from execution, and the shorter job is allocated CPU cycle □ Prevents a newer ready state process from holding the completion of an older process. □ Applied in batch environments □ Helps to schedule the process with the shortest possible time 				
Disadvantages				
 □ Job completion time must be known earlier □ Long turnaround times or starvation □ Elapsed time should be recorded, that results in more overhead on the processor. □ Hard to know the length of the upcoming CPU request. 				



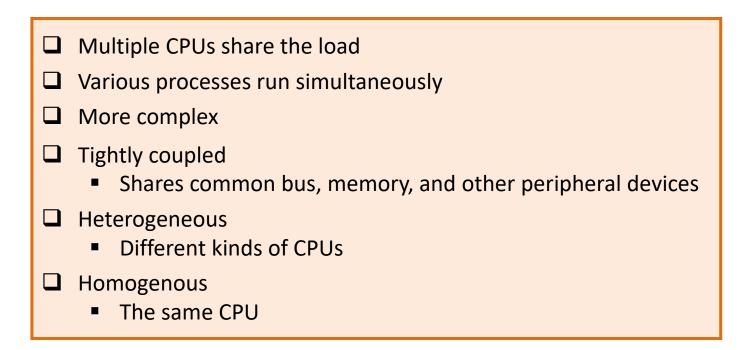
	Round Robin			
□ Implemented pre-emptive □ The oldest, easiest, fairest and simplest □ Each ready task runs turn by turn only in a cyclic queue for a limited time slice □ Preempted process added to the end of the queue. □ Clock-driven □ Used for multitasking □ Starvation free □ Without priority □ Performance in average response time				
Disadvantages				
Spends more time on contePerformance depends on ti	the processor output will be reduced ext switching me quantum. Its in higher the context switching or			



Multiprocessor Scheduling

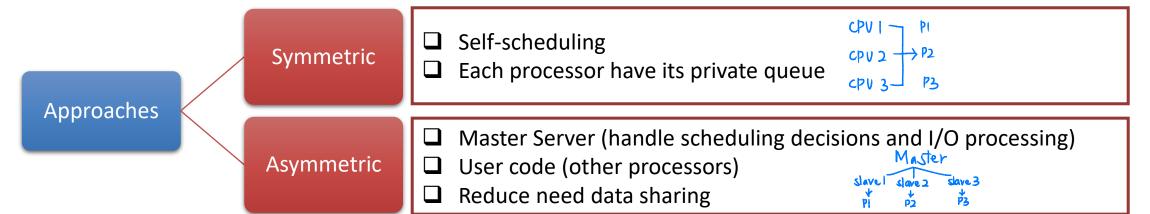


Introduction



Scheduling Algorithms

- Processor Affinity
- Load Balancing
- Multi-core Processors
- Symmetric Multiprocessor
- Master-Slave Multiprocessor
- Virtualization and Threading



Process	Arrival Time	Burst Time
P1	2 🚱	6
P2	5 🕏	2
P3	1 🕗	8
P4	0 (1)	3
P5	4 🖐	4

Calculate the Average Turnaround Time (TAT) and Average Waiting Time (AWT), if First-Come First-Serve (FCFS) algorithm is followed.

V 13 ડ 17 PI PS PY P3 Completion Time Turnariund Time Process ID Waiting Time p 12 PZ 16 73 10 \mathcal{V} PY 0 17 PS γI 13 Average Turnaround Time = 63 - 5 =17.6 units Average Waiting Time = 40 - 5 = 8 units

Gantt

Chart

Process	Arrival Time	Burst Time
P1	0 (1)	21
P2	1	3 🕟
P3	2	6 ③
P4	3	7 🏵

Calculate the Average Turnaround Time (TAT) and Average Waiting Time (AWT), if Shortest-Job-First(SJF) algorithm is followed.

Gantt chart

0 ~1 ~4 &0 &7

P1 | PV | P3 | P4

Process 40	Completion 7ime	Turnariund Time	Waiting Time
pj	ν •	V -4	0
PZ	W 1	γ3 -3	70
P3	20 1	7	\frac{1}{2}
РЧ	37 -3	34 -7	74

Kverage Turnaround Time = 106 - 4 = 76.5 mults

Average Waiting Time = 69 = + = 17.25 units

Revision Questions

Consider the set of five processes whose arrival time and burst time are given below:

Process	Arrival Time	Burst Time	Priority
P1	0	4	2
P2	1 ν	3	3
Р3	2 3	1	4
P4	3 +	5	5
P5	4 5	2	5

Calculate the Average Turnaround Time (TAT) and Average Waiting Time (AWT), if Priority Based algorithm is followed.

Gartt Chart PY P3 PS 74 Completion 7ime Turnariund Time Process ID Waiting Time PI 4 PZ (3 73 6 74 DI 25 Π 」 Average 77 = 37 55 Average WT = 22 5 5

Process	Arrival Time	Burst Time
P1	2	6
P2	5	2
P3	1	8
P4	0	3

Calculate the Average Turnaround Time (TAT) and Average Waiting Time (AWT), if Shortest Remaining Time algorithm is followed.

Process	Arrival Time	Burst Time
P1	0	5
P2	1	3
P3	2	1
P4	3	2
P5	4	3

If the CPU scheduling policy is Round Robin with time quantum = 2, calculate the Average Turnaround Time and Average Waiting Time