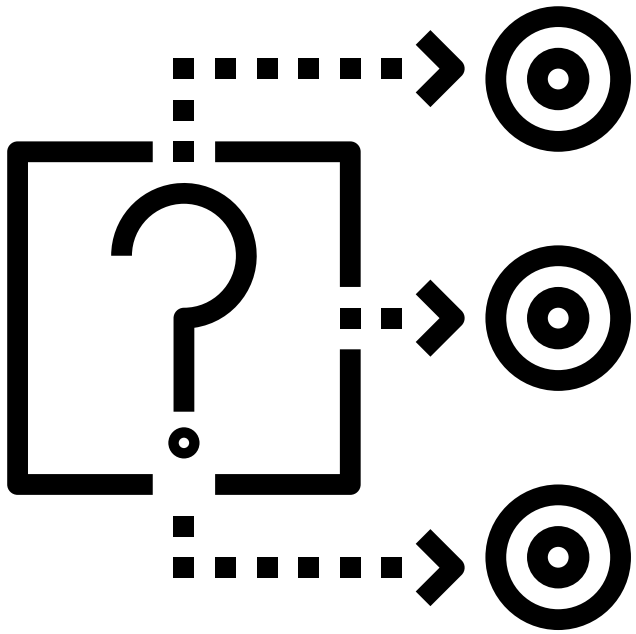




# DEADLOCK AND STARVATION

DCS4103 Operating System

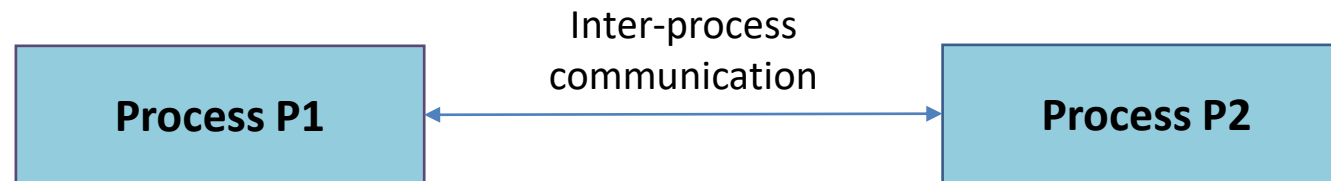


# Inter-Process Communication



# Introduction

Allows processes to communicate with each other and letting another process know that some event has occurred or the transferring of data from one process to another



## Advantages

- ☐ Helps to speedup modularity
- ☐ Computational
- ☐ Privilege separation
- ☐ Convenience
- ☐ Communication



# Synchronization

## 信号量 Semaphore

- Variable that controls the access to a common resource by multiple processes
  - Binary
  - Counting

## 互斥 Mutual Exclusion

- Only one process thread can enter the critical section at a time

## 障码 Barrier

- Does not allow individual process to process until all the process reach it

## Spinlock

- Acquire the spinlock waits or stays in a loop while checking that the lock is available or not
- Busy waiting
  - The process is active, but does not perform any functional operation



执行

# Implementation

## Pipe

- Unidirectional
- Moved in only a single direction at a time
- POSIX systems
- Windows OS

## Message Passing

- Allows processes to synchronize and communicate with each other without restoring the shared variables

## Direct Communication

- Create or establish link between two communication processes
- Only one link can exist

## Indirect Communication

- Only exist or be established when processes share a common mailbox
- Share multiple communication links



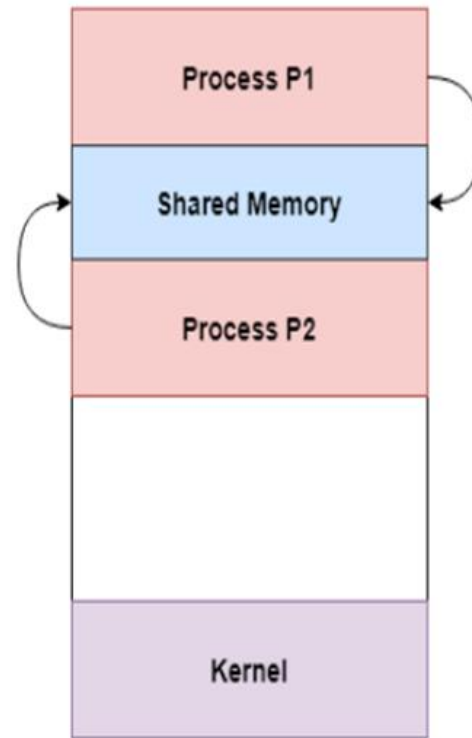
# Implementation

## Shared Memory

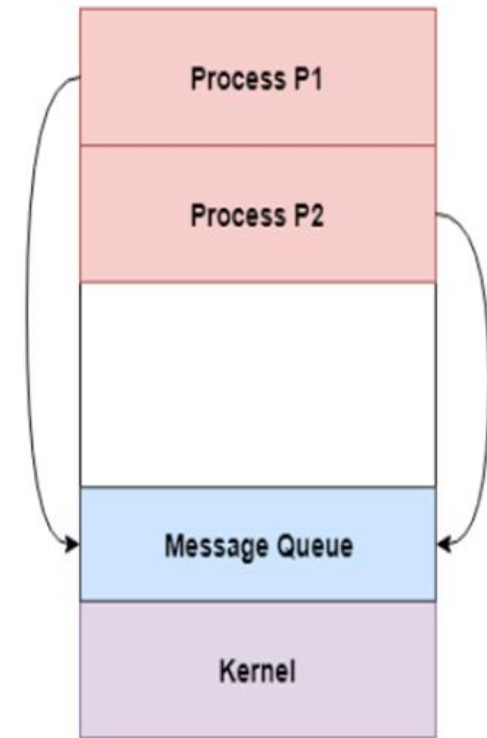
- Which can be used or accessed by multiple processes simultaneously
- Processes can communicate with each other

## Message Queue

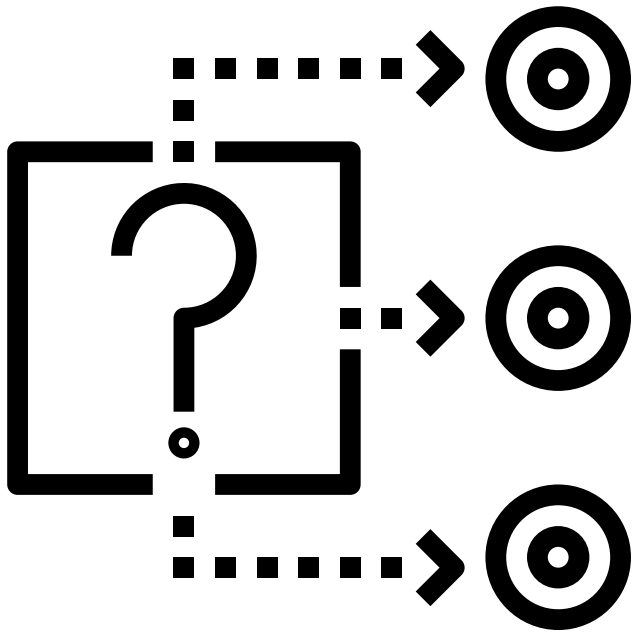
- Allows different messages to read and write the data
- The message are stored or stay in the queues unless their recipients retrieve them



Shared Memory



Message Queue

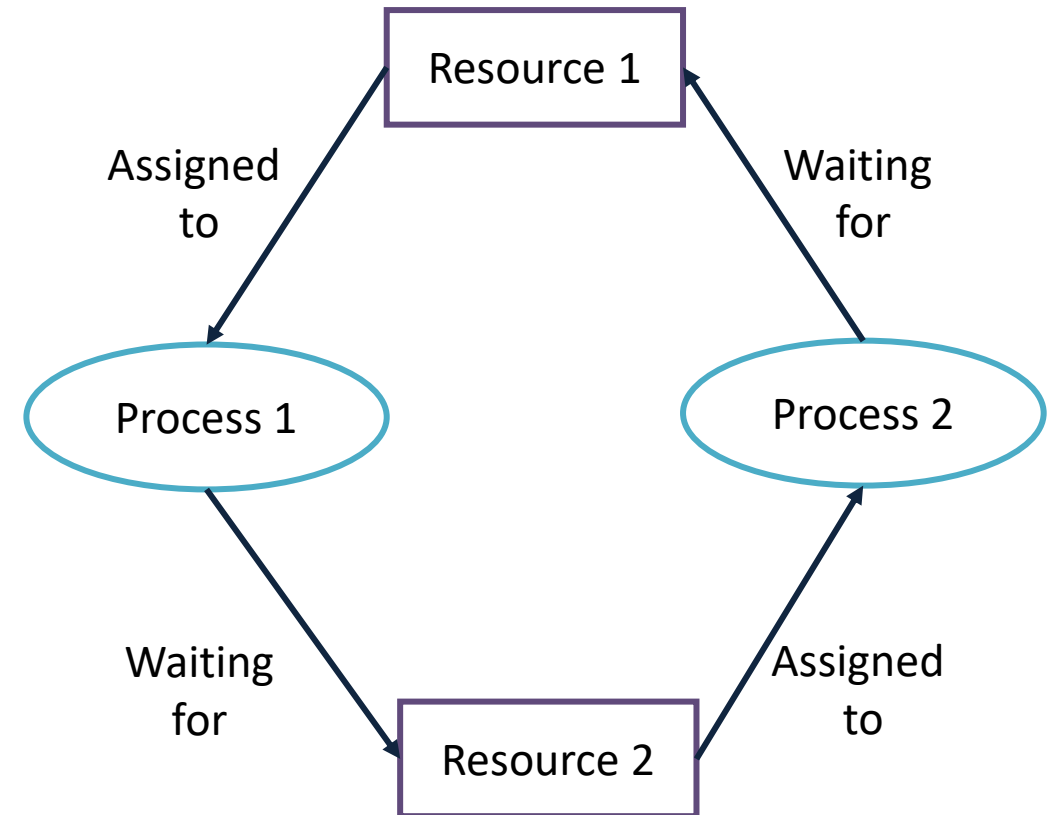
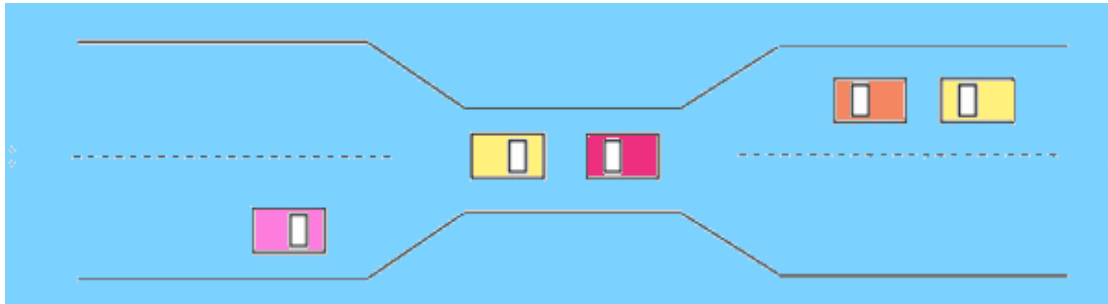


# Deadlock



# Introduction

- ❑ A situation that occurs in OS when any process enters a waiting state
- ❑ Another waiting process holding the demanded resource
- ❑ Examples
  - Traffic
  - Bridges







## Principles / Causes

### Mutual Exclusion

- Two process cannot use the same resource at the same time

### Hold and Wait

- A process waits for some resources while holding another resource at the same time.

### No preemption

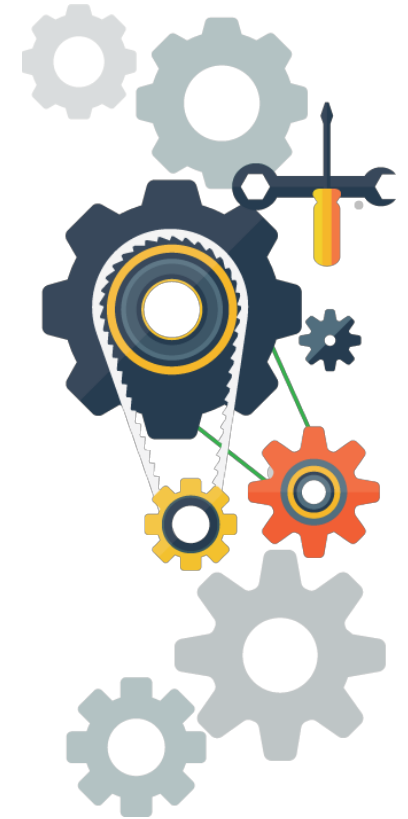
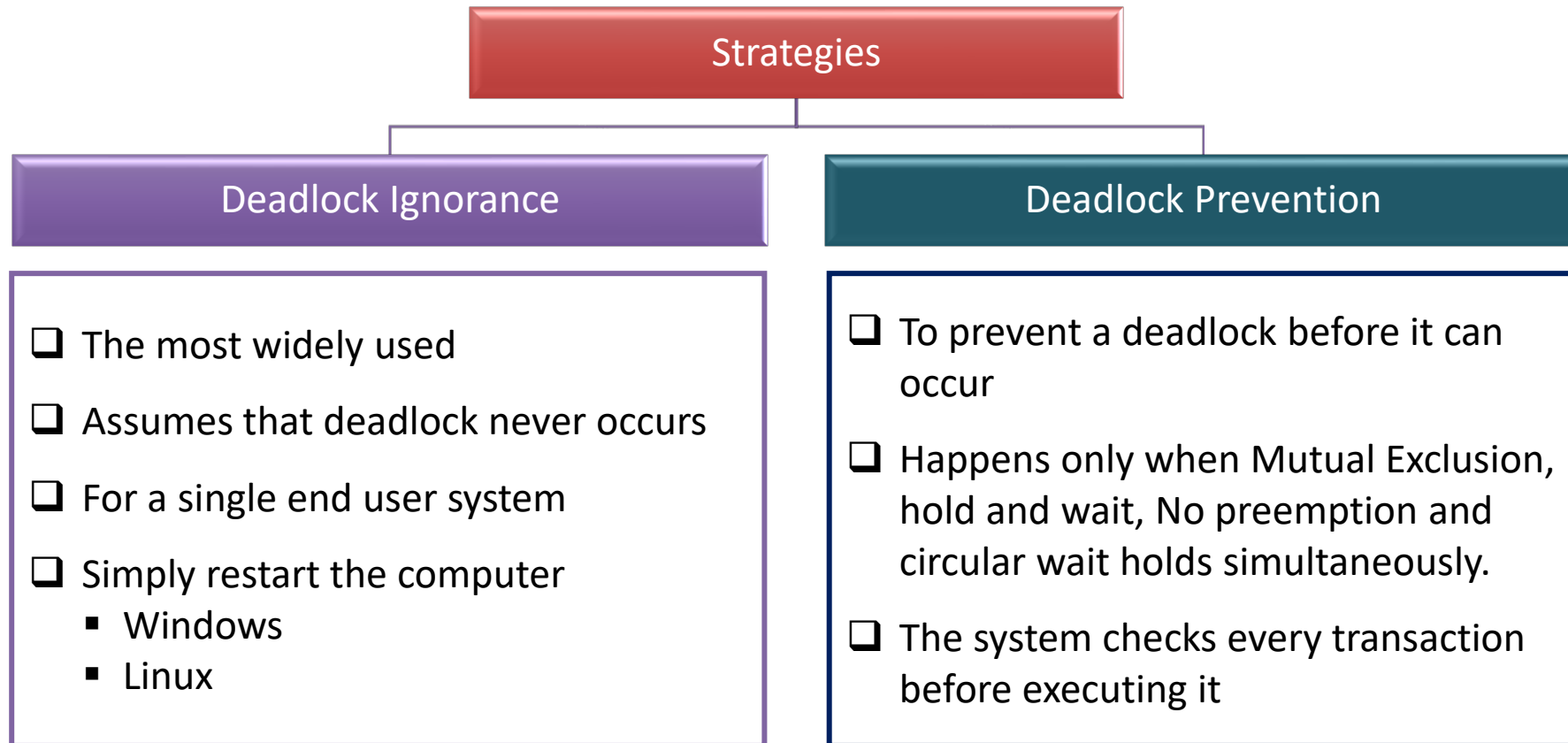
- The process which once scheduled will be executed till the completion.
- No other process can be scheduled by the scheduler meanwhile.

### Circular Wait

- All the processes must be waiting for the resources in a cyclic manner so that the last process is waiting for the resource which is being held by the first process.

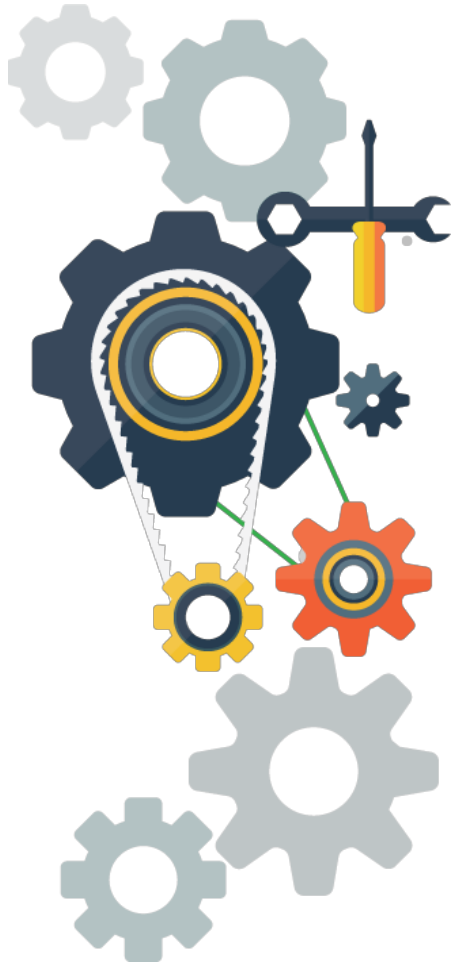


# Strategies





# Strategies



## Deadlock Avoidance

- ☐ The simplest and most useful
- ☐ Declare the maximum number of resources a process can request to complete its execution.

## Avoidance Algorithm

- ☐ Resource-allocation graph
  - Pictorial representation of the state of a system

## Banker's Algorithm

- ☐ If the customer's request does not cause the bank to leave a safe state, the cash will be allocated, otherwise the customer must wait until some other customer deposits enough
- ☐ To test for safely simulating the allocation for determining the maximum amount available for all resources
- ☐ Checks for all the possible activities before determining whether allocation should be continued or not.



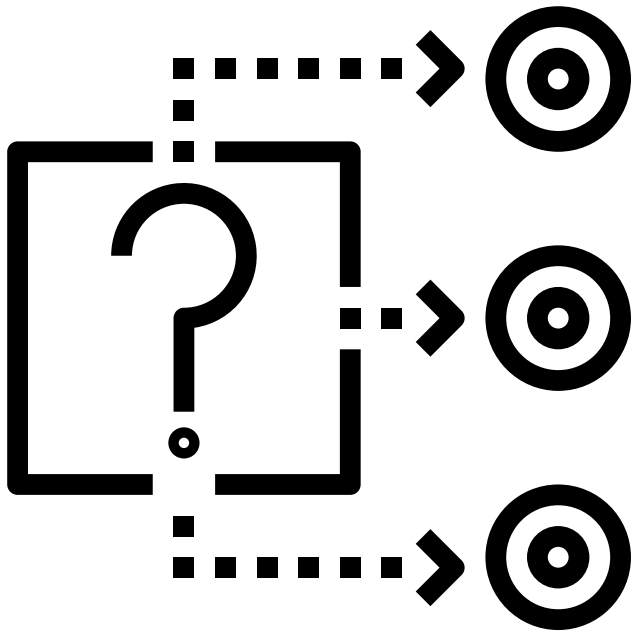
# Banker's Algorithm

## Characteristics

- ☐ Keep many resources that satisfy the requirement of at least one client
- ☐ Whenever a process gets all its resources, it needs to return them in a restricted period.
- ☐ When a process requests a resource, it needs to wait
- ☐ The system has a limited number of resources
- ☐ Advance feature for max resource allocation

## Disadvantages

- ☐ Does not allow the process to change its Maximum need while processing
- ☐ It allows all requests to be granted in restricted time, but one year is a fixed period for that.
- ☐ All processes must know and state their maximum resource needs in advance.



**Starvation**



# Introduction

- ☐ Happens when a low priority program requests a system resource but cannot run
- ☐ Higher priority program has been employing that resource for a long time
- ☐ Prevent the lower priority process from obtaining the requested resource.
  - Priority Scheduling Algorithm

## Causes

- ☐ Lack of resources
- ☐ Faulty resource allocation decisions
- ☐ Higher priority operations monopolize the processor

## Solutions

### Aging

- The priority of a process increases the longer it waits.

### Manager

- Independently distributes resources properly

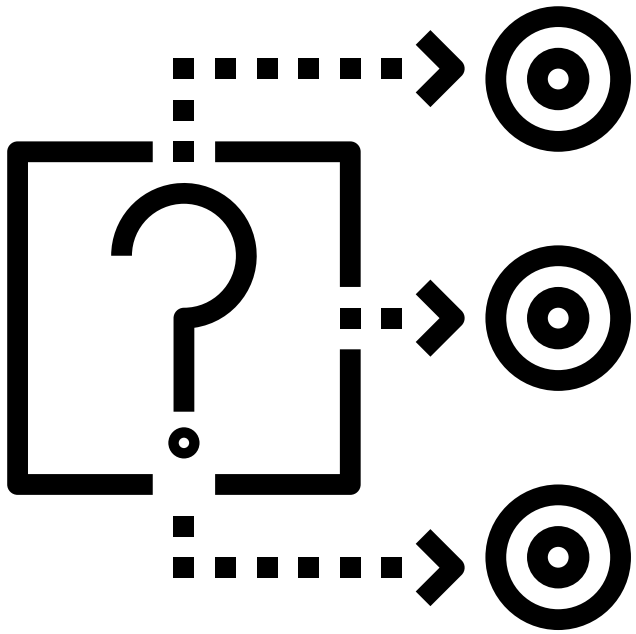
### Avoidance

- Random process selection



# Deadlock vs Starvation

| Features          | Deadlock  | Starvation  |
|-------------------|---|---|
| Definition        | When every process holds a resource and waits for another process to hold another resource.               | When a low priority program requests a system resource but cannot run because a higher priority program has been employing that resource for a long time. |
| Basic             | Occurs when no process can proceed and becomes blocked.   | Occurs when low priority procedures are blocked while high priority operations proceed.   |
| Other names       | Circular wait.  | Lived Lock  |
| Resources         | Other processes block requested resources while a process is deadlocked.                                  | High-priority processes continue to use the requested resources.  |
| Arising Condition | Mutual exclusion's occurrence, Hold and wait, No preemption, and Circular wait all happen simultaneously. | Uncontrolled resource management, enforcement of priorities.  |
| Prevention        | Avoidance   | Aging   |



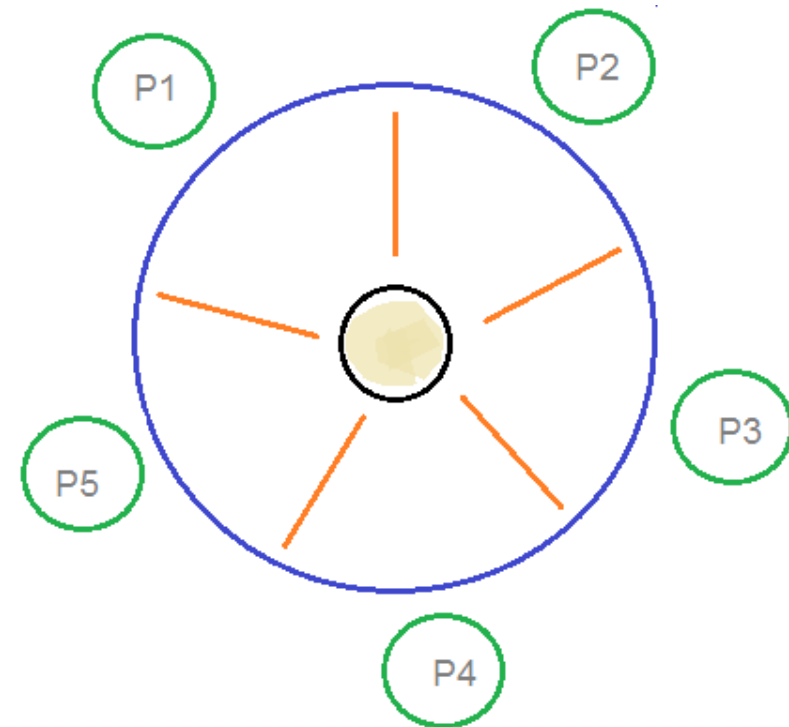
# Dining Philosophers Problem





# Introduction

- ❑ Evaluate situations where there is a need of allocating multiple resources to multiple processes.
- ❑ When a philosopher wants to eat, he uses two chopsticks
  - one from their left
  - one from their right
- ❑ When a philosopher wants to think, he keeps down both chopsticks at their original place.





# Semaphore

- ☐ Wait for the chopstick at the left and picks up that chopstick.
- ☐ Then wait for the right chopstick to be available, and then picks it too.
- ☐ After eating, put both the chopsticks down.
- ☐ Deadlock
  - Two neighboring philosophers want to eat at the same time

## Solution

- ☐ One-to-one relationship of user-level thread to the At most four philosophers on the table
- ☐ An even philosopher should pick the right chopstick and then the left chopstick while an odd philosopher should pick the left chopstick and then the right chopstick.
- ☐ Only be allowed to pick chopstick if both are available at the same time.



## Revision Questions

1. List approaches to implement inter-process communication.
2. Describe the cause deadlock to arise
3. Explain starvation and the causes.

EXERCISE





## Revision Questions

Consider the following set of processes:

| Process | Allocation |    |    | Max |    |    | Work |    |    |
|---------|------------|----|----|-----|----|----|------|----|----|
|         | R1         | R2 | R3 | R1  | R2 | R3 | R1   | R2 | R3 |
| $P_0$   | 1          | 1  | 2  | 4   | 3  | 3  | 2    | 1  | 0  |
| $P_1$   | 2          | 1  | 2  | 3   | 2  | 2  |      |    |    |
| $P_2$   | 4          | 0  | 1  | 9   | 0  | 2  |      |    |    |
| $P_3$   | 0          | 2  | 0  | 7   | 5  | 3  |      |    |    |
| $P_4$   | 1          | 1  | 2  | 1   | 1  | 2  |      |    |    |

Compute whether the system is safe or not using the Banker's algorithm.  
Determine the sequence if it is safe.



## Revision Questions

Consider the following set of processes:

| Process | Allocation |    |    |    | Max |    |    |    | Work |    |    |    |
|---------|------------|----|----|----|-----|----|----|----|------|----|----|----|
|         | R1         | R2 | R3 | R4 | R1  | R2 | R3 | R4 | R1   | R2 | R3 | R4 |
| $P_0$   | 1          | 2  | 1  | 0  | 3   | 3  | 2  | 3  | 6    | 4  | 6  | 4  |
| $P_1$   | 2          | 1  | 2  | 0  | 3   | 3  | 2  | 2  |      |    |    |    |
| $P_2$   | 1          | 0  | 1  | 2  | 2   | 1  | 3  | 3  |      |    |    |    |
| $P_3$   | 1          | 0  | 0  | 1  | 3   | 2  | 2  | 2  |      |    |    |    |

Compute whether the system is safe or not using the Banker's algorithm.  
Determine the sequence if it is safe.