This documentation provides an illustration (and also a proof) of Metropolis algorithm of sampling. The proof follows Metropolis et al (1953), but is modified and generalized in its morden face.

Notation 1. [Discrete Version]¹ Let \mathcal{X} state-space of a given stochastic system. Suppose \mathcal{X} is discrete or has been discretized, within which state is denoted by x_r , or simply r. Let p(r) denotes the target distribution to be mimicked. Let $q(r \to s)$ (or q(s|r) by statistics) the proposed transition-distribution of Markov process from state r to s. Suppose, for any $r, s \in \mathcal{X}$, $q(r \to s) \neq 0$ and $p(r) \neq 0$ (thus they are positive). That is, in \mathcal{X} any r and s are "connected". And for Metropolis algorithm, $q(r \to s) = q(s \to r)$ for $\forall r, s \in \mathcal{X}$ is supposed.

Algorithm 1

```
[Metropolis]
#!/usr/bin/env python3
# -*- coding: utf-8 -*-
import random
class MetropolisSampler:
   Args:
      iterations: int
      initialize state: (None -> State)
      markov process: (State -> State)
      burn in: int
   Attributes:
      accept ratio: float
         Generated only after calling MetropolisSampler.
   Methods:
      sampling:
         Do the sampling by Metropolis algorithm.
   Remarks:
      The "State" can be any abstract class.
   def __init__(self, iterations, initialize_state, markov_process, burn_in):
      self.iterations = iterations
      self.initialize \ state = initialize\_state
      self.markov\_process = markov\_process
      self.burn \ \ in = burn\_in
   def sampling(self, target distribution):
      Do the sampling.
      Args:
         target distribution: (State -> float)
```

^{1. [}Continuum Version] Let $\mathcal X$ state-space of a given stochastic system. Let p(x) denotes the target distribution to be mimicked. Let $q(x \to y)$ (or q(y|x) by statistics) the proposed transition-distribution of Markov process from state x to y. Suppose, for any $x, y \in \mathcal X$, $q(x \to y) \neq 0$ and $p(x) \neq 0$ (thus they are positive). That is, in $\mathcal X$ any x and y are "connected". And for Metropolis algorithm, $q(x \to y) = q(y \to x)$ for $\forall x, y \in \mathcal X$ is supposed.

```
Returns:
  list of State, with length being iterations - burn in.
init state = self.initialize state()
assert target distribution(init state) > 0
chain = [init\_state]
accepted = 0
for i in range(self.iterations):
   next state = self.markov process(init state)
   alpha = target distribution(next state) / target distribution(init state)
   u = random.uniform(0, 1)
   if alpha > u:
      accepted += 1
      chain.append(next state)
      init state = next state.copy()
   else:
      chain.append(init state)
self.accept ratio = accepted / self.iterations
print('accept-ratio = \{0\}'.format(self.accept ratio))
return chain[self.burn in:]
```

Lemma 2. Define, for $\forall r \in \mathcal{X}$, $J(r) := {\forall s \in \mathcal{X} : p(r) \ge p(s)}$. Let $h_i(r)$ the distribution of $\forall r \in \mathcal{X}$ at the ith Markov epoch. For any given h_i , we have

$$h_{i+1}(r) - h_i(r) = \sum_{s \in \mathcal{X}} \left\{ q(s \to r) \left[\frac{\delta_{s \in J(r)}}{p(r)} + \frac{\delta_{s \notin J(r)}}{p(s)} \right] \right\} \left\{ h_i(s) \ p(r) - h_i(r) \ p(s) \right\}.$$

Note that the first $\{...\}$ are symmetric for s and r, while the second is anti-symmetric.

Proof. Directly from Metropolis algorithm, we have,

$$h_{i+1}(r) - h_i(r) = \sum_{s \in J(r)} h_i(s) \ q(s \to r) + \sum_{s \notin J(r)} h_i(s) \ q(s \to r) \frac{p(r)}{p(s)} - \sum_{s \in J(r)} h_i(r) \ q(r \to s) \frac{p(s)}{p(r)} + \sum_{s \notin J(r)} h_i(r) \ q(r \to s).$$

Since, in Metropolis algorithm, $q(s \rightarrow r) = q(r \rightarrow s)$ is supposed,

$$h_{i+1}(r) - h_{i}(r) = \sum_{s \in J(r)} h_{i}(s) q(s \to r) + \sum_{s \notin J(r)} h_{i}(s) q(s \to r) \frac{p(r)}{p(s)} - \sum_{s \in J(r)} h_{i}(r) q(s \to r) \frac{p(s)}{p(r)} + \sum_{s \notin J(r)} h_{i}(r) q(s \to r);$$

then by direct simplification, we reach

$$h_{i+1}(r) - h_i(r) = \sum_{s \in \mathcal{X}} \left\{ q(s \to r) \left[\frac{\delta_{s \in J(r)}}{p(r)} + \frac{\delta_{s \notin J(r)}}{p(s)} \right] \right\} \{ h_i(s) \ p(r) - h_i(r) \ p(s) \}.$$

Corollary 3. If $h_i = p$, then $h_{i+1} = p$.

Theorem 4. [Metropolis] Samples generated by algorithm 1 approximately obeys the target distribution p(x). That is, algorithm 1 creates a sampler of p(x).

Proof. [Intuitive]

Let $\forall r_1 \in \mathcal{X}$ given. We have a Markov chain generated by Metropolis algorithm starting at r_1 , say $\{r_1, r_2, ..., r_N\}$. We want to prove that $\{r_1, r_2, ..., r_N\}$ approximately obeys distribution p(x). This proof has three steps.

- S1) Define h by $\{r_1, r_2, ..., r_{N-1}\} \sim h$, and h' by $\{r_2, r_3, ..., r_N\} \sim h'$. That is, the histogram of $\{r_1, r_2, ..., r_{N-1}\}$, after normalization, can be fitted by h, and the same for h'. We will proof first that h = h' + o(1) on \mathcal{X} as $N \to +\infty$. Indeed, if N is large enough, $\{r_1, r_2, ..., r_{N-1}\} \sim h$, since dropping a single element will affect little on the fitting of histogram (will be reviewed later in remark 5). So, h = h' + o(1) on \mathcal{X} as $N \to +\infty$.
- S2) Next is to proof that $h' = h \Rightarrow h = p$ on \mathcal{X} . (This proof temporally employs the Perron–Frobenius theorem.)
- S2.1) Indeed, we shall not forget that $\{r_1, r_2, ..., r_N\}$ are generated by a Markov chain obeying Metropolis algorithm. Thus, the Metropolis algorithm as a Markov process brings $r_i \to r_{i+1}$ for $\forall i = 1, ..., N-1$, s.t. $\{r_1, r_2, ..., r_{N-1}\} \to \{r_2, r_3, ..., r_N\}$. That is to say, h' can be regarded as being generated by Metropolis algorithm from h. So, h and h' are related by lemma 2.
- S2.2) Since h and h' are related so, as in the proof of lemma 2 before inserting the condition $q(s \to r) = q(r \to s)$,

$$\begin{array}{lll} h'(r) & = & h(r) \\ & + & \displaystyle \sum_{s \in J(r)} \; h(s) \; q(s \mathop{\rightarrow} r) + \sum_{s \not\in J(r)} \; h(s) \; q(s \mathop{\rightarrow} r) \frac{p(r)}{p(s)} \\ & - & \displaystyle \sum_{s \in J(r)} \; h(r) \; q(r \mathop{\rightarrow} s) \frac{p(s)}{p(r)} - \sum_{s \not\in J(r)} \; h(r) \; q(r \mathop{\rightarrow} s) \\ & := & \displaystyle \sum_{s \in \mathcal{X}} K(r,s) \; h(s), \end{array}$$

where

$$K(r,s) = \delta_{s \in J(r)} q(s \to r) + \delta_{s \notin J(r)} q(s \to r) \frac{p(r)}{p(s)} + \delta_{r,s} \times \left[1 - \sum_{t \in J(r)} q(r \to t) \frac{p(t)}{p(r)} - \sum_{t \notin J(r)} q(r \to t) \right].$$

Or in matrix form h = Kh. That is, h is the eigen-vector of K with eigen-value 1.

Since, for $\forall t \in J(r), \ 0 < p(t)/p(r) \leq 1$, we have

$$\begin{split} &1 - \sum_{t \in J(r)} \ q(r \rightarrow t) \, \frac{p(t)}{p(r)} - \sum_{t \notin J(r)} \ q(r \rightarrow t) \\ \geqslant & 1 - \sum_{t \in J(r)} \ q(r \rightarrow t) - \sum_{t \notin J(r)} \ q(r \rightarrow t) \\ = & 1 - 1 = 0; \end{split}$$

also since, for $\forall r, s \in \mathcal{X}$, both $q(s \to r)$ and p(r) are positive, we thus conclude that, for $\forall r, s \in \mathcal{X}$

that is, K is a positive real square matrix. Recall Perron–Frobenius theorem for positive matrices states that "given positive matrix" A, the Perron–Frobenius eigenvector 3 is the only (up to multiplication by constant) non-negative eigenvector for A". As in corollary 3, by letting h = p, we have gained an eigen-value of K such that all components are real and non-negative. So, as a distribution (thus all components have to be real and non-negative), h has no choice but be p, which is what we want to prove, that is $h' = h \Rightarrow h = p$ on \mathcal{X} .

Remark 5. [Burn-in]

Within this intuitive proof, we have to ensure that dropping r_1 from $\{r_1, r_2, ..., r_{N-1}\}$ affects little on the fitting of h. This does affect h if r_1 happens to be the state where $p(r) \ll 1$, so that causes a "Poisson error". After all, r_1 is initialized randomly. For this reason, "burn-in" mechanism is introduced in. While adding r_N will affects little on h' - h, since the probability of being in the "important region" of p for r_N is large, after all, r_N is not initialized randomly as r_1 .

^{2.} I.e. positive real squre matrix.

^{3.} I.e. that unique eigen-vector (up to multiplication by constant) of the Perron-Frobenius eigen-value, which is defined herein.

^{4.} Recall that this corollary requires the Metropolis condition $q(s \to r) = q(r \to s)$ for $\forall r, s \in \mathcal{X}$. This is where this condition is employed in the proof of Metropolis algorithm.