

## Network Security Assignment - ABCs of Digital Certificates

cs23mtech14018 - Yash Shukla

### PART - A

Field Name	Subject (CN) of certificate holder (website)  = nytimes.com	Subject (CN) of certificate holder (intermediate)	Subject (CN) of certificate holder (root)  = DigiCert Global Root CA	Remarks/obse rvations
Issuer	Thawte RSA CA 2018	DigiCert Global Root CA	DigiCert Global Root CA	NA
Version No.	Version 3	Version 3	Version 3	All are using Version 3
Signature Algo	PKCS #1 SHA-256 With RSA Encryption	PKCS #1 SHA-256 With RSA Encryption	PKCS #1 SHA-1 With RSA Encryption	This information reveals the cryptographic algorithm employed to create the unique signature for the certificate by hashing and encrypting its contents.
Size of digest that is signed to generate Cert Sign	256 bits	256 bits	160 bits	Hash Length
Size of Cert Signature	2048 bits	2048 bits	2048 bits	Implies strength and efficiency
Validity period	22/03/2023 to 22/04/2024	06/11/2017 to 06/11/2027	10/11/2006 to 10/11/2031	More validity from top to bottom

Is Subject field (CN), FQDN?	Nytimes.com , yes	CN = DigiCert Global Root CA , no	CN = Thawte RSA CA 2018 , No	FQDN stands for Fully Qualified Domain Name Allows us to see the location
Certificate type: DV, IV, OV or EV? Tell also how you are able to determine the type!	DV	OV	EV	Provides valuable insights into validation and trustworthiness and intended use of certificate
Subject Alternative Name(s) (SAN/UCC), if any	DNS Name: www.homedelivery.nytimes.com DNS Name: *.api.dev.nytimes.com DNS Name: *.api.nytimes.com	NA	NA	Domain Names
Certificate category: Single domain, wildcard, Multi-domain SAN/UCC cert?	Wildcard	Single Domain Certificate	Single Domain Cert	The end user is having many SANs whereas the intermediate and the root has a single domain certificate.
Public Key Info like key algo, key length, public exponent (e) in case of RSA	RSA , 2048 bits, Public Exponent (17 bits): 01 00 01	RSA , 2048 Bits, Public Exponent (17 bits): 01 00 01	RSA ,2048 Bits , Public Exponent (17 bits): 01 00 01	Information about the public key algorithm with its parameters
Public key or modulus (n) in case of RSA	E2 3B E1 11 72 DE A8 A4 D3 A3 57 AA 50 A2 8F 0B 77 90 C9 A2	CA 08 5E E5 53 8A 97 1C 1E 43 2F B6 8A A7 56 E9 8B 84 43 A8	B7 C8 EC BD 69 FE 11 63 D1 CF BF 82 3D 07 26 0E 89 F4 0D AC	The Public key

A5 EE 12 CE 96 5B 01 09 20 CC 01 93 A7 4E 30 B7 53 F7 43 C4 69 00 57 9D E2 8D 22 DD 87 06 40 00 81 09 CE CE 1B 83 BF DF CD 3B 71 46 E2 D6 66 C7 05 B3 76 27 16 8F 7B 9E 1E 95 7D EE B7 48 A3 08 DA D6 AF 7A 0C 39 06 65 7F 4A 5D 1F BC 17 F8 AB BE EE 28 D7 74 7F 7A 78 99 59 85 68 6E 5C 23 32 4B BF 4E C0 E8 5A 6D E3 70 BF 77 10 BF FC 01 F6 85 D9 A8 44 10 58 32 A9 75 18 D5 D1 A2 BE 47 E2 27 6A F4 9A 33 F8 49 08 60 8B D4 5F B4 3A 84 BF A1 AA 4A 4C 7D 3E CF 4F 5F 6C 76 5E A0 4B 37 91 9E DC 22 E6 6D CE 14 1A 8E 6A CB FE CD B3 14 64 17 C7 5B 29 9E 32 BF F2 EE FA D3 0B 42 D4 AB B7 41 32 DA 0C D4 EF F8 81	AC 9D 7A 55 82 7A 14 4B 86 B7 2F 8F 52 9F 1C CA B1 20 5B 6F BA 22 DD A6 9C 2D 78 DA E9 06 08 4E BE 13 A6 EB CB BB 3E B9 05 0C 3E 4A E1 F0 32 1F 13 4E F5 06 C5 47 73 89 3E 80 A3 8B F1 01 24 9B A3 99 66 92 6B 68 AD 0D 2D B4 CD 72 A2 F4 F9 38 5A 65 A6 B4 8C 53 C1 08 1A 84 F8 FD 2E F3 11 75 6E DC 6A 31 29 AC 0D 87 CC 93 60 78 DF 25 BA 26 59 91 C6 83 52 35 A6 CA 9C B8 28 1A CE D7 1C EE 14 BF 76 5C 65 AB 38 1E 79 E9 7C CC 49 23 26 A2 52 50 66 D0 59 61 FF A0 FE 9A 4C 0C 9F F9 E8 8E DE 09 8B B8 15 C1 A4 08 4C 26 9C 7B 06 DB FD 8A 74 5B 58 7E CD 63 A4 91 2F 45 F0 7A 3C 94 0B 8A 7C B2 05 A9 67 93 9F 68 E5	19 3A 1C 20 2F A5 75 61 19 1A 0A B2 07 65 8A A6 62 A8 24 EF 0B 8A AA E7 65 3D F7 50 00 4C 55 56 18 AD EF DC 59 20 9B 7A 57 93 89 C2 77 1C 03 8A 6C A5 BC 01 F3 48 46 0C 7F 4F 3C 00 3A 9C A0 47 4F A2 E2 1B 89 50 95 3D F7 15 81 D2 33 FB 98 3B E7 7D E7 01 6F EA 7E 29 15 EB 25 D6 0E 26 64 08 3F 25 7A AC C5 E4 FF 83 EA E4 B9 C7 54 E9 0F 8A 78 34 9E 01 40 E4 A2 4D 52 FF 55 87 35 00 80 D7 19 E6 61 D3 39 03 B0 B1 5F F0 4B 39 65 BF 6B 51 FE BC AA 25 14 98 87 A3 C7 46 82 6A 6A 18 4C B9 F4 2D 4E 07 B5 E4 C5 D5 EC CE 44 14 3F 81 95 5C E3 77 6C E3 A5 32 C6 A8 34 D3 C5 E7 3A A0 FE 9B F7 63 5C E1 07 36 6F 52 29 3F
--	---	--

	D5 BB 8D 58 3F B5 1B E8 49 28 A2 70 DA 31 04 DD F7 B2 16 F2 4C 0A 4E 07 A8 ED 4A 3D 5E B5 7F A3 90 C3 AF 27	95 63 60 D8 58 95 5F E0 55 EF 93 A7 11 3B 7C E6 92 D8 66 44 E0 AB BD A7 8F CD A4 85 78 41 24 54 E7 D8 03	8E 59 9E 0B E7 E6 71 51 67 A4 C3 8E CD 40 73 ED 88 D2 35 84 94 90 7A E7 34 72 89 2C 46 2A 5C ED 5D 23	
Key usages; how do they vary in the chain, mention in the remarks?	Signing Key Encipherment	Signing Certificate Signer CRL Signer	Signing Certificate Signer CRL Signer	Key difference is just Certificate Signer and Signer.
Basic constraints, how do they vary in the chain?	Not Critical Is not a Certification Authority	Subject Type=CA Path Length Constraint=0	Subject Type=CA Path Length Constraint=None	CA: No for end user cannot provide certificate for others
Size of the certificate	2.84 KB	1.62 KB	1.32 KB	NA
URI of CRL	URI: <a href="http://cdp.thawte.com/ThawteRSACA2018.crl">http://cdp.thawte.com/ThawteRSACA2018.crl</a>	<a href="http://crl3.digicert.com/DigiCertGlobalRootCA.crl">http://crl3.digicert.com/DigiCertGlobalRootCA.crl</a>	NA	Location where the Certificate Revocation List (CRL) associated with the certificate can be found
URI of OCSP Responder	URI: <a href="http://status.thawte.com">http://status.thawte.com</a>	URI: <a href="http://ocsp.digicert.com">http://ocsp.digicert.com</a>	NA	Location where the OCSP responder service associated

				<b>with the certificate</b>
<b>Any other parameters that you found interesting?</b>	<b>NA</b>	<b>NA</b>	<b>NA</b>	<b>NA</b>

Answer the following queries after filling out the above table:

**Q1.** Which certificate type (DV/OV/IV/EV) is more trustable, secure, and expensive?

**Answer 1 ) :** EV certificates are the most trusted and have high security but they are expensive.

**Q2.** What is the role of the Subject Alternative Name (SAN) field in X.509 certificates?

**Answer 2 ) :** they just provide more flexibility and security and compatibility and secure the diverse environments efficiently while considering limitations.

**Q3.** Why are key usages and basic constraints different for root, intermediate and end certificates? What could go wrong if all of them have the same values?

**Answer 3 ) :** The root, intermediate, and end certificates have distinct key usages and essential limits based on the hierarchical structure and security assurance. Intermediates are limited in depth, root certificates cannot sign other certificates that further certify each other, and allowing the same values could compromise the PKI's trust model and lead to misuse, interoperability issues, hierarchy collapse, and security vulnerabilities.

**Q4.** What is the difference between Signature value and Thumbprint aka Fingerprint of a digital certificate?

**Answer 3 ) :** The CA creates signature value, a cryptographic value, certify the contents of the certificate and guarantee its validity and integrity. confirms the authenticity of the certificate.

On the other hand, the fingerprint or thumbprint is a condensed version of the certificate that is produced through the use of a hash function. It functions as a distinct identification that can be compared and verified. The thumbprint helps with identification and comparison.

**Q5.** Why do RSA key lengths increase over the years? Why is ECDSA being preferred over RSA now-a-days?

**Answer 5 ) :** RSA Key lengths increase over the years because To keep security against advances in computing power and cryptographic assaults, and ECDSA is preferred Because it requires shorter key lengths for equal security levels than RSA, it is preferable because it facilitates faster cryptographic operations and lowers computing cost and ECDSA is a more effective option for contemporary cryptographic applications since it also performs better in limited settings.

**Q6.** What are pros and cons of pre-loading root and intermediate certificates in the root stores of browsers and OSes?

**Answer 6 ) :** pro is pre-loading root and intermediate certificates in root stores improves security. This expedites the SSL/TLS handshake procedure and avoids security alerts. Cons is the root store larger and more complicated, which could make maintenance more difficult and update more slowly. Furthermore, it centralises judgements about trust, which might be dangerous if a compromised certificate is on the pre-loaded list.

**Q7.** Why are root CAs kept offline? How do they issue certificates to intermediate CAs?

**Answer 7 ) :** To reduce the possibility of compromise, root CAs are maintained offline Because to reduce the possibility of compromise.and they are issued through safe offline methods like hardware security modules (HSMs) or secure signing servers, they provide certificates to intermediate CAs. This lessens the possibility of unauthorised certificate issuing by guaranteeing the extreme security of the root CAs' private keys.

**Q8.** Why are root and intermediate certificates of new CAs cross-signed by the legacy CAs? Answer this by taking Let's Encrypt and its parent organization as root and intermediate CAs

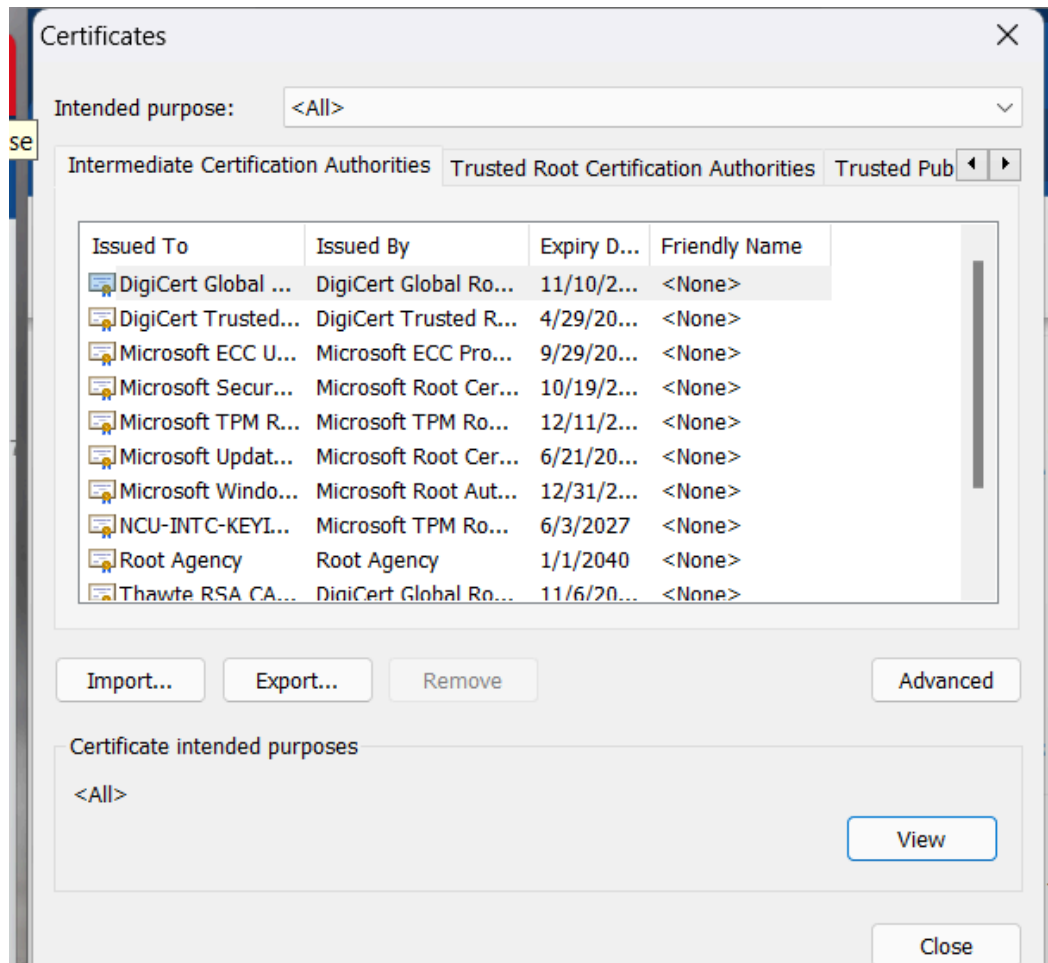
**Answer 8 ) :** In order to quickly build confidence with clients and systems that may not yet be aware of the new CA, traditional CAs cross-sign the root and intermediate certificates of new CAs, such as Let's Encrypt. The Internet Security Research Group (ISRG), Let's Encrypt's parent organisation, cross-signs its intermediate certificates, and the IdenTrust root CA does the same.

**Q9.** What challenge is posed to the certificate seekers (Alice) by Let's Encrypt CA before issuing wildcard certificates? How does Alice respond to it so that she passes it?

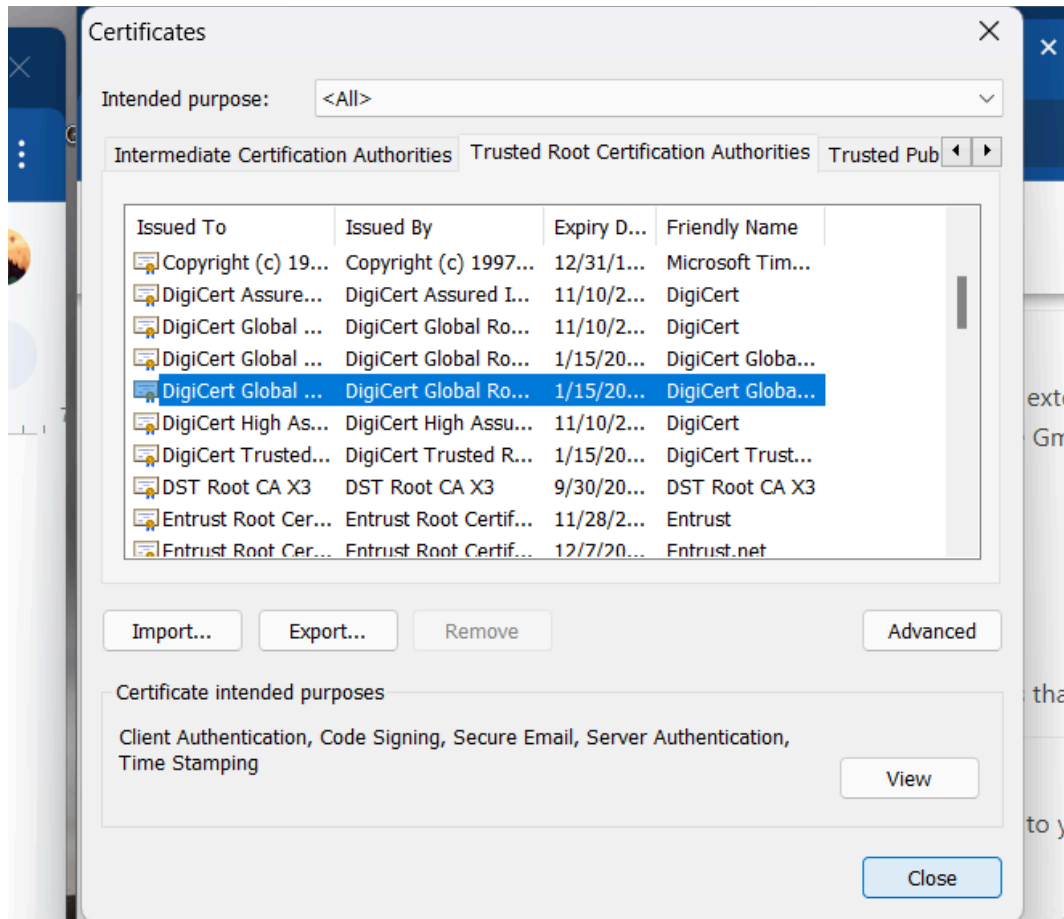
**Answer 9 ) :** To Answer this question before providing wildcard certificates, Let's Encrypt challenges certificate seekers like Alice to validate their domains. By completing task, such as adding a TXT record to the domain's DNS or putting a file with a specific content on her web server, and Alice must prove that she is the owner of the domain. so in order to successfully complete the challenge, Alice must demonstrate her control over the domain by responding to the validation request by doing the designated action within the allotted time.

**Q10.** List out names of OS/Browser/Company whose root stores were pre-populated with Root and Intermediate CA certificates of the website #N?

**Answer 10 ) :**







## PART B

**Q1** . A browser X has received the digital certificate of the website #N over a TLS connection. How does it verify whether the certificate is valid? Write a psuedo-code of browser X's verifier function named myCertVerifier( ) and explain how it works by picking the entire chain of trust of an end-user cert (of the website #N) in PART-A of this assignment.

**Answer 1 )** : A Browser has Received the Digital Certificate of the Website [nytimes.com](https://www.nytimes.com) over a TLS Connection.

\* To Verify whether the Certificate is valid or not : this are the potential reasons to declare invalid :

- Expired Certificate
- Invalid Signature
- Untrusted Certificate Authority (CA)
- Certificate Chain Issues (Incomplete Chain,Invalid Chain,Self-Signed Certificate)

- Domain Name Mismatch
- Revoked Certificate
- Critical Extensions Handling
- Fingerprint or Public Key Mismatch
- Revocation Status Check Failure

```

** Examination flow : Start -> myCertVerifier()
                        Signature Verification ->
                        Check Certificate Expiration ->
                        Verify Issuer Trustworthiness ->
                        Extract Certificate Chain ->
                        Chain Verification ->
                        Additional Checks ->
                        Overall Decision (Valid/Invalid)

```

Pseudo Code mentioned here as respective to the flow.

```
{.....
```

```
function myCertVerifier(cert):
```

```
    if not check_certificate_signature(cert):
```

```
        return False
```

```
    if not check_certificate_expiration(cert):
```

```
        return False
```

```
    if not check_trusted_CA(cert):
```

```
        return False
```

```
    if not verify_certificate_chain(cert):
```

```
        return False
```

```
    if not additional_checks(cert):
```

```
        return False
```

```
    return True
```

```
function check_certificate_signature(cert):
```

```
    # Verify the cryptographic signature of the certificate
```

```
    # using the public key of the issuer
```

```
    return signature_verified
```

```
function check_certificate_expiration(cert):
```

```
    # Check if the certificate has not expired
```

```
    return not expired
```

```
function check_trusted_CA(cert):
```

```

    # Check if the issuer (CA) of the certificate is trusted
    return trusted_CA

function verify_certificate_chain(cert):
    # Verify the entire certificate chain
    chain = extract_certificate_chain(cert)
    for cert in chain:
        if not check_certificate_signature(cert):
            return False
        if not check_certificate_expiration(cert):
            return False
        if not check_trusted_CA(cert):
            return False
    # Optionally, perform additional checks here
    return True

function extract_certificate_chain(cert):
    # Function to extract the certificate chain from the server certificate
    # This may involve retrieving additional certificates from the server
    return certificate_chain

function additional_checks(cert):
    # Perform additional checks such as domain name verification,
    # revocation status check, extension critical flag check, etc.
    return additional_checks_passed

.....}

```

**Q2.)** Consider the scenario in which evil Trudy has used the domain validated (DV) digital certificate of the website (Bob) named Bob.com to launch her own web server with the domain name, xyz.com. Does your function myCertVerify( ) returns valid or invalid for this when someone like Alice (browser) tries to access Trudy's website xyz.com?

**Answer 2 ) :** DV certificates are primarily used to verify control of domain, rather than the identity behind the domain, they are susceptible to misuse in scenarios where a certificate issued for one domain is used for another domain, this will return invalid.

**Q3. )** Consider another scenario in which evil Trudy has used the digital certificate of Bob's website Bob.com to launch her own web server with the domain name, xyz.com.

When a web client (Alice) tries to connect with Bob's website by sending a DNS query, Trudy responds with her IP address by DNS cache poisoning ([What is DNS cache poisoning? | DNS spoofing | Cloudflare](#)) Does your function myCertVerifier( ) returns valid or invalid for this and what are the consequences? What kind of attacks can Trudy launch in this scenario?

**Answer 3 ) :** My Function checks during the verification of the entire "Chain of Trust," the function will detect that the certificate presented by Trudy does not belong to the domain for which the DNS query was made (Bob.com), but instead, it's being used for xyz.com. Trudy can launch Phishing ,Man in the middle attacks ,Data Theft ,Session Hijacking.

### **About : 7-zip**

7-Zip Encryption:

- 7-Zip uses passwords to encrypt files
- Passwords are converted into encryption keys using a SHA.
- The keys are used with symmetric encryption algorithm (like AES) to encrypt the file data.

Role of Password Length in Brute Force Attacks:

- Longer passwords make brute force attacks harder.
- Each additional character exponentially increases the search space for possible passwords.
- Strong, lengthy passwords greatly improve the security of encrypted files.

### **References:**

1. <https://crt.sh/>
2. <https://ahrefs.com/blog/most-visited-websites/>
3. <http://lapo.it/asn1js/#>
4. <http://phpseclib.sourceforge.net/x509/decoder.php>
5. <https://www.ssl.com/article/dv-ov-and-ev-certificates/>
6. <https://www.ccadb.org/>
7. [DV, OV, IV, and EV Certificates - SSL.com](#)
8. [7-Zip \(7-zip.org\)](#)

## **PLAGIARISM STATEMENT**

*I certify that this assignment/report is my own work, based on my personal study and/or research and that I have acknowledged all material and sources used in its preparation, whether they be books, articles, reports, lecture notes, and any other kind of document, electronic or personal communication. I also certify that this assignment/report has not previously been submitted for assessment in any other course, except where specific permission has been granted from all course instructors involved, or at any other time in this course, and that I have not copied in part or whole or otherwise plagiarized the work of other students and/or persons. I pledge to uphold the principles of honesty and responsibility at CSE@IITH. In addition, I understand my responsibility to report honor violations by other students if I become aware of it.*

Name: Yash Shukla

Date: 06 / 02 / 2024

Signature: Yash Shukla