

PROGRAMMING AND DATA STRUCTURES

BINARY TREES

HOURLIA OUDGHIRI

FALL 2021

OUTLINE

- ◆ Characteristics of binary trees
- ◆ Binary Search Trees (BST)
- ◆ Operations on BSTs
- ◆ Implementation of the BST class
- ◆ Heap
- ◆ Operations on the Heap
- ◆ Implementation of the Heap class

STUDENT LEARNING OUTCOMES

At the end of this chapter, you should be able to:

- ▶ Describe the properties of binary trees/BST/Heap
- ▶ Trace operations on binary trees/BST/Heap
- ▶ Implement BST/Heap generic data structures
- ▶ Use the BST/Heap data structures
- ▶ Evaluate the complexity of the operations on BST/Heap

What is a binary tree?

- ◆ Data organized in a binary tree structure
- ◆ Easy and efficient access and update in large collections of data
- ◆ Used for efficient search operations
- ◆ Wide range of applications: mathematical expressions, game strategies, decision trees, data compression, ...

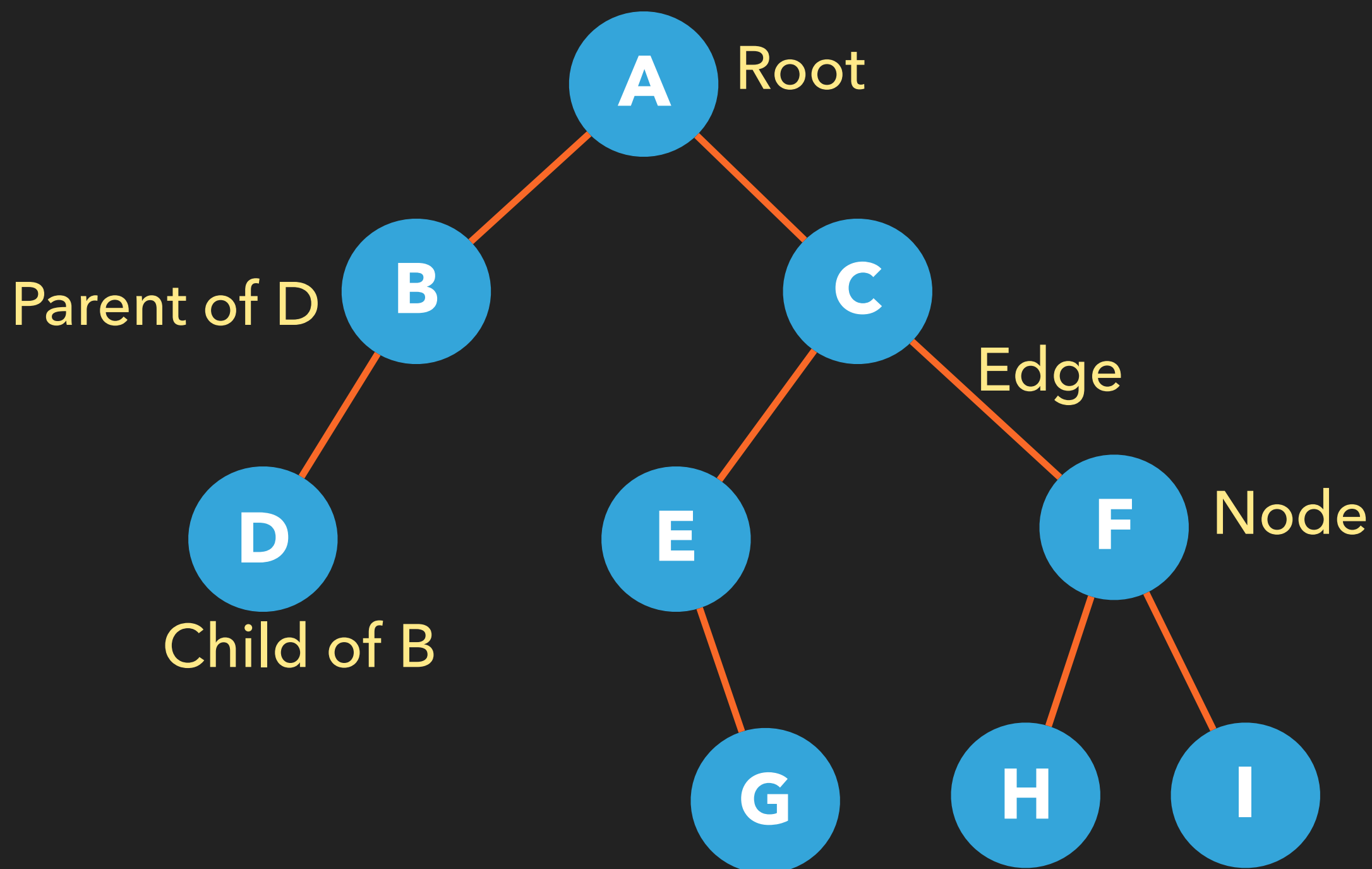
Properties

- ◆ Set of elements called **nodes** (vertices) interconnected with **edges** (arcs)
- ◆ The first node is called the **root**
- ◆ The root is connected to two binary trees (**left subtree** and **right subtree**)

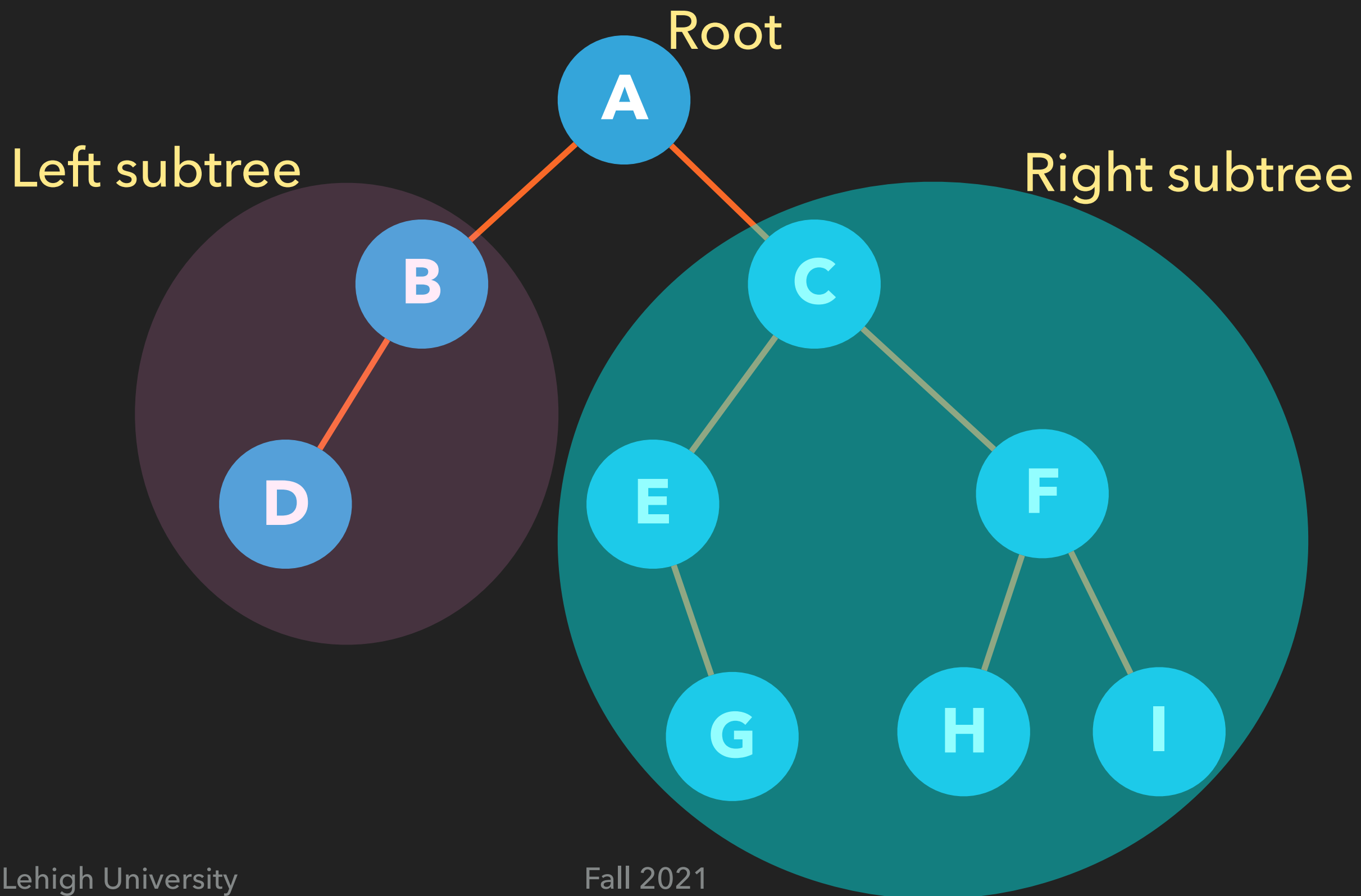
Properties

- ◆ Every node has a parent (except the root) and may have one child or two children at most
- ◆ The root is the ancestor of all the nodes in the tree

Properties



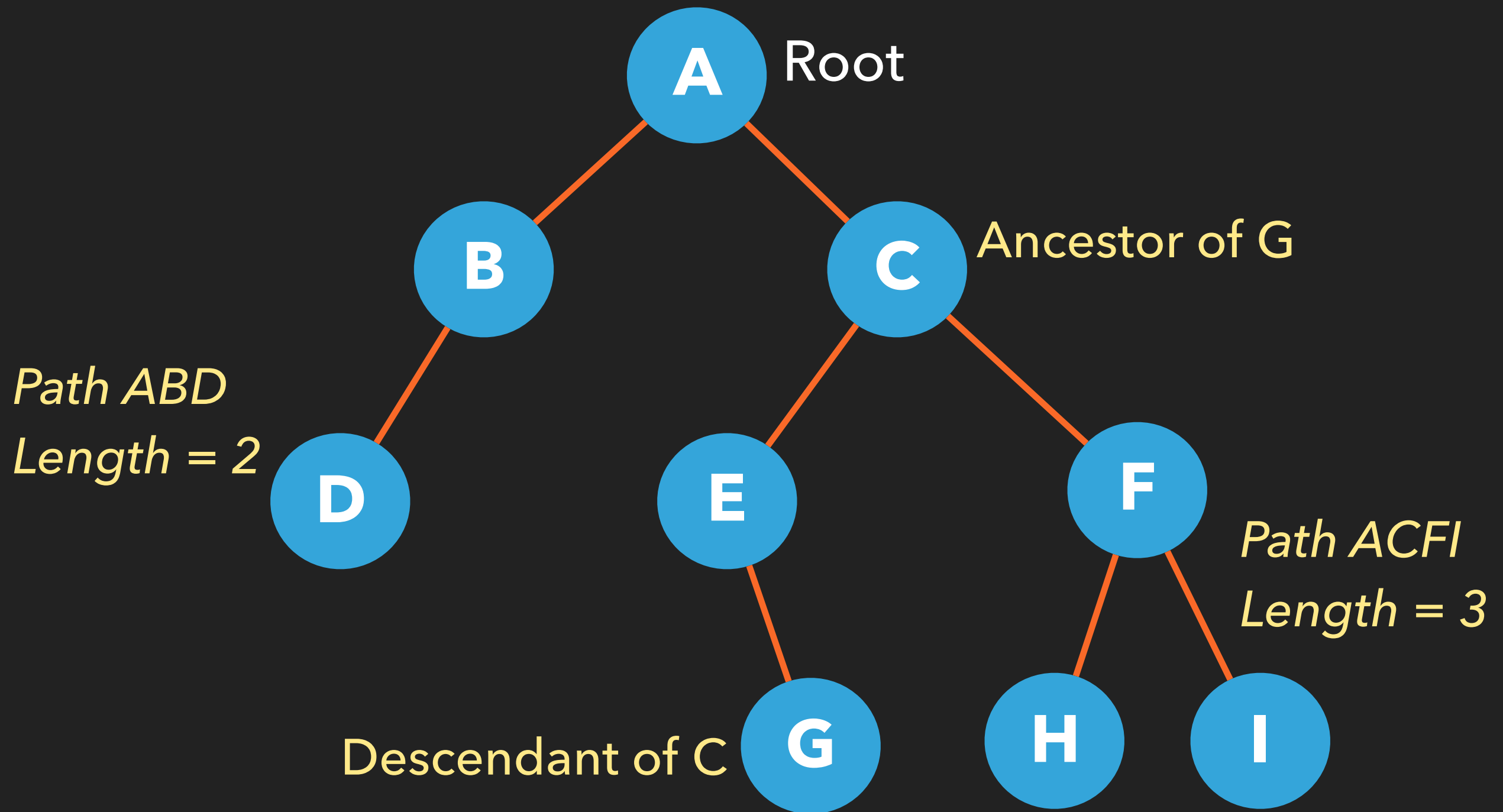
Properties



Properties

- ◆ **Path**: Sequence of connected nodes starting at any level of the tree
- ◆ **Length of a path**: the number of nodes in the sequence - 1 (number of edges)
- ◆ If there is a path from node P to node Q , Q is the **descendant** of P and P is an **ancestor** of Q

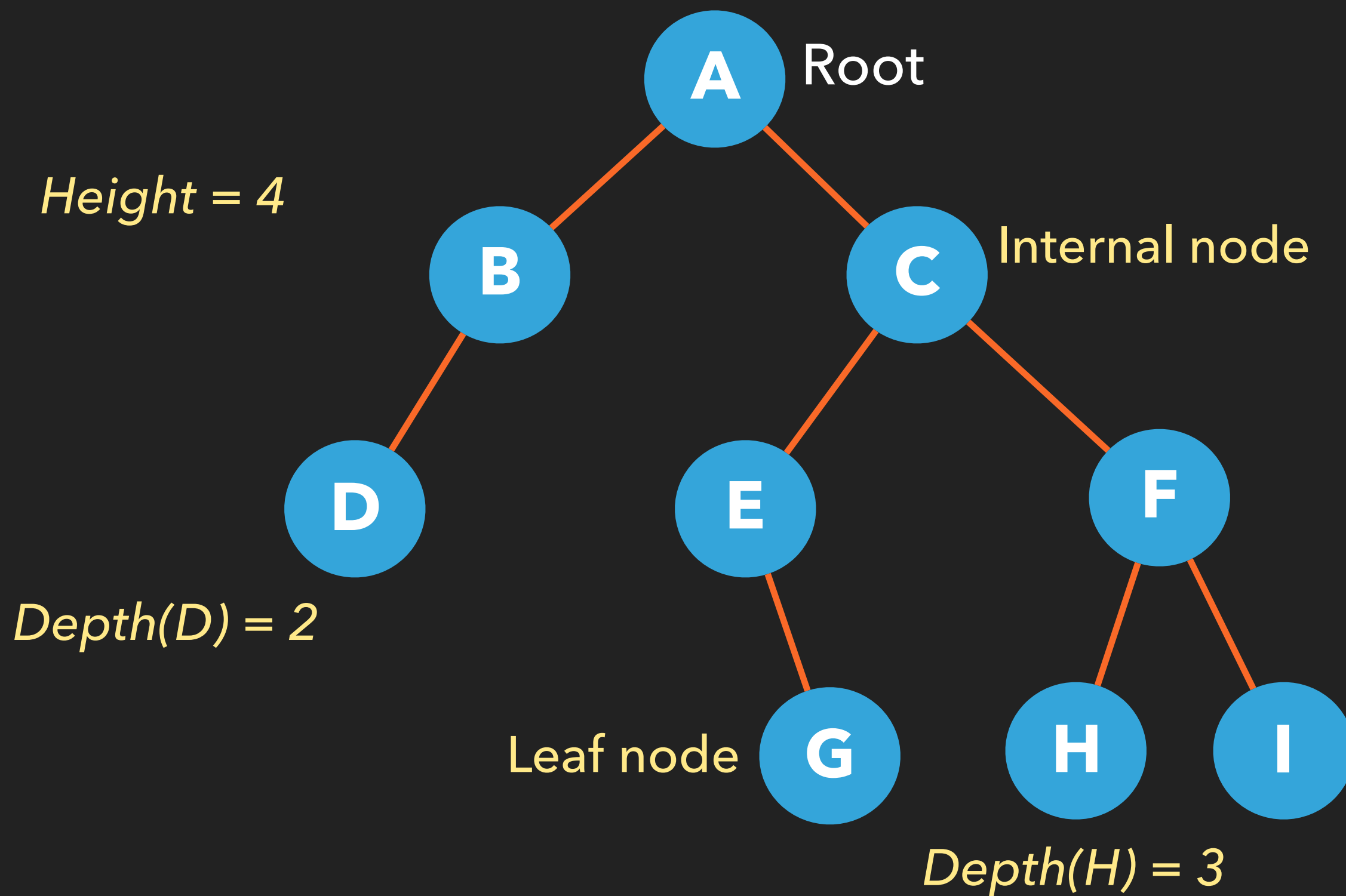
Properties



Properties

- ◆ **Depth of a node:** Length of the path from the root to the node
- ◆ **Height of a tree:** the depth of the deepest node + 1
- ◆ **Leaf node:** node that has no children
- ◆ **Internal node:** node that has at least one child

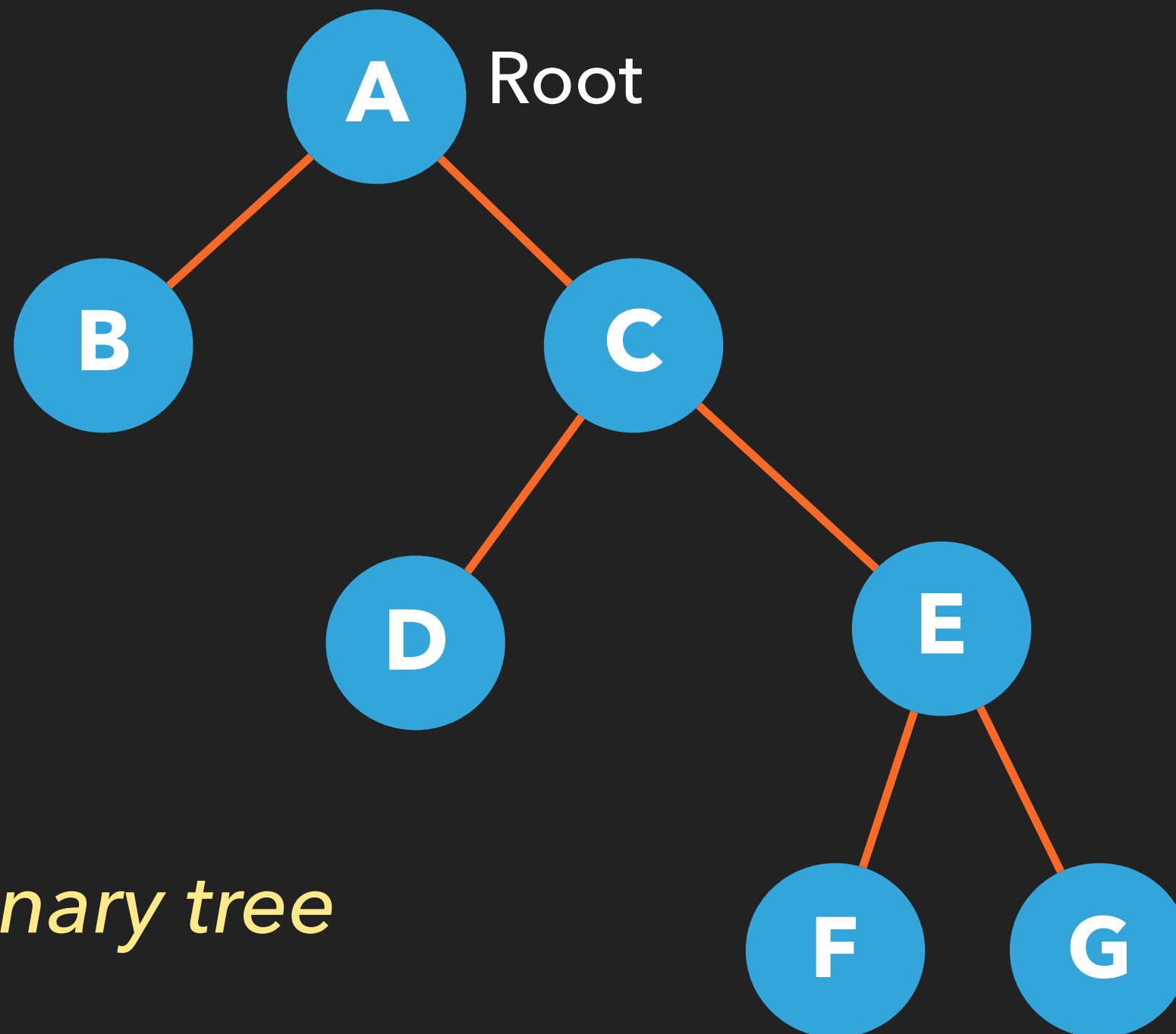
Properties



Properties

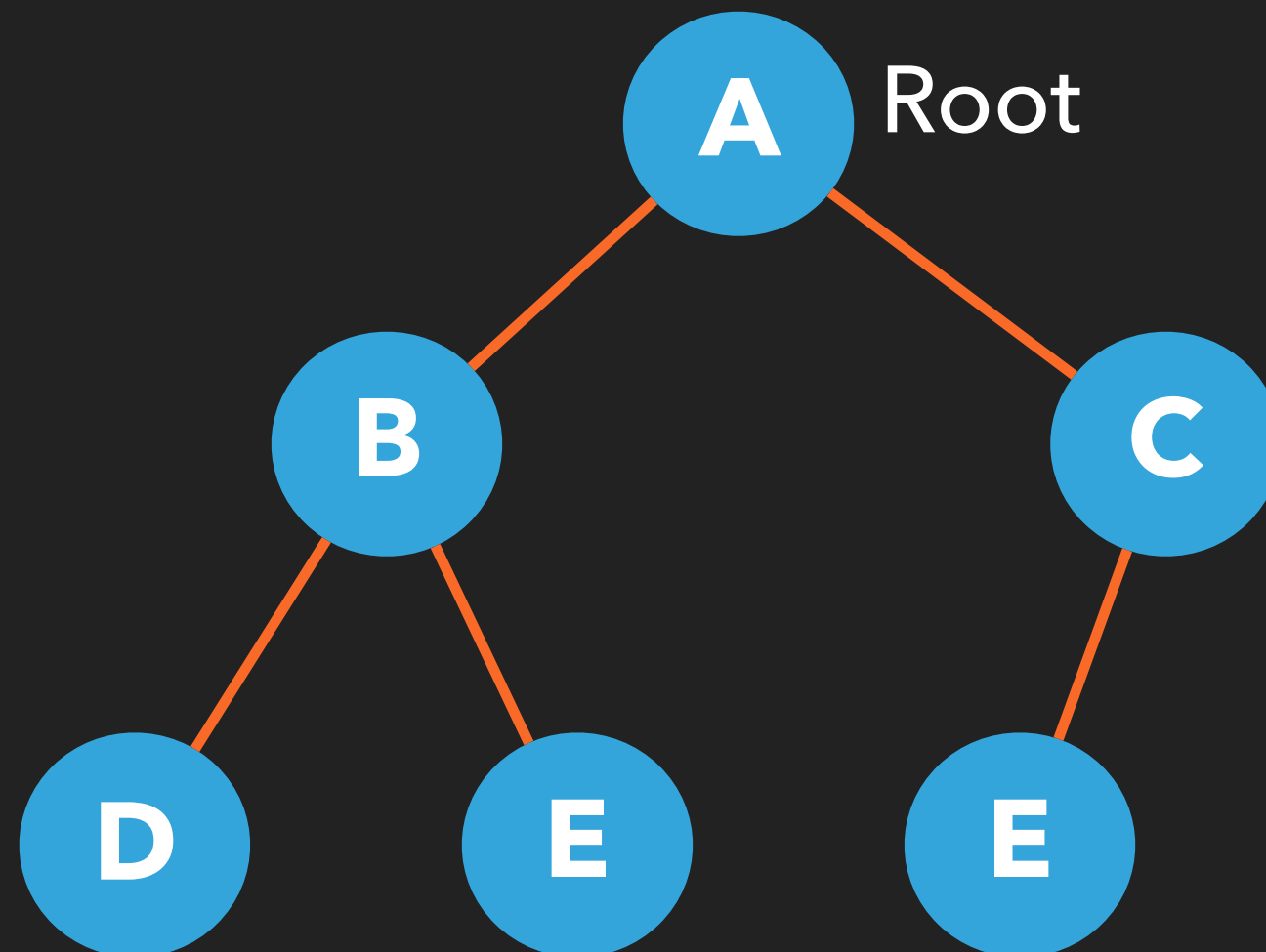
- ◆ **Full Binary Tree:** each node is a leaf or an internal node with exactly two children
- ◆ **Complete Binary Tree:** Every level is filled except the last level, and the leaves on the last level are placed leftmost

Properties



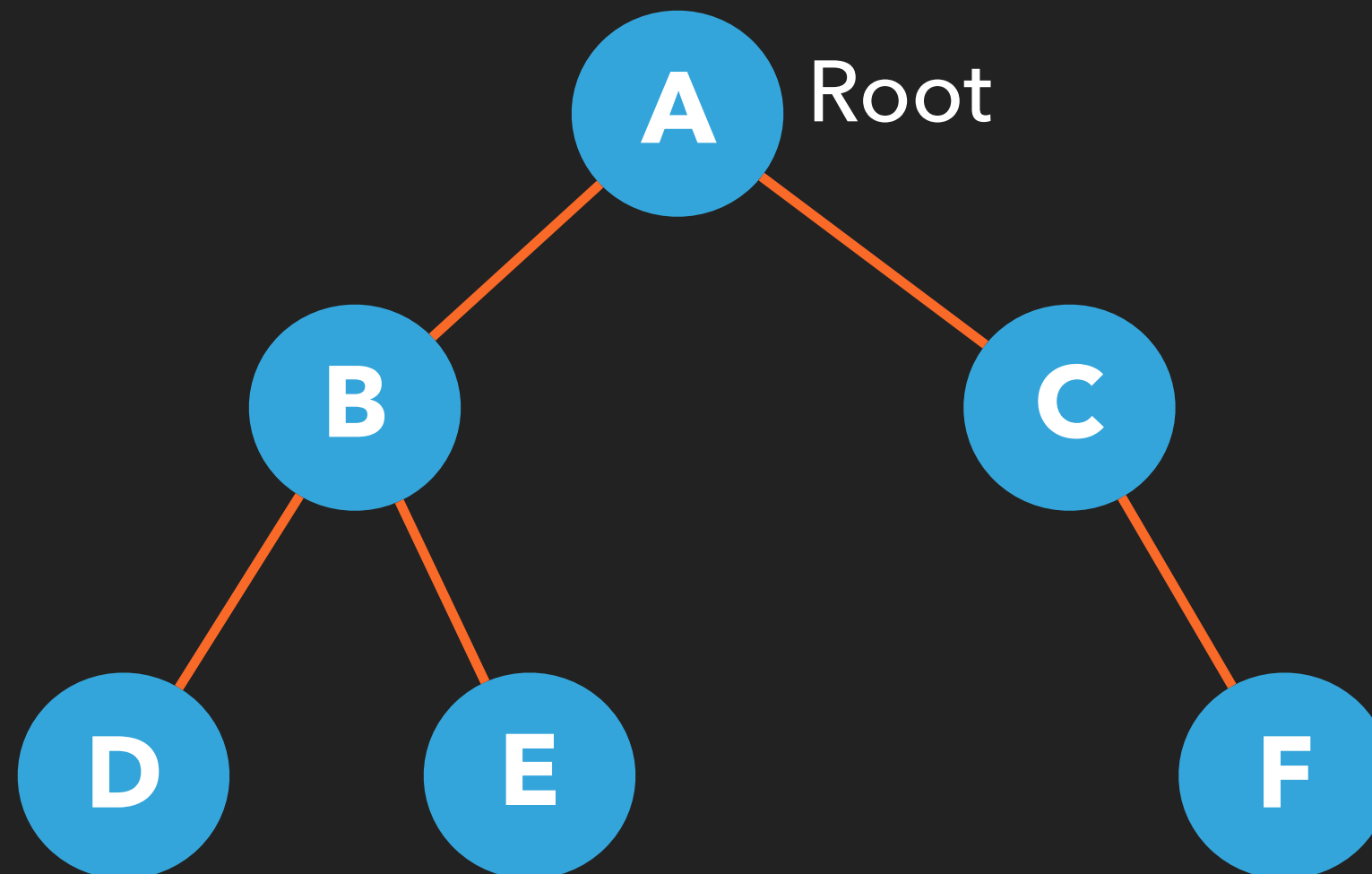
Full binary tree

Properties



Complete binary tree but not full

Properties

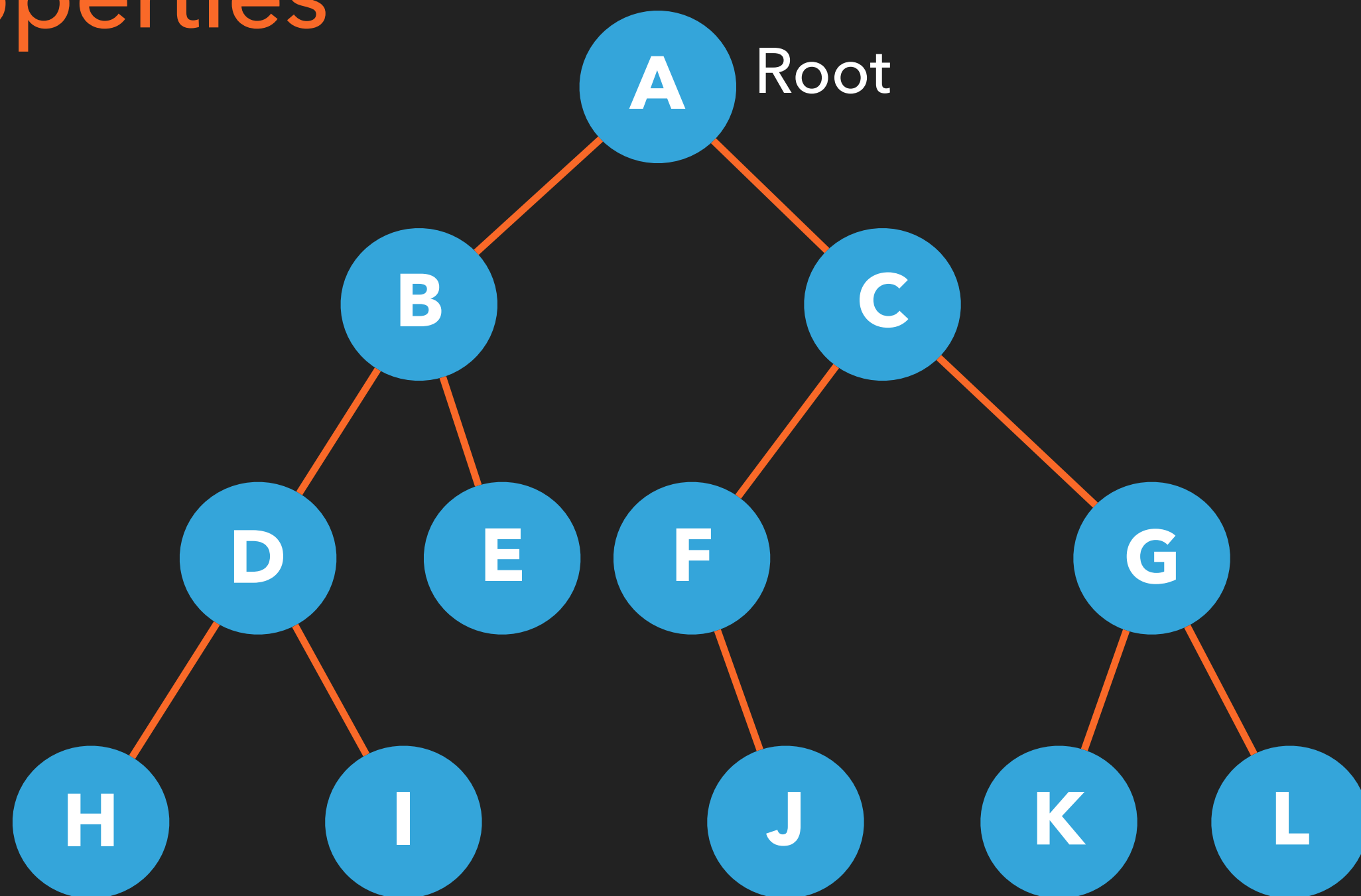


Not Complete - not full

Properties

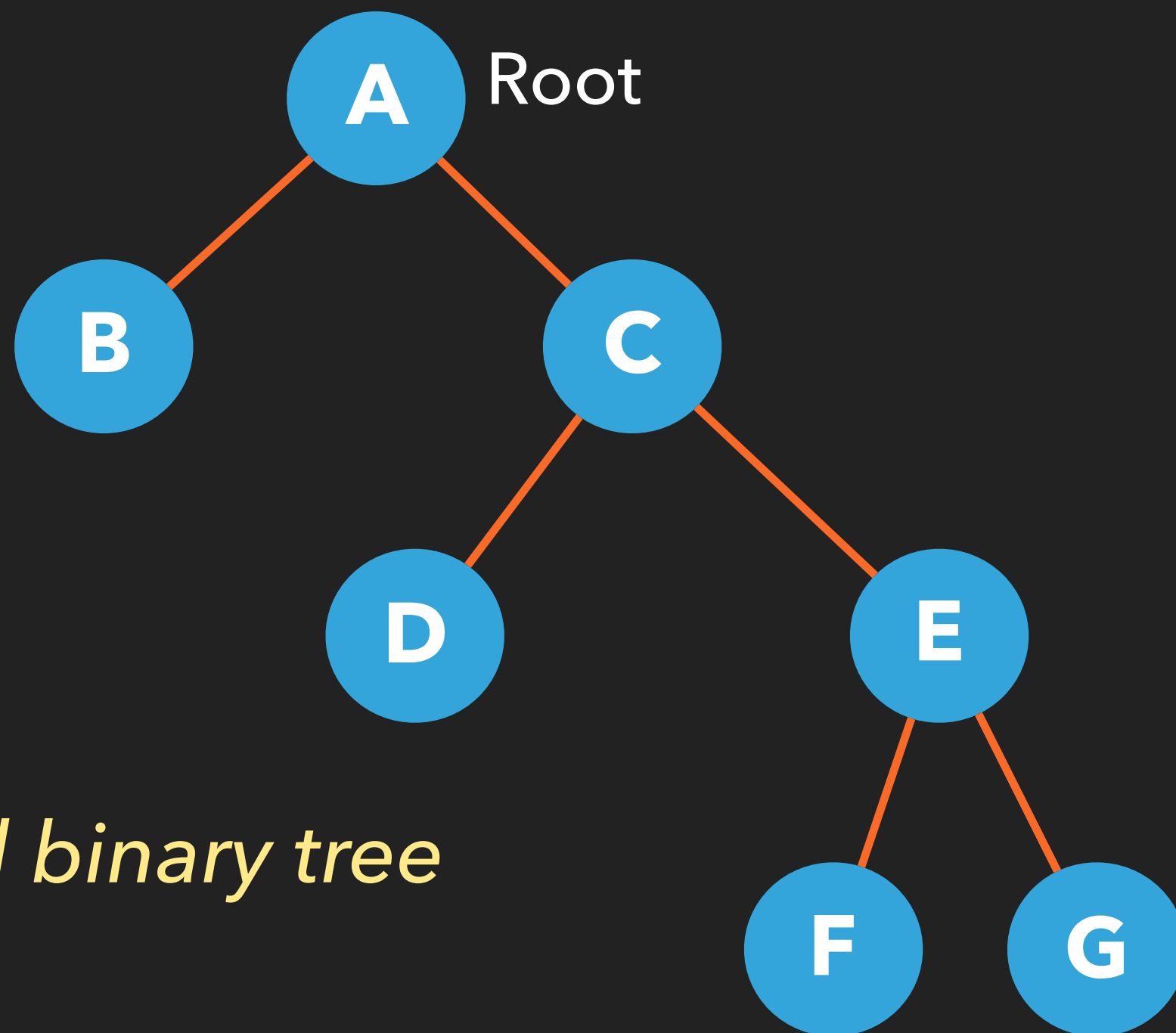
- ◆ **Balanced Binary Tree:** for each node, the height of the left subtree and the height of the right subtree differ by at most 1

Properties



Balanced binary tree

Properties



Unbalanced binary tree

Practice

Root?

Depth(35)?

Path from 70 to 72?

Height of the tree?

Leaf nodes?

Internal nodes?

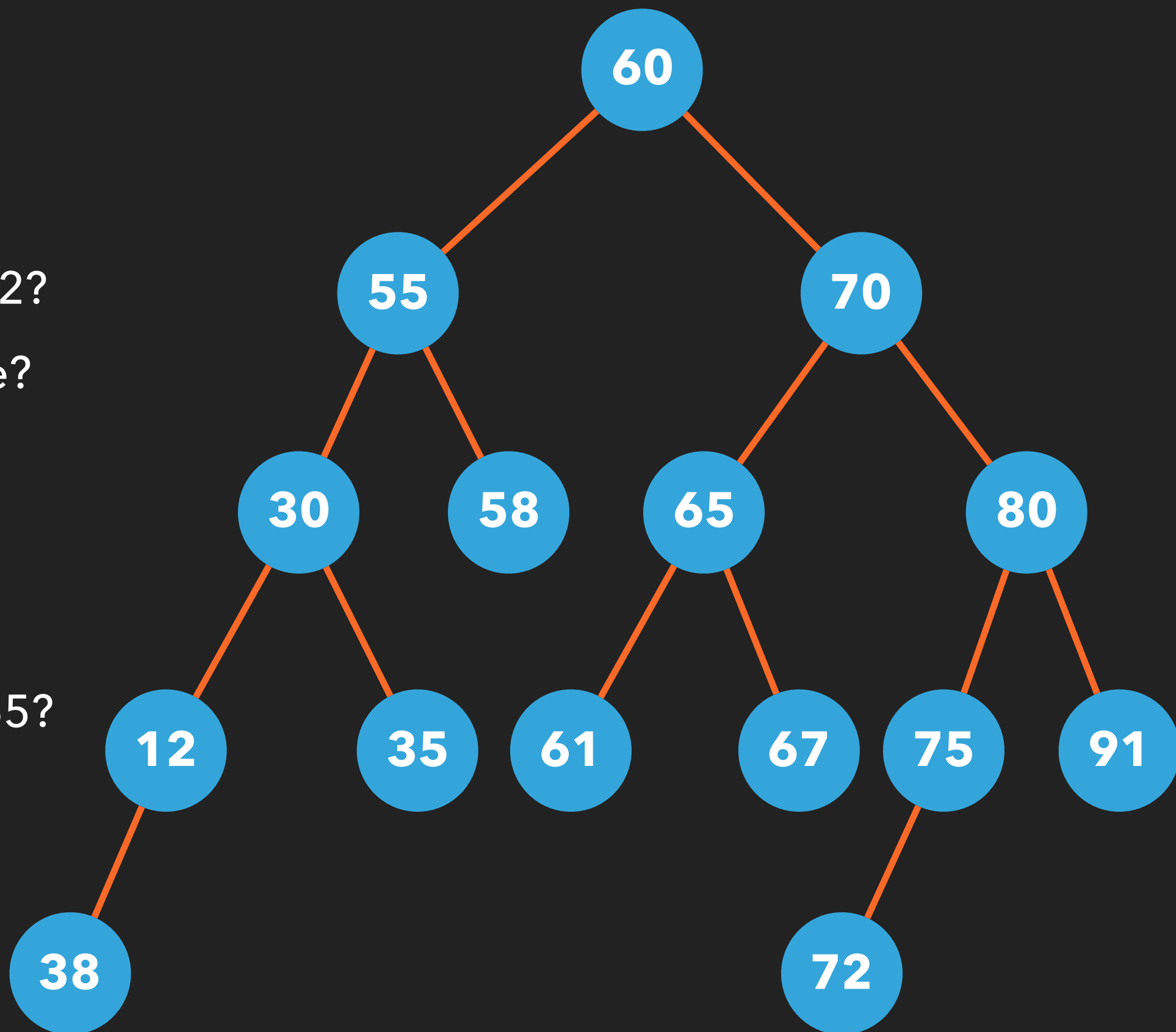
Children of 55?

Descendants of 55?

Full?

Complete?

Balanced?



Binary Tree Traversal

- ◆ Any process of visiting all the nodes in the tree is called **traversal**
- ◆ Three common traversals
 - ◆ **Preorder**
 - ◆ **Inorder**
 - ◆ **Postorder**

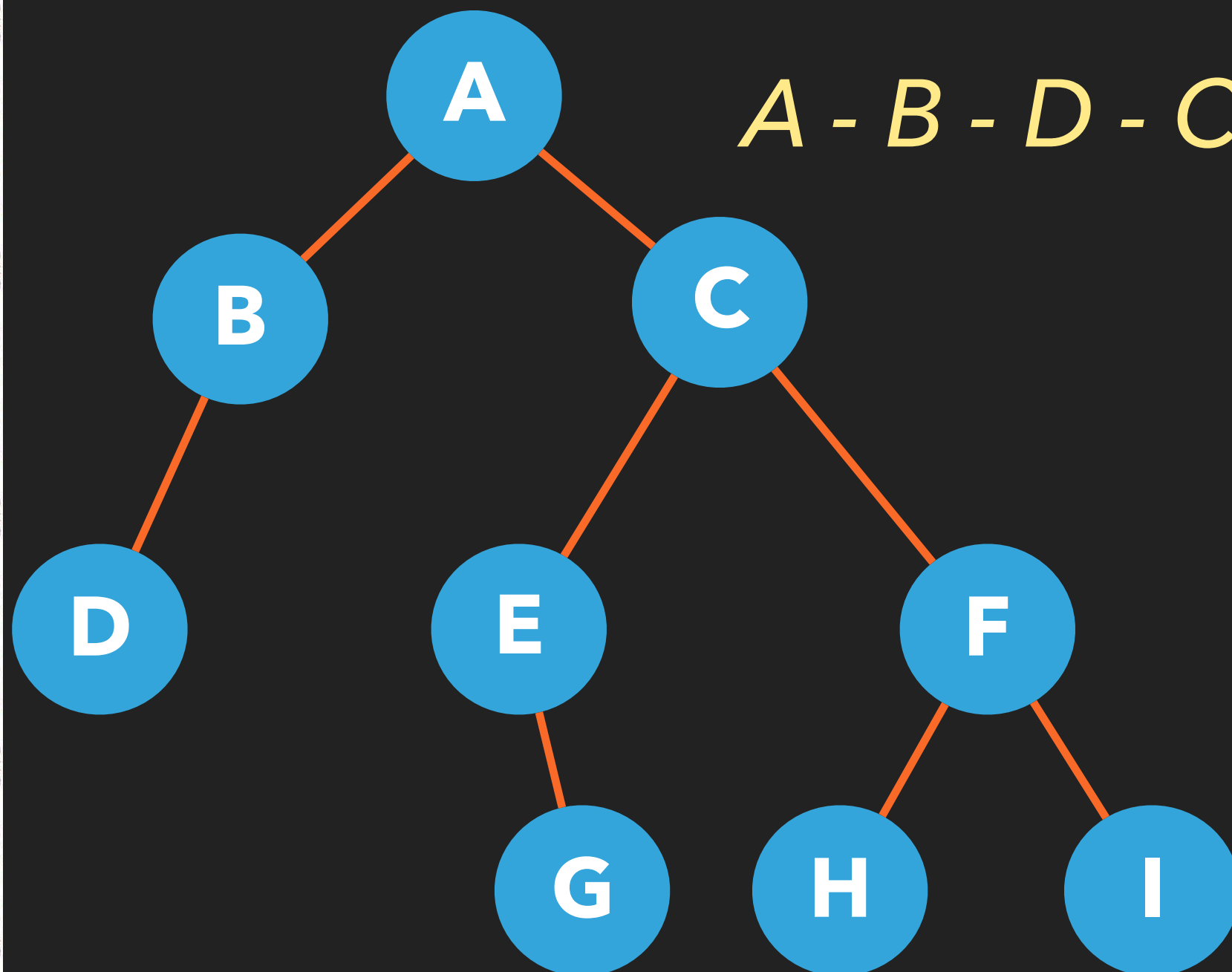
Preorder Traversal

- ◆ Any node is visited before its children
- ◆ **V-L-R**: Visit Node - Go Left - Go Right
 - ◆ Visit the node
 - ◆ Traverse the left subtree
 - ◆ Traverse the right subtree

Preorder Traversal

♦ V-L-R: Visit Node - Go Left - Go Right

A - B - D - C - E - G - F - H - I



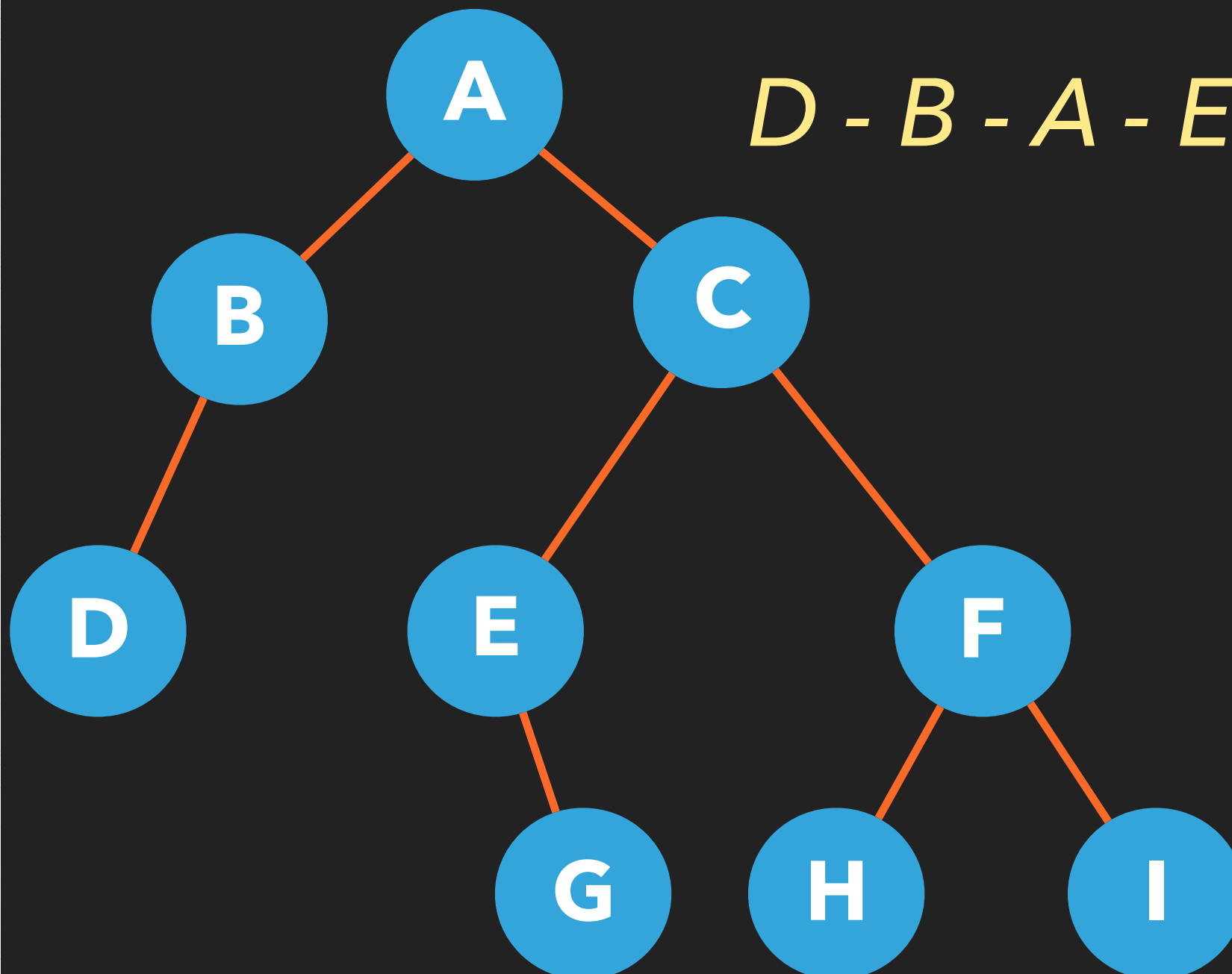
Inorder Traversal

- ◆ Any node is visited after its left subtree and before its right subtree
- ◆ **L-V-R**: Go Left - Visit Node - Go Right
 - ◆ Traverse the left subtree
 - ◆ Visit the node
 - ◆ Traverse the right subtree

In order Traversal

♦ L-V-R: Go Left - Visit Node - Go Right

D - B - A - E - G - C - H - F - I



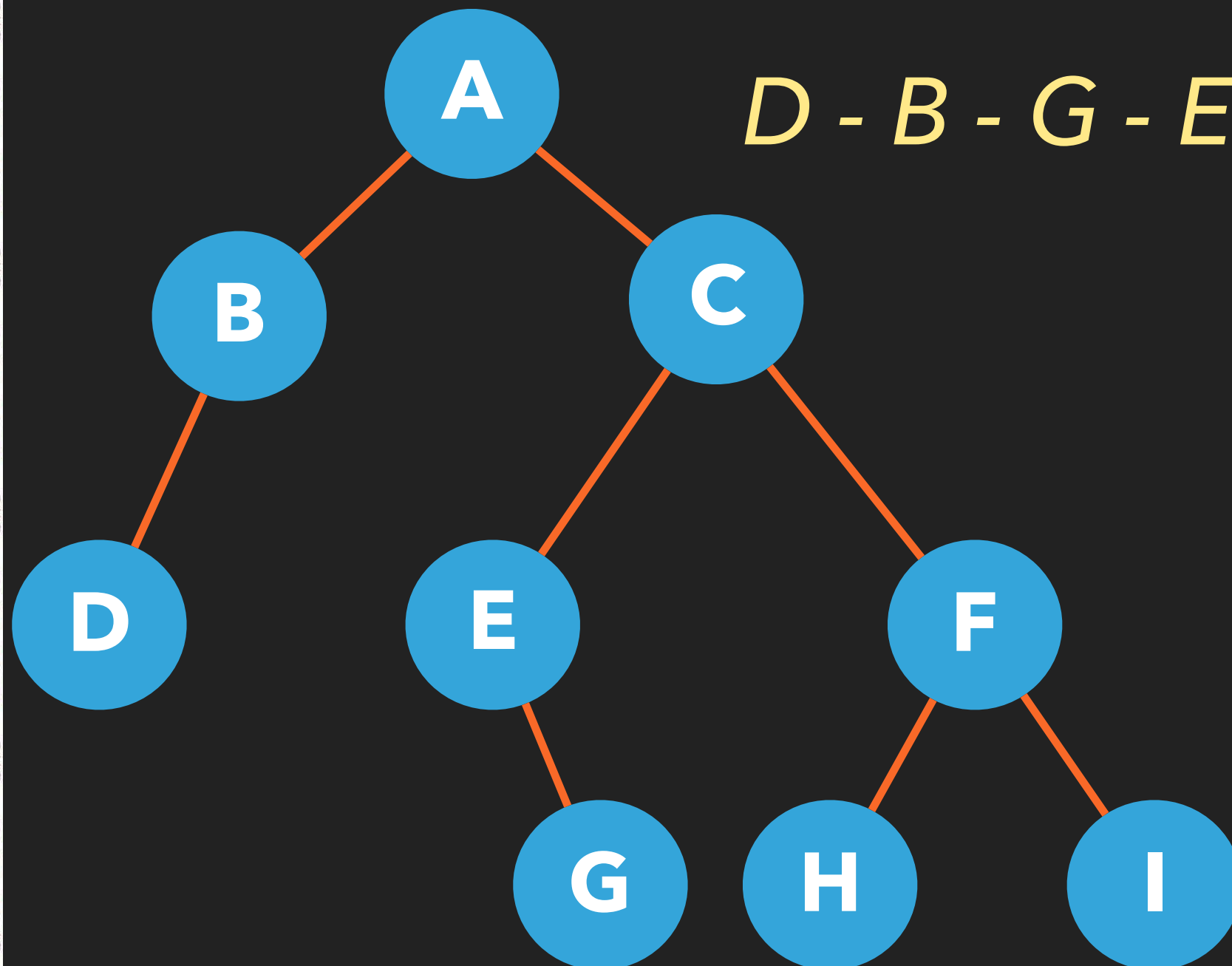
Postorder Traversal

- ◆ Any node is visited after its left subtree and right subtree
- ◆ **L-R-V**: Go Left - Go Right - Visit Node
 - ◆ Traverse the left subtree
 - ◆ Traverse the right subtree
 - ◆ Visit the node

Postorder Traversal

♦ **L-R-V**: Go Left - Go Right - Visit Node

D - B - G - E - H - I - F - C - A

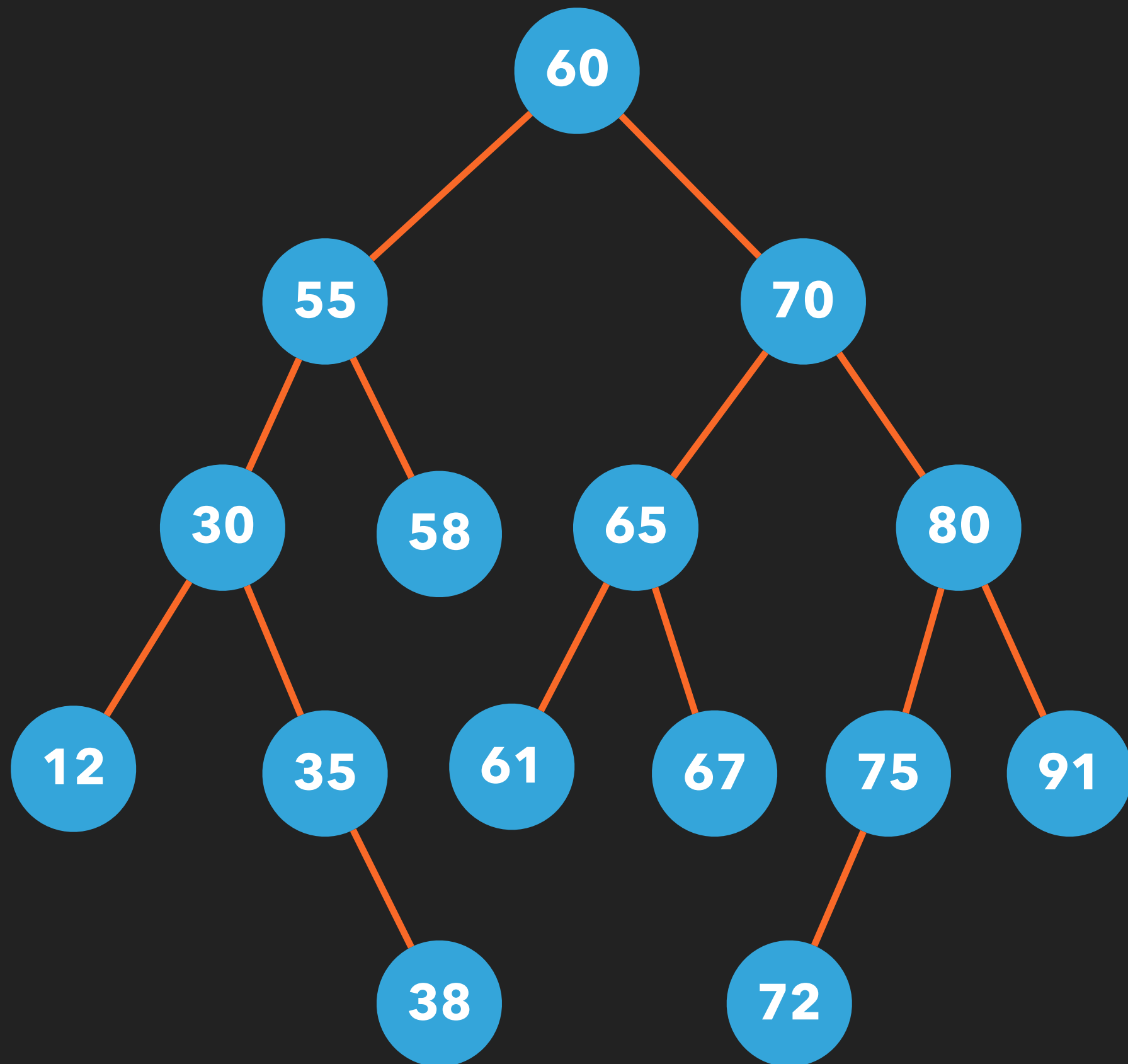


Practice

Preorder:

Inorder:

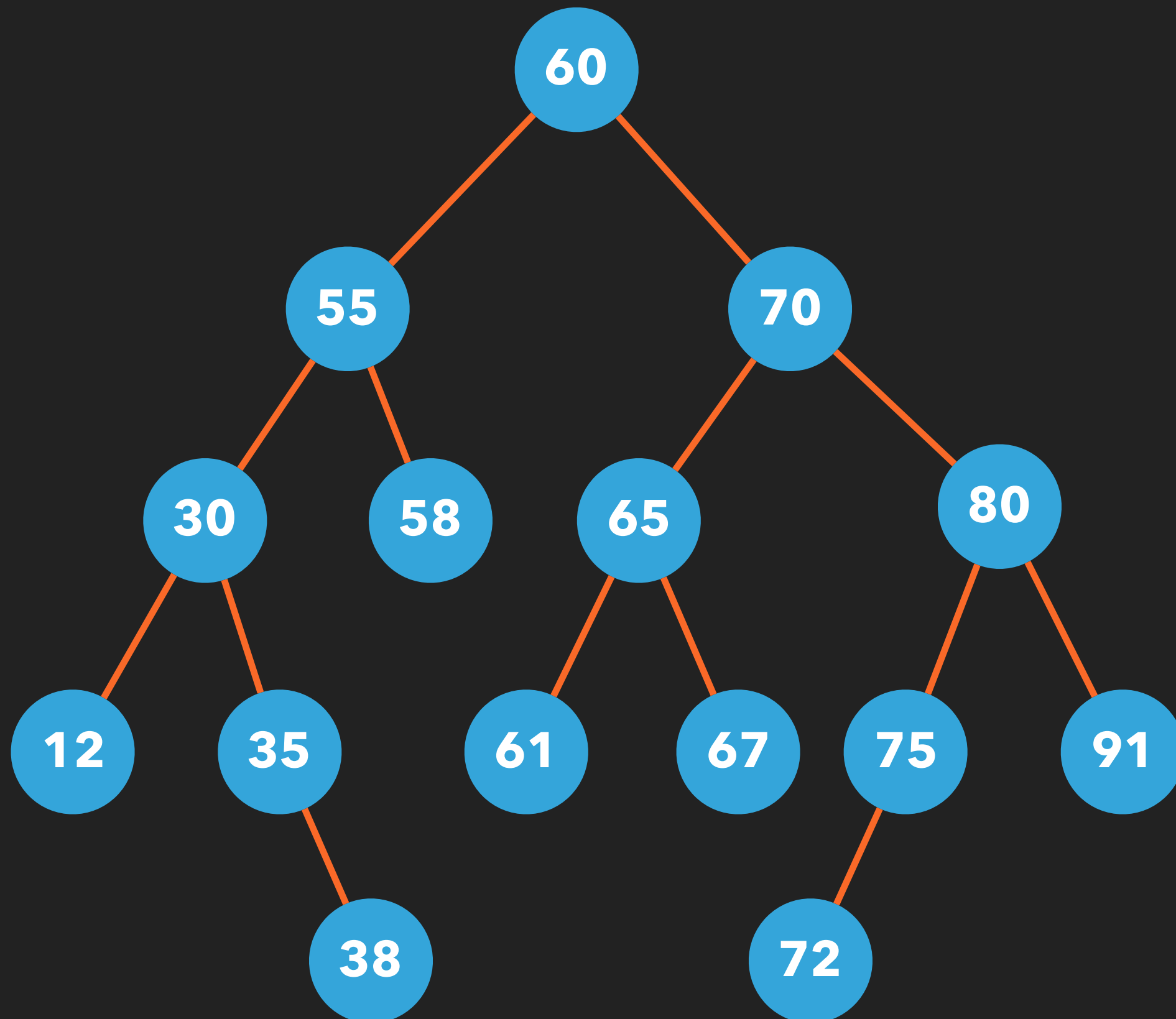
Postorder:



Binary Search Tree (BST)

- ◆ Special binary tree
 - ◆ BST has a root, a left subtree (L) and a right subtree (R)
 - ◆ The value of the root is greater than the value of every node in L
 - ◆ The value of the root is less than the value of every node in R
 - ◆ L and R are also BSTs
 - ◆ Used for efficient search in large data sets

BST

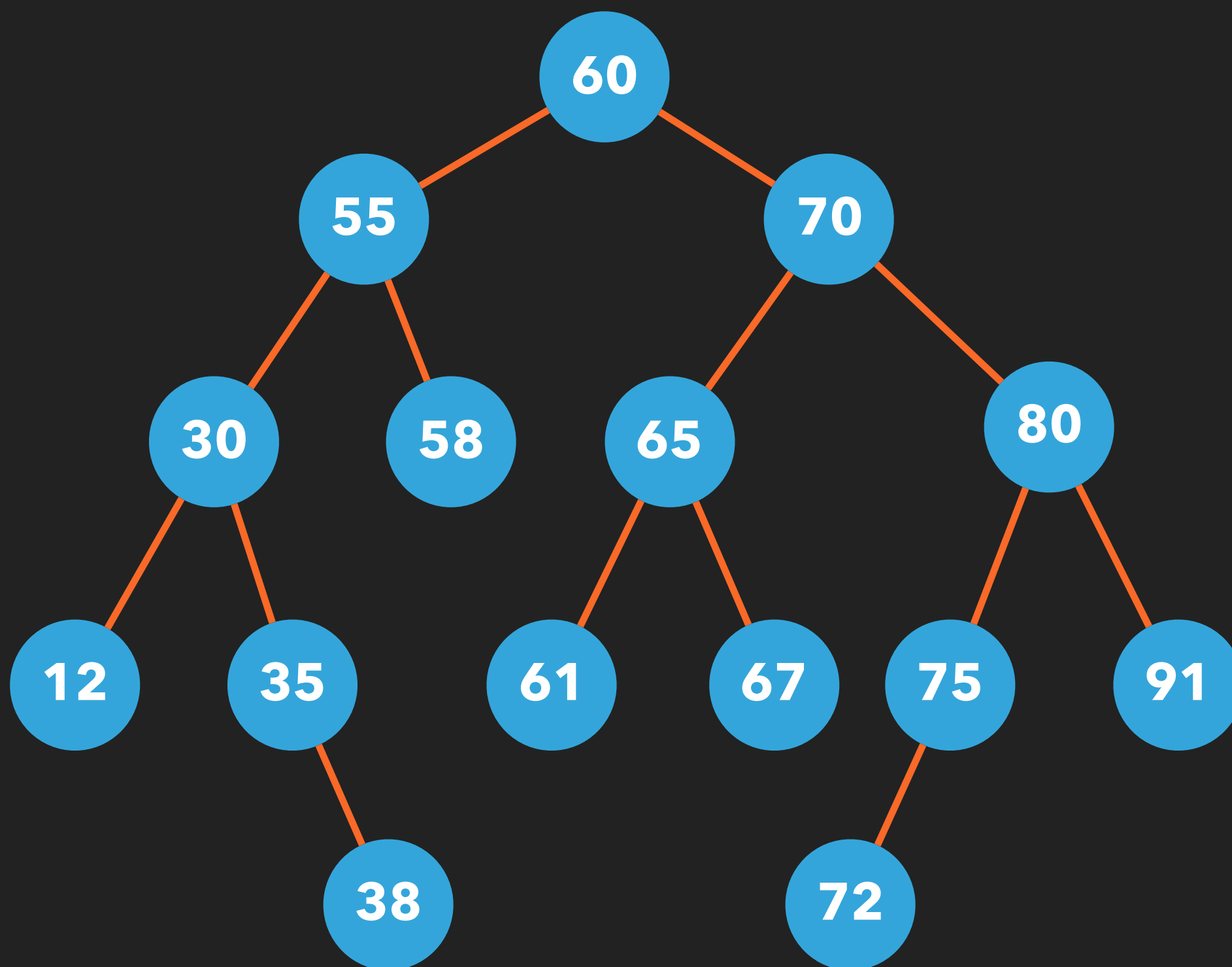


BST

- ◆ Common operations on BST
 - ◆ **Search** for a specific value in the BST
 - ◆ **Add** a node to the BST while keeping the BST properties
 - ◆ **Remove** a node from the BST while keeping the BST properties
 - ◆ **Traverse** the BST (preorder, inorder, postorder)

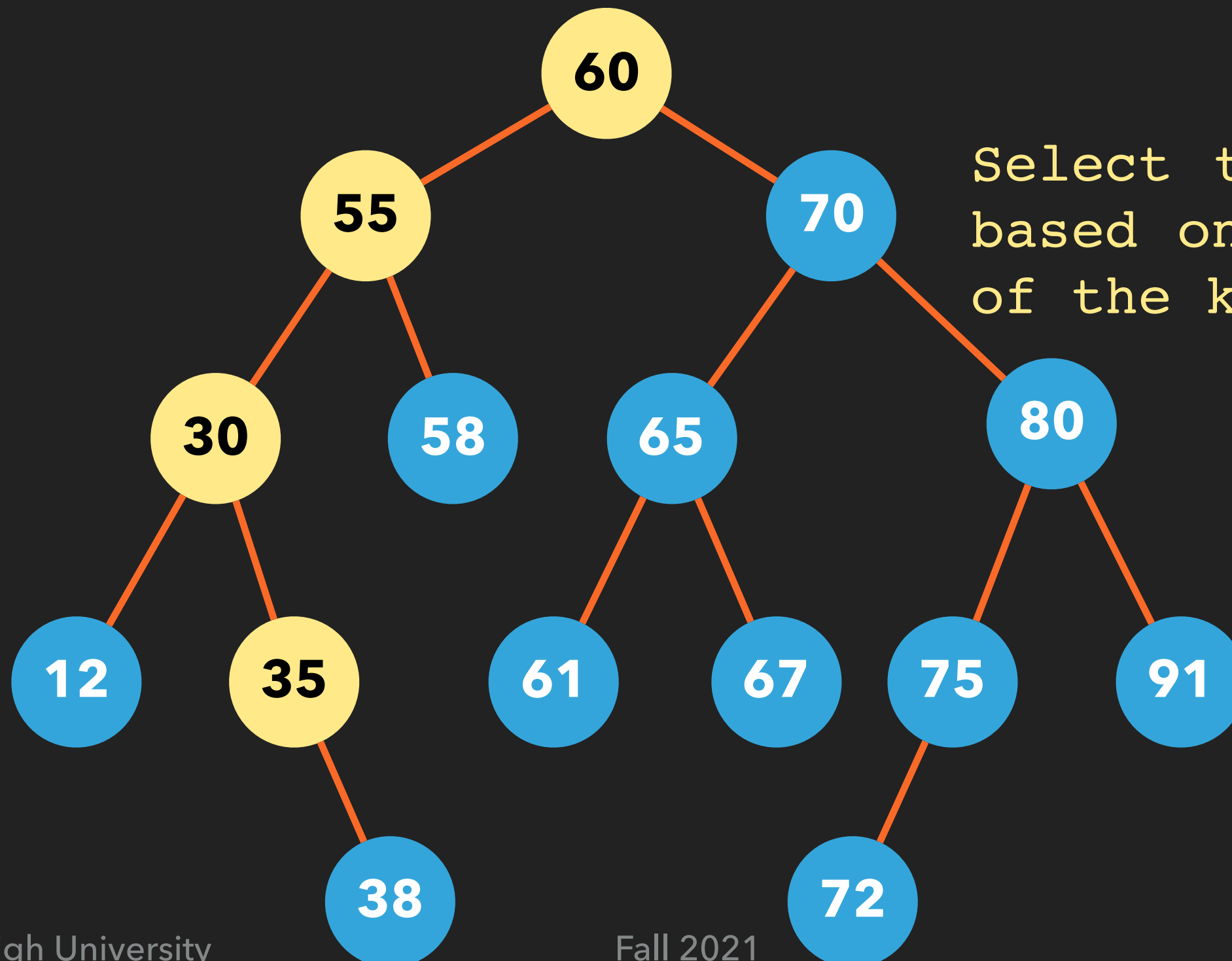
BST - Search

Search for the value 35



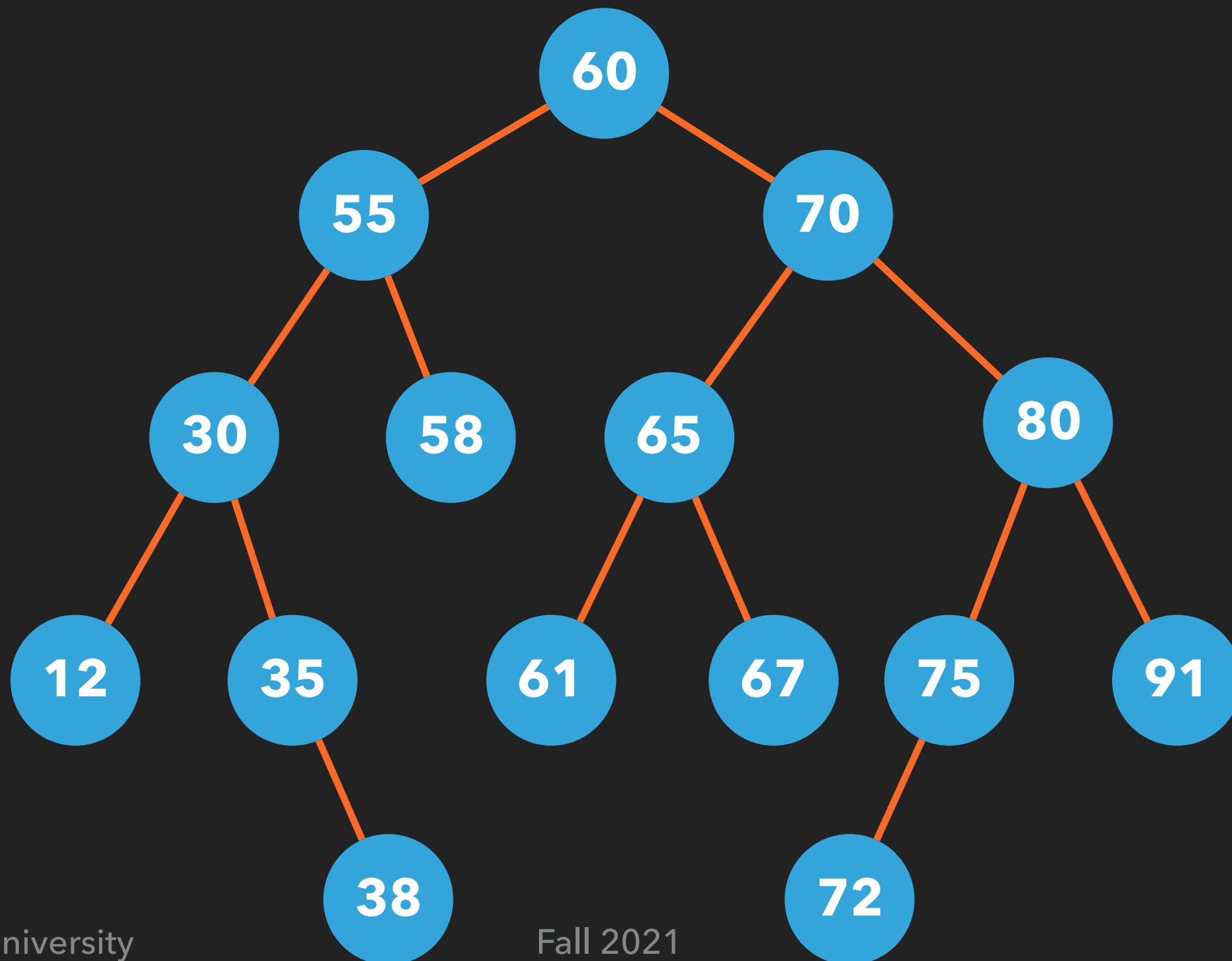
BST - Search operation

Search for the value 35



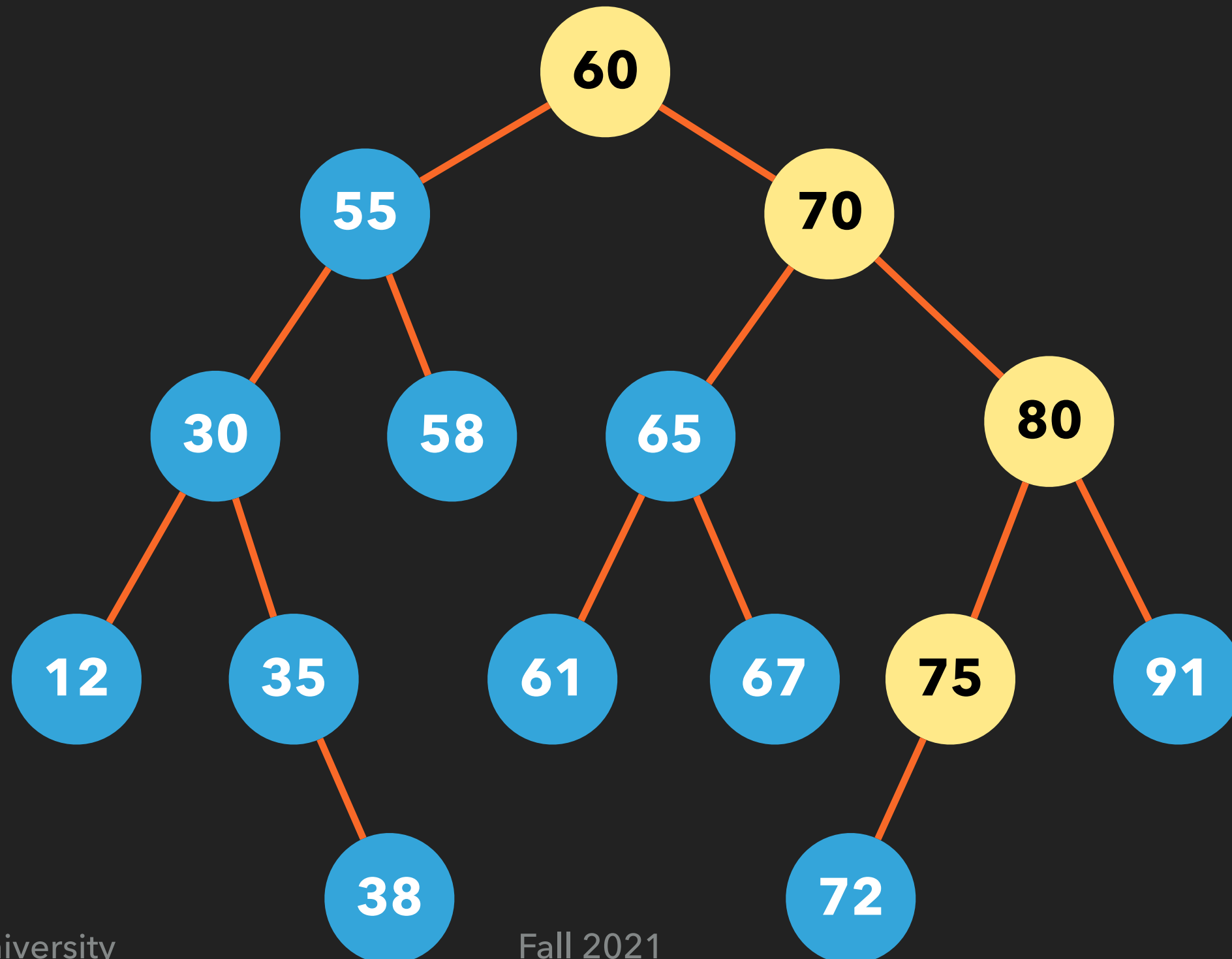
BST - Search operation

Search for the value 75



BST - Search operation

Search for the value 75



BST - Search algorithm

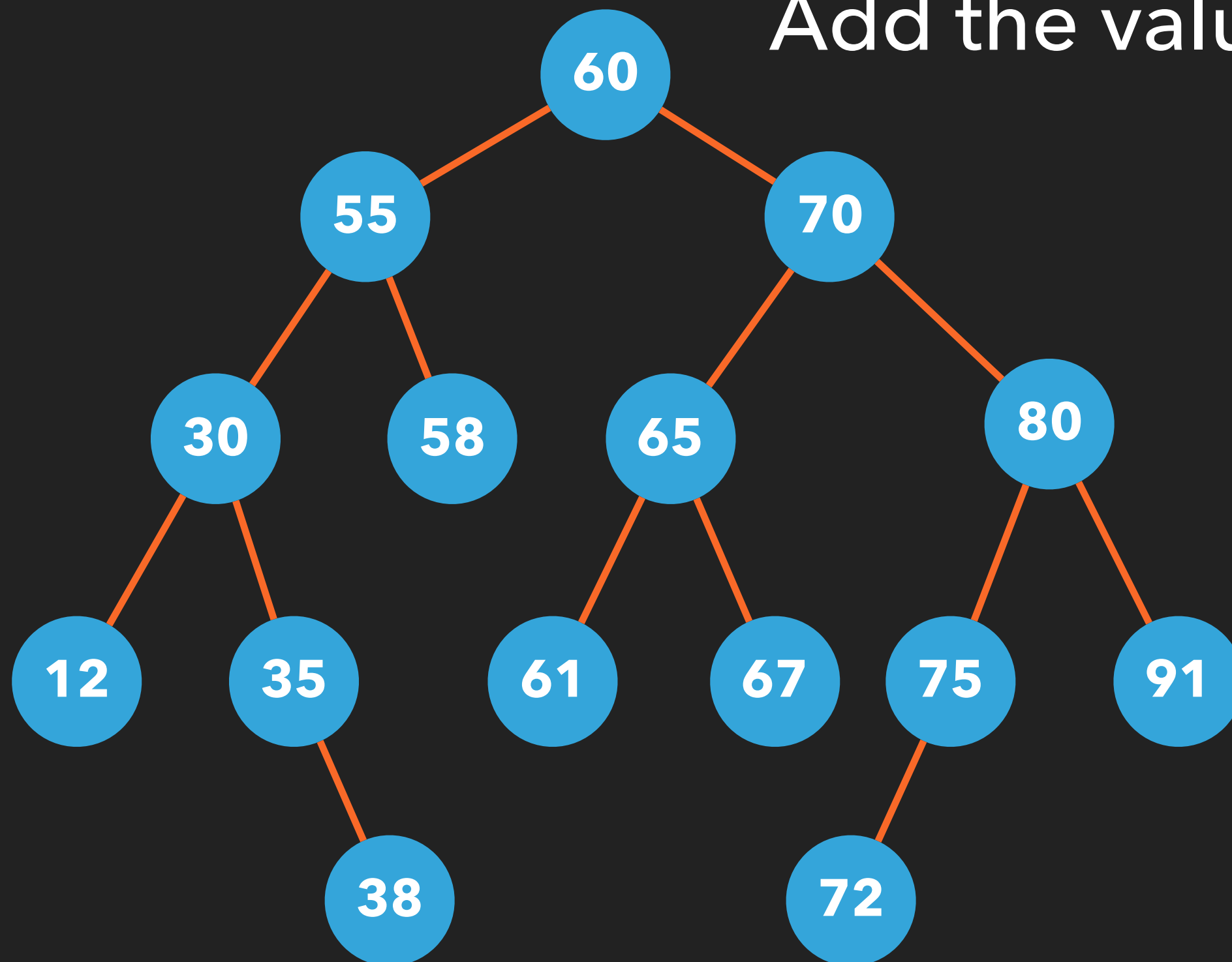
contains**boolean contains (item)**

```
current node = root // start from the root
while(current node is not null){
    if(the value of the current node == item)
        return true
    else if (value of the current node > item)
        current node is set to its left child
    else
        current node is set to its right child
}
return false
```

end contains

BST - Add operation

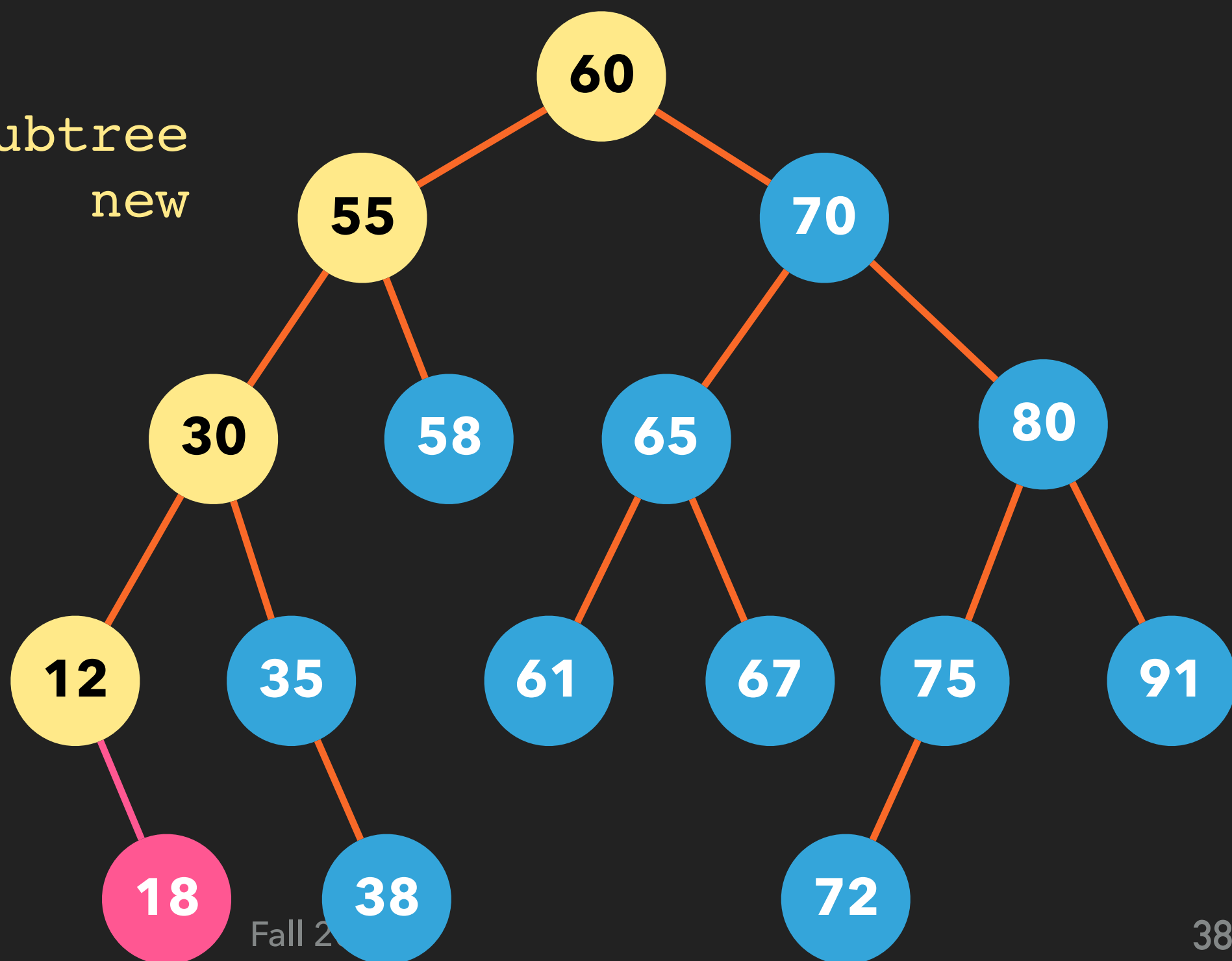
Add the value 18



BST - Add operation

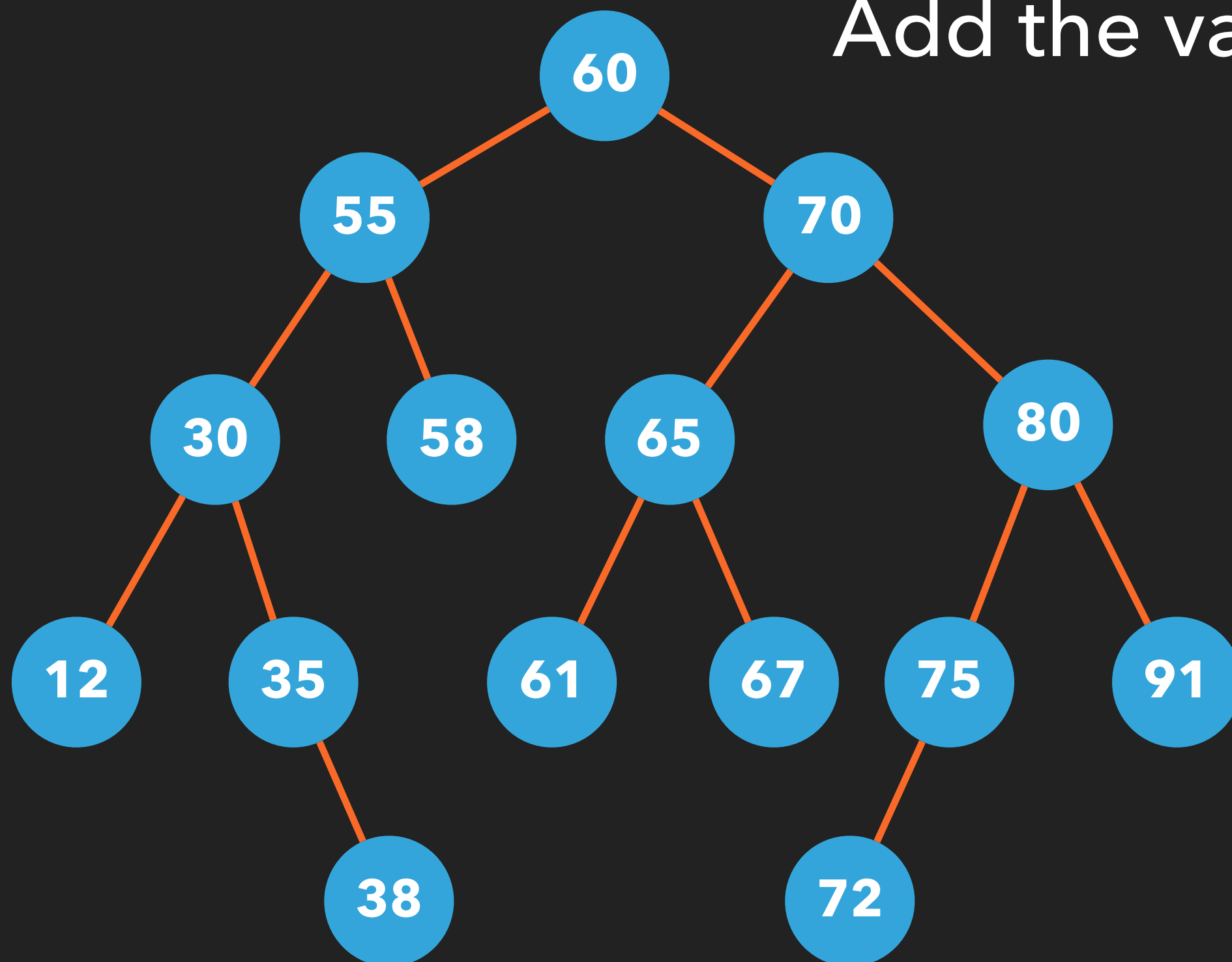
Add the value 18

Find the subtree
where the new
value fits



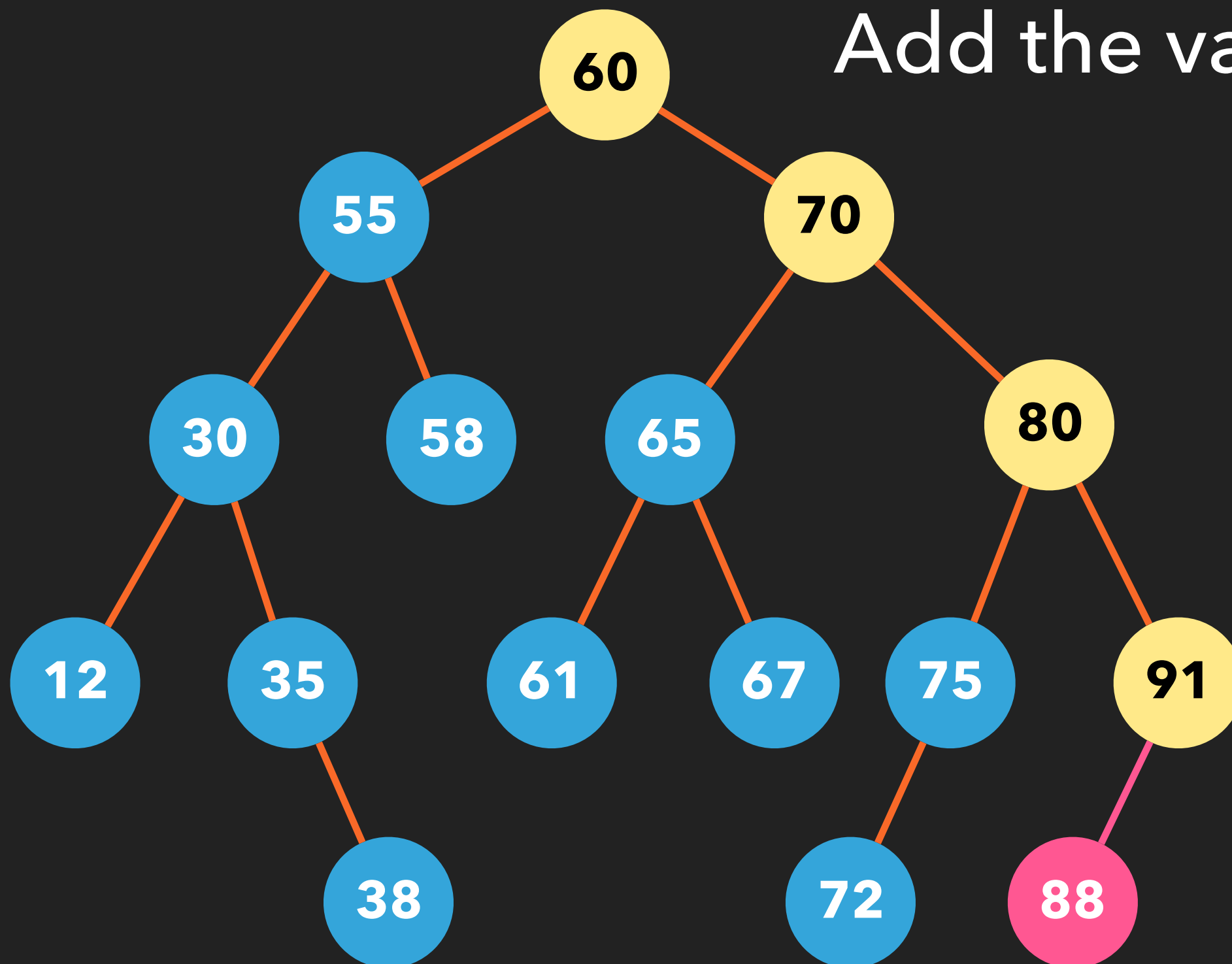
BST - Add operation

Add the value 88



BST - Add operation

Add the value 88



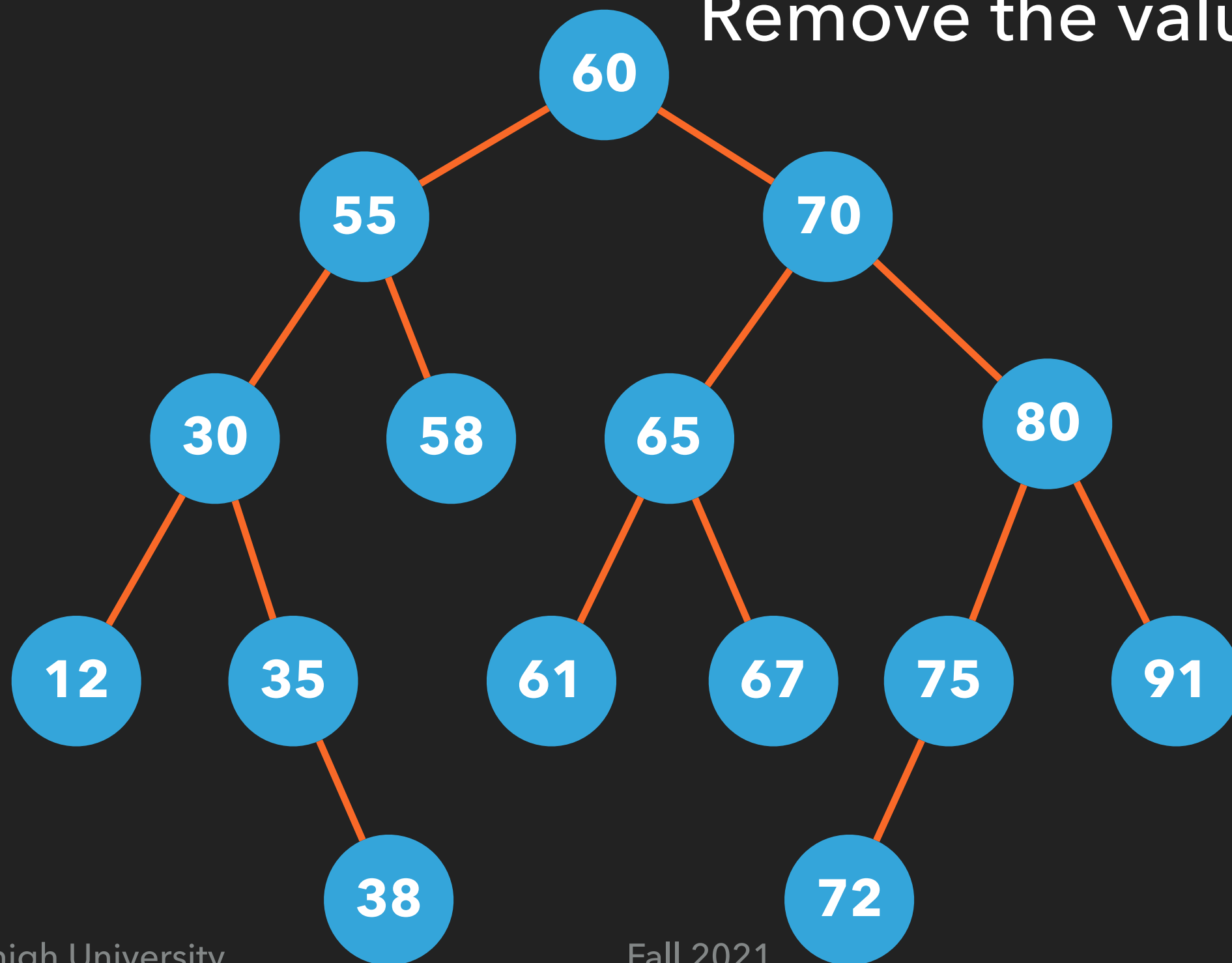
BST - Add operation

add

```
boolean add (item)
    current node = root
    while(current is not null){
        parent = current node
        if( the value of the current node == item)
            return false (duplicates are not allowed)
        else if (value of the current node > item)
            current is set to the left child
        else
            current node is set to the right child
    }
    if (the value of the parent node > item)
        Create a left child to parent (value=item)
    else
        Create a right child to parent (value=item)
    end if
    return true
end add
```

BST - Remove operation

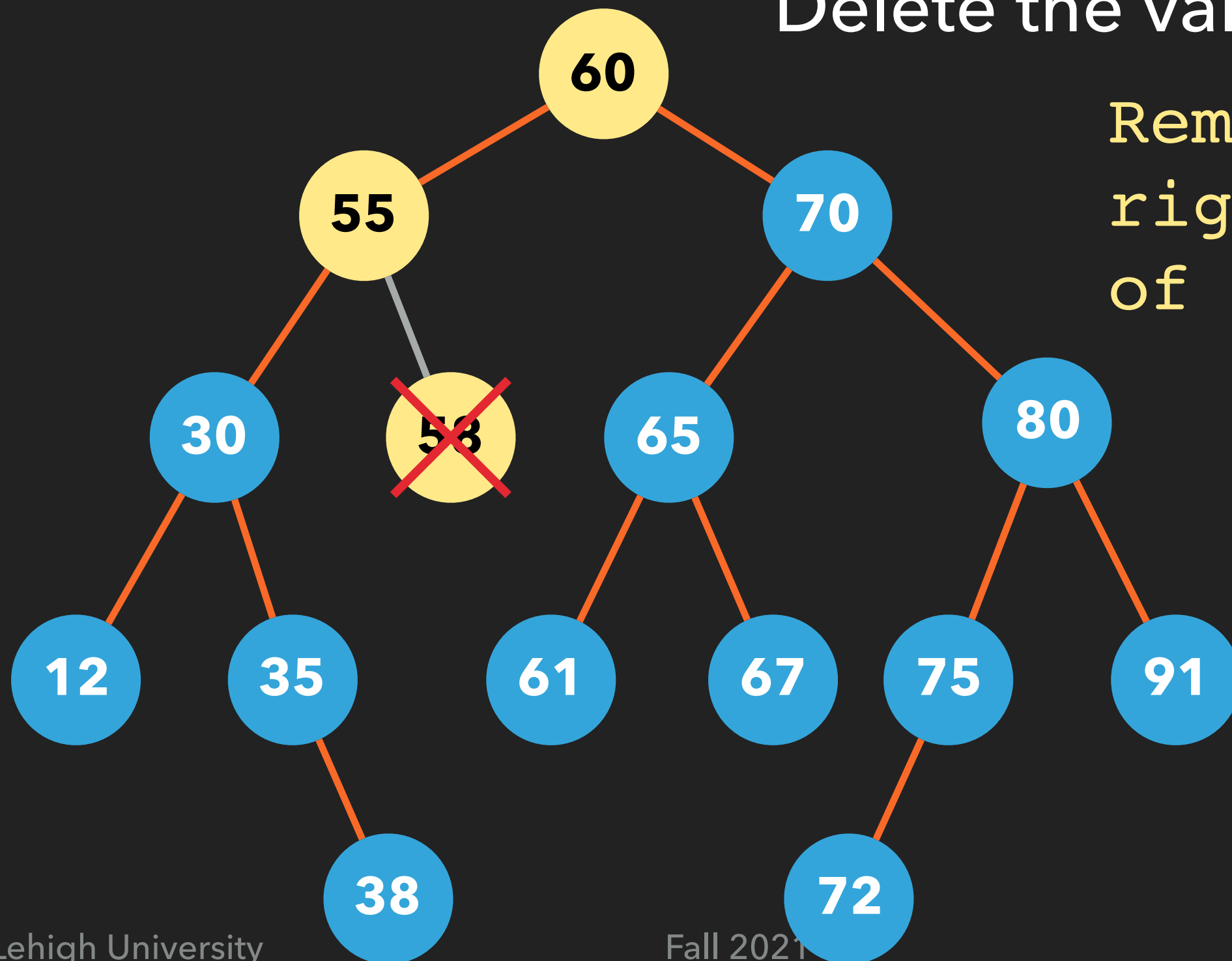
Remove the value 58 (Leaf)



BST - Remove operation

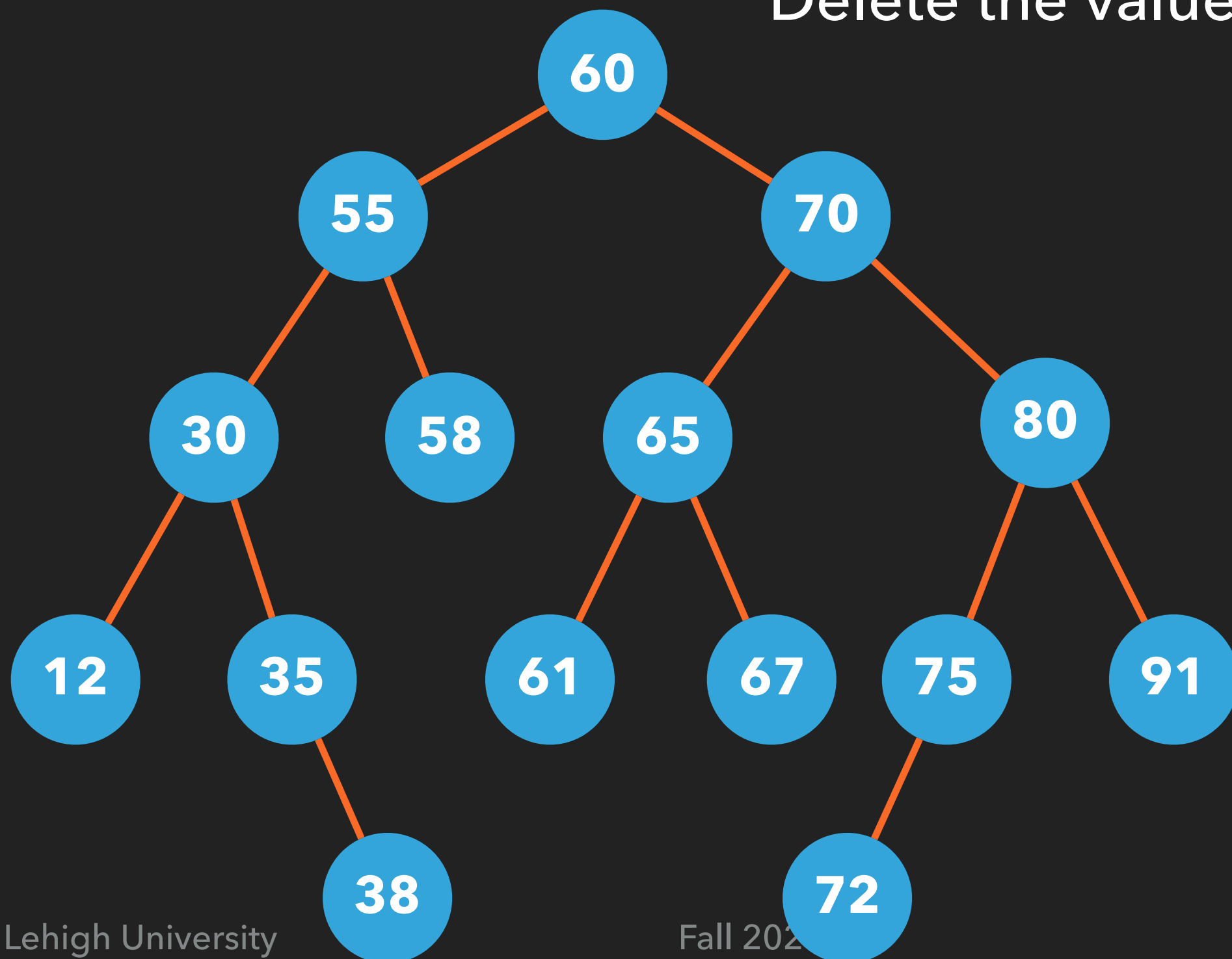
Delete the value 58 (Leaf)

Remove the
right child
of 55



BST - Remove operation

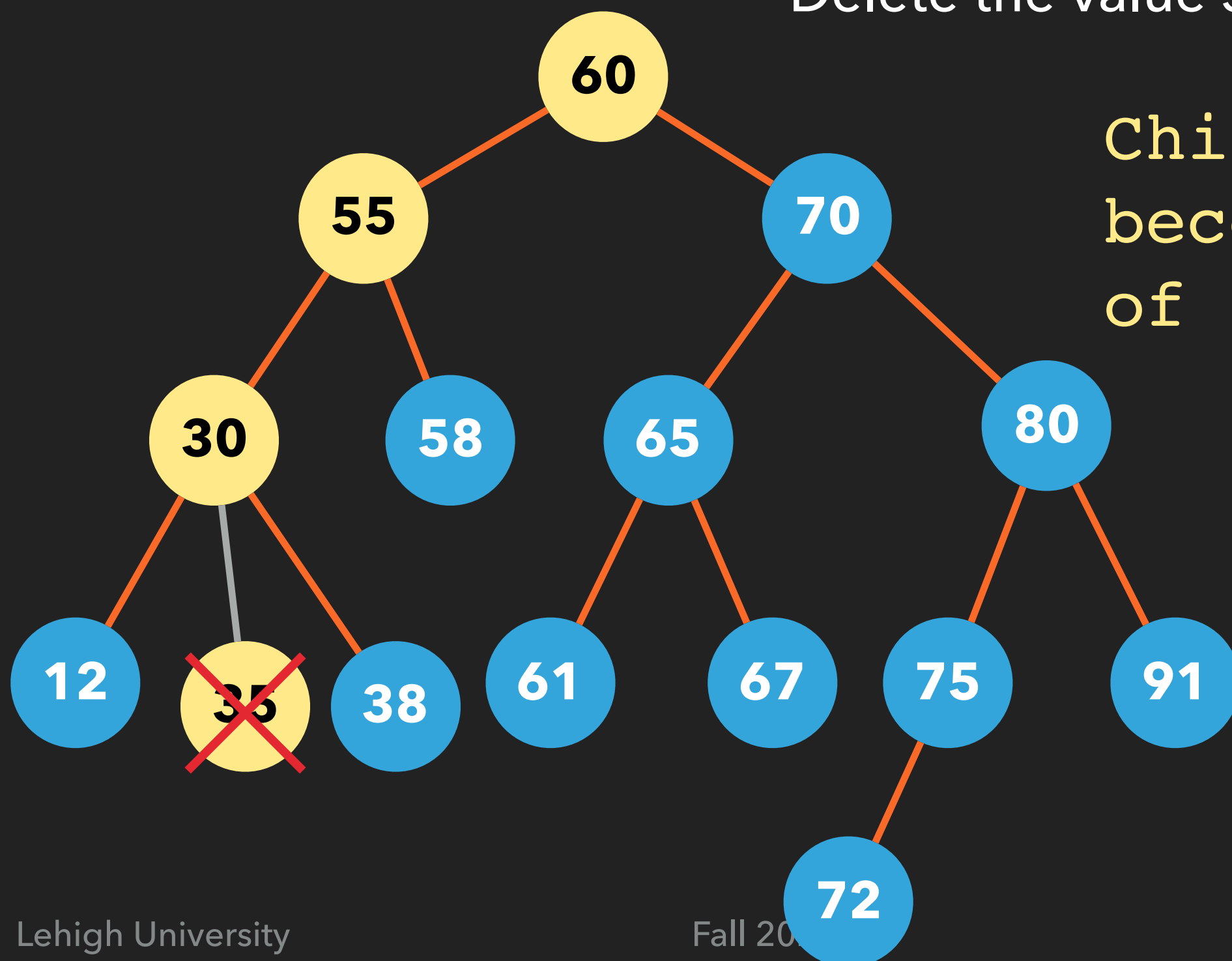
Delete the value 35 (one child)



BST - Remove operation

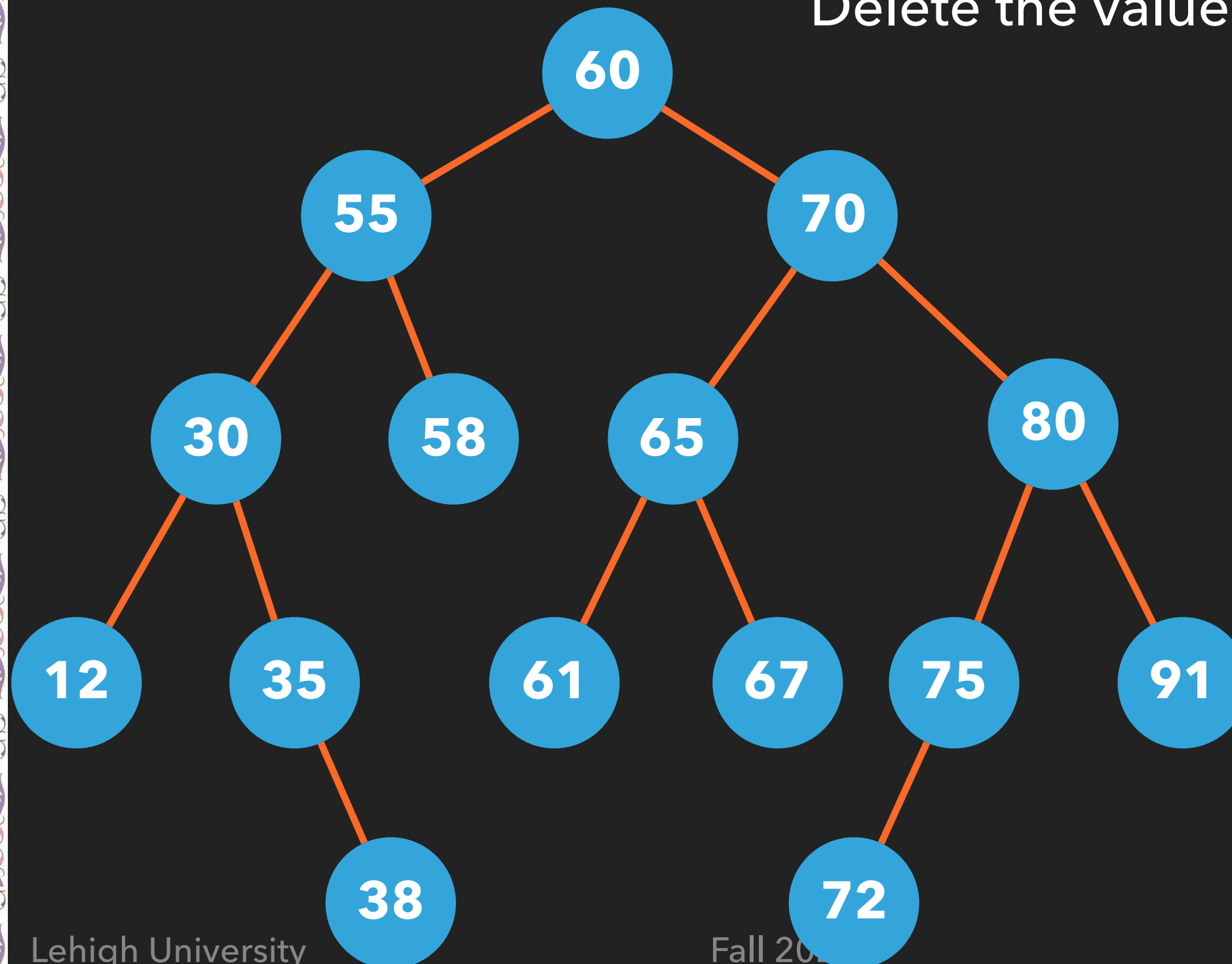
Delete the value 35 (has one child)

Child of 35
becomes child
of 30



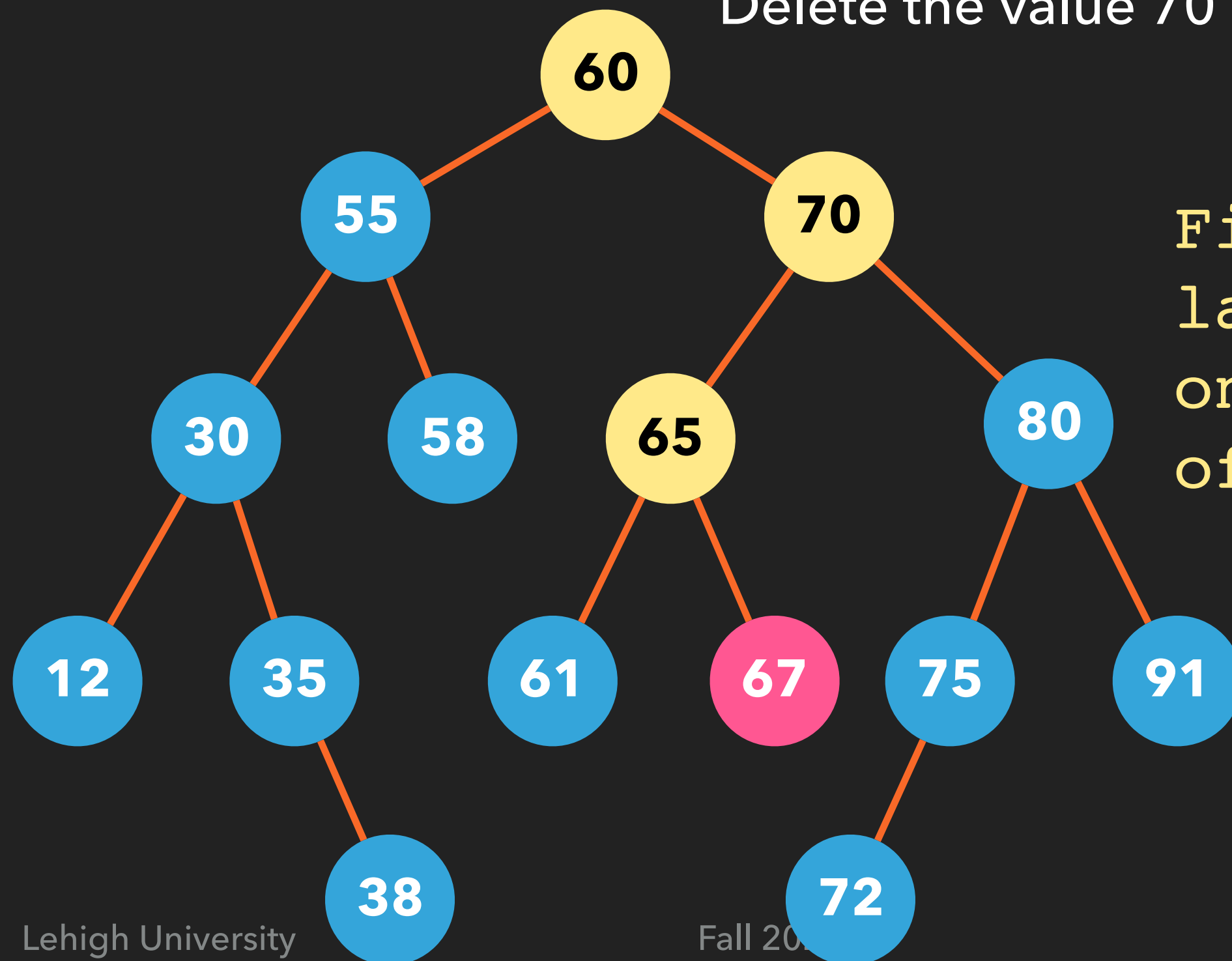
BST - Remove operation

Delete the value 70 (two children)



BST - Remove operation

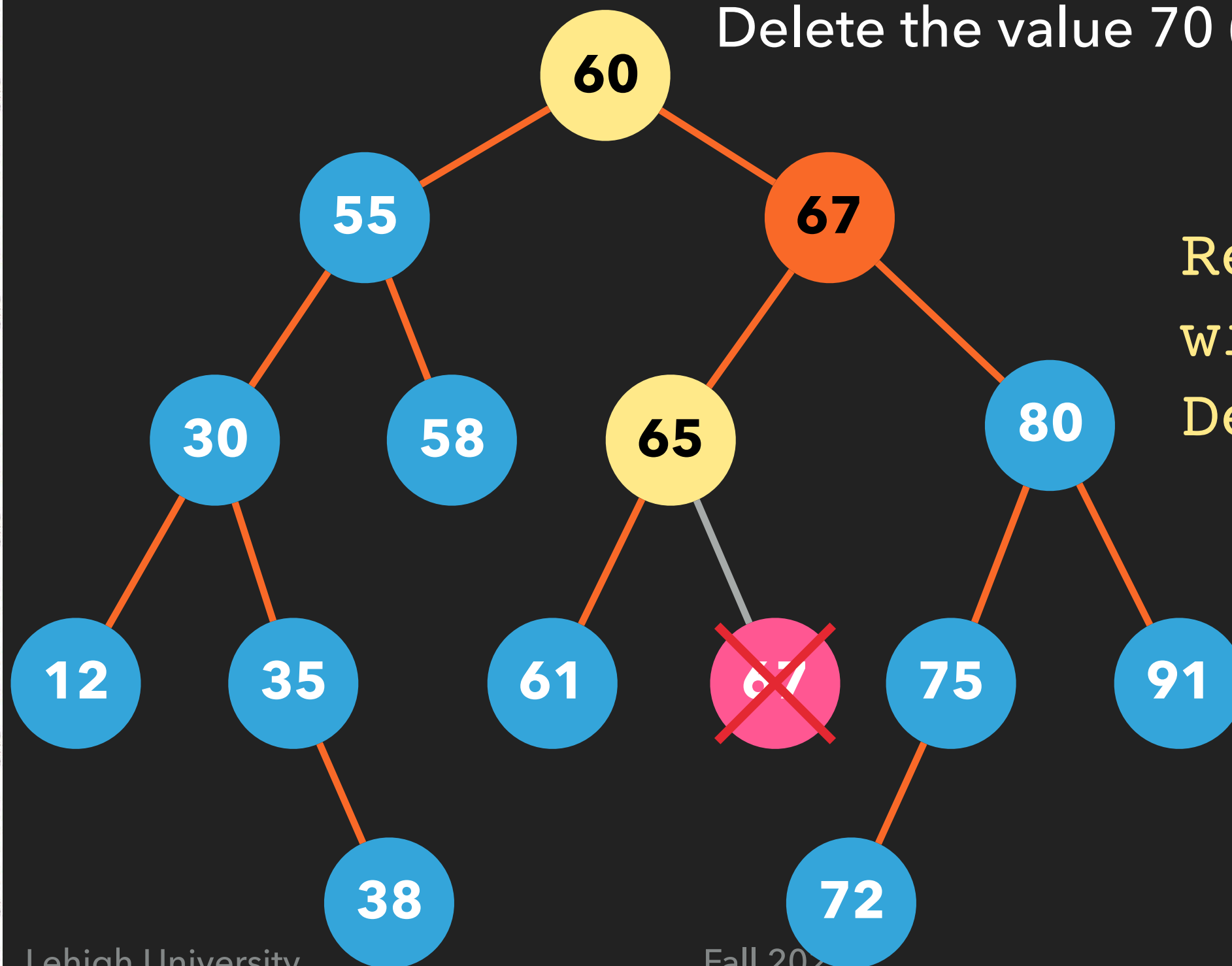
Delete the value 70 (has two children)



Find the
largest node
on the left
of 70

BST - Remove operation

Delete the value 70 (has two children)



Replace 70
with 67 -
Delete 67

BST - Remove operation

remove

```
boolean remove (item)
    node = search(item) // find node with value item first
    if (node == null)
        return false (item not found in the BST)
    else
        if (node has no children)
            remove link to node (parent points to null)
        else if (node has one child)
            replace node with node's child
        else if (node has two children)
            find the largest node on the left subtree of node
            copy the value of the largest node to node
            remove the largest node
        end if
    end if
    return true
end remove
```

Traversals

```
preorder() {  
    preorder(root) }  
preorder(node) {  
    print node  
    preorder(left child of node)  
    preorder(right child of node) }
```

```
inorder() {  
    inorder(root) }  
inorder(node) {  
    inorder(left child of node)  
    Print node  
    inorder(right child of node) }
```

```
postorder() {  
    postorder(root) }  
postorder(node) {  
    postorder(left child of node)  
    postorder(right child of node)  
    print node    }
```

BST - Implementation

- ◆ BST may be implemented in two ways
 - ◆ Array Based BST
 - ◆ Linked BST

BST - Implementation

- ◆ Nodes of the tree are stored in an array
- ◆ Children of a node follow the node (at specific indices)

BST - Implementation

- ◆ Nodes of the tree are linked
- ◆ Every node has two references in addition to its value: left child and right child

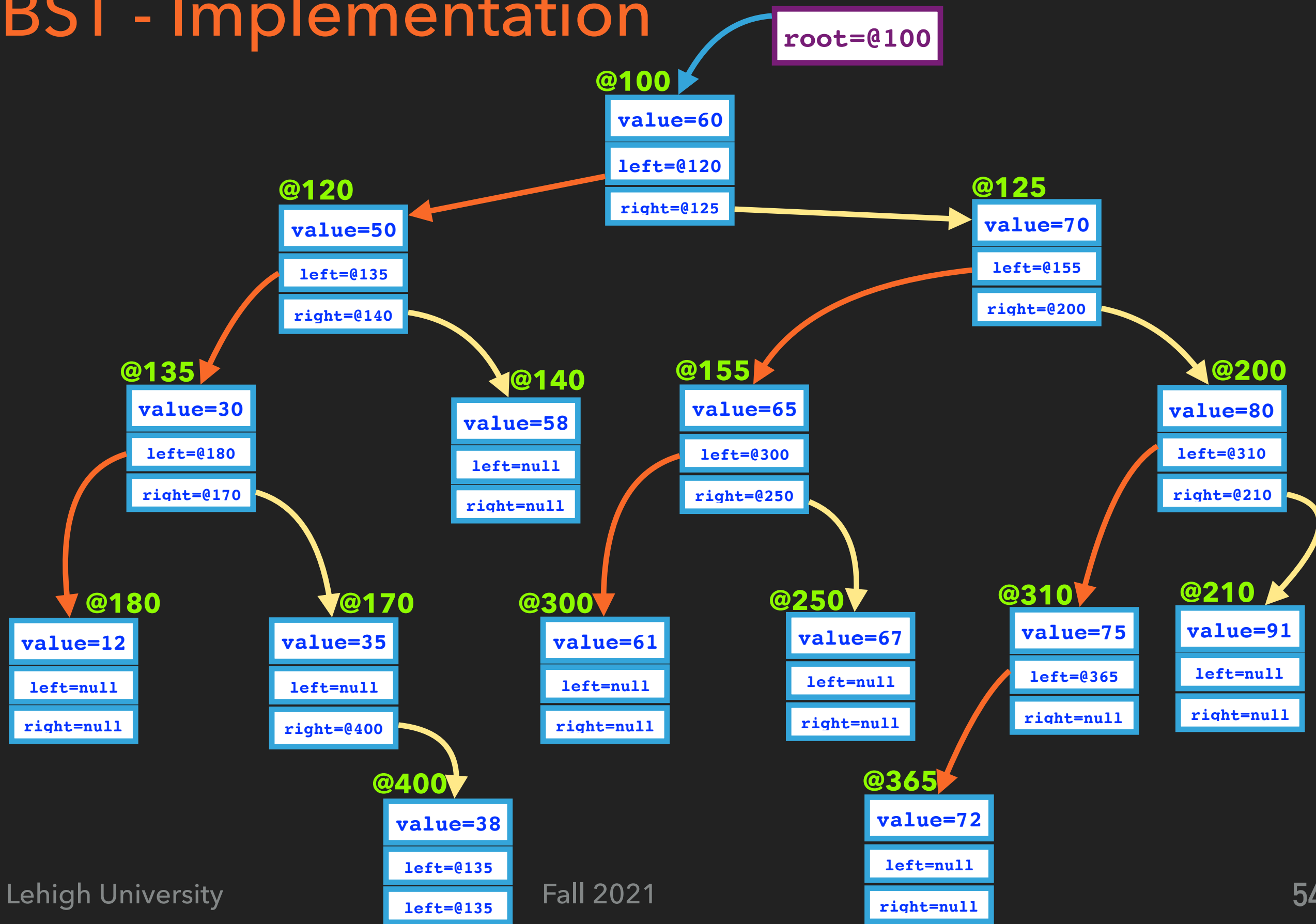
TreeNode

value

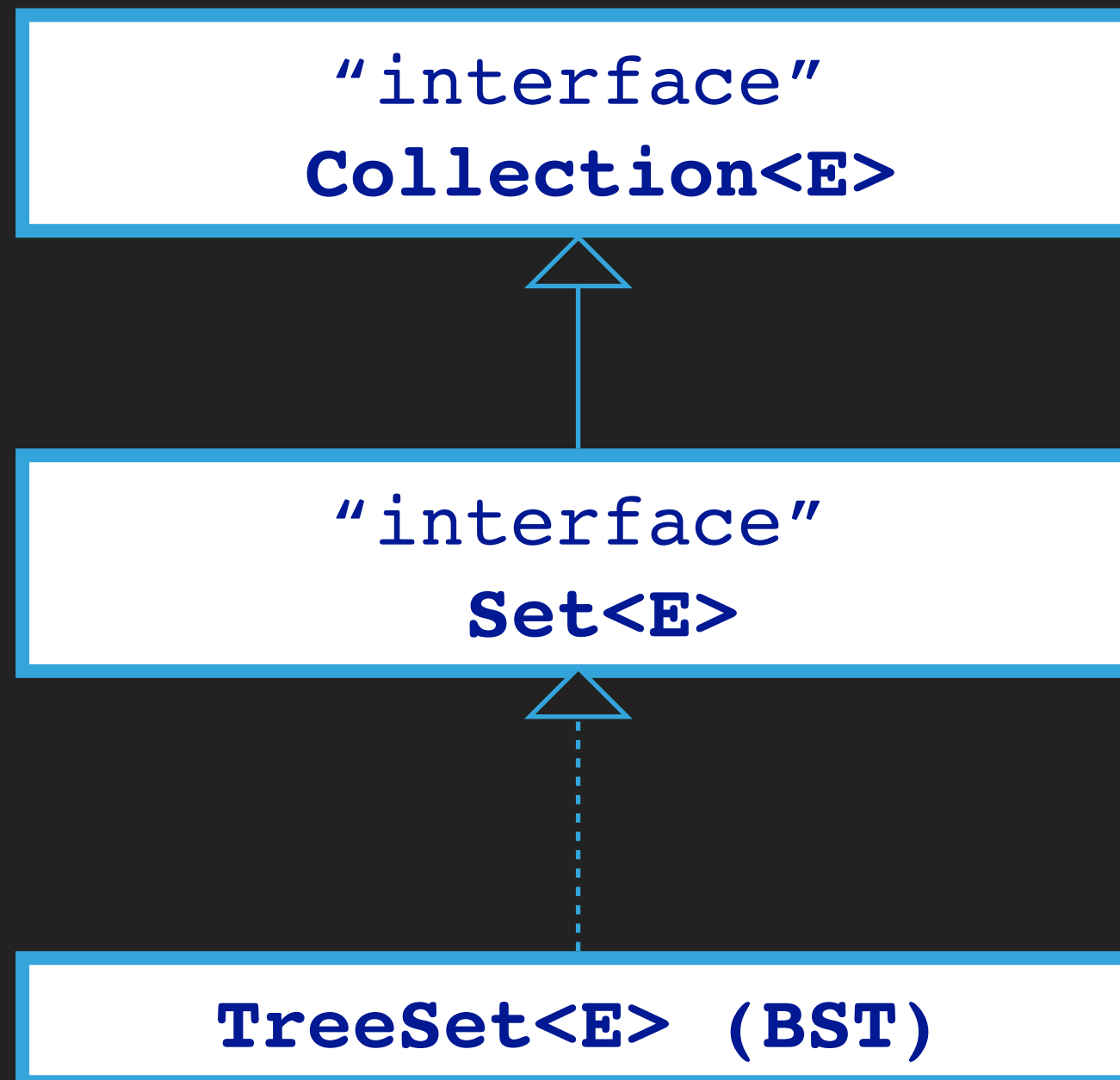
left

right

BST - Implementation



BST - Implementation



BST - Implementation

has

BST<E extends Comparable<E>>

-root: TreeNode

-size: int

+BST()

+size(): int

+isEmpty(): boolean

+clear(): void

+contains(E): boolean

+add(E): boolean

+remove(E): boolean

+inorder(): void

+preorder(): void

+postorder(): void

TreeNode

value: E

Left: TreeNode

Right: TreeNode

TreeNode(E val)

BST - Implementation

BST.java

```
public class BST<E extends Comparable<E>> {  
    private TreeNode root;  
    private int size;  
  
    private class TreeNode{  
        E value;  
        TreeNode left;  
        TreeNode right;  
        TreeNode(E val){  
            value = val;  
            left = right = null;  
        }  
    }  
}
```


BST - Implementation

BST.java

```
BST() {  
    root = null;  
    size = 0;  
}  
public int size() {  
    return size;  
}  
public boolean isEmpty() {  
    return (size == 0);  
}  
public void clear() {  
    root = null;  
}
```


BST - Implementation

BST.java

// Search method

```
public boolean contains(E item) {  
    TreeNode node = root;  
    while (node != null) {  
        if( item.compareTo(node.value) < 0)  
            node = node.left;  
        else if (item.compareTo(node.value) > 0)  
            node = node.right;  
        else  
            return true;  
    }  
    return false;  
}
```

```
// Method add()
public boolean add(E item) {
    if (root == null) // first node to be inserted
        root = new TreeNode(item);
    else {
        TreeNode parent, node;
        parent = null; node = root;
        while (node != null) { // Looking for a leaf node
            parent = node;
            if (item.compareTo(node.value) < 0) {
                node = node.left; }
            else if (item.compareTo(node.value) > 0) {
                node = node.right; }
            else
                return false; // duplicates are not allowed
        }
        if (item.compareTo(parent.value) < 0)
            parent.left = new TreeNode(item);
        else
            parent.right = new TreeNode(item);
    }
    size++;
    return true;
}
```

BST - Implementation

BST.java

```
// Method remove()
public boolean remove(E item) {
    TreeNode parent, node;
    parent = null; node = root;
    // Find item first
    while (node != null) {
        if (item.compareTo(node.value) < 0) {
            parent = node;
            node = node.left;
        }
        else if (item.compareTo(node.value) > 0) {
            parent = node;
            node = node.right;
        }
        else
            break; // item found
    }
}
```

BST - Implementation

BST.java

```
if (node == null) // item not in the tree
    return false;

// Case 1: node has no children
if (node.left == null && node.right == null) {
    if (parent == null) { // delete root
        root = null;
    }
    else {
        changeChild(parent, node, null);
    }
}
```


BST - Implementation

BST.java

```
// case 2: one right child
else if (node.left == null) {
    if (parent == null) { // delete root
        root = node.right;
    }
    else {
        changeChild(parent, node, node.right);
    }
}

// case 2: one left child
else if (node.right == null) {
    if (parent == null) { // delete root
        root = node.left;
    }
    else {
        changeChild(parent, node, node.left);
    }
}
```


BST - Implementation

BST.java

```
// Case 3: node has two children
else {
    TreeNode rightMostParent = node;
    TreeNode rightMost = node.left;
    // go right on the left subtree
    while (rightMost.right != null) {
        rightMostParent = rightMost;
        rightMost = rightMost.right;
    }
    // copy the value of rightMost to node
    node.value = rightMost.value;
    //delete rightMost
    changeChild(rightMostParent, rightMost,
                rightMost.left);
}
size--;
return true;
}
```

BST - Implementation

BST.java

```
private void changeChild(TreeNode parent,  
    TreeNode node, TreeNode newChild){  
    if(parent.left == node)  
        parent.left = newChild;  
    else  
        parent.right = newChild;  
}
```

BST - Implementation

BST.java

```
// Recursive method inorder()
public void inorder() {
    inorder(root);
}
private void inorder(TreeNode node) {
    if (node != null) {
        inorder(node.left);
        System.out.print(node.value + " ");
        inorder(node.right);
    }
}
```

BST - Implementation

BST.java

```
// Recursive method preorder()
public void preorder() {
    preorder(root);
}
private void preorder(TreeNode node) {
    if (node != null) {
        System.out.print(node.value + " ");
        preorder(node.left);
        preorder(node.right);
    }
}
```


BST - Implementation

BST.java

```
// Recursive method postorder()
public void postorder() {
    postorder(root);
}
private void postorder(TreeNode node) {
    if (node != null) {
        postorder(node.left);
        postorder(node.right);
        System.out.print(node.value + " ");
    }
}
```

BST - Testing

Test.java

```
// Testing the class BST
public static void main(String[] args){
    BST<String> bst = new BST<>();
    bst.add("Kiwi");
    bst.add("Strawberry");
    bst.add("Apple");
    bst.add("Banana");
    bst.add("Orange");
    bst.add("Lemon");
    bst.add("Watermelon");
    bst.inorder();
    bst.remove("Banana");
    System.out.println(bst.contains("Banana"));
    bst.inorder();
    bst.remove("Orange");
    bst.inorder();
    bst.remove("Kiwi");
    bst.inorder();
}
```

BST - Implementation

◆ Performance of the BST operations

Method	Complexity
BST()	$O(1)$
size()	$O(1)$
clear()	$O(1)$
isEmpty()	$O(1)$
contains(E)	$O(n)$ / $O(\log n)$
add(E)	$O(n)$ / $O(\log n)$
remove(E)	$O(n)$ / $O(\log n)$
inorder()	$O(n)$
preorder()	$O(n)$
postorder()	$O(n)$

BST - Testing

Test.java

```
// Testing the class BST
```

```
public static void main(String[] args){
```

```
    BST<String> bst = new BST<>();
```

```
    bst.add("Apple");
```

```
    bst.add("Banana");
```

```
    bst.add("Kiwi");
```

```
    bst.add("Lemon");
```

```
    bst.add("Orange");
```

```
    bst.add("Strawberry");
```

```
    bst.add("Watermelon");
```

```
    bst.inorder();
```

```
}
```

Order of the
added values
affects the shape
of the tree

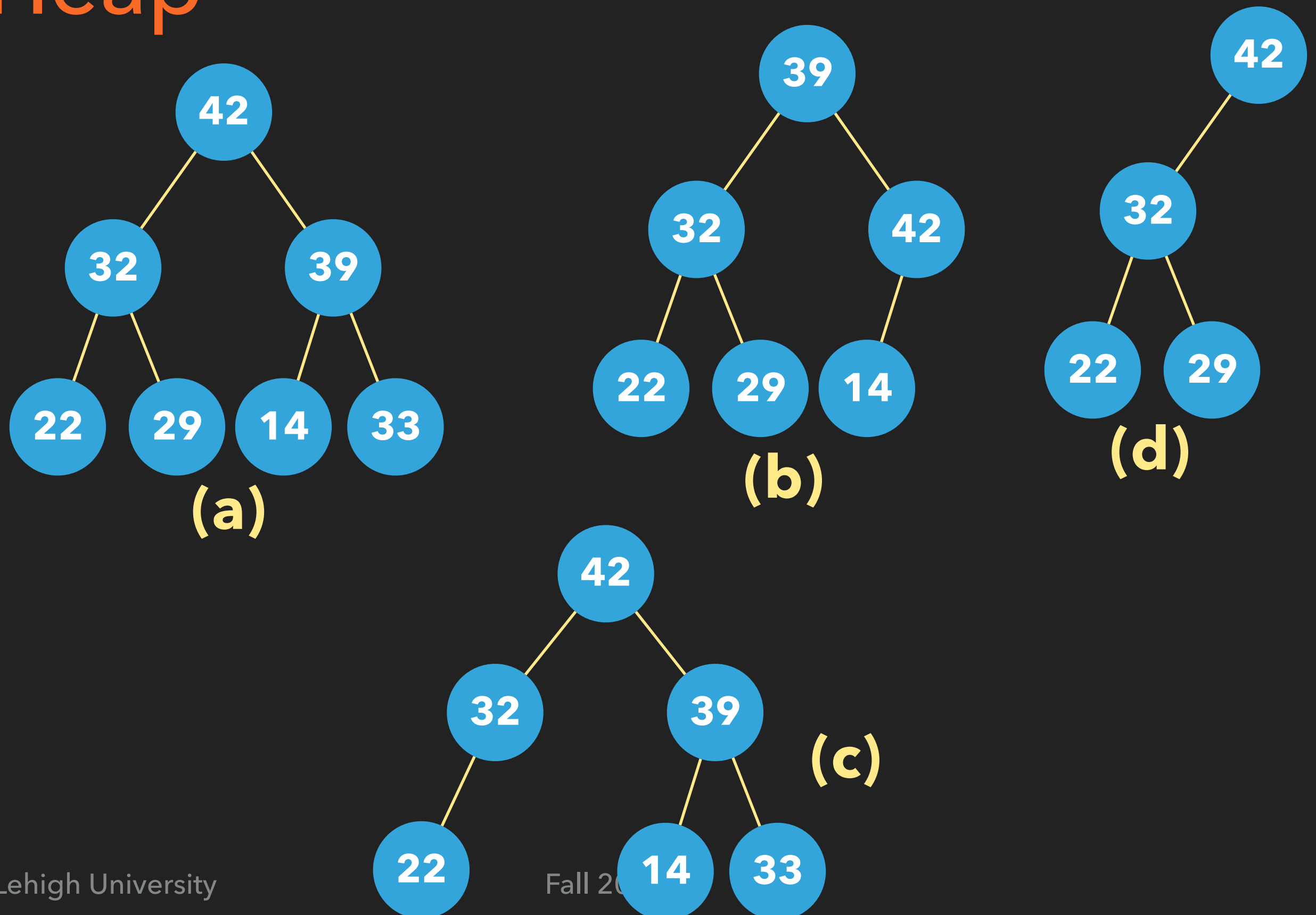
Summary

- ◆ Binary trees and binary search trees
- ◆ Operations: Search, Add, Remove, Traversals
- ◆ Implementation - Linked Nodes
- ◆ Order of add operations has an effect on the shape of the BST

Heap

- ◆ Special binary tree
 - ◆ Complete binary tree - All the levels are filled except the last level
All leaves on the last level are placed leftmost
- ◆ Every node is greater than or equal to any of its children (**Max Heap**) [Min Heap: less than or equal]
- ◆ Used for efficient sorting

Heap

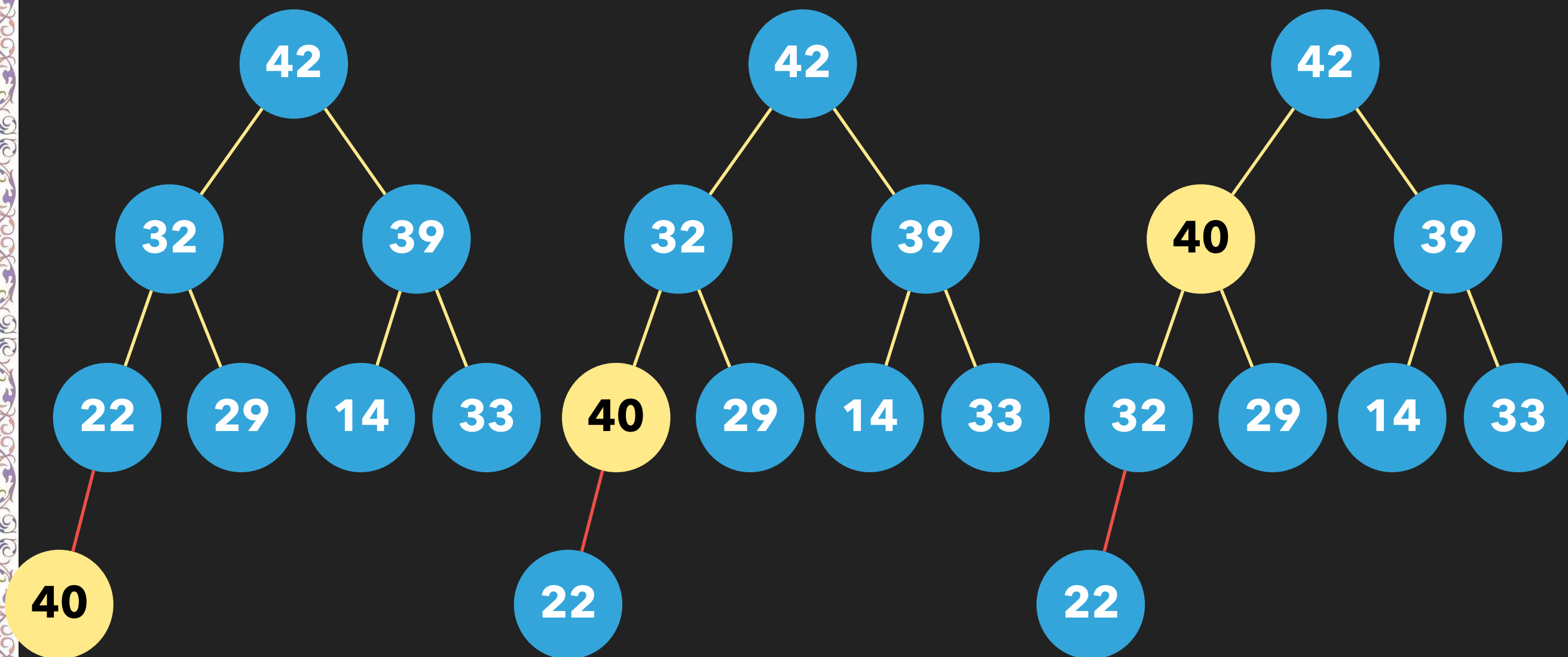


Heap

- ◆ Two main operations on the Heap
 - ◆ Adding a new node while keeping the heap properties
 - ◆ Removing a node while keeping the heap properties

Heap

- ◆ Adding a new node to the heap (40)



Heap

◆ Adding a new node to the heap

Algorithm **add**

Add the new node at the end of the heap

Current node = added node

While (current node > its parent)

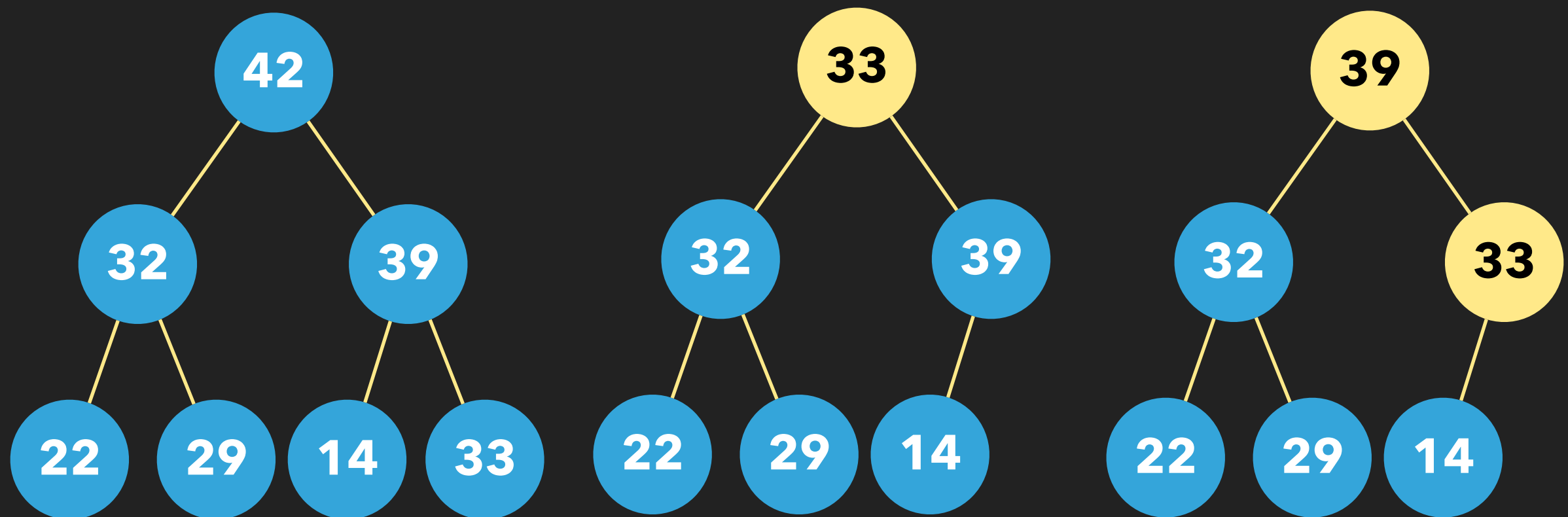
 Swap current node with its parent

 Current node becomes the parent

End

Heap

◆ Removing a node from the heap (42)



Heap

◆ Removing a node from the heap (root)

Algorithm **remove**

Move the last node to replace the root

Current node = root

While (current node < its children)

 Swap current node with the largest
 of its children

 Current node becomes the largest child

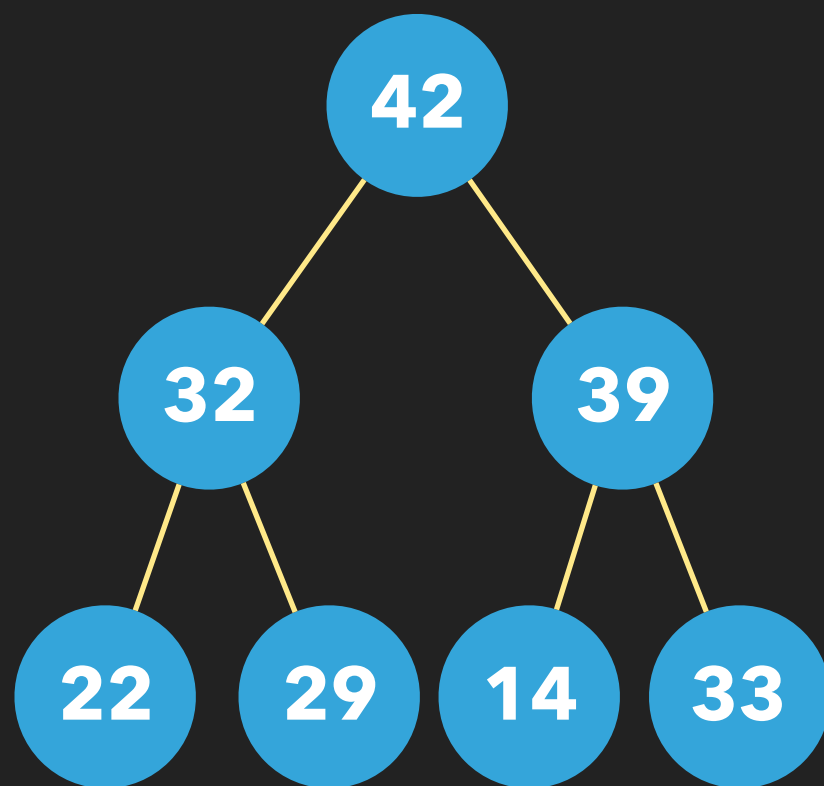
End

Heap

- ◆ Heap implementation
 - ◆ ArrayList to store the heap nodes
 - ◆ Easy access to children and parent

Heap

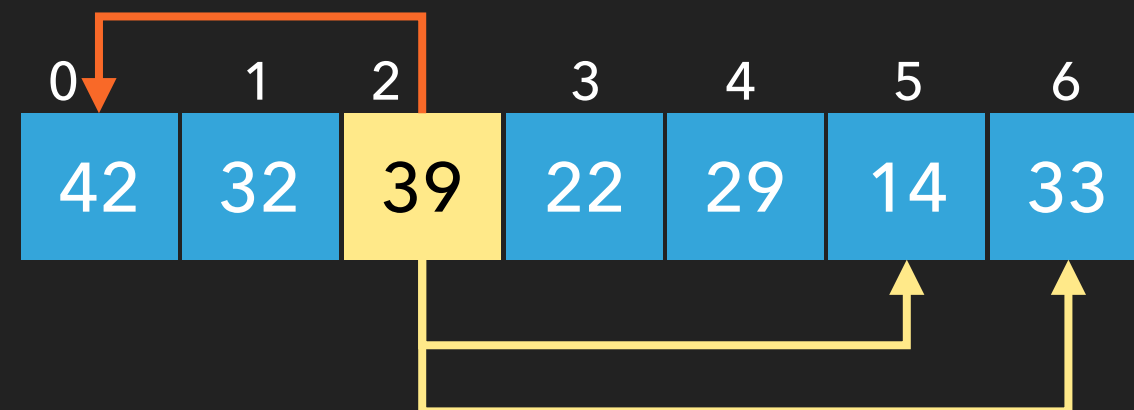
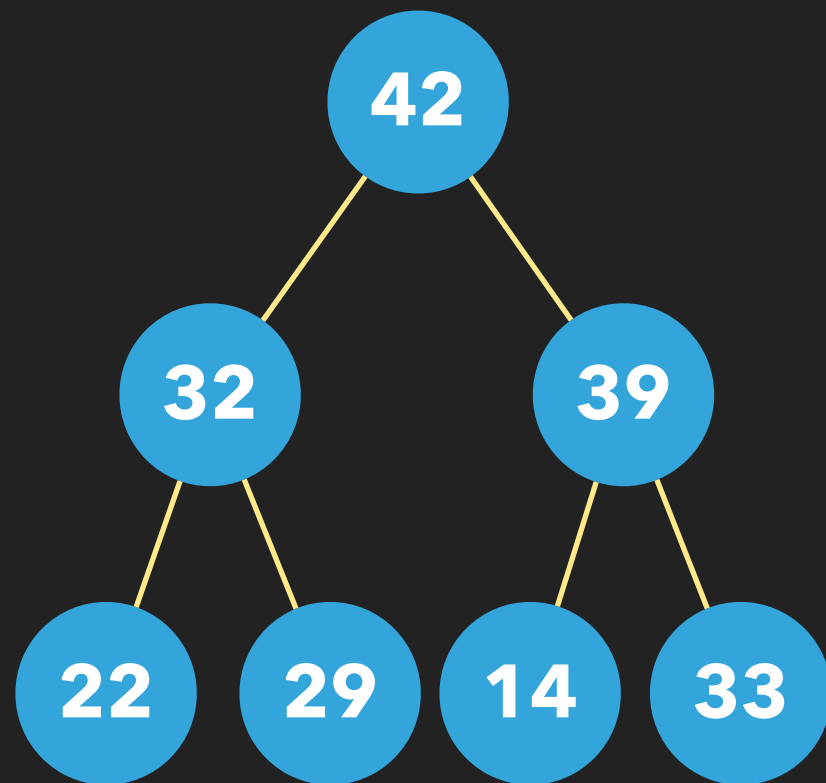
◆ Heap implementation



0	1	2	3	4	5	6
42	32	39	22	29	14	33

Heap

◆ Heap data structure



$$\text{IndexOf(Parent)} = (\text{IndexOf(current)} - 1) / 2$$

$$\text{IndexOf(Left child)} = 2 * \text{IndexOf(current)} + 1$$

$$\text{IndexOf(Right child)} = 2 * \text{IndexOf(current)} + 2$$

Heap

```
Heap<E extends Comparable<E>>
```

```
-list: ArrayList<E>
```

```
+Heap()
```

```
+add(E): void
```

```
+remove(): E
```

```
+contains(E): boolean
```

```
+size(): int
```

```
+isEmpty(): boolean
```

```
+clear(): void
```

```
+toString(): String
```


Heap

Heap.java

```
public class Heap<E extends Comparable<E>> {  
    private ArrayList<E> list;  
    public Heap(){  
        list = new ArrayList<>();  
    }  
    public int size(){  
        return list.size();  
    }  
    public boolean isEmpty(){  
        return list.isEmpty();  
    }  
    public void clear(){  
        list.clear();  
    }  
    public String toString(){  
        return list.toString();  
    }  
}
```

Heap

Heap.java

```
public boolean contains(E item) {  
    for(int i=0; i<list.size(); i++) {  
        if(list.get(i).equals(item))  
            return true;  
    }  
    return false;  
}
```

```
public void add(E item) {  
    list.add(item); //append item to the heap  
    int currentIndex = list.size()-1;  
    //index of the last element  
    while(currentIndex > 0) {  
        int parentIndex = (currentIndex-1)/2;  
        //swap if current is greater than its parent  
        E current = list.get(currentIndex);  
        E parent = list.get(parentIndex);  
        if(current.compareTo(parent) > 0) {  
            list.set(currentIndex, parent);  
            list.set(parentIndex, current);  
        }  
        else  
            break; // the tree is a heap  
        currentIndex = parentIndex;  
    }  
}
```

Heap

```
public E remove() {  
    if(list.size() == 0) return null;  
    //copy the value of the last node to root  
    E removedItem = list.get(0);  
    list.set(0, list.get(list.size()-1));  
    //remove the last node from the heap  
    list.remove(list.size()-1);  
    int currentIndex = 0;
```

Heap.java

Heap

Heap.java

```
while (currentIndex < list.size()) {  
    int left = 2 * currentIndex + 1;  
    int right = 2 * currentIndex + 2;  
    //find the maximum of the left and right nodes  
    if (left >= list.size())  
        break; // no left child  
    int maxIndex = left;  
    E max = list.get(maxIndex);  
    if (right < list.size()) // right child exists  
        if (max.compareTo(list.get(right)) < 0)  
            maxIndex = right;
```

Heap

Heap.java

```
// swap if current is less than max
E current = list.get(currentIndex);
max = list.get(maxIndex);
if(current.compareTo(max) < 0){
    list.set(maxIndex, current);
    list.set(currentIndex, max);
    currentIndex = maxIndex;
}
else
    break; // the tree is a heap
}
return removedItem;
}
```

Heap

Test.java

```
public class TestBST {  
    public static void main(String[] args) {  
        Heap<String> heap = new Heap<>();  
        heap.add("Apple");  
        heap.add("Banana");  
        heap.add("Kiwi");  
        heap.add("Lemon");  
        heap.add("Orange");  
        heap.add("Strawberry");  
        heap.add("Watermelon");  
        System.out.println("Heap: " + heap.toString());  
        System.out.println("Removed: " + heap.remove());  
        System.out.println("Heap: " + heap.toString());  
        System.out.println("Heap contains Pear?: "  
                            + heap.contains("Pear"));  
    }  
}
```

Heap

◆ Performance of the Heap operations

Method	Complexity
Heap()	$O(1)$
size()	$O(1)$
clear()	$O(1)$
isEmpty()	$O(1)$
add(E)	$O(\log n)$
remove(E)	$O(\log n)$
contains(E)	$O(n)$
toString()	$O(n)$

Summary

- ◆ Binary trees - Special binary trees: BST, Heap
- ◆ Operations: Search, Add, Remove, Traversals
- ◆ Implementation - Linked Nodes for BST, ArrayList for Heap
- ◆ Performance of the operations on BST/Heap
- ◆ BST height determined by the order of insertion of the nodes
- ◆ Heap is a balanced binary tree
(height = $\log(\text{number of nodes})$)