

# Programs and Processes

*The concept of process is fundamental to the structure of modern computer operating systems. Its evolution in analyzing problems of synchronization, deadlock, and scheduling in operating systems has been a major intellectual contribution of computer science.*

*WHAT CAN BE AUTOMATED?:  
THE COMPUTER SCIENCE AND  
ENGINEERING RESEARCH STUDY,  
MIT Press, 1980*

# Objectives

- Define the term process and explain the relationship between processes and process control block
- Explain the concept of a process state and discuss the state transition the process undergoes
- List and describe the purpose of the data structures and data structure elements used by an OS to manage processes
- Assess the requirements for process control by the OS
- Understand the issues involved in the execution of OS code

- multiprogramming
  - one process at a time
- multiprocessing
  - many processes at a time
- process
  - execution environment
    - provided by operating system
  - communicate with OS and with each other
  - system processes
  - user processes

- library files
  - collection of functions
  - available to programs
  - library function
    - runtime library functions
- static libraries
  - during linking
- shared object libraries
  - at load time
  - shared by many processes

- library files

- .a
- static
- ar utility
- examine library functions
- create library functions

```
linux$ cc -c change_case.c
linux$ cc -c ascii.c
linux$ ar cr libmy_demo.a ascii.o change_case.o
```

```
FILE: main.c
#include <iostream>
#include "my_demo.h"
...
```

```
linux$ cc -o main main.c -L. -lmy_demo
```

```
FILE: ascii.c
char *
ascii( in start, int finish ) {
char *b = new cha(finish-start+1);
for ( int i=start; i <= finish; ++i)
    b( i-start ) = char ( i );
return b;
}

FILE: change_case.c
#include <ctype.h>
char *
change_case ( char *s ) {
char *t = &s[0];
while ( *t ) {
    if ( isalpha(*t) 0
        *t += islower(*t) ? -32 : 32;
    ++t;
}
return s;
}

FILE: my_demo.h
/*
prototypes for my_demo library functions
*/
#ifndef MY_DEMO_H

#define MY_DEMO_H
char * ascii ( int, int );
char * change_case ( char * );

#endif
```

- **system calls**
  - request service from operating system
  - OS performs work on behalf of caller
    - kernel
  - user mode
    - subset of instructions
  - system (privileged) mode
    - entire instruction set

- object code
  - library files and compiled code
    - combined at compile time
  - files not in standard library
    - specified at compile time

```
linux$ cc prgm.c -lm
```

- include math library libm.a

- system calls / library function return codes
  - external global **errno**
  - defined constants in **<sys/errno.h>**
  - always examine return codes!!!!
  - **perror**
    - to produce error message

```
void perror ( const char *s );
```

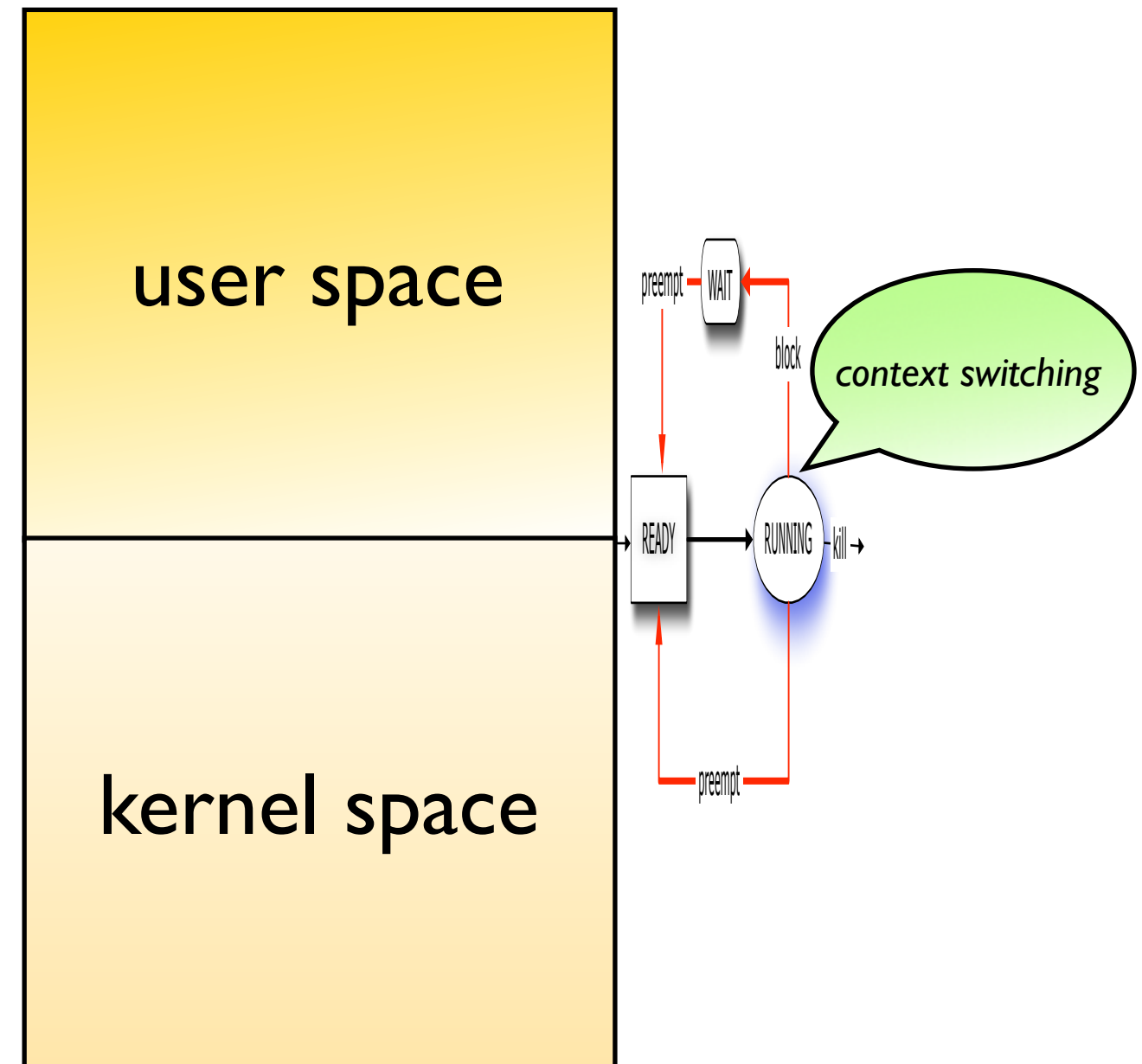
- review **strerror**



- **ELF (executable and linking format)**
  - marked as executable
    - program text
    - data
    - relocation information
    - symbol table
    - string table

old format: `a.out`  
Assembler OUtpuT Format

- System memory
  - user space
    - user processes
    - protected
    - run in user mode
  - kernel space
    - kernel execution
    - system processes
    - privileged mode



- process memory
  - private memory
  - divided into three segments
    - text
      - code
    - data
      - initialized variables
      - uninitialized variables
    - stack
      - automatic identifiers
      - function call information

- text segment
  - read-only segment
  - instruction segment
  - constant data
  - can be shared among processes
    - different instantiations of same code

- data segment
  - virtually contiguous to text segment
  - initialized/uninitialized data
  - may be expanded at execution time
    - `new`, `malloc`, `calloc`
- stack segment
  - automatic identifiers
  - register variables
  - function call information
  - grows towards uninitialized data segment

- u-area
  - maintained by operating system
    - one per process
  - specific process information
    - open files
    - current directory
    - signal actions
    - accounting information
    - a system stack segment
      - system calls
  - process can access information
    - via system calls

## 1.8 Process Memory

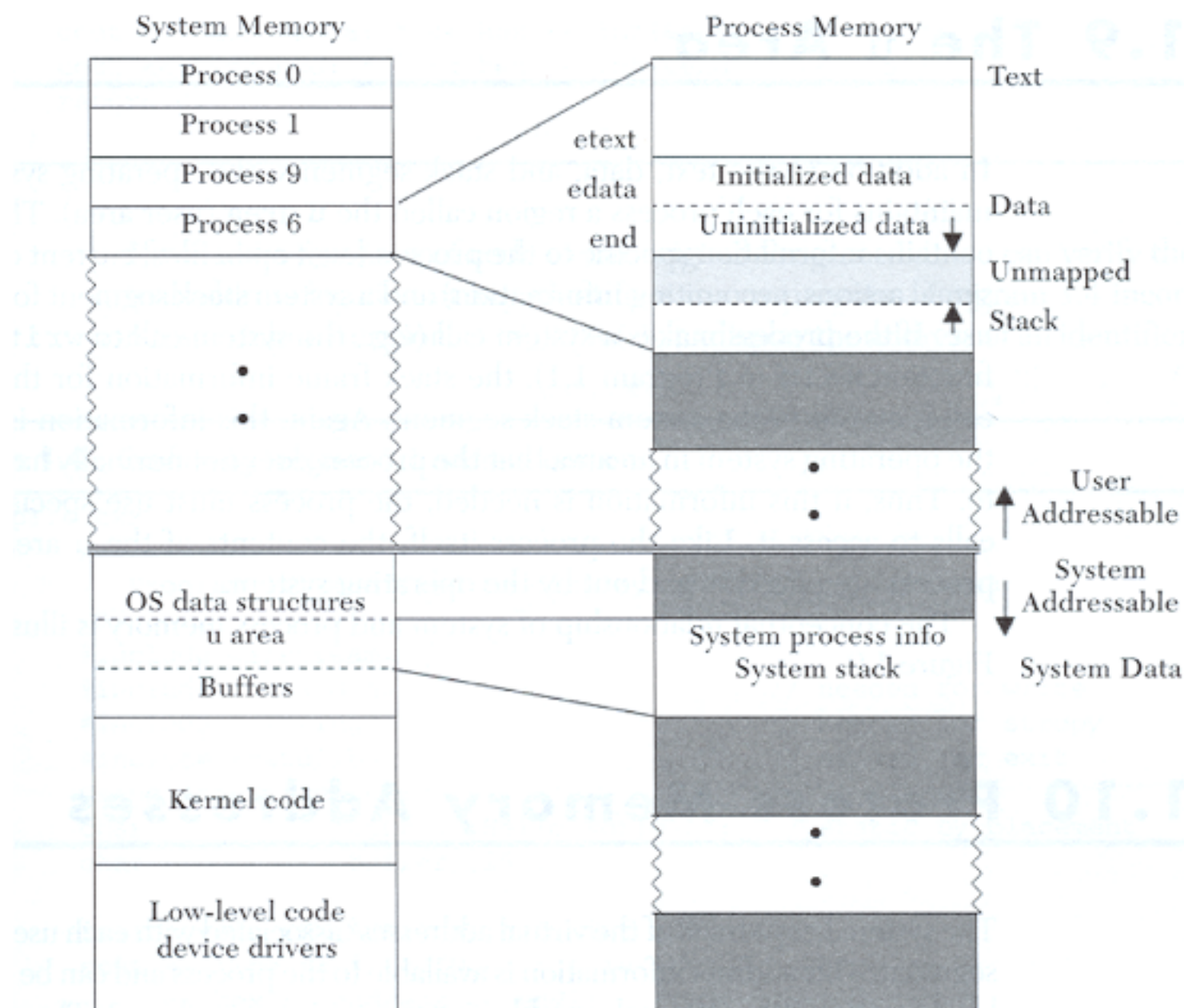


Figure 1.9 System and process memory.

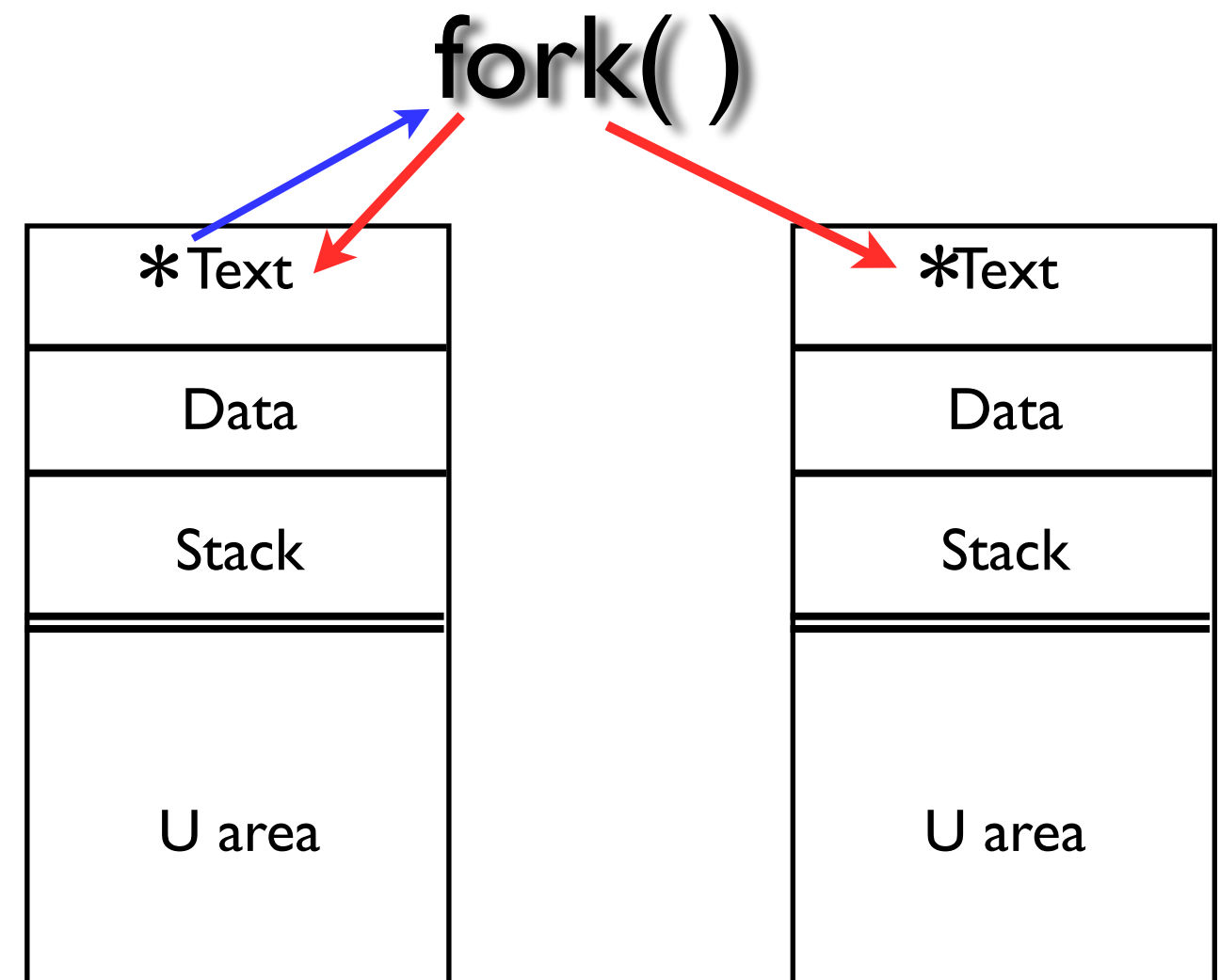
- `fork ( )`
  - parent process
    - process executing the `fork ( )`
  - child process
    - created by `fork ( )` system call

```
pid_t fork ( void );
```

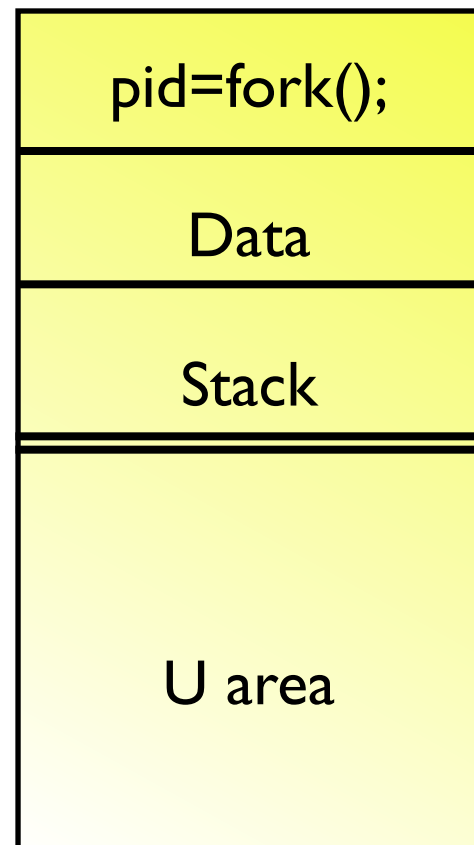
Constant	error Message	Explanation
EAGAIN	Resource temporarily unavailable	No memory to copy parents page table and allocate task structure
ENOMEM	Cannot allocate memory	Insufficient swap space



- returns twice
  - in parent
    - returns child PID
  - in child
    - returns 0
- process can identify itself
  - parent or child

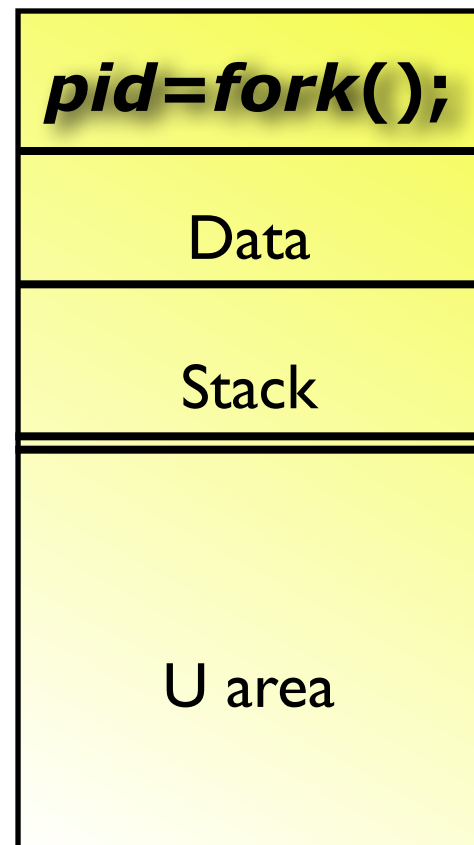


process id=789

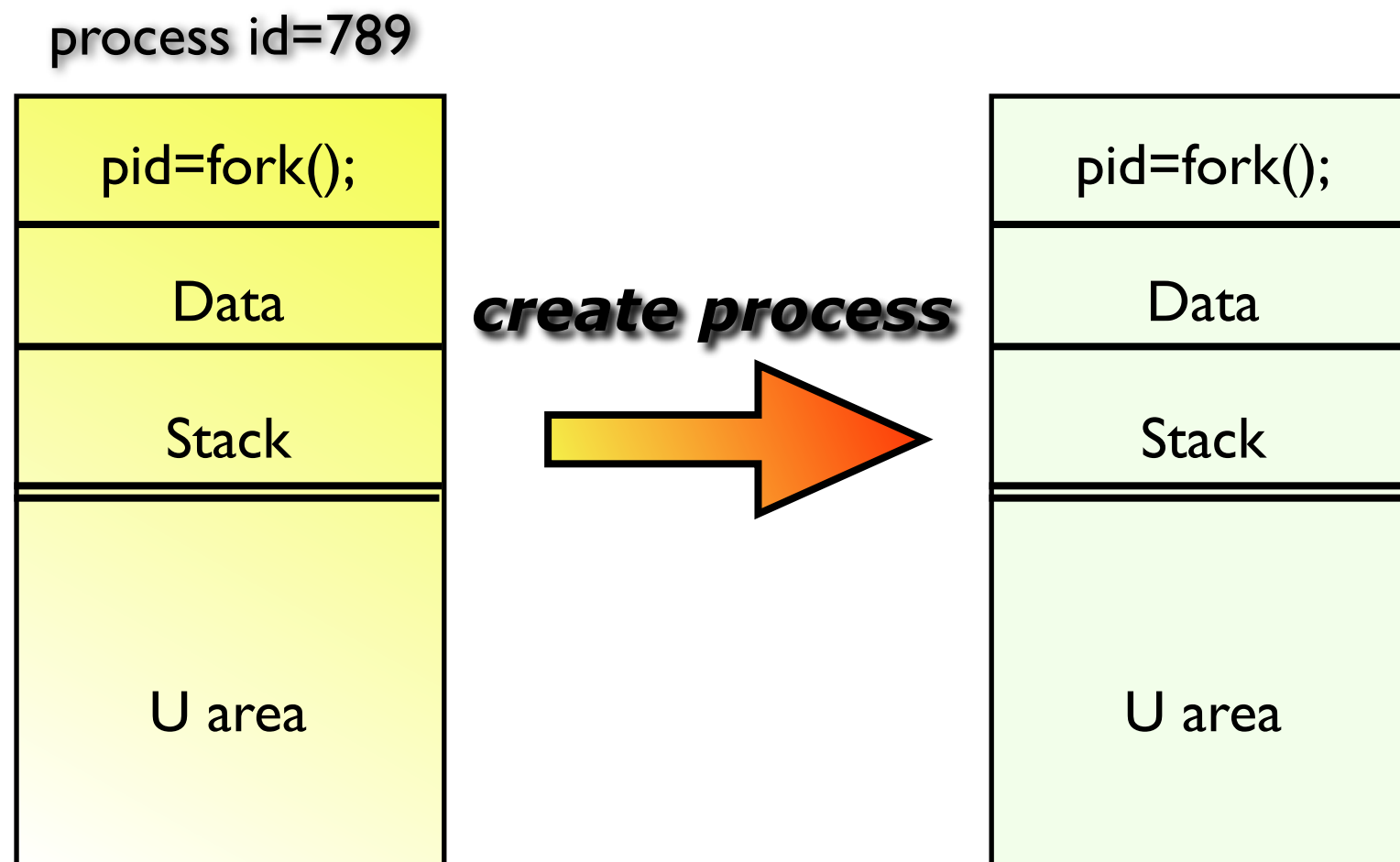


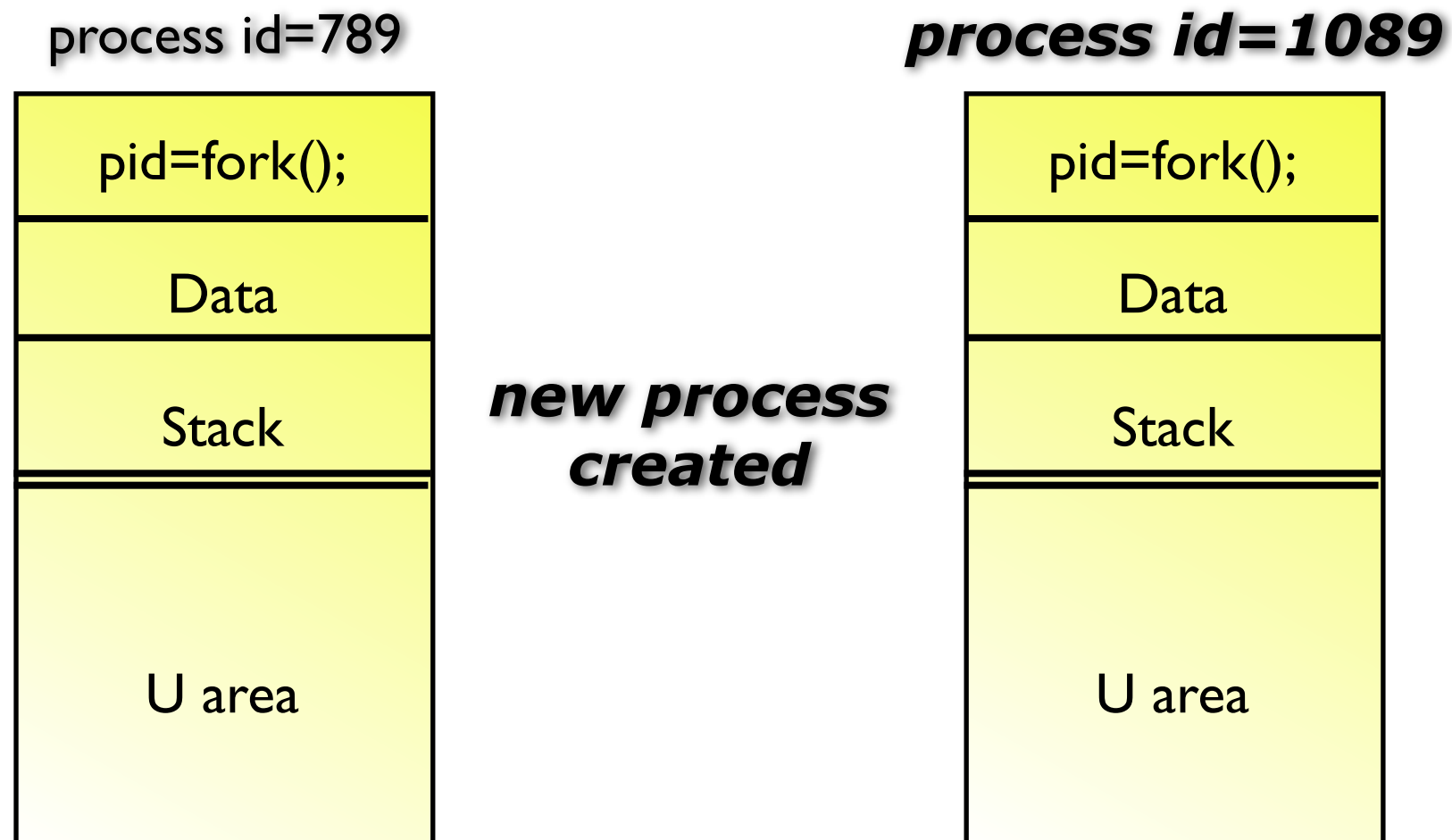
process executes a fork():

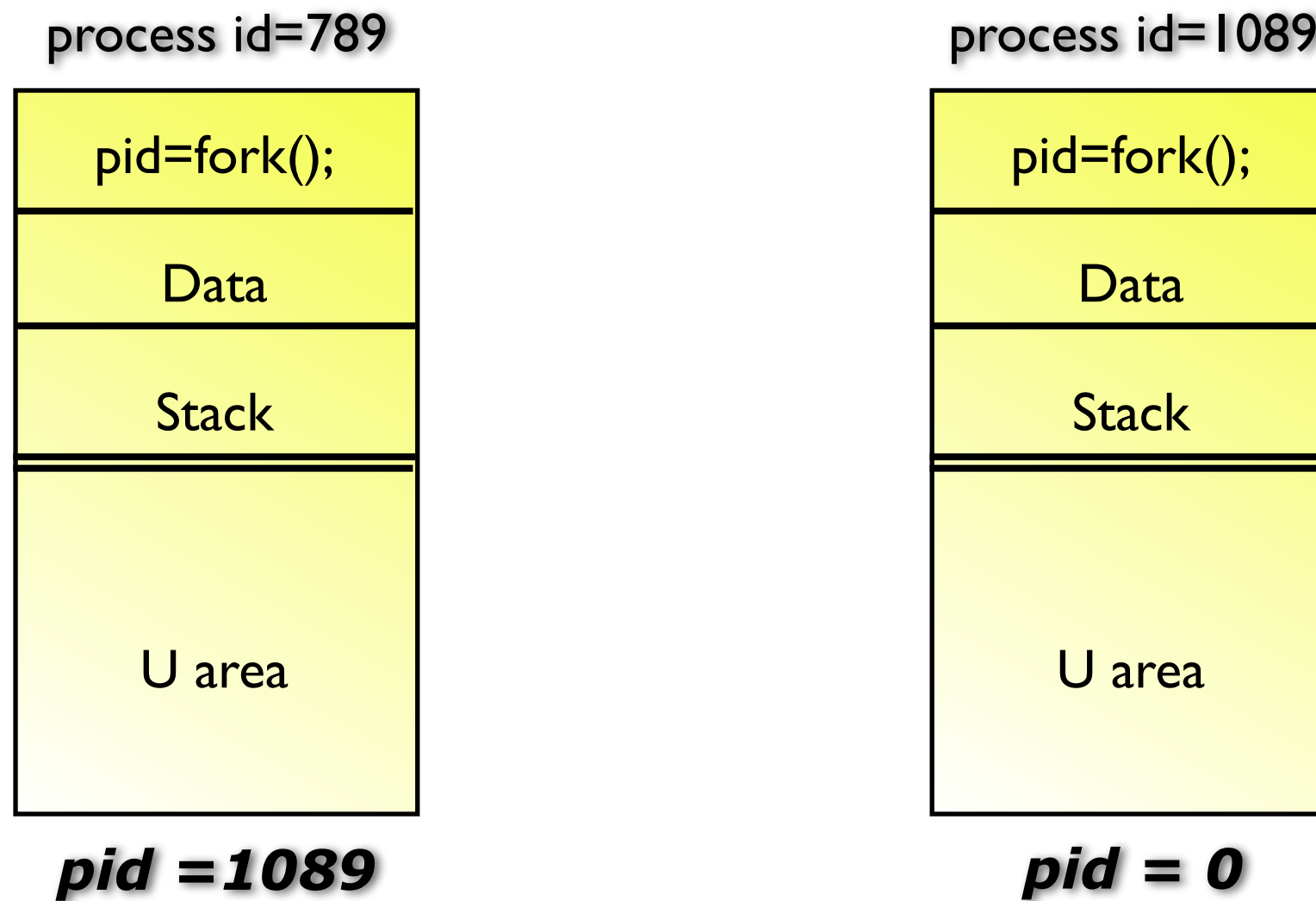
process id=789



fork system call







```
pid = fork ();  
if ( pid == 0 ){  
    // in child process  
}  
// in parent process
```

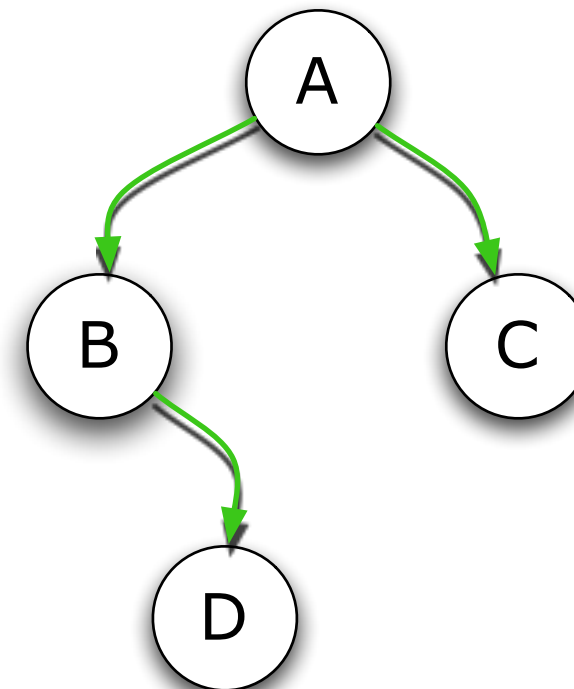
- Unique to process
  - PID
  - PPID
  - System imposed process limits are reset
    - CPU allocation time
  - File record locks are reset
  - Signal actions



## ● Share

- real UID
- real GID
- effective UID
- effective GID
- process GUI
- controlling terminal
- set-user-id flag and set-group-id flag
- current working directory
- root directory
- file mode creation mask
- signal mask and dispositions
- the close-on-exec flag for any open file descriptors
- attached shared memory segments
- open file descriptors
- Environment information

- processes may create other processes
  - parent - child relation
    - A parent of B and C
    - B parent of D



- **Process States**

- **Ready**

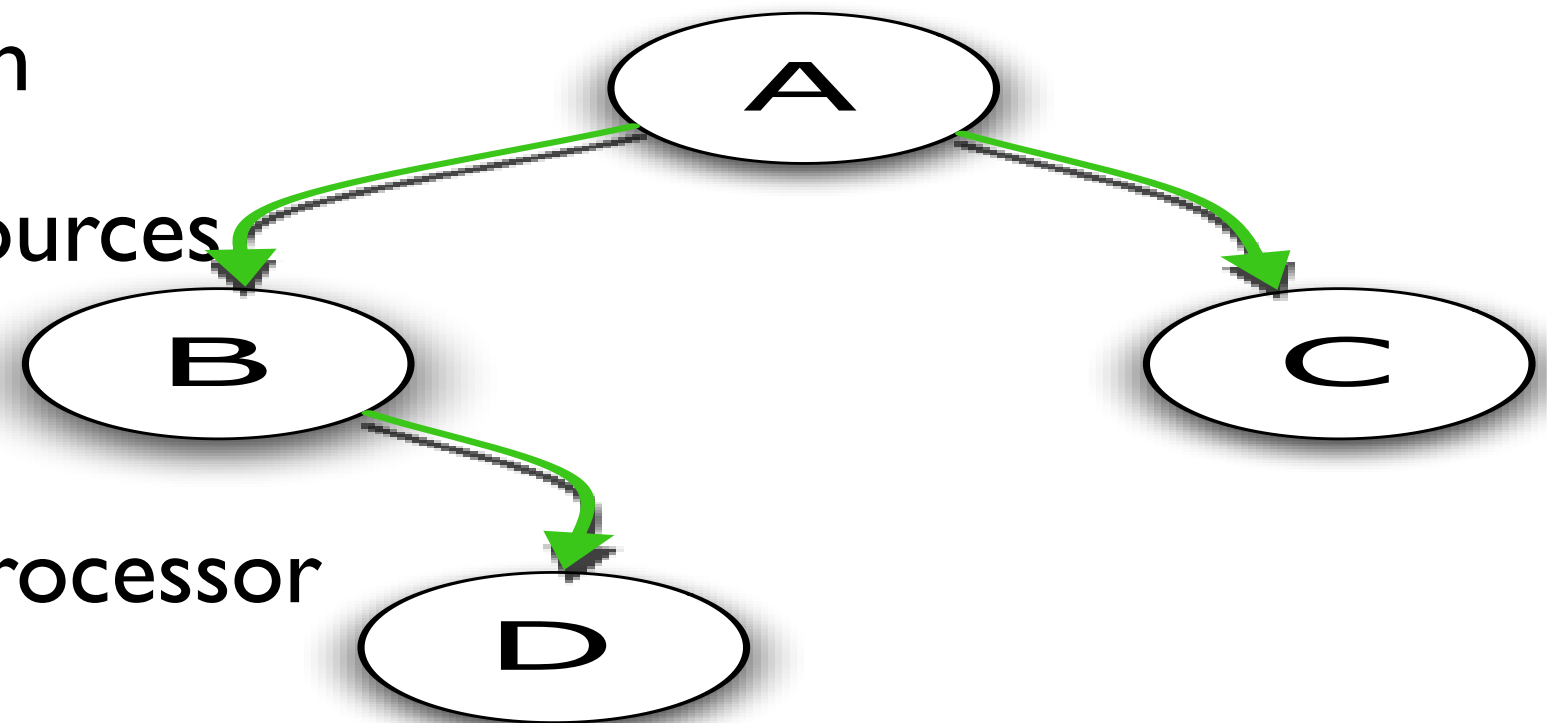
- ready for execution
- has all needed resources

- **Running**

- in control of the processor

- **Wait**

- waiting for an event to occur
  - waiting for a resource



- **Process Control Block (PCB)**
  - data structure created by the Operating System
    - one entry per process
      - create a process  $\longrightarrow$  entry in PCB
  - all information about a process
    - process id (PID)
    - dispatching information
      - priority
        - CPU time
    - resources
      - allocated
      - outstanding

A process is known through its PCB