Pipes

```
Pipes
 interprocess communication
 system provided facilities
 passing information
 read/write
"special file" with limited capacity
 PIPE SIZE
          <linux/limits.h>
FIFO
```

Pipes

write end

read end

system maintains pointer to current location

system manages current location

no seek permitted



- Synchronization
 - done by system
 - atomic read/write
 - no seek
 - blocks on
 - write full
 - empty read
 - open for read with no open for write
 - use read/write system call
 - unbuffered I/O

Synchronization

```
ssize_t write (int fd, const void *buf, size_t count);
```

write count bytes from buffer to file descriptor fd returns number of bytes written

-1 on error

sets errno

errors

- i/o error
- bad file pointer
- resource temporarily unavailable
- bad address
- file too large
- no space left on device
- numerical result out of range

write

- append to end of pipe
- writes of PIPE_BUF size or less are guaranteed to not be interleaved with other write requests to same pipe
- write to a pipe not opened for read generates a SIGPIPE signal (default action is to terminate)
- if both ON_NONBLOCK and O_NDELAY flags are clear, write blocks if device is busy
 - set in fcntl system call
- if both ON_NONBLOCK and O_NDELAY flags are not clear write will not block

• read

```
ssize_t read (int fd, void *buf, size_t count);
reads count bytes
unbuffered I/O read system call
returns
  actual number of bytes read
  zero if at end of pipe
  -1 on error
  sets errno
```

- possible errors
 - I/O error
 - No such file or address)file descriptor)
 - Bad file descriptor (not opened for reading)
 - bad address (buffer)

- read
 - initiated from current position
 - no seeks
 - If both ON_NONBLOCK and O_NDELAY flags are clear read blocks on empty until data is written or pipe is closed
 - if pipe is not opened for writing returns a 0

- unnamed pipes
 - used only by related processes
 - parent/child
 - child/child
 - exists as long as they are in use
- named pipes
 - exists as directory entries
 - have file access permissions
 - can be used by unrelated processes

- unnamed pipes
 - pipe system call

```
int pipe( int filedes[2]);
```

- Returns a pair of integer file descriptors
 - filedes[0]
 - filedes[1]
 - they reference two data streamsLinux
 - half duplex: unidirectional
 - Solaris
 - full duplex: bidirectional

- unnamed pipes
 - half duplex
 - filedes[0]
 - reading
 - filedes[1]
 - writing

- full duplex
 - agree on use of read/write ends
 - both opened for reading and writing

possible errors

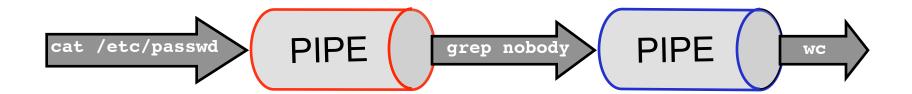
Operating Systems

- file table overflow
- too many opened files
- bad address (invalid file descriptor)

```
int main(int argc, char *argv[]) {
                 f des[2];
  int
  static char
                 message[BUFSIZ];
  if (argc != 2) {
   cerr << "Usage: " << *arqv << " message\n";</pre>
   return 1;
  if (pipe(f des) == -1) {
                                         // generate the pipe
   perror("Pipe"); return 2;
  switch (fork( )) {
  case -1:
   perror("Fork"); return 3;
                                       // In the child
  case 0:
   close(f des[1]);  // close write end of pipe
    if (read(f des[0], message, BUFSIZ) != -1) {
      cout << "Message received by child: [" << message</pre>
           << "1" << endl:
      cout.flush();
    } else {
     perror("Read"); return 4;
    break;
  default:
                                       // In the Parent
    close(f des[0]); // close read end of pipe
    if (write(f des[1], argv[1], strlen(argv[1])) != -1) {
      cout << "Message sent by parent : [" <<</pre>
           argv[1] << "]" << endl;
                                                       $ ./a.out Hello
      cout.flush();
                                                       Message sent by parent : [Hello]
    } else {
                                                       Message received by child: [Hello]
     perror("Write"); return 5;
                            pipe to send first argument from parent to child
 return 0;
```

- I/O redirection
 - associate standard input/output to pipe

cat /etc/passwd grep nobody wc



dup and dup2

```
int dup ( int oldfd );
```

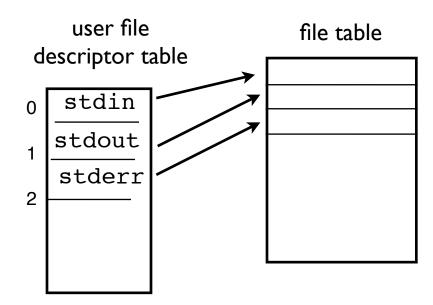
- duplicates an opened file descriptor
- new descriptor
 - references file system table entry for next available nonnegative file descriptor
 - always returns next lowest available file descriptor
 - share the same file pointer (offset) as original
 - same access control as original
 - share locks
 - remain open across exec
 - do not share close-on_exec flag

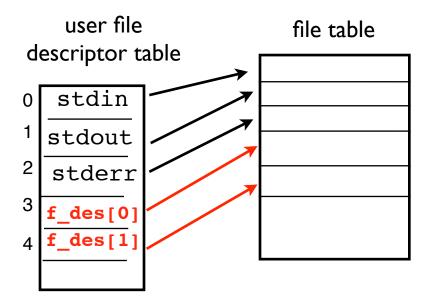
dup

*always return the next lowest available descriptor

race condition between close and dup

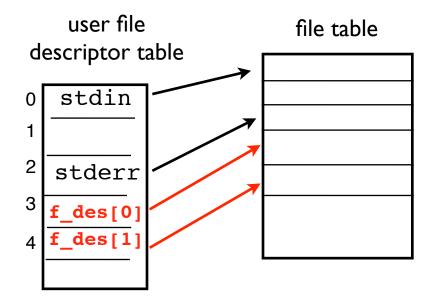
for example a signal-catcher that closes a file



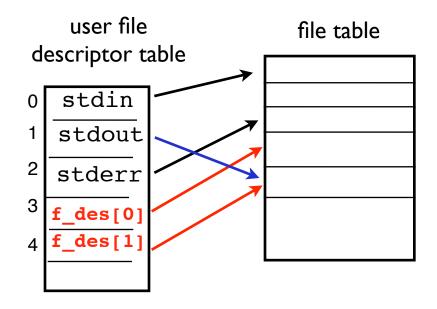


```
int f_des_[2];
pipe(f_des);
```

pipes

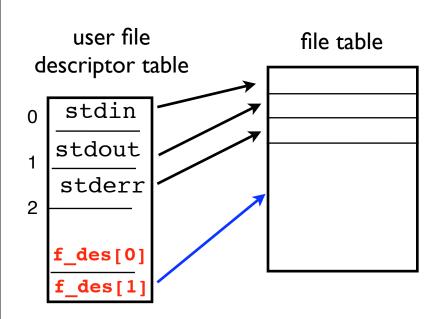


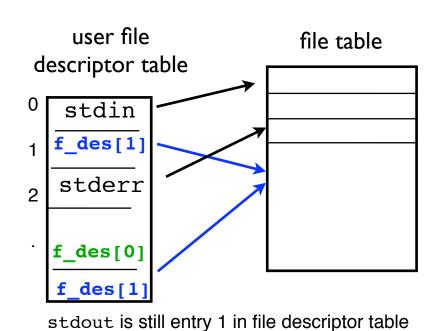
```
int f_des_[2];
pipe(f_des);
close(fileno(stdout) ); // close standard output
```



stdout is still entry 1 in file descriptor table

pipes

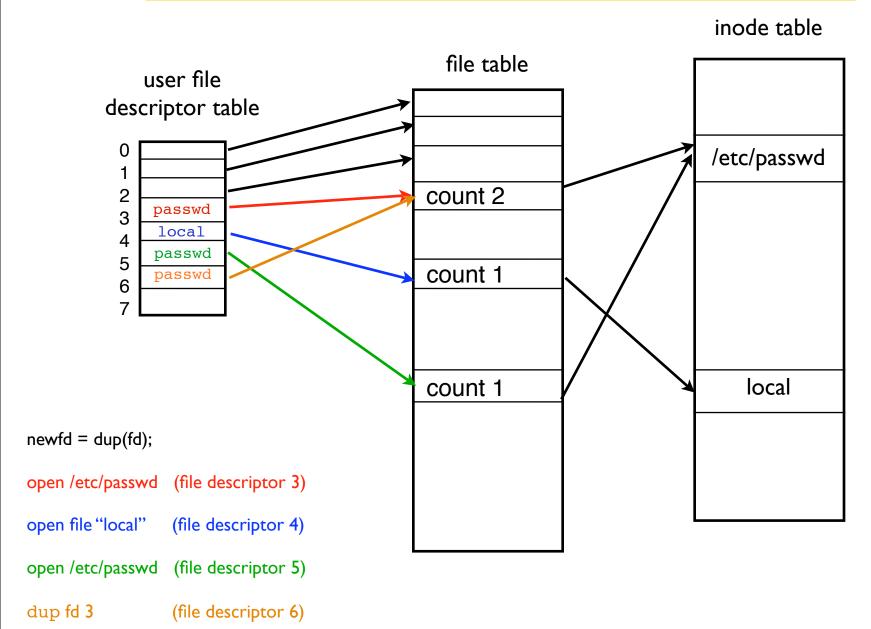




- newfd = dup(fd)
 - open /etc/passwd (file descriptor 3)
 - open file "local" (file descriptor 4)
 - open /etc/passwd (file descriptor 5)
 - dup fd 3 (file descriptor 6)

user file descriptor table

0	
1	
2	
3	passwd
4	local
5	passwd
6	passwd
7	



• dup2

```
int dup2 ( int oldfd, int newfd );
```

Closes and duplicates file descriptor atomic operation

if **newfd** is already open closed before duplicating

5

child

```
/* A home grown last | sort cmd pipeline. */
enum { READ, WRITE };
int
                                              parent
main( ) {
                                           stdin
                                                           stdin
  int
           f des[2];
                                           stdout
                                                           stdout
  if (pipe(f des) == -1) {
                                           stderr
                                                           stderr
    perror("Pipe");
                                                           f des[0]
                                           f des[0]
                                                           f des[1]
    return 1;
                                            f des[1]
                                        5
  switch (fork( )) {
  case -1:
    perror("Fork");
    return 2;
  case 0:
                                         // In the child
    dup2( f des[WRITE], fileno(stdout));
    close(f des[READ] );
    close(f des[WRITE]);
    execl("/usr/bin/last", "last", (char *) 0);
    return 3:
  default:
                                         // In the parent
    dup2( f des[READ], fileno(stdin));
    close(f des[READ] );
    close(f des[WRITE]);
    execl("/bin/sort", "sort", (char *) 0);
    return 4;
  return 0;
```

- Summary unnamed pipes
- Create pipes
- Generate child process
- Close/duplicate file descriptors
- Close unneeded ends of pipe
- Communicate

popen pclose system calls

```
FILE *popen ( const char *command, const char *type);
int pclose ( FILE *stream);
```

fork child process duplicate file descriptors pass command

```
popen
```

duplicate file descriptors fork child process which will exec /bin/sh /bin/sh executes passed shell command duplicate file descriptors

W: parent process can write to standard input of new shell file pointer referenced by popen

R: parent process can read from standard output of new shell

fully buffered pass command argument

poclose

close data stream penned by popen returns exit status of shell command

```
/* Using the popen and pclose I/O commands*/
                                                                       write
                                                         read
int
main(int argc, char *argv[]) {
  FILE
        *fin, *fout;
          buffer[PIPE BUF];
  char
                                                                      #a.out last more
  int
           n;
  if (argc < 3) {
    cerr << "Usage " << arqv << "cmd1 cmd2" << endl;</pre>
    return 1;
                                        forks child that executes command argv[1] output read by parent
                                        with file pointer fin
       = popen(argv[1], "r");
  fin
                                        forks child that executes command argv[2] input read from
  fout = popen(argv[2], "w");
                                        parent's output
  fflush(fout);
  while ((n = read(fileno(fin), buffer, PIPE BUF)) > 0)
    write(fileno(fout), buffer, n);
  pclose(fin);
  pclose(fout);
  return 0;
```

- named pipes (FIFO)
 - directory entry
 - permissions
 - can be used by unrelated processes
 - data stored in kernel not in file system
 - creation
 - mknod name p
 - mkfifo name p

mknod: generate special files (block character device files in /dev) nonpriviledged users can only generate named pipes

%mkfifo APIPE p

prw-r--r--

1 dcanas dcanas

0 Sep 16 11:21 APIPE

```
%mkfifo APIPE p
```

prw-r--r- 1 dcanas dcanas 0 Sep 16 11:21 APIPE

% cat copy.c > APIPE & [1] 287

% cat < APIPE

this is the contents of copy.c a test file for pipes last line

[1]+ Done

System call

int mknod(const char *pathmane, mode_t mode, dev_t dev);

- creates a file descriptor by pathname
- file types and permissions are ORed
- dev argument only for character, block devices
 - use 0 for FIFO
- File types;

```
S_FIFO FIFO special (non privileged users)
S_IFCHR character special
S_IFDIR directory
S_IFBLK block special
S_IFREG ordinary file
```

System call

```
int mkfifo(const char *pathmane, mode_t mode );
```

uses mknod to create a FIFO