

# Discrete Optimization

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Non-Linear Optimization

April 18, 2013



# Framework

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- A general introduction to discrete optimization
- Discrete Optimization in Radiation Therapy
- Discrete Optimization in Medical Imaging
- Bipartite Matching

# Definition of discrete optimization

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- **Discrete optimization** is a branch of optimization in applied mathematics and computer science, as opposed to continuous optimization, the variables used in the mathematical program (or some of them) are restricted to assume only **discrete values**, such as the integers
- Two notable branches of discrete optimization are:
  - combinatorial optimization, which refers to problems on graphs, matroids and other discrete structures
  - integer programming

# Definition of discrete optimization

- **combinatorial optimization** is a topic that consists of finding an optimal object from a finite set of objects
- An **integer programming** problem is a mathematical optimization in which some or all of the variables are restricted to be integers

# Some examples in discrete optimization

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- Integer linear programming
- Set cover problem
- Knapsack problem
- Graph theory
  - Minimum spanning tree
  - Vertex cover problem
  - Traveling salesman problem (Hamiltonian circuit)
  - Shortest path problem
- Scheduling problem
- Maximum flow problem

# Integer linear programming (NP-hard)

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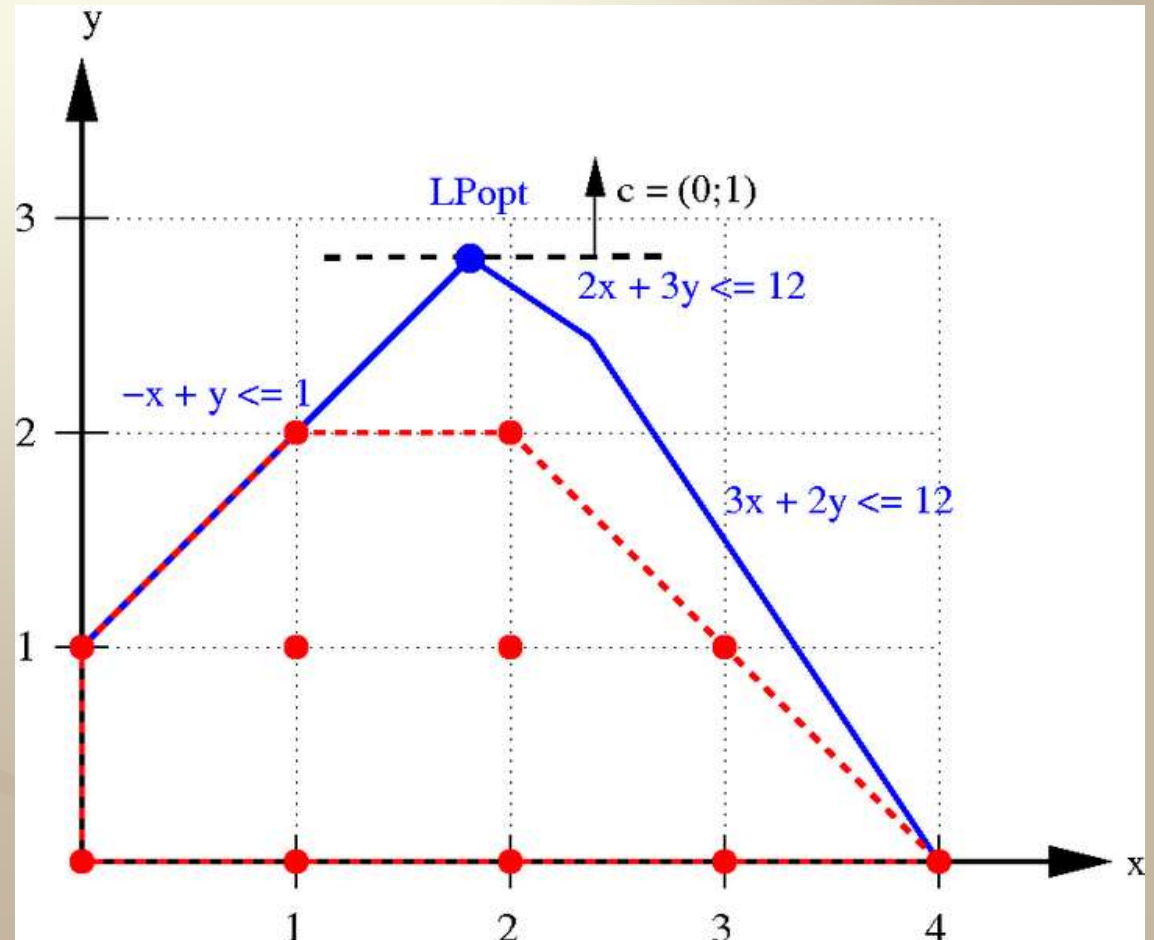
- An ILP in canonical form is expressed as:

$$\begin{array}{ll}\text{maximize} & c^T x \\ \text{subject to} & Ax \leq b \\ & x \geq 0 \\ & \text{and } x \text{ integer}\end{array}$$

# Integer linear programming (NP-hard)

- An example:

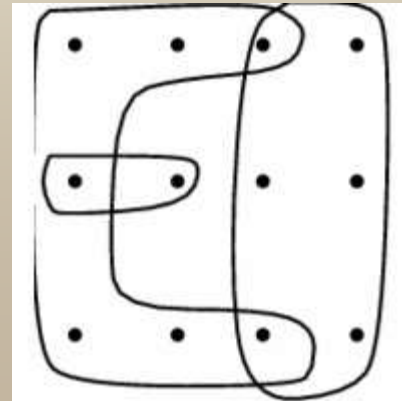
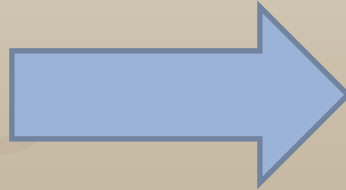
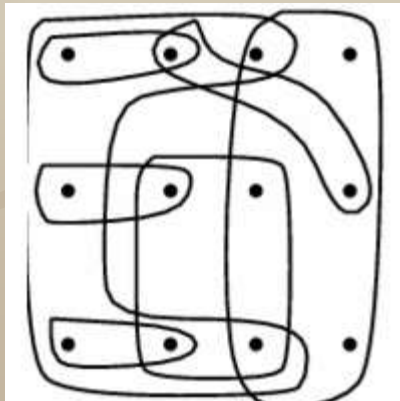
$$\begin{aligned} \max \quad & y \\ \text{s.t.} \quad & -x + y \leq 1 \\ & 3x + 2y \leq 12 \\ & 2x + 3y \leq 12 \\ & x, y \geq 0, \text{ integer} \end{aligned}$$



# Set cover problems (NP-hard)

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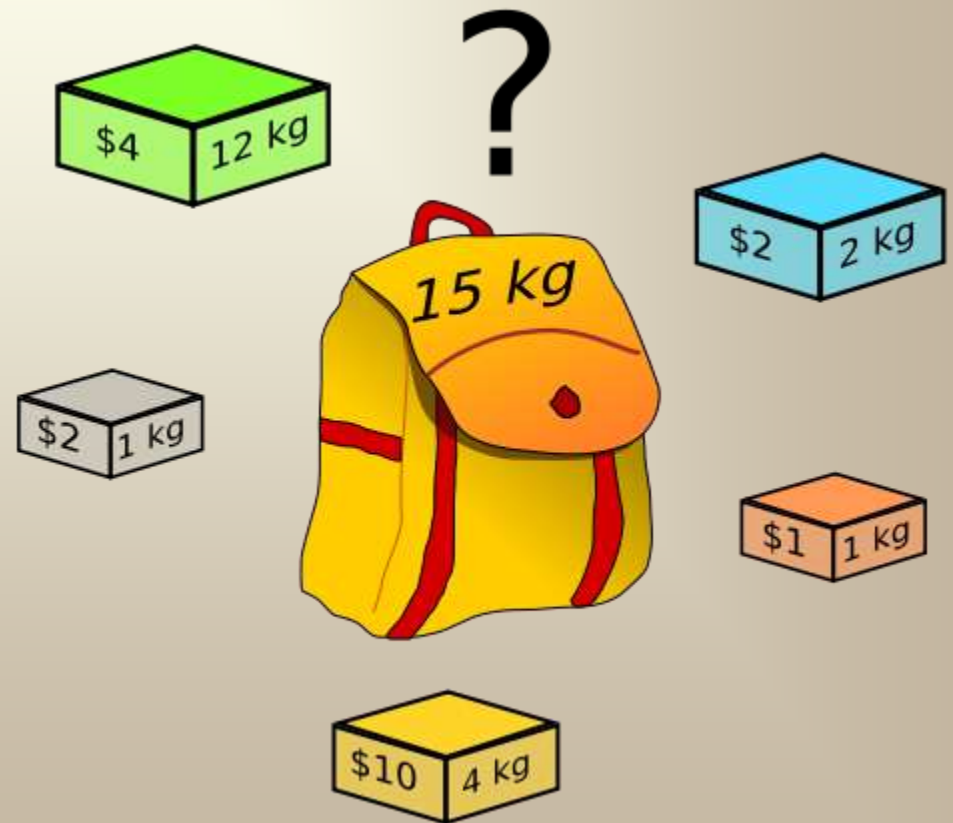
- Given a set of elements  $\{1, 2, \dots, n\}$  (called the universe) and a set  $S$  of  $m$  sets whose union equals the universe, the set cover problem is to identify the smallest subset of the union of which contains all elements in the universe.





# Knapsack problems (NP-hard)

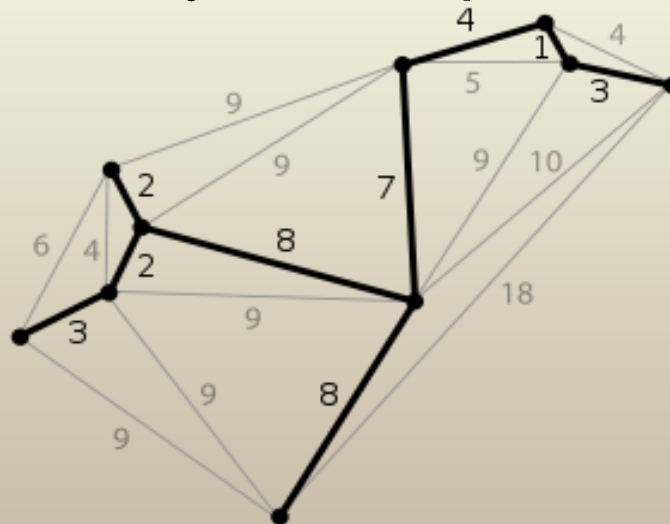
- Given a set of items, each with a weight and a value, determine the number of each item to include in a collection so that the total weight is less than or equal to a given limit and the total value is as large as possible.
- Backtracking Algorithm



# Minimum spanning tree( $O(n^2)$ )

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- A **minimum spanning tree (MST)** is then a spanning tree with weight less than or equal to the weight of every other spanning tree.

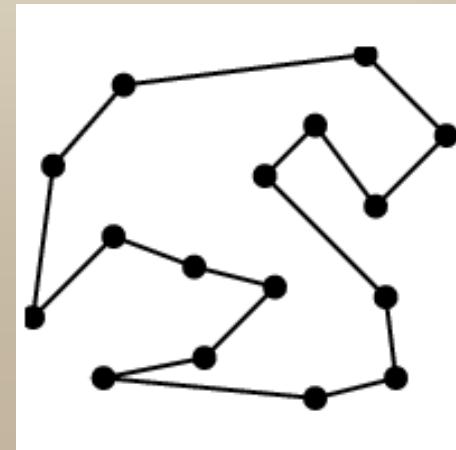
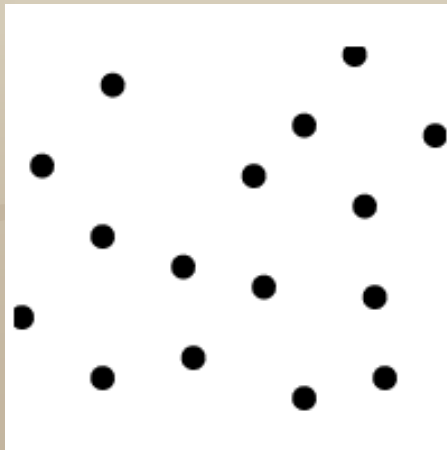


- Kruskal's algorithm (greedy Alg.)

# Traveling salesman problem (NP-hard)

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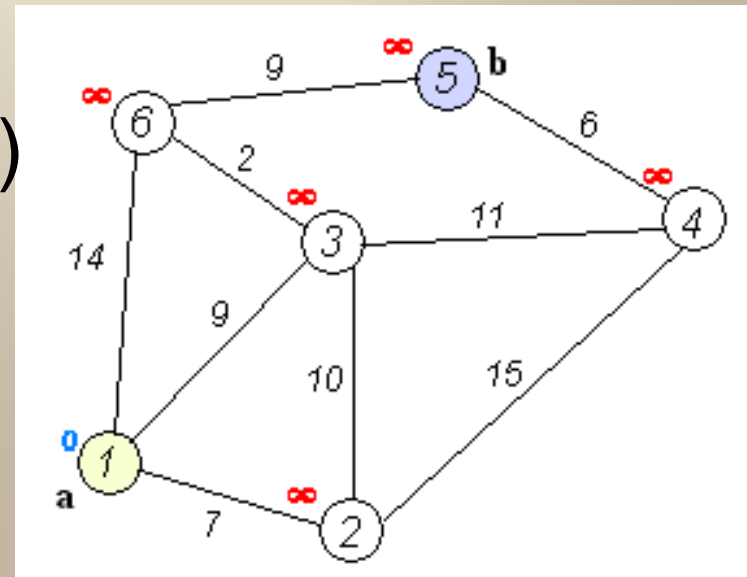
- The **travelling salesman problem (TSP)** asks the following question: given a list of cities and the distances between each pair of cities, what is the shortest possible route that visits each city exactly once and returns to the origin city?



# Shortest path problem

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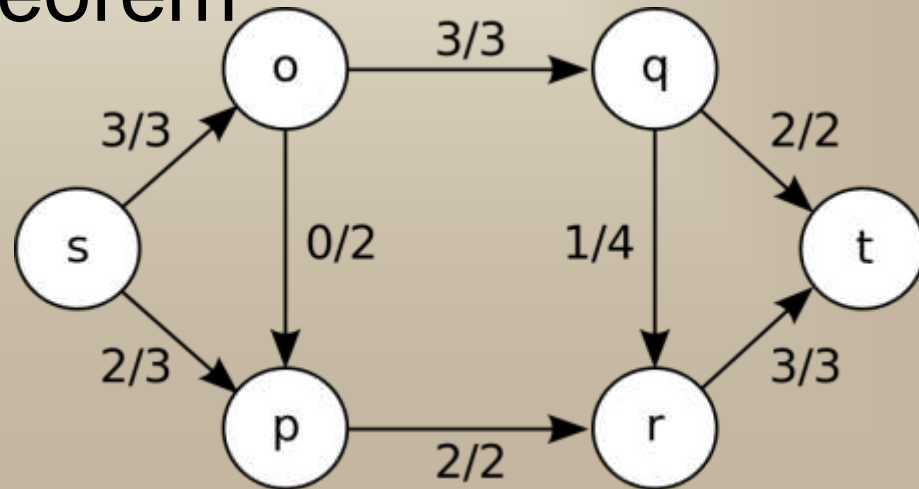
- The shortest path problem is the problem of finding a path between two vertices (or nodes) in a graph such that the sum of the weights of its constituent edges is minimized
- Dijkstra's algorithm( $O(n^2)$ )
- Floyd's algorithm( $O(n^3)$ )



# Maximum flow problem (NP-hard)

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- **Maximum flow problem** is to find a feasible flow through a single-source, single-sink flow network that is maximum
- Max-flow min-cut theorem



# Uses of Discrete Optimization

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- Resource management
- Supply Chain
- Off-shore oil drilling infrastructure planning
- Medical
  - Computational molecular biology
  - Pharmaceutical testing schedules
  - Radiation Treatment Planning
  - Medical Imaging

# Radiation Treatment for Cancer

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- Types
  - Photons (x-rays and  $\gamma$ -rays)
  - Protons
- Delivery
  - External beam
  - Intensity Modulated Radiation Treatment (IMRT)
  - Gamma Knife
  - Brachytherapy (seeds)

# Goals of Radiation Therapy

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- 50% isodose line coverage of the target volume
- Minimize the non-target volume covered by external beams
- Minimize dose to near-by sensitive structures
- Minimize the number of shots → maximize the number of treated patients
- Simplicity



# Optimization

- Calculate Dose (Gy) to each voxel (i,j,k) in the target volume and surrounding volume

$$\begin{aligned} Dose(i, j, k) &\leq U_{\mathcal{R}} + M * Exceed(i, j, k) \quad \forall (i, j, k) \in \mathcal{R} \\ \sum_{(i,j,k) \in \mathcal{R}} Exceed(i, j, k) &\leq \beta_{\mathcal{R}} * card(\mathcal{R}). \end{aligned}$$

$U_{\mathcal{R}}$  – Threshold dose

$M$  – constant

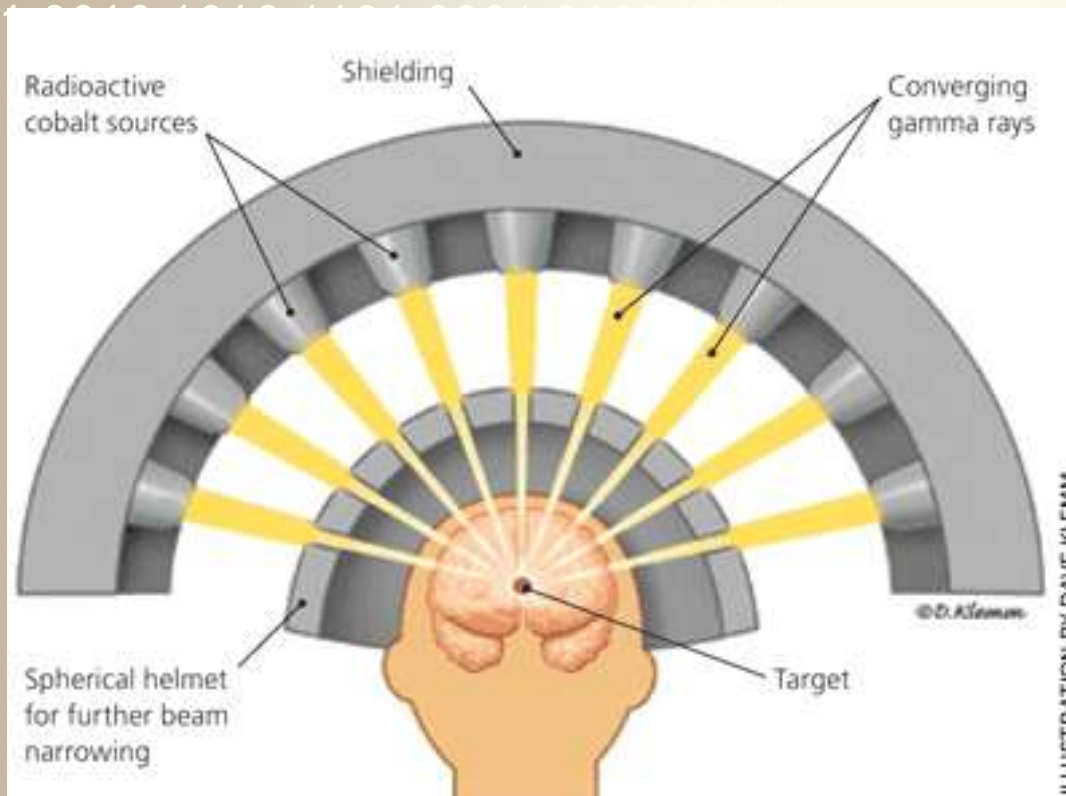
$\mathcal{R}$  – target region

$\beta_{\mathcal{R}}$  – percent over-dose allowed (tolerance)

$Exceed$  – binary variable

Mixed Integer Programming to solve

# Gamma Knife Radio-surgery



Co-60 source emits gamma rays by radioactive decay

Allows for much fewer fractions of radiation

Main application to brain tumors

Delivered in several “shots” – time and weight are controlled

# Gamma Knife Optimization

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- Two steps to optimization
  - Continuous variables optimized using quassi-Newton method
    - position and weight of shots
  - Discrete variables optimized using simulated annealing
    - number of shots and collimator size

# Gamma Knife Optimization

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- Mixed Integer Programming
  - Very large solution space

$$\min \sum_{(i,j,k) \in \mathcal{N}} Dose(i,j,k)$$

$$\text{subject to } Dose(i,j,k) = \sum_{(s,w) \in \mathcal{S} \times \mathcal{W}} t_{s,w} D_w(x_s, y_s, z_s, i, j, k)$$

$$\Theta \leq Dose(i,j,k) \leq 1, \quad \forall (i,j,k) \in \mathcal{T}$$

$$n = \sum_{(s,w) \in \{1, \dots, n\} \times \mathcal{W}} H_{\alpha}(t_{s,w})$$

$$t_{s,w} \geq 0.$$

# IMRT

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# IMRT

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- Mixed Integer Programming
  - Optimize beam angles, wedge orientations, beam intensities
  - Long run time to solve

$$\begin{aligned} \min_w \quad & \lambda_t \|D_T - \theta e_T\|_\infty + \lambda_s \frac{\|(D_S - \phi e_S)_+\|_1}{\text{card}(\mathcal{S})} + \lambda_n \frac{\|D_N\|_1}{\text{card}(\mathcal{N})} \\ \text{s.t.} \quad & D_\Omega = \sum_{A \in \mathcal{A}, F \in \mathcal{F}} w_{A,F} \mathcal{D}_{A,F,\Omega}, \quad \Omega \in \mathcal{T} \cup \mathcal{S} \cup \mathcal{N}, \\ & \frac{u}{\rho_A} \psi_A \geq w_{A,0} + \tau_1 \sum_{F \in \mathcal{F} \setminus 0} w_{A,F} \\ & K \geq \sum_{A \in \mathcal{A}} \psi_A, \\ & w_{A,F} \geq 0, \quad \forall A \in \mathcal{A}, \\ & \psi_A \in \{0, 1\}, \quad \forall A \in \mathcal{A}. \end{aligned}$$

# Brachytherapy Optimization

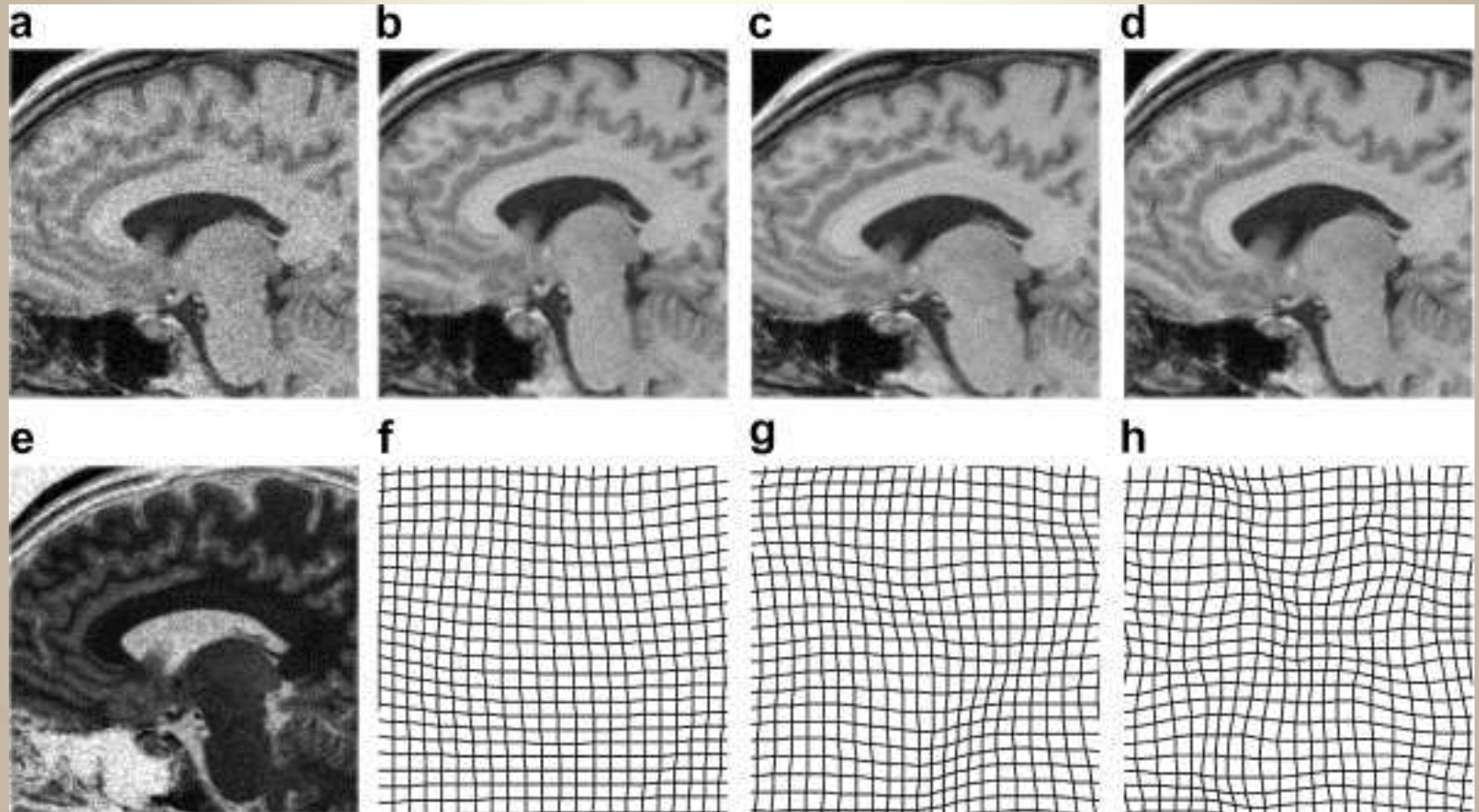
$$Dose(i, j, k) = \sum_{(r,s,t) \in T} D_{r,s,t}(i, j, k) * Seed(r, s, t).$$

$$\begin{aligned} \min \quad & \alpha * \sum_{(i,j,k) \in T} Underdose(i, j, k) + \beta * \\ & \sum_{(i,j,k) \in OAR} Overdose(i, j, k) + \gamma * \sum_{(i,j) \in template} Needle(i, j) \end{aligned}$$

- Subject to constraints related to limiting needles to place the seed, limited dose to sensitive organs
- Mixed integer programming and branch-and-bound methods



# Discrete Optimization in Medical Imaging



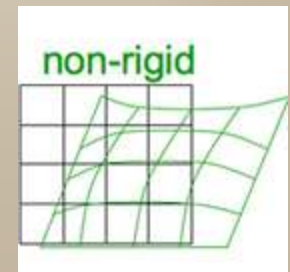
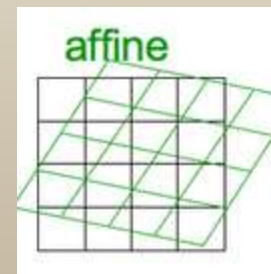
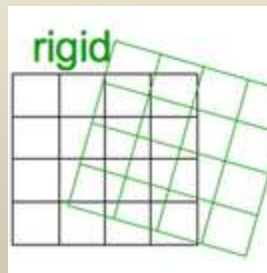


# Basics of Image Registration

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- Images must be registered in order to compare multiple subjects in a group
- From the most simple to the most complex:

- Rigid
- Affine
- Non-rigid
  - Varying DOF



# Why does medical imaging need optimization?

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- Curse of Dimensionality
  - 1 million + voxels
- Curse of non-convexity
  - Ill-posed problem (more parameters than constraints)
- Curse of non-linearity
  - Not all registrations are affine
- Curse of modularity
  - Solutions for very specific problems

# What have they already tried?

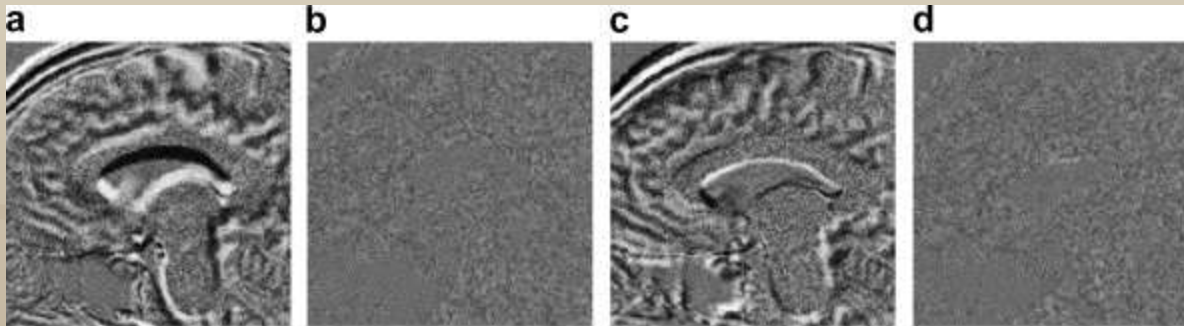
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- Powell's (conjugate direction) method
- Downhill Simplex
- Levenberg-Marquardt
  - Currently used by SPM8
- Newton-Raphson
- Stochastic search methods
- Gradient descent methods
- Quasi-exhaustive search methods

# What's going on now?

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- Glocker et al. recently published a paper using discrete optimization techniques for image registration
  - Model registration as Markov Random Field
  - MRF is optimized using the primal-dual schema
  - Performs well on MRI brain image registration



# The Primal-Dual Schema

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- Main goal: Minimize gap between the primal and dual costs

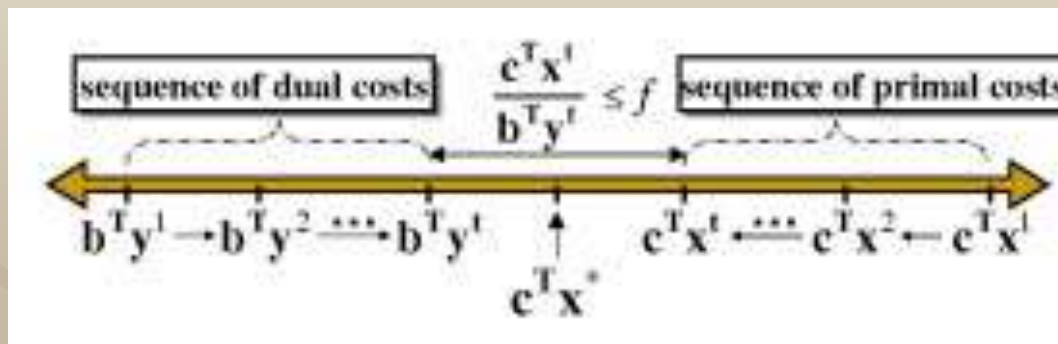
- Primal problem is NP hard

Primal :  $\min c^T x$

s.t.  $Ax = b, x \geq 0$

Dual :  $\max b^T y$

s.t.  $A^T y \leq c$



# Estimating Biomarkers

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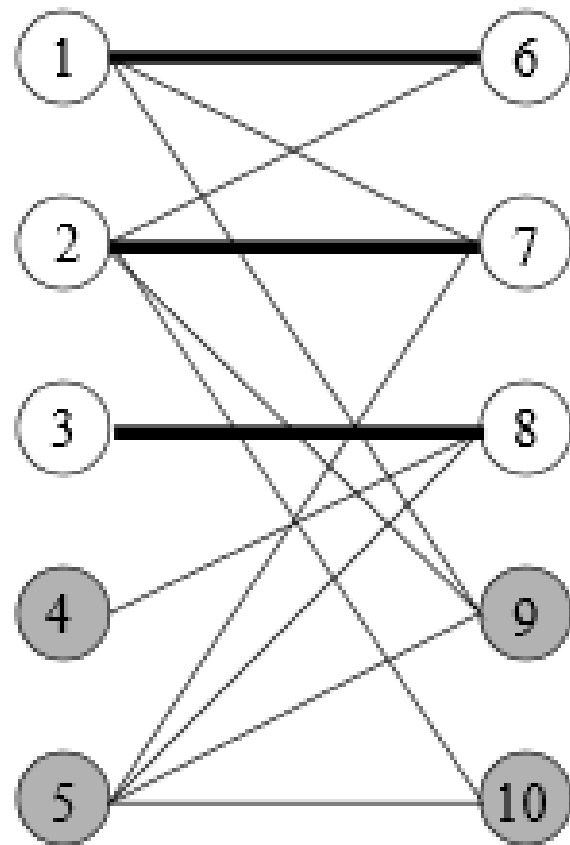
- Once the images are registered discrete optimization is also being used to find biomarkers related to the imaging data
- Biomarker = an indicator of a biological state
- Optimization helps with the curse of non-convexity
- Not much literature on this yet
  - Jobs?

# Bipartite Matching

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- A graph  $G = (V, E)$  consists of a set  $V$  of *vertices* and a set of  $E$  of pairs of vertices called *edges*.
  - An edge  $e = (u, v)$
- A graph is bipartite if the vertex set  $V$  can be partitioned into two sets  $A$  and  $B$  such that no edge in  $E$  has both endpoints in the same set of the bipartition.
  - A matching is *perfect* if no vertex is exposed

# Bipartite Matching



There are three edges

1. (1,6)
2. (2,7)
3. (3,8)



# Bipartite Matching

- Consider a rectangular matrix  $A$  of random numbers

```
A = rand(4,2)
```

```
A =
```

0.4242	0.8010
0.5079	0.0292
0.0855	0.9289
0.2625	0.7303

- Determine the matching

```
[val mi mj] = bipartite_matching(A);
```

- $mi = [2;3]$
- $mj = [1;2]$

# Bipartite Matching

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- Rearrange the matrix based on the outputs of the previous problem

```
A1 = A(mi,mj)
```

```
A1 =
```

```
0.5079    0.0292  
0.0855    0.9289
```

# Bipartite Matching

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- Sum the diagonal of the new matrix

```
Ans = sum(diag(A1))
```

```
val
```

```
% The answer should verify that the maximum weighted matching is obtained  
% when using the bipartite_matching command.
```

```
Ans =
```

```
1.4367
```

```
val =
```

```
1.4367
```

# References

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