

Heap sort

分成四頁

第一頁文字說明

第二-四頁流程圖

(製造成max heap 讓陣列由小到大排序) ->最大的放在最後,剩下再次heapify,依序處理,最後排成一個由小到大的list

概念如下：

step1:是將陣列排序到完全二元樹裡面,

因為要製造max heap所以

step2-heapify:將每一個二元樹裡面的左子樹和右子樹和該二元樹的父節點比較,如果左子樹有大於該二元樹的父節點的話,則左子樹換到父節點(因為要滿足父節點大於子節點的條件),同理,如果右節點也大於父節點的話,那麼就將右子樹也換到父節點,然後再繼續向下處理,最後排序成一個max heap(滿足所有父節點都大於子節點的條件)!

step3-heapsort:把heapify後的二元樹當中位在 $i=0$ (整個完全二元樹的最上面的節點)與 $i=\text{len}(\text{nums})-1$ (整個完全二元樹的最後一個節點)交換位置

重複步驟二:繼續heapify除了最後一個節點以外的二元樹。

重複步驟三

依此類推,完成整個陣列有小到大的排序。

時間複雜度 $O(n\log n)$:

將陣列做maxheapify:需要花 $\log n * n$ 次->因為每次對調完,會再對剩下的node進行maxheapify所以要乘上n

將root與last node對調: $n-1$ 次

total: $n-1+n\log n$ 省略常數,所以 $O(n\log n)$

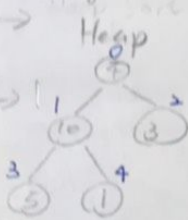
What: Heap Sort

Step 1 $\frac{1}{2}$ input data

ex:

4	10	3	5	1
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Binary Tree



def heap_sort(A, R)

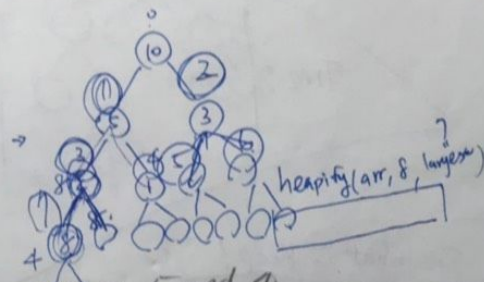
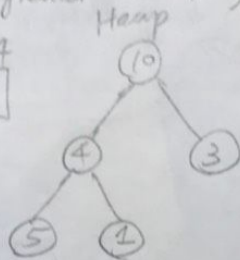
```
R = 1
l = 2 * l
r = 2 * l + 1
if A[l] < A[R]:
    A[l], A[R] = A[R], A[l]
    R = l
if A[r] < A[R]:
    A[r], A[R] = A[R], A[r]
    R = r
if R != l:
    heapify(A, R)
```

Step 2 Create Max Heap (Parent is greater than or equal to child nodes)

step 2-1 $\frac{10}{10}$ is greater than 4, So we swap 4 and 10

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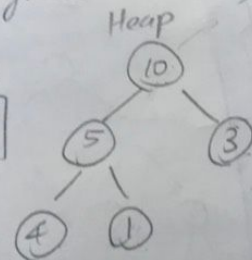
0	1	2	3	4
10	4	3	5	1



Step 2-2 5 is greater than 4, So we swap 5 and 4

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0	1	2	3	4
10	5	3	4	1



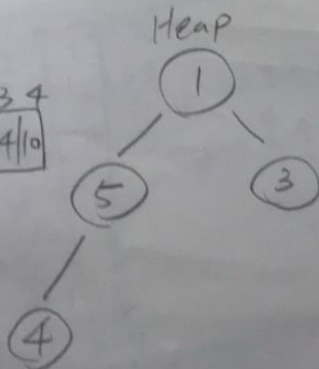
$arr[3], arr[7] = arr[7], arr[3]$

(when all parent node is greater than or equal to child nodes)

Step 2-3 remove the node: Swap first and last node and delete the last node from heap

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0	1	2	3	4
1	5	3	4	10

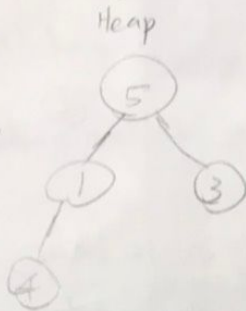


Why swap first & last node:
→ Return to the Parent node is not greater than or equal to child nodes's situation and Create Max Heap AGAIN!

Step 2-4 5 is greater than 1, So we swap 5 and 1

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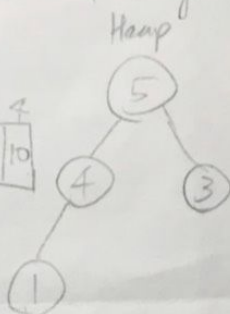
0	1	2	3	4
5	1	3	4	10



Step 2-5 4 is greater than 1, so we swap 4 and 1.

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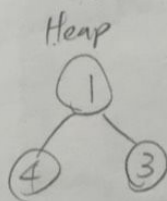
0	1	2	3	4
5	4	3	1	10



Step 2-6 remove the node: Swap first and last node and delete the last node from heap

Index

0	1	2	3	4
1	4	3	5	10

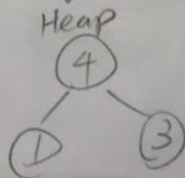


Swap 1 and 5

Step 2-7 4 is greater than 1, so swap 4 and 1

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0	1	2	3	4
4	1	3	5	10

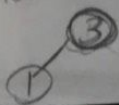


last node First node and delete last node

Step 2-8 remove the node: Swap ³1 and 4 ⇒

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0	1	2	3	4
3	1	4	5	10



Step 2-9 remove the node

Index 0 1 2 3 4

3	4	5	10	
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swap first node and last node
and delete the last node

Step 2-10

Index 0 1 2 3 4

1	3	4	5	10
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min

max



Only one element left in the heap,
Algorithm ends.

Heap