High Performance Computing and modern Architectures

Comparison analysis of parallel sorting algorithms

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algorithms Sorting

Merge Sort

Divide the unsorted list into *n* sublists, (a list of one element is considered sorted). Repeatedly merge sublists to produce new sorted sublists until there is only one sublist remaining.

- \blacktriangle Linear Complexity O(nlogn)
- ▲ Parallel Complexity $O(\frac{n}{p}log(n))$

Bucket Sort

- Set up an array of initially empty buckets
- Scatter: Go over the original array, putting each object in its bucket.
- Sort each non-empty bucket.
- Gather: Visit the buckets in order and put all elements back into the original array.

Merge Sort

- **✓ Vector**: 100000000
- **◀ Number of threads: 2**

Sequential version

101

MPI

82.75

OMP

70

Results

Bucket Sort

- Vector: 100000000
- **◀ Number of threads: 2**

Sequential version

108.596

MPI

87.76

OMP

72

Results

Bucket Sort

Results

Sort	Time
Seq Merge	158
OMP	128

Experiments

- 1. Sequential Merge
- 2. OMP Merge

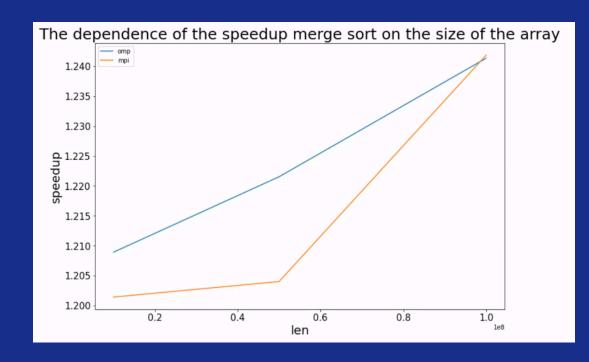
▼ Vector: 100000000

Number of threads: 2

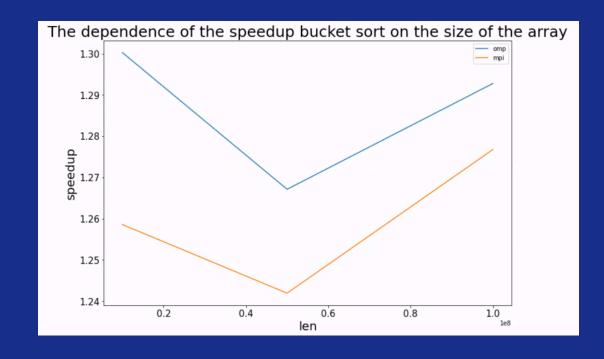
Results Comparison

Vector size

Merge Sort

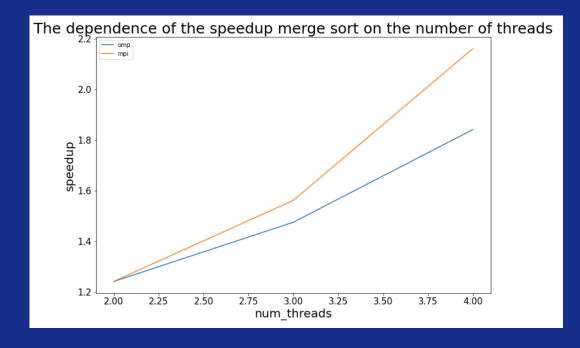


Bucket Sort



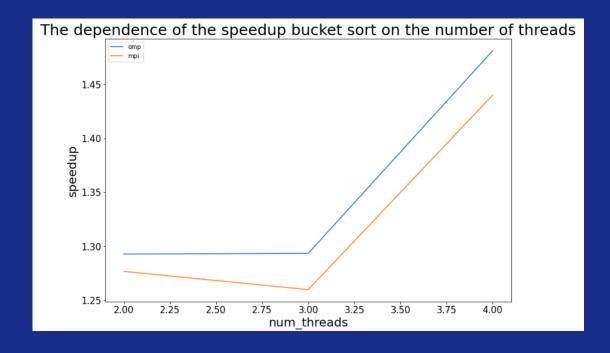
Results Comparison

Merge Sort



Number of threads

Bucket Sort



Conclusion

Two versions of parallel algorithms are implemented

√ a comparison of results is made

problems with memory problems have been reached and analysed

✓ algorithm optimization is performed

Conclusion

not enough memory on my computer

▲ MPI needs more memory then OpenMP

OMP is faster for my implementation

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