

## 2 Relative Entropy and Mutual Information [17 Points]

We define **relative entropy** between two discrete distributions  $\mathbf{p} \in \{p_1, \dots, p_n\}$  and  $\mathbf{q} = \{q_1, \dots, q_n\}$  as:

$$D(\mathbf{p} \parallel \mathbf{q}) = \sum_{i=1}^n p_i \log \frac{p_i}{q_i}$$

The above quantity is equal to 0 if and only if  $\mathbf{p} = \mathbf{q}$ . Otherwise, this quantity is positive. We can think of **relative entropy** as a measure of "distance" between two distributions. This quantity is widely known as the **Kullback-Leibler (KL) divergence** between two distributions (you can assume that none of the  $q_i$  and  $p_i$  are 0).

1. [5 Points] Prove that the **KL divergence** is always non-negative. Show all steps of your work

*Hint:* You may use this inequality without proof:  $x - 1 \geq \log(x)$ , where  $x \in \mathbb{R}$  with equality if and only if  $x = 1$ .

$$\begin{aligned}
 D(\mathbf{p} \parallel \mathbf{q}) &= \sum_{i=1}^n p_i \log \frac{p_i}{q_i} \\
 \xrightarrow{(-1)} -D(\mathbf{p} \parallel \mathbf{q}) &= - \sum_{i=1}^n p_i \log \frac{p_i}{q_i} \\
 &= - \sum_{i=1}^n p_i (\log p_i - \log q_i) \\
 &= \sum_{i=1}^n p_i (\log q_i - \log p_i) \\
 &= \sum_{i=1}^n p_i \log \frac{q_i}{p_i} \leq \sum_{i=1}^n p_i \left( \frac{q_i}{p_i} - 1 \right) \\
 &\quad \downarrow \times \\
 &= \sum_{i=1}^n p_i \times \frac{q_i}{p_i} - \sum_{i=1}^n p_i \\
 &= \sum_{i=1}^n q_i - \sum_{i=1}^n p_i = 1 - 1 = 0 \\
 \Rightarrow - \sum_{i=1}^n p_i \log \frac{p_i}{q_i} \leq 0 &\Rightarrow \sum_{i=1}^n p_i \log \frac{p_i}{q_i} \geq 0 \quad \times
 \end{aligned}$$

2. [5 Points] True or False: **KL-divergence** is symmetric (i.e.  $D(p||q) = D(q||p)$ ).  
If you think the statement is true, please write down a sketch of proof. Otherwise, give a counterexample. Show all steps of your work

False, counter example :

$$P \in \{0.2, 0.8\}$$

$$q \in \{0.6, 0.4\}$$

$$D(p||q) = \sum_{i=1}^n p_i \log \frac{p_i}{q_i} = 0.2 \log \frac{0.2}{0.6} + 0.8 \log \frac{0.8}{0.4} \approx 0.1454$$

$$D(q||p) = \sum_{i=1}^n q_i \log \frac{q_i}{p_i} = 0.6 \log \frac{0.6}{0.2} + 0.4 \log \frac{0.4}{0.8} \approx 0.1659$$

$$\rightarrow D(p||q) \neq D(q||p)$$

3. [7 Points] Let  $X$  and  $Y$  be random variables taking the values  $\{1, \dots, n\}$ . Recall the definition of mutual information between  $X$  and  $Y$ :

$$I(X; Y) = H(Y) - H(Y|X)$$

Using the conclusion of part 1), show that the mutual information is always non-negative. Show all steps of your work.

$$H(Y) = - \sum_{y \in Y} P(y) \log P(y)$$

$$= - \sum_{y \in Y} \sum_{x \in X} P(x, y) \log P(y)$$

$$H(Y|X) = \sum_{i=1}^n P(X=i) H(Y|X=i)$$

$$= \sum_{i=1}^n P(X=i) \times \left( - \sum_{j=1}^n P(Y=j|X=i) \log_2 P(Y=j|X=i) \right)$$

$$= - \sum_{i=1}^n \sum_{j=1}^n \underbrace{P(X=i) P(Y=j|X=i)} \log_2 P(Y=j|X=i)$$

$$= - \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 P(Y=j|X=i)$$

$$I(X; Y) = H(Y) - H(Y|X)$$

$$= - \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 P(Y=j) + \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 P(Y=j|X=i)$$

$$= \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \left[ \log_2 P(Y=j|X=i) - \log_2 P(Y=j) \right]$$

$$= \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 \frac{P(Y=j|X=i)}{P(Y=j)}$$

$$= \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 \frac{P(Y=j | X=i) P(X=i)}{P(Y=j) P(X=i)}$$

$$= \sum_{i=1}^n \sum_{j=1}^n P(X=i, Y=j) \log_2 \frac{P(Y=j, X=i)}{P(Y=j) P(X=i)}$$

$$= D(p(X=i, Y=j) \| p(X=i) P(Y=j))$$

Based on the conclusion of part 1)  $= D(p \| q) \geq 0$

$$I(X; Y) = D(p(X=i, Y=j) \| p(X=i) P(Y=j)) \geq 0 \quad \star$$