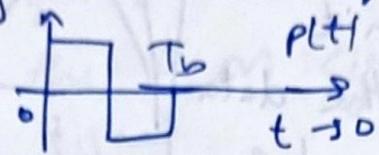


4th November, 2023

Tut-11 comm sys

Q1. Derive $S_y(\omega)$, the PSD of the Manchester (split phase) signal, assuming 1, 0 are equally likely. Sketch the PSD

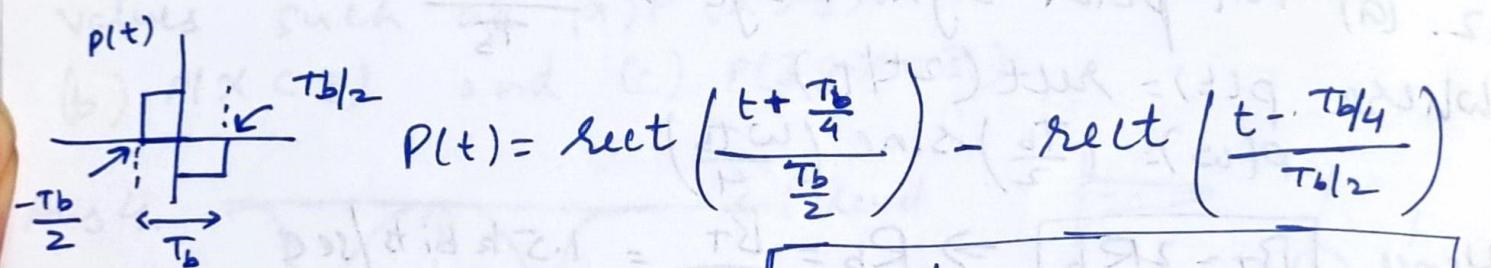


Ans1. $S_y(\omega) = \frac{|P(\omega)|^2}{T_b} \left(R_0 + 2 \sum_{n=1}^{\infty} R_n \cos n\omega T_b \right)$

Polar Code: "1" is transmitted by $p(t)$ and "0" is transmitted by $-p(t)$. Hence,

$$R_0 = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_k a_k^2 = \lim_{N \rightarrow \infty} \frac{1}{N} (N) = 1$$

$$R_n = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{k=1}^N a_k a_{k+n} = \lim_{N \rightarrow \infty} \frac{1}{N} \left[\frac{N}{2}(1) + \frac{N}{2}(-1) \right] = 0$$



we know that $\text{rect} \frac{t}{T} \leftrightarrow T \sin(\omega T) \frac{1}{2}$

Here, $T = 0.5T_b$

$$\Rightarrow \text{rect} \left(\frac{t}{0.5T_b} \right) \leftrightarrow 0.5T_b \sin(0.25\omega T_b)$$

We also know that $n(t-t_0) \leftrightarrow X(w) e^{-j\omega t_0}$

$$n(t+0.25T_b) \leftrightarrow X(w) e^{-j\omega(0.25(-T_b))} = X(w) e^{j\omega(0.25T_b)}$$

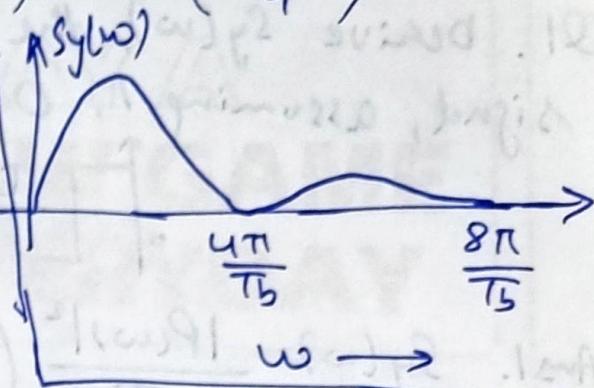
$$\therefore \text{rect} \left(\frac{t+0.25T_b}{0.5T_b} \right) \leftrightarrow 0.5T_b \sin(0.25\omega T_b) e^{0.25j\omega T_b}$$

$$\text{rect} \left(\frac{t-0.25T_b}{0.5T_b} \right) \leftrightarrow 0.5T_b \sin(0.25\omega T_b) e^{-0.25j\omega T_b}$$

$$\therefore P(\omega) = jT_b \sin \left(\frac{\omega T_b}{4} \right) \sin \left(\frac{\omega T_b}{4} \right)$$

$$S_y(\omega) = \frac{|P(\omega)|^2}{T_b} = T_b \operatorname{sinc}^2\left(\frac{\omega T_b}{4}\right) \sin^2\left(\frac{\omega T_b}{4}\right)$$

Q2. A leased telephone of bandwidth 3 kHz is used to transmit binary data. Calculate the data rate (in bps) that can be transmitted if we use:



- (a) Polar signal with rectangular half-width pulses.
- (b) Polar signal with rectangular full-width pulses.
- (c) Polar signal ~~with~~ using Nyquist criterion pulses of $\alpha = 0.25$.
- (d) Bipolar signal with rectangular half-width pulses.
- (e) Bipolar signal with rectangular fullwidth pulses.

Ans2. (a) For polar signal: $S_y(\omega) = \frac{|P(\omega)|^2}{T_b}$

where $p(t) = \operatorname{rect}(2t/T_b)$

$$P(\omega) = \left(\frac{T_b}{2}\right) \operatorname{sinc}\left(\frac{\omega T_b}{4}\right)$$

Hence, $B_T = 2R_b \Rightarrow R_b = \frac{B_T}{2} = 1.5 \text{ k bits/sec}$

(b) $p(t) = \operatorname{rect}(t/T_b); P(\omega) = T_b \operatorname{sinc}\left(\frac{\omega T_b}{2}\right)$

Essential $BW(B_T) = R_b$, hence transmission = 3 kbits/sec bit rate

(c) $B_T = (1+\alpha)R_b = 3000$, hence $R_b = \frac{6000}{1.25} = 4800 \frac{\text{bit}}{\text{sec}}$

(d) $S_y(\omega) = \frac{T_b}{4} \operatorname{sinc}^2\left(\frac{\omega T_b}{4}\right) \sin^2\left(\frac{\omega T_b}{2}\right)$

Essential $BW = R_b$: Bit rate possible is 3 kbit

(e) Bipolar $\Rightarrow S_y(\omega) = \frac{|P(\omega)|^2}{T_b} \sin^2\left(\frac{\omega T_b}{2}\right)$

$$P(\omega) = T_b \operatorname{sinc}\left(\frac{\omega T_b}{2}\right)$$

$$S_y(\omega) = T_b \operatorname{sinc}^2\left(\frac{\omega T_b}{2}\right) \sin^2\left(\frac{\omega T_b}{2}\right)$$

Q3. A 64 kbps binary PCM Bipolar NRZ signal is passed through a communication system with a raised cosine filter with $\alpha = 0.25$. Determine the bandwidth of the PCM signal and the bandwidth of the filtered PCM signal.

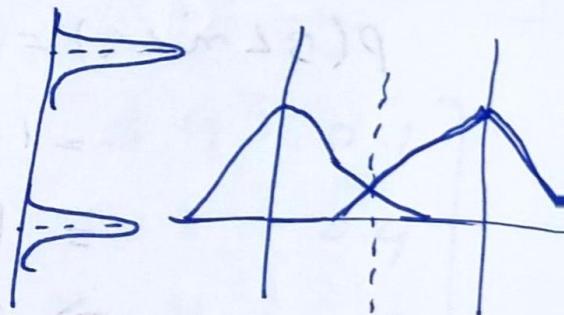
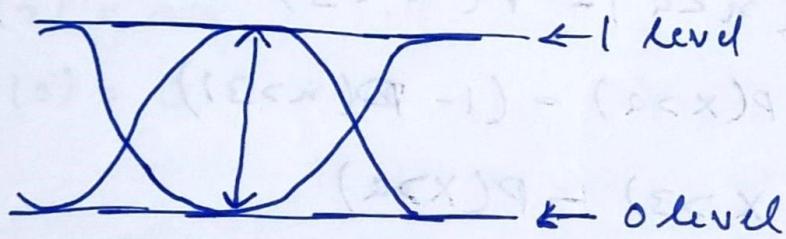
Ans 3. $BW = (1+\alpha) \frac{R_b}{2}$ = Bandwidth of PCM signal

Filtered signals

$$\begin{aligned} &= (1+0.25) \frac{64}{2} \text{ kHz} = 40 \text{ kHz} \end{aligned}$$

Q4. For the Gaussian random variable, X with mean and standard deviation values are given as $\mu = 5$ and $\sigma = 2$, respectively. Find the probability values such as (a) $P(X > 8)$
 (b) $P(X < 8)$ and (c) $P(3 < X < 8)$

Ans 4.



$$\begin{aligned} (a) P(X > 8) &= Q\left(\frac{x-\mu}{\sigma}\right) \\ &= Q\left(\frac{8-5}{2}\right) = Q(1.5) = 0.0668 \quad (\text{From the } Q \text{ table}) \end{aligned}$$

$$\begin{aligned} (b) P(X < 8) &= 1 - P(X > 8) \\ &= 1 - 0.0668 = 0.9332 \end{aligned}$$

$$Q(-x) = 1 - Q(x)$$

$$\begin{aligned} (c) P(3 < X < 8) &= P(X < 8) - P(X < 3) \\ &= P(X < 8) - (1 - P(X > 3)) \\ &= 0.9332 - 1 + Q\left(\frac{3-5}{2}\right) \\ &= \cancel{0.9332} + 0.774533 \end{aligned}$$

Q5. For the Gaussian random variable, X with mean and standard deviation values are given as $\mu=5$ and $\sigma=3$, respectively. Find the factor a such that $P(X < a) = 0.9$

Ans 5. $P(X < a) = 0.9$

$$P(X < a) = 1 - P(X > a) = 1 - Q\left(\frac{x-\mu}{\sigma}\right) = 1 - Q\left(\frac{a-5}{3}\right)$$

$$\Rightarrow Q\left(\frac{a-5}{3}\right) = 0.1 = 0.9$$

$$\Rightarrow \frac{a-5}{3} = 1.3 \Rightarrow \boxed{a = 5 + 3(1.3) = 8.9}$$

↓ take either 1.25 or 1.3

Q6. For the Gaussian random variable, X with mean and standard deviation values are given as $\mu=5$, $\sigma=3$ respectively. Find the factor such that

$$P(3 < X < a) = 0.7$$

Ans 6. $P(3 < X < a) = 0.7$

$$P(3 < n < a) = P(n < a) - P(n < 3)$$

$$= 1 - P(X > a) = (1 - Q(n > 3))$$

$$= P(X > 3) - P(X > a)$$

$$\Rightarrow Q\left(\frac{3-5}{3}\right) - Q\left(\frac{a-5}{3}\right) = 0.7$$

$$\boxed{a = 9.65}$$

11th November, 2025

Tutorial-12

Q1. The received signal input to the zero forcing equaliser (ZFE), is given as $n(n) = \{0, 0.1, 0.4, 0.8, 0.4, 0.1, 0\}$. The centre value $n(0) = 0.8$. Design a 3-tap ZFE and find the equaliser coefficients. Find the equalised output. Check for the reduction of ISI in the filtered signal.

Ans1. $n(n) = \{0, 0.1, 0.4, 0.8, 0.4, 0.1, 0\} \Rightarrow \text{ISI} \rightarrow$

$$\Rightarrow n(-1) = 0.4 \quad \text{and} \quad n(+1) = 0.4.$$

$$n(-3) = 0$$

$$n(-2) = 0.1$$

$$n(-1) = 0.4$$

$$n(0) = 0.8$$

$$n(1) = 0.4$$

$$n(2) = 0.1$$

$$n(3) = 0$$

3 tap
equaliser matrix

3 Tap ZFE

$$X = \begin{bmatrix} n(0) & n(-1) & n(-2) \\ n(+1) & n(0) & n(-1) \\ n(+2) & n(+1) & n(0) \end{bmatrix}$$

$$X = \begin{bmatrix} 0.8 & 0.4 & 0.1 \\ 0.4 & 0.8 & 0.4 \\ 0.1 & 0.4 & 0.8 \end{bmatrix}$$

coefficient matrix = $C = \boxed{X^{-1} Z}$

Now, ~~$X^{-1} Z$~~

$$\left[\begin{array}{ccc|ccc} 0.8 & 0.4 & 0.1 & 1 & 0 & 0 \\ 0.4 & 0.8 & 0.4 & 0 & 1 & 0 \\ 0.1 & 0.4 & 0.8 & 0 & 0 & 1 \end{array} \right]$$

$$X^{-1} = \left[\begin{array}{ccc|ccc} 0.8 & 0.4 & 0.1 & 1 & 0 & 0 \\ 0.4 & 0.8 & 0.4 & 0 & 1 & 0 \\ 0.1 & 0.4 & 0.8 & 0 & 0 & 1 \end{array} \right]$$

$$C = \begin{bmatrix} 1 \\ 2.25 \\ -1 \end{bmatrix}$$

$$Z = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \text{ for zero forcing equaliser}$$

Equaliser output = $x^2(n) = x(n) * C(n)$

$$y(n) = (-1)x(n+1) + 2.25x(n) + 1x(n-1) \quad \text{Convolution.}$$

$$\underline{x^2(n)} = (0, -0.01, -0.0175, \underline{\underline{0, 1, 0}}, -0.0175, \cancel{0}, -0.01, 0)$$

Reduction

ISI required for ~~$x(n+1), x(n+2)$~~ sample values

Q2 The sampled values of unequalised pulses are given by 0.4, 0.6 and 0.8 for $k=1, 0$, and -1 respectively. Design the zero forcing equaliser. What will be equaliser output for $k=-3, -2, +2$ and +3 respectively?

$$\text{Ans. } P_x[+1] = 0.4$$

$$P_x[0] = 0.6$$

$$P_x[-1] = 0.8$$

$$P_x[2] = 0 \quad \left. \begin{array}{l} \text{Assume since} \\ P_x[-2] = 0 \quad \text{not given.} \end{array} \right\}$$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} P_x(0) & P_x(-1) & P_x(-2) \\ P_x(+1) & P_x(0) & P_x(+1) \\ P_x(+2) & P_x(+1) & P_x(0) \end{bmatrix} \begin{bmatrix} c_{-1} \\ c_0 \\ c_1 \end{bmatrix}$$

$$\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0.6 & 0.8 & 0 \\ 0.4 & 0.6 & 0.8 \\ 0 & 0.4 & 0.6 \end{bmatrix} \begin{bmatrix} c_1 \\ c_0 \\ c_1 \end{bmatrix}$$

$$C_1 = \frac{20}{7} = 2.85$$

$$C_0 = -\frac{15}{7} = -2.142$$

$$C_1 = \frac{10}{7}$$

$$= 1.4$$

By using

$$\begin{aligned} y(n) &= \frac{10}{7} n(n-1) + \left(-\frac{15}{7}\right) n(n) + \frac{20}{7} (n(n+1)) \\ y(n-3) &= \frac{20}{7} (n-4) + \left(-\frac{15}{7}\right) (n-3) + \frac{10}{7} (n-2) = 0 \end{aligned}$$

$$\begin{bmatrix} P_0(-3) \\ P_0(-2) \\ 0 \\ 1 \\ 0 \\ P_0(2) \\ P_0(3) \end{bmatrix} = \begin{bmatrix} 0.6 & 0.8 & 0 & 0 & 0 & 0 & 0 \\ 0.4 & 0.6 & 0.8 & 0 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & 0.4 & 0.6 & 0.8 \\ 0 & 0 & 0 & 0 & 0 & 0.4 & 0.6 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 20/7 \\ -15/7 \\ 10/7 \\ 0 \\ 0 \end{bmatrix}$$

$$\text{For } k=-3, P_0(k) = 0$$

$$k = -2, P_0(k) = 0.8 \times \frac{20}{7} = \frac{16}{7} = 2.28$$

$$k=2, P_0(k) = 0.4 \times \frac{10}{7} = \frac{4}{7} = 0.571$$

$$k=3, P_0(k) = 0$$

Q3. Given $N_0 = 10^{-10} \frac{W}{Hz}$, $A = 50mV$, and the bit rate is $R_b = \frac{1}{T} = 5Mbps$, find P_E .

$$\text{Ans 3. } P_E = \frac{1}{2} \operatorname{erfc}(\sqrt{\gamma}) \quad \text{where} \quad \boxed{\gamma = \frac{E_b}{N_0} = \frac{A^2 T}{N_0}}$$

$$\therefore P_E = \frac{1}{2} \operatorname{erfc}(\sqrt{\gamma}) \quad \gamma = \frac{(50 \times 10^{-3})^2 (1s)(10^6)}{10^{-10}}$$

1. Get the value from the error function table $\gamma = 5$

$$2. \operatorname{erfc}(x) \cong \frac{1}{x\sqrt{\pi}} e^{-x^2} \Rightarrow \text{true only for } x > 3$$

$$\operatorname{erfc}(x) \cong \frac{2x}{\sqrt{\pi}} \Rightarrow \text{true for } |x| \ll 1$$

We can use the first formula despite $\sqrt{S} < 3$

$$\Rightarrow \operatorname{erfc}(m) = \frac{1}{2\sqrt{\pi}} e^{-m^2} = 1.7 \times 10^{-3}$$

$$\therefore P_E = \frac{1}{2} \operatorname{erfc}(\sqrt{S}) = 0.768 \times 10^{-3} = Q\left(\sqrt{\frac{2E}{N_0}}\right) = Q\sqrt{2r}$$

Q4. Prove that the maximum signal to noise ratio of the matched filter is

$$\left(\frac{S}{N}\right)_{\text{max}} = \frac{2E}{N_0}.$$

$$\text{Ans 4. } \left(\frac{S}{N}\right)_o = \int_{-\infty}^{\infty} \frac{|x(f)|^2 df}{S_N(f)} = \int_{-\infty}^{+\infty} \frac{|x(f)|^2 df}{N_0/2}$$

For system assume

$$\text{AWGN} = \frac{2}{N_0} \int_{-\infty}^{\infty} |x(f)|^2 df = \frac{2E}{N_0}$$

$$\text{Noise PSD} = \frac{N_0}{2} = S_N(f)$$

$$\text{Rayleigh's Theorem, } E = \int_{-\infty}^{+\infty} n^2(t) dt = \int_{-\infty}^{\infty} |x(f)|^2 df$$

Q5. In binary transmission, one of the messages is represented by a rectangular pulse $n(t)$. Another message is transmitted by the absence of the pulse.

Evaluate the signal to noise ratio $t=T$. Assume the PSD of Noise is $N_0/2$. Sketch the impulse response and output of match filter.

Ans 5.

$$n(t) = \begin{cases} A & ; 0 \leq t \leq T \rightarrow \text{bit 1} \\ 0 & ; \text{otherwise} \rightarrow \text{bit 0} \end{cases}$$

Noise is AWGN or PSD = $\frac{N_0}{2}$

Input

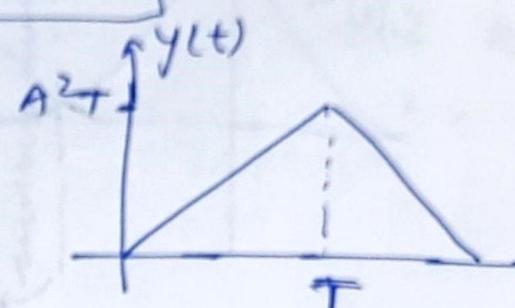
$$h(t) = \frac{2A}{T} e^{-\frac{|t|}{T}} \quad \text{assume}$$

$$y(t) = n(t) * h(t)$$

Now, $n(-t) = \begin{cases} A & -T \leq t \leq 0 \\ 0 & \text{otherwise} \end{cases}$

$$\Rightarrow n(T-t) = \begin{cases} A & 0 \leq t \leq T \\ 0 & \text{otherwise} \end{cases} = h(t)$$

$$SNR = \frac{2E_b}{N_0} = \frac{2A^2 T}{N_0}$$



18th November, 2025

Tut-B

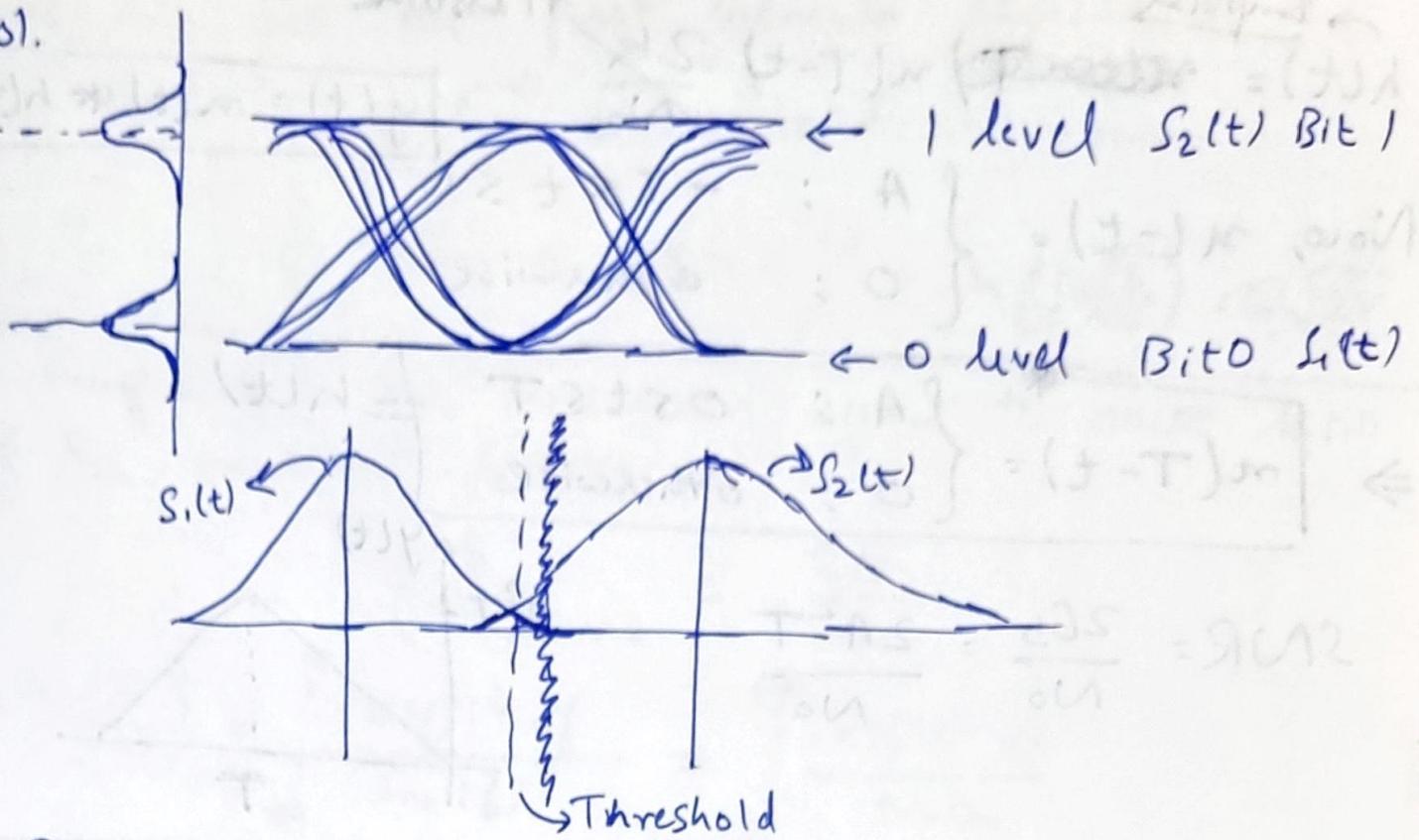
Q1. A binary communication system transmits signals $s_i(t) \quad i=1, 2$. The receiver test statistic $z(T) = a_i + n_0$, where the signal component a_i is either $a_1 = +1$ or $a_2 = -1$ and the noise component n_0 is uniformly distributed, yielding the conditional density functions $p(z/s_i)$ given by:

~~Q1.~~ $p(z/s_1) = \begin{cases} 1/2 & \text{for } -0.2 \leq z \leq 1.8 \\ 0 & \text{otherwise} \end{cases} \quad \text{and}$

$p(z/s_2) = \begin{cases} 1/2 & \text{for } -1.8 \leq z \leq 0.2 \\ 0 & \text{otherwise} \end{cases}$

Find the probability of a bit error, P_B , for the case of equally likely signalling and the use of an optimum decision threshold.

Ans.



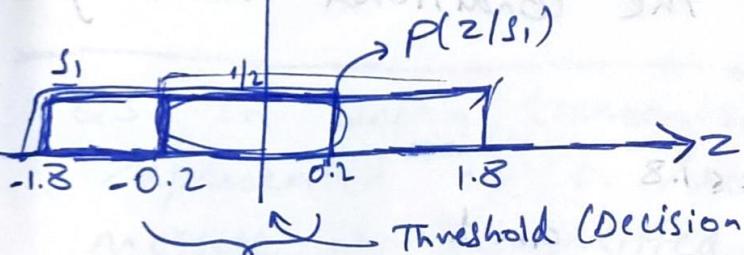
$$P_e = P[1_D, 0_T] + P[0_D, 1_T]$$

$$= \underline{P[1_D | 0_T]} P[0_T] + P[0_D | 1_T] P[1_T]$$

$$P(z)$$

H_1 : $S_1(t)$ is detected/received

H_2 : $S_2(t)$ is detected/received



$$\frac{P(z|a_2)}{P(z|a_1)} \stackrel{H_2}{>} \frac{P(S_1(t))}{P(S_2(t))}$$

changes with respect to
modulation formats

The error will be in this range.

$$\therefore \frac{P(z|a_2)}{P(z|a_1)} \stackrel{H_2}{>} 1 \stackrel{H_1}{<}$$

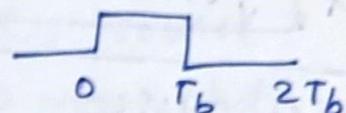
$$\text{Probability of error} = P_B = 0.5 \int_{-0.2}^0 0.5 dz + 0.5 \int_0^{0.2} 0.5 dz$$

$$= 0.25 \left[\int_{-0.2}^0 dz + \int_0^{0.2} dz \right] = \frac{1}{4} \times 0.4 = \boxed{0.1}$$

Q2. Consider that NRZ binary pulses are transmitted along a cable that attenuates the signal power by 3dB (from transmitter to receiver). The pulses are coherently detected at the receiver, and the data rate is 56 kbps/sec. Assume Gaussian noise with $N_0 = 10^{-6} \text{ W/Hz}$. What is the ~~minimum~~ amount of power needed at the transmitter in order to maintain a bit error probability of $P_B = 10^{-3}$?

{ Qtable }

Ans 2. Binary transmission \rightarrow NRZ



$$\text{Attenuation} = 3 \text{ dB}$$

$$\text{Data Rate} = R_b = 56 \text{ kbps}$$

$$P_B = \frac{10^{-3}}{2} = Q\left(\sqrt{\frac{2E_b}{N_0}}\right) = 0.001 \Rightarrow \sqrt{\frac{2E_b}{N_0}} = 3.08$$

$$N_0 = 10^{-6} \text{ W/Hz}$$

$$E_b = A^2 T_b$$

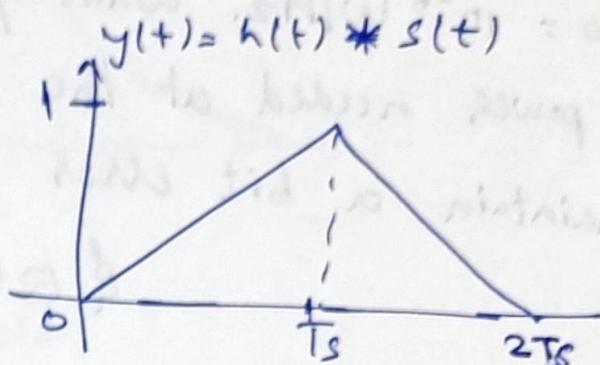
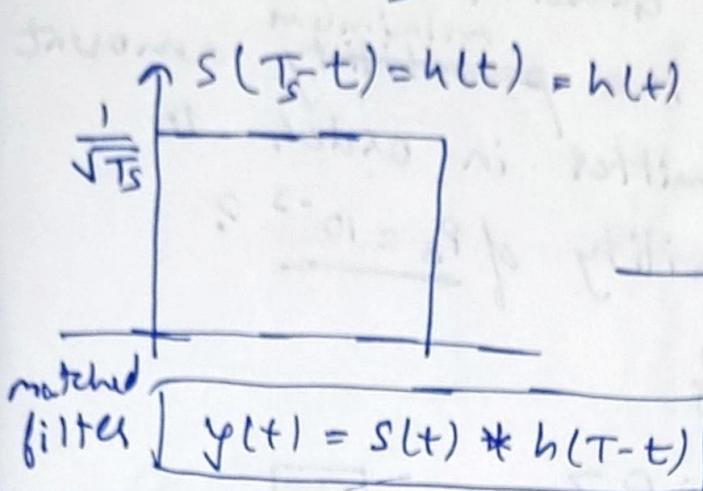
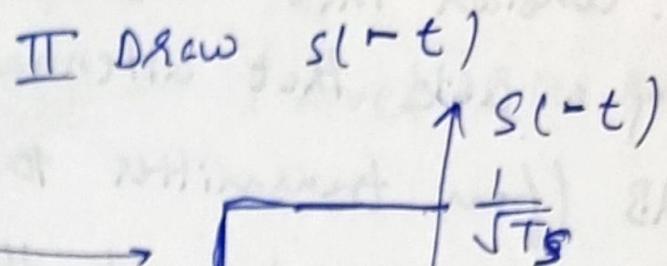
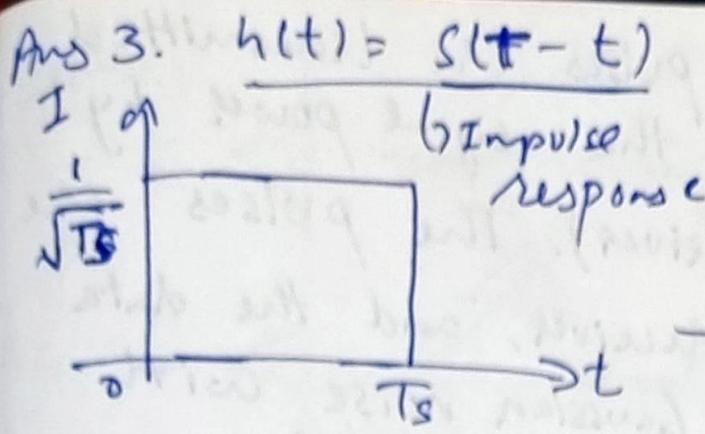
$$A^2 T_b = 4.74 \times 10^{-6}$$

$$\Rightarrow A^2 = (4.74 \times 10^{-6} \times 56 \times 10^3) \text{ W} = 265 \text{ mW}$$

$$\begin{aligned} \text{Now, the required power at Tx} &= 2 P_{Rx} \text{ (on 3dB loss)} \\ &= (2 \times 265) \text{ mW} \\ &\approx 0.531 \text{ W} \end{aligned}$$

Q3. A binary digital communication system employs the signal $s(t) = \begin{cases} \frac{1}{\sqrt{T}} & 0 \leq t \leq T \\ 0 & \text{otherwise} \end{cases}$

- Determine the impulse response of a filter matched to this signal and sketch it as a function of time.
- Plot the matched filter output as a function of time.



Q4. A source emits seven messages with probabilities $1/2, 1/4, 1/8, 1/16, 1/32, 1/64$, and $1/64$ respectively. Find the entropy of the source.

Ans 4. $H = \sum_i p_i \log_2 \left(\frac{1}{p_i} \right) = \frac{1}{2} \log_2 \left(\frac{1}{1/2} \right) + \frac{1}{4} \log_2 \left(\frac{1}{1/4} \right) + \dots$
 $\cong 1.96 \text{ bits/msg.}$

Q5. Assume that normal speech rate is 250 words/minute. Also assume that the average word contains six characters and the characters within a word are statistically independent. Based on these assumptions, compute the information rate considering the English alphabet.
 (Note: Information rate = $(R_s \times H(m)) \text{ per sec}$)

Ans 5. Per min words = 250
 per msg characters = 6

$$\text{Total ch. per min} = 1500$$

$$\text{Total ch. per sec} = \frac{1500}{60} = 25 \text{ ch/msg/sec}$$

$$H = \left[\frac{1}{26} \log_2 \left(\frac{1}{1/26} \right) \right] \times 26 = 4.7 \text{ bits/ch}$$

$$I = 4.7 \text{ bits/ch} \times 25 \\ = (4.7 \times 25) \text{ bps} = 117.5 \text{ bps.}$$

25th November, 2025

Tut-14

Q6. An analog signal having 4kHz bandwidth is sampled at 1.25 times the Nyquist rate, and each sample is quantised into one of equally likely levels. Assume that the successive samples are statistically independent.

- What is the information rate of this source?
- Can the output of this source be transmitted without error over an AWGN channel with a bandwidth of 10kHz and an S/N ratio of 20dB?
- Find the S/N ratio required for error-free transmission for part (i).
- Find the bandwidth required for an AWGN channel for error free transmission of the output of this source if the S/N ratio is 20 dB.

$$\text{No. of symbols} = 256$$

$$\text{Ans} \quad (i) f_m = 4\text{kHz}, \quad f_s = 2f_m = 8\text{kHz}$$

$$f_s^{\text{actual}} = 1.25f_s = (1.25 \times 8) = 10\text{kHz}$$

$$\text{No. of symbols} = 256, \quad \text{each symbol probability} = \frac{1}{256}$$

$$H(m) = - \sum_{i=1}^m P(m_i) \log_2 \left(\frac{1}{P(m_i)} \right) \rightarrow \text{entropy}$$

$$= \log_2(256) = 8 \text{ bits/sample.}$$

Information rate $R = H(n) f_s^{\text{actual}}$

$$= 8 \frac{\text{bits}}{\text{sample}} \times 10^4 \frac{\cancel{\text{bits samples}}}{\text{second}}$$

$$R = 80 \text{ kbps}$$

(ii) $C = B \log_2 (1 + SNR)$

$$B = 10 \text{ kHz}$$

$$SNR = 20 \text{ dB} = 10^{\frac{20}{10}} = 100$$

$$C = (10 \times 10^3) \times \log_2 (1 + 100) = 66.58 \text{ kbits/s}$$

$\Rightarrow R > C \Rightarrow$ no error free Transmission

$R < C \Rightarrow$ error free transmission.

(iii) $R = 80 \text{ kbps}$

for error free transmission, the required condition is $C \geq R$

$$B \log_2 (1 + SNR) \geq R$$

$$\Rightarrow (10 \times 10^3) (\log_2 (1 + SNR_{\text{req}})) \geq (80 \times 10^3)$$

$$\Rightarrow SNR_{\text{req}} \geq 2^8 - 1 = \boxed{24.06 \text{ dB}} = 255$$

(iv) $B \log_2 (1 + SNR) \geq R$

$$\Rightarrow B_{\text{req}} \log_2 (1 + 100) \geq 80 \times 10^3$$

$$\Rightarrow \boxed{B_{\text{req}} \geq 12.01 \text{ kHz}}$$

Q5. Given an AWGN channel with 4 kHz bandwidth and the noise power spectral density $n/2 = 10^{-12}$

The signal power required at the receiver = 0.1 mW
Calculate the capacity of this channel.

Ans. $P_{Tx} = 0.1 \text{ mW}, n/2 = 10^{-12} \frac{\text{W}}{\text{Hz}}, B = 4 \text{ kHz}$

$$C = B \log_2 (1 + SNR) \Rightarrow SNR = \frac{\text{Signal Power}}{\text{Noise Power}} = \frac{0.1 \times 10^3}{2 \times 10^{-1} \times 4 \times 10^{-2}}$$

$$\Rightarrow C = B \log_2 (1 + 12500) = (4 \times 10^3) (\log_2(12501))$$

$C = 54.439 \text{ kbps}$

Q4. A DMS X has five symbols n_1, n_2, n_3, n_4 , and n_5 with the $P(n_1) = 0.4, P(n_2) = 0.19, P(n_3) = 0.16, P(n_4) = 0.15, P(n_5) = 0.1$.

- (i) Construct a Shannon-Fano code for X , and calculate the efficiency of the code.
- (ii) Repeat for the Huffman code and compare the results.

write in descending order only

	n_i	$P(n_i)$	<u>Step 1</u>	<u>Step 2</u>	<u>Step 3</u>
	n_1	0.4	0	0	0
	n_2	0.19	0	0	0
	n_3	0.16	1	0	0
	n_4	0.15	1	0	0
	n_5	0.1	1	1	1

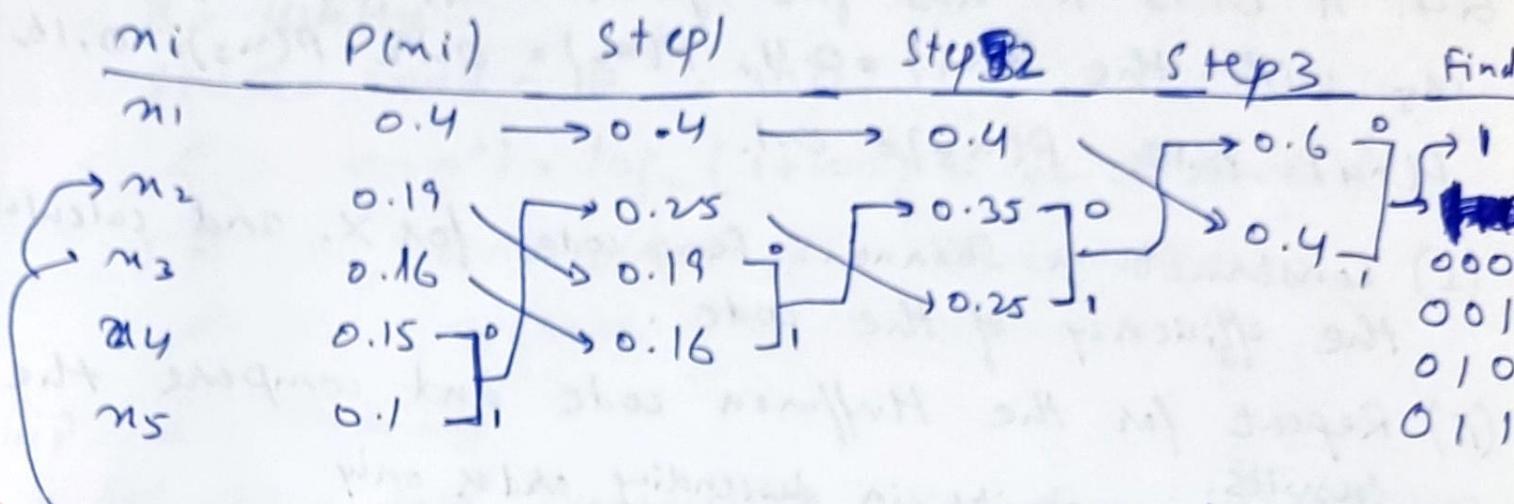
Find Code

n_1	<u> </u>	0 0	$L = (0.4/2) + (0.19 \times 2)$
n_2	<u> </u>	0 1	$+ (0.16 \times 2) + (0.15 \times 3)$
n_3	<u> </u>	1 0	$+ 0.1 \times 3$
n_4	<u> </u>	1 1 0	
n_5	<u> </u>	1 1 1	$L = 2.25$

$$H(m) = -\sum_{i=1}^5 P(m_i) \log_2 [P(m_i)] = 2.15$$

$$\eta = \frac{H(m)}{L} = \text{Coding Efficiency} = \frac{2.15}{2.25} = 95.55\%$$

(ii) Huffman Coding (not unique)



Since $0.25 \geq 0.19$, we have to interchange, else keep it there.

$$L = (0.4 \times 1) + (0.19 \times 3) + (0.16 \times 3) + (0.15 \times 3) + (0.1 \times 3) \\ = 2.2$$

$$H(m) = 2.15 \Rightarrow \eta = \frac{H(m)}{L} = \frac{2.15}{2.2} = 97.73\%$$

Q3. A DMS X has 5 equally likely symbols.

(i) Construct a Shannon-Fano code for X , and calculate the efficiency of the code.

(ii) Construct another Shannon-Fano code and compare the results.

(iii) Repeat for the Huffman code and compare the results.

$$\text{Ans 3. } P(X_i) = \frac{1}{5} = 0.2$$

(i)

<u>$n_i P(n_i)$</u>	<u>Step 1</u>	<u>Step 2</u>	<u>Step 3</u>	<u>Step 4 (Final)</u>
$n_1 0.2$	0	0		0 0
$n_2 0.2$	0	1	0	0 1 0
$n_3 0.2$	0	1	1	0 1 1
$n_4 0.2$	1	0		1 0
$n_5 0.2$	1	1		1 1

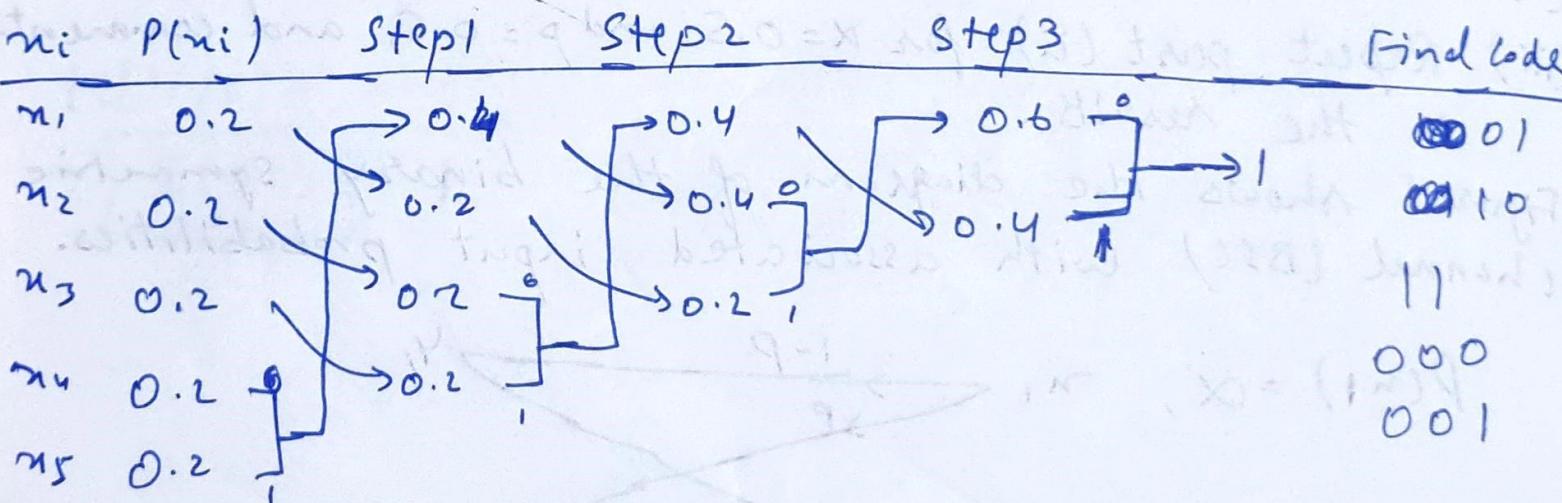
(ii)

<u>$n_i P(n_i)$</u>	<u>Step 1</u>	<u>Step 2</u>	<u>Step 3</u>	<u>Find Code</u>
$n_1 0.2$	0	0		00
$n_2 0.2$	0	1		01
$n_3 0.2$	1	0		10
$n_4 0.2$	1	1	0	11 0
$n_5 0.2$	1	1	1	11 1

$$L = (0.2)(2) + (0.2)(2) + (0.2)(2) + (0.2)(3) + (0.2)(3)$$

$$= \cancel{0.2} 1.2 + 1.2 = \boxed{2.4}$$

(iii) Huffman Coding



Q2. A DMS X has 4 symbols m_1, m_2, m_3 and m_4 with $P(m_1) = \frac{1}{2}$, $P(m_2) = \frac{1}{4}$ and $P(m_3) = P(m_4) = \frac{1}{8}$. Construct a Shannon-Fano code for X ; show that this code has the optimum property that $n_i = I(m_i)$ and that the code efficiency is 100%.

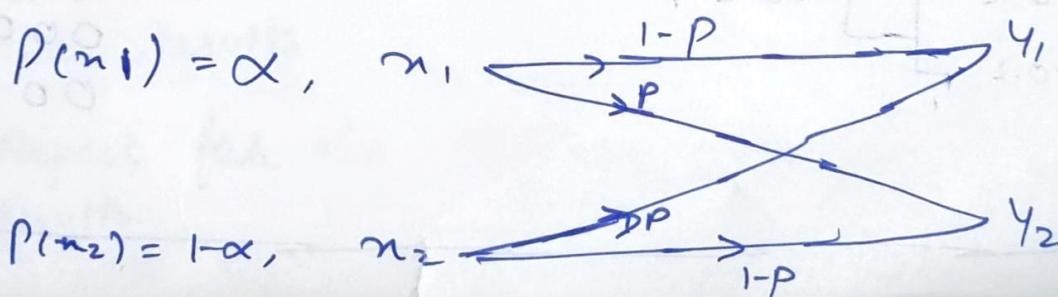
<u>Ans2.</u>	<u>m_i</u>	<u>$P(m_i)$</u>	<u>Step1</u>	<u>Step2</u>	<u>Step3</u>	<u>Final Code</u>
	m_1	0.5	0			0
	m_2	0.25	1	0		10
	m_3	0.125	1	1	0	110
	m_4	0.125	1	1	1	111

$$L = (0.5 \times 1) + (0.25 \times 2) + (0.125 \times 3) \times 2 = 1.75$$

$$H(m) = 1.75 \quad \therefore \eta = \frac{H(m)}{L} = 100\%$$

- RQ1. Given a binary symmetric channel (BSC) (Figure 1) with $P(m_i) = \alpha$,
- Show that the mutual information $I(X; Y)$ is given by $I(X; Y) = H(Y) + p \log_2 p + (1-p) \log_2 (1-p)$
 - Calculate $I(X; Y)$ for $\alpha = 0.5$ and $p = 0.1$.
 - Repeat part (ii) for $\alpha = 0.5$ and $p = 0.5$, and comment on the results.

Figure 1 shows the diagram of the binary symmetric channel (BSC) with associated input probabilities.



Ans1.

$$[P(m_i, y_j)] = \begin{bmatrix} \alpha & 0 \\ 0 & 1-\alpha \end{bmatrix} \begin{bmatrix} 1-p & p \\ p & 1-p \end{bmatrix}$$

$$= \begin{bmatrix} \alpha(1-p) & \alpha p \\ p(1-\alpha) & (1-\alpha)(1-p) \end{bmatrix} = \begin{bmatrix} P(m_1, y_1) & P(m_1, y_2) \\ P(m_2, y_1) & P(m_2, y_2) \end{bmatrix}$$

$$\Rightarrow H(Y/x) = - \sum_{j=1}^n \sum_{i=1}^m P(m_i, y_j) \log_2 (y_j/m_i)$$

$$= -p \log_2 p - (1-p) \log_2 (1-p)$$

$$\begin{aligned} I(x; Y) &= H(Y) - H(Y/x) \\ &= H(Y) + p \log_2 p + (1-p) \log_2 (1-p) \end{aligned}$$

$$\alpha = 0.5, \quad p = 0.1$$

~~$$H(x) = - \sum_{i=1}^m P(m_i) \log_2 (P(m_i)) = 1$$~~

$$H(Y) = - \sum_{i=1}^m P(y_i) \log_2 (P(y_i)) = 1$$

$$H(Y/x) = 0.469$$

$$\therefore I(x; Y) = 1 - 0.469 = 0.5301$$