

Solutions of Non-Linear Equations :-

In science and engineering, mathematical models for different problems can be formulated into the equations of the form,

$$f(n) = 0 \quad \text{--- (1)}$$

where, n and $f(n)$ may be real, complex. The solution process often involves finding the values of n that would satisfy $f(n) = 0$. These values are called roots of the equation. The equation (1) may belong to one of the following types of equations.

- ① Algebraic equation
- ② Polynomial equation
- ③ Transcendental equation

A Non-linear Function :

Any function of one variable which does not graph as a straight line in two dimensions or any function of two variables which does not graph as a plane in three dimensions can be said to be non linear.

Consider the function :

$y = f(n)$, $f(n)$ is said to be non linear if the response of the dependent variable y is not direct or exact proportion to the changes in the independent variable n .

For example : $y = n^2 + 1$

Linear Equation:

Consider the function,

$y = f(n)$, $f(n)$ is a linear function if the dependent variable y changes in direct proportion to the change in independent variable n .

$$\text{eg: } y = 2n + 3$$

Types of Equations:

1) Algebraic Equation: An equation of the form $y = f(n)$ is said to be algebraic if it can be expressed as:

$f_n y^n + f_{n-1} y^{n-1} + \dots + f_1 y^1 + f_0 = 0$,
where, f_i is an i^{th} order polynomial in n (say).

Eg:- $2n + 4y - 19 = 0$, $n^3 - ny - 4y^3 = 0$
In general form it can also be represented as $f(n, y) = 0$

2) Polynomial Equations:

Polynomial equations are simple class of algebraic equations that are represented as:-

$$a_n n^n + a_{n-1} n^{n-1} + \dots + a_1 n + a_0 = 0$$

This is called n^{th} degree polynomial and has n roots. The roots may be

- i) real and distinct
- ii) real and repeated
- iii) complex

3) Transcendental Equation:

A non-algebraic equation is called a transcendental equation. These include the trigonometric, exponential and logarithmic functions.

e.g. i) $2 \sin x - x = 0$,

ii) $e^x \sin x - \frac{1}{2}x = 0$

iii) $\log x^2 - 1 = 0$

Root: The value 'a' of x which satisfies $f(x)=0$, is called root of $f(x)=0$.

Geometrically, a root of $f(x)=0$ is that value of x where the graph of $y=f(x)$ intersects the x-axis.

The process of finding the roots of an equation is known as the solution of that equation.

Q. Solve the equation:

$$2x^3 + x^2 - 13x + 6 = 0$$

$$\Rightarrow 2x^3 + x^2 - 13x + 6 = 0$$

By inspection we can find $x=2$, which satisfies the given equation.

$\therefore 2$ is its root, so $x=2$ is a factor of $x-2 | 2x^3 + x^2 - 13x + 6$

Dividing this polynomial by $x-2$, we get the quotient $2x^2 + 5x - 3$ and remainder 0.

$$x-2 | 2x^3 + x^2 - 13x + 6$$

The quotient is in the form of $an^2 + bn + c$

where, $a = 2$, $b = 5$, $c = -3$

so,

$$n = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$= \frac{-5 \pm \sqrt{25 + 24}}{4}$$

$$= \frac{-5 \pm 7}{4}$$

Taking +ve

$$n = \frac{-5 + 7}{4}$$

$$= \frac{1}{2}$$

Taking -ve sign

$$n = \frac{-5 - 7}{4}$$

$$= -3$$

Errors: The deviation from the actual value to the observed value is termed as error.

Types of errors:

1. Inherent errors: Error which are already present in the statement of a problem before its solution, are called Inherent error. Such errors arise either due to the given data being approximate or due to the limitations of mathematical tables, calculators, or the digital computer.

2) Numerical Errors:

Numerical errors (also known as procedural errors) are introduced during the process of implementation of numerical method. They are classified into two classes as:-

a) Roundoff Errors: This type of errors arise from the process of rounding off the numbers during the computation. Such errors are unavoidable in most of the calculations due to the limitations of the computing aids. Roundoff errors can however be reduced:-

- i) By changing the calculation procedure so as to avoid subtraction of nearly equal number or division by small number.
- ii) By retaining at least one more significant figure at each step than that given in the data and rounding off at the last step.

b) Truncation Errors: These errors are caused by using approximate results or replacing an infinite process by a finite one. If we are using a decimal computer having a fixed word length, of 4 digits, rounding off of 13.658 gives 13.66 whereas truncation gives 13.65

Example :

$$\text{If } e^m = 1 + m + \frac{m^2}{2!} + \frac{m^3}{3!} + \dots + \infty = n \text{ (say)}$$

is replaced by

$$1 + m + \frac{m^2}{2!} + \frac{m^3}{3!} = x'; \text{ then the}$$

truncation error is $(x - x')$

If Absolute, Relative and Percentage error

If u is the true value of a equation u' is its approximate value, then $|u - u'|$ is called absolute (E_a) error.

The relative error is given as:

$$Er = \left| \frac{u - u'}{u} \right|$$

and the percentage error is

$$Ep = 100 \left| \frac{u - u'}{u} \right| \%$$

Taylor's Theorem : The Taylor series of a real valued function $f(m)$ that is infinitely differentiable at a real number 'a' is :

$$f(a) + \frac{f'(a)}{1!}(m-a) + \frac{f''(a)}{2!}(m-a)^2 + \frac{f'''(a)}{3!}(m-a)^3 + \dots$$

which can be written in the more compact form as

$$\sum_{n=0}^{\infty} \frac{f^{(n)}(a)}{n!} (m-a)^n$$

where, $n!$ denotes the factorial of n and $f^{(n)}(a)$ denotes the n^{th} derivative of f' evaluated at the point a . The derivative of order zero of f' is defined to be f itself and $(n-n)^0$ and $0!$ both are defined to be 1. When $a=0$ the series is also called MacLaurin series.

Q. Derive the MacLaurin series of $\sin(x)$

$$\Rightarrow \sin(x) =$$

$$f(n) = \sin x$$

$$f'(n) = \cos n$$

$$f''(n) = -\sin n$$

$$f'''(n) = -\cos n$$

$$f''''(n) = \sin n$$

$$f''''''(n) = \cos n$$

$$f(0) = 0$$

$$f'(0) = 1$$

$$f''(0) = 0$$

$$f'''(0) = -1$$

$$f''''(0) = 0$$

$$f''''''(0) = 1$$

MacLaurin series is :-

$$f(n) = n - \frac{n^3}{3!} +$$

$$f(0) + \frac{f'(0)}{1!} + \frac{f''(0)}{2!} n^2 + \frac{f'''(0)}{3!} n^3 + \frac{f''''(0)}{4!} n^4 + \dots$$

$$0 + \frac{1}{1!} n + \frac{0}{2!} n^2 + \frac{1}{3!} n^3 + \frac{0}{4!} n^4 + \frac{1}{5!} n^5 + \dots$$

$$\Rightarrow n - \frac{n^3}{3!} + \frac{n^5}{5!} - \frac{n^7}{7!} + \dots$$

4 Methods of solution:

There are number of methods to find the roots of non-linear eqn as:

- 1) Direct analytical method
- 2) Trial and error method
- 3) Graphical method
- 4) Iterative method

1. Direct Analytical Method:

Equations are solved by using derived formulas. we just need to supply the value of input parameters in the eqn.

Solution of quadratic equation

$$an^2 + bn + c = 0 \text{ is given by}$$

$$n = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

Q. Solve the equation $2n^2 + 5n + 2 = 0$

$$a = 2, b = 5, c = 2$$

$$n = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$= \frac{-5 \pm \sqrt{25 - 4 \times 2 \times 2}}{2 \times 2}$$

$$= \frac{-5 \pm \sqrt{25 - 16}}{4}$$

$$= \frac{-5 \pm 3}{4}$$

Taking +ve

$$\frac{-5 + 3}{4} = \frac{-2}{4} = -\frac{1}{2}$$

Taking -ve

$$\frac{-5 - 3}{4} = \frac{-8}{4} = -2$$

2. Trial and Error method:-

In this method, we make a series of guess for n and ^{we} evaluate the function $f(n)$ at that n , if it is close to zero, it is one of the approximate root of the given non linear equation, otherwise make another guess for n and repeat the same process again until the n for which $f(n)=0$ is found.

Example:

Solve the eqⁿ $f(n) = 2n^2 + n - 6 = 0$

\Rightarrow Sol:

Step 1: Guess $n=0 \Rightarrow f(n) = -6$

Step 2: Guess $n=1 \Rightarrow f(n) = -3$

Step 3: Guess $n=2 \Rightarrow f(n) = 4$

Step 4: Guess $n=-2 \Rightarrow f(n) = 0$

Step 5: Guess $n=\frac{3}{2} \Rightarrow f(n) = 0$

3. Iterative Method:- There are number of iterative methods that have been tried and use successfully in various problem. All these methods, typically generate a sequence of estimators of the solution, which is expected to converge to the true solution.

They are categorized into two classes:-

- 1) Bracketing methods (Interpolation method)
- a. Bisection method (half interval)
- b. False position method

- ??) Open-end method (Extrapolation method)
- Newton-Raphson method
 - Secant Method
 - Muller's method
 - Fixed point method
 - Bairstow's method

Bisection Method :-

The bisection method is one of the simplest and most reliable iterative methods for the solutions of non-linear equation. It is also known as half-interval method, relies on the fact that if $f(n)$ is continuous in the interval $a < n < b$, and $f(a)$ and $f(b)$ are of opposite signs i.e $f(a) \times f(b) < 0$,

then, there is at least one real root in the interval between a and b .

Let $n_1 = a$ and $n_2 = b$, let us also define another point n_0 to the mid-point bet " a and b ", i.e

$$n_0 = \frac{a+b}{2}$$

Now, there exists the following three conditions :-

1. If $f(n_0) = 0$, we have a root at n_0 .
2. If $f(n_0) \cdot f(n_1) < 0$, there exist a root betⁿ n_0 and n_1 .
3. If $f(n_0) \cdot f(n_2) < 0$, there is a root betⁿ n_0 and n_2 .

Pictorially, it can be represented as:

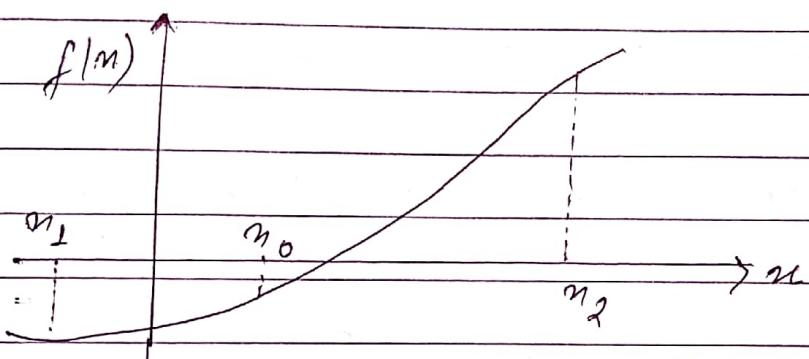
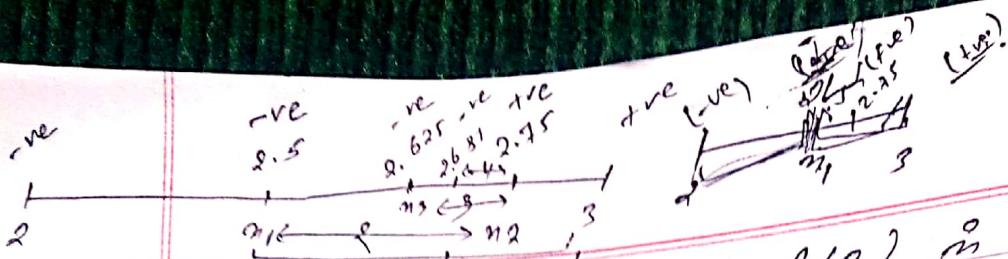


fig: Illustration of Bisection Method.

From figure $f(n_0)$ and $f(n_1)$ are of opposite sign, a root lies betⁿ n_0 and n_1 , we can further divide this subinterval into two halves to locate a new subinterval containing the root. This process can be repeated until the interval containing the root.

Q Example: Find the root of the equation $n^3 - 4n - 9 = 0$, using the bisection method within the range $2 < n < 3$ correct upto three decimal places.

\Rightarrow Sol'n: // Let $f(n) = n^3 - 4n - 9$.
 $f(2) = 8 - 8 - 9 = -9$ i.e. -ve
 $f(3) = 27 - 12 - 9 = 6$ i.e. +ve



Since $f(2)$ is -ve and $f(3)$ is +ve, so root lies between 2 and 3.

\therefore First approximation to the root is

$$n_1 = \frac{2+3}{2} = 2.5$$

$$\begin{aligned} f(n_1) &= f(2.5) = (2.5)^3 - 4 \times 2.5 - 9 \\ &= -3.375 \text{ (ie -ve)} \end{aligned}$$

\therefore The root lies betⁿ $n_1 = 2.5$ and 3

Second approximation to the root is

$$n_2 = \frac{2.5+3}{2} = 2.75$$

$$\begin{aligned} f(n_2) &= f(2.75) = (2.75)^3 - 4 \times 2.75 - 9 \\ &= 0.79 \text{ ie +ve} \end{aligned}$$

\therefore The root lies betⁿ $n_1 = 2.5$ and $n_2 = 2.75$

Third approximation to the root is

$$n_3 = \frac{2.5+2.75}{2} = 2.625$$

$$f(n_3) = (2.625)^3 - 4 \times 2.625 - 9 = -1.413 \text{ ie -ve}$$

\therefore The root lies betⁿ $n_2 = 2.75$ and $n_3 = 2.625$

Fourth approximation to the root is

$$n_4 = \frac{2.75+2.625}{2} = 2.6875$$

$$f(n_4) = (2.6875)^3 - 4 \times 2.6875 - 9 = -0.399 \text{ ie -ve}$$

\therefore The root lies betⁿ $n_3 = 2.625$ and $n_4 = 2.6875$

Fifth approximation to the root is : 97185

$$n_5 = \frac{2.6875 + 2.75}{2} = 2.7185$$

$$f(1.85) = (2.7185)^3 - 4 \times 2.7185 - 9 = 0.2209 \text{ re}$$

The root lies betw' my and me'

Sixth approximation to the root is 1.

$$n_6 = \frac{2.6875 + 2.7185}{2} = 2.7031$$

$$f(176) = (2.7031)^3 - 4 \times 2.7031 - 9 = -0.0610 \text{ ie -ve}$$

\therefore The root lies between $m_6 = 2.7031$ and $m_5 = 2.7185$

Seventh approximation to the root is 1.

$$m_7 = \frac{2.7031 + 2.7185}{2} = 2.7109$$

$$f(1.7) = (2.7109)^3 - 4 \times (2.7109) - 9 \\ = 0.076 \text{ re + re}$$

\therefore The root lies betw $n_6 = 2.7031$ and $n_7 = 2.7109$

Eighth approximation to the root is:

$$m_8 = \frac{2.7031 + 2.7109}{2} = 2.7071$$

$$f(1.8) = (2.7071)^3 - 4 \times 2.7071 - 9 = 9.04 \times 10^{-3} \text{ re-arrange}$$

The root error had $m_8 = 2.7097$ and $m_6 = 2.7031$

Ninth approximation to the root is :-

$$Mg = \frac{2.7031 + 2.7071}{2} \\ = 2.7050$$

$$a_n n^n + a_{n-1} n^{n-1} + \dots + a_0 = 0$$

$$(m_{\max}) = \sqrt{\left(\frac{a_{n-1}}{a_n}\right)^2 - 2\left(\frac{a_{n-2}}{a_n}\right)}$$

$$f(m_1) = -0.026 \text{ is -ve}$$

\therefore The root lies between $m_1 = 2.7071$ and $m_2 = 2.7050$

Tenth approximation to the root :-

$$m_{10} = \frac{2.7050 + 2.7071}{2} = 2.70605$$

$$f(m_{10}) = -8.5 \times 10^{-3} \text{ is -ve}$$

\therefore The root lies between $m_{10} = 2.7060$ and $m_8 = 2.7071$

Eleventh approximation to the root is :-

$$m_{11} = \frac{2.70605 + 2.7071}{2} = 2.70695$$

$$f(m_{11}) = 2.69 \times 10^{-4} \text{ is +ve}$$

\therefore The approximate root is 2.706

Algorithm :-

1. Start

2. Decide initial values for m_1 and m_2 and stopping criteria (E_{toler})

3. Calculate $f_1 = f(m_1)$ and $f_2 = f(m_2)$

4. If $f_1 \neq f_2 > 0$ and m_1 and m_2 does not contain any root and goto step 8. otherwise continue.

5. Calculate $m_0 = (m_1 + m_2)/2$ and compute $f_0 = f(m_0)$

6. If $f_0 \neq f_2 < 0$ then set $m_2 = m_0$ } $f_1 \neq f_0 < 0$ otherwise

Set $m_2 = m_0$

Set $f_2 = f_0$

Set $m_2 = m_0$

otherwise

Set $m_1 = m_0$

Set $f_1 = f_0$

$$|a-b| < 0.0005 \text{ or } (a+b)/2 < 0.0005$$

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7. If absolute value of $(n_2 - n_1)/n_2$ is less than error ϵ , then

$$\text{root} = (n_1 + n_2)/2$$

print the value of root.
goto step 8.

else

goto step 5

8. STOP

- Q. Find the root of equation $x^2 - 4x - 10 = 0$ using bisection method within interval $-2 \leq x \leq -1$

Solution:

$$\text{Let, } f(x) = x^2 - 4x - 10$$

$$f(-2) = (-2)^2 - 4(-2) - 10 \\ = 2 \text{ ie +ve}$$

$$f(-1) = (-1)^2 - 4(-1) - 10 \\ = -5 \text{ ie -ve}$$

So, root lies betw -2 and -1.

$$\therefore n_1 = \frac{-2 - (-1)}{2} = -1.5$$

$$f(n_1) = -1.75 \text{ ie -ve}$$

The root lies betw -1.5 and -2.

$$n_2 = \frac{-1.5 - (-2)}{2} = -1.75$$

$$f(n_2) = 0.0625 \text{ ie +ve}$$

\therefore The root lies betw -1.5 and -1.75

$$\therefore n_3 = \frac{-1.5 - 1.75}{2}$$

$$f(n_3) = -0.85937 \text{ re-ve}$$

\therefore Root lies betⁿ -1.75 and -1.625

$$n_4 = \frac{-1.75 + 1.625}{2} = -1.6875$$

$$f(n_4) = -0.402 \text{ re-ve}$$

\therefore Root lies betⁿ -1.75 and -1.6875

$$n_5 = \frac{-1.75 - 1.6875}{2} = -1.71875$$

$$f(n_5) = -1.017 \text{ re-ve}$$

\therefore Root lies betⁿ -1.75 to -1.71875.

$$n_6 = \frac{-1.758 - 1.71875}{2} = -1.734$$

$$n_7 = -1.742$$

$$n_8 = -1.738$$

$$n_9 \approx -1.740$$

$$n_{10} \approx -1.741$$

$$n_{11} \approx -1.742$$

Secant Method:

Secant method, like bisection method, uses two initial estimates but does not require that they must bracket the root.

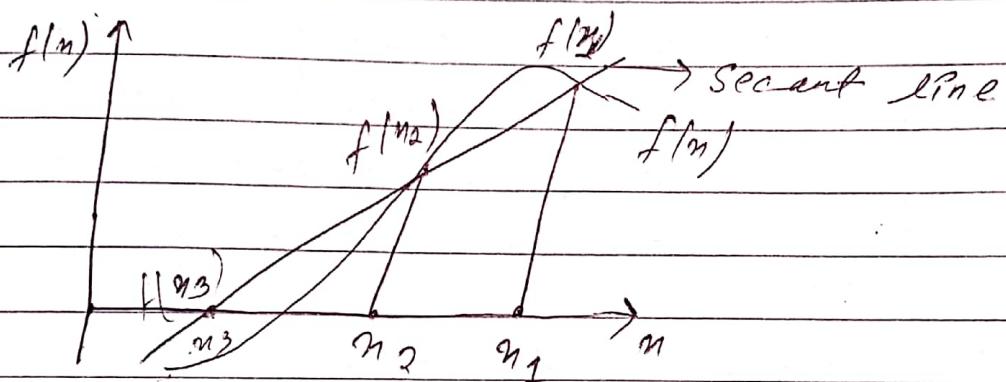


Fig:- Graphical representation of secant method.

Let n_1 and n_2 as starting values they do not bracket the root. slope of the secant line passing through n_1 and n_2 is given by :

$$\frac{f(n_1)}{n_1 - n_3} = \frac{f(n_2)}{n_2 - n_3}$$

$$\therefore f(n_1)(n_2 - n_3) = f(n_2)(n_1 - n_3)$$

$$\Rightarrow n_3 [f(n_2) - f(n_1)] = f(n_2)n_1 - f(n_1)n_2$$

$$\Rightarrow n_3 = \frac{f(n_2)n_1 - f(n_1)n_2}{f(n_2) - f(n_1)}$$

By adding and subtracting $f(n_2)/n_2$ to the numerator and rearranging the terms we get,

$$n_3 = n_2 - \frac{f(n_2)(n_2 - n_1)}{f(n_2) - f(n_1)}$$

This formula is known as secant formula.

General form for i^{th} root,

$$n_i^o = n_{i-1} - \frac{f(n_{i-1}^o)(n_i^o - n_{i-2}^o)}{f(n_{i-1}^o) - f(n_{i-2}^o)}$$

The approximate value of root can be defined by repeating this procedure by replacing n_1 and n_2 by n_2 and n_3 respectively.

The next approximate value is given as:

$$n_4 = n_3 - \frac{f(n_3)(n_3 - n_2)}{f(n_3) - f(n_2)}$$

$$n_{i+1} = n_i - \frac{f(n_i)(n_i - n_{i-1})}{f(n_i) - f(n_{i-1})}$$

\checkmark Q 1st q direction

#include < stdio.h >

#include < conio.h >

#define eqn $x^4 + 2x^3 - 2x^2 + x - 5$

float f (float n)

{

 float ans;

 ans = eqn;

 return (ans);

}

main()

Q. Use the Secant method to find the root of the equation $n^2 - 4n - 10 = 0$ with initial estimates of $n_1 = 4$ and $n_2 = 2$, calculate upto 6th iterations.

\Rightarrow Sol'n:

Let given eqn is $n^2 - 4n - 10 = f(n)$

$$f(n) = n^2 - 4n - 10$$

$$n_1 = 4, n_2 = 2$$

Now,

$$f(n_1) = 4^2 - 4 \times 4 - 10 = -10$$

$$f(n_2) = 2^2 - 4 \times 2 - 10 = -14$$

: The first approximation to the root

$$\therefore n_3 = n_2 - \frac{f(n_2)(n_2 - n_1)}{f(n_2) - f(n_1)}$$

$$= 2 - \frac{-14 \times 2}{-14 + 10}$$

$$= 9$$

$$f(n_3) = 9^2 - 4 \times 9 - 10 = 35$$

: The second approximation to the root is

$$n_4 = n_3 - \frac{f(n_3)(n_3 - n_2)}{f(n_3) - f(n_2)}$$

$$= 9 - \frac{35 \times (8.9 - 2)}{35 + 14}$$

$$= 4$$

$$f(n_4) = 4^2 - 4 \times 4 - 10 = -10$$

\therefore The third approximation of root

$$\begin{aligned} n_5 &= n_4 - \frac{f(n_4)}{f(n_4) - f(n_3)} \\ &= 4 - \frac{-10}{-10 - 9} \\ &= 5.1111 \end{aligned}$$

$$f(n_5) = (5.1111)^2 - 4 \times 5.1111 - 10 = -4.3210$$

\therefore The fourth approximation to the root :

$$\begin{aligned} n_6 &= n_5 - \frac{f(n_5)}{f(n_5) - f(n_4)} \\ &= 5.1111 - \frac{-4.3210}{-4.3210 + 10} \\ &= 5.9565 \end{aligned}$$

$$f(n_6) = 1.6538$$

\therefore The fifth approximation to the root is :

$$\begin{aligned} n_7 &= n_6 - \frac{f(n_6)}{f(n_6) - f(n_5)} \\ &= 5.9565 - \frac{1.6538}{1.6538 + 4.3210} \\ &= 5.7224 \end{aligned}$$

$$f(n_7) = -0.1437$$

\therefore The sixth approximation to the root :

$$\begin{aligned} n_8 &= n_7 - \frac{f(n_7)}{f(n_7) - f(n_6)} \\ &= 5.7224 - \frac{-0.1437}{-0.1437 - 1.6538} \\ &= 5.7411 \end{aligned}$$

$$f(n_8) = -4.27079 \times 10^{-5}$$

Algorithm for secant method:

1. Start
2. Decide two initial points n_1 and n_2 , accuracy level E .
3. Compute $f_1 = f(n_1)$ and $f_2 = f(n_2)$.
4. Compute $n_3 = n_2 - \frac{f_2(n_2 - n_1)}{f_2 - f_1}$
5. Test for accuracy of n_3 if $\left| \frac{n_3 - n_2}{n_3} \right| > E$, then
 - set $n_1 = n_2$ and $f_1 = f_2$
 - set $n_2 = n_3$ and $f_2 = f(n_3)$
 - goto step 4
 otherwise,
 - set $x_{\text{root}} = n_3$
 - print the value of x_{root} . goto step 6
6. STOP
7. Find the root of the equation $n e^n = \cos n$ using the secant method. Correct to four decimal places. Take initial estimates are $n_1 = 0$ and $n_2 = 1$.

\Rightarrow for n :

The given eqn is $f(n) = n e^n - \cos n$

$$f(n_1) = 0 \cdot e^0 - \cos 0 = 0 - 1 = -1$$

$$f(n_2) = 1 \cdot e^1 - \cos 1 = 2.71828 - 0.99984 \\ = 1.71844$$

First approximation to the root is:

$$n_3 = n_2 - \frac{f(n_2)(n_2 - n_1)}{f(n_2) - f(n_1)}$$

$$= 1 - \frac{1.71844(1-0)}{1.71844+1}$$

$$= 0.36785$$

$$f(n_3) = 0.36785 e^{(0.36785)} - \cos(0.36785)$$

$$= -0.46857$$

∴ Second approximation to the root is

$$n_4 = n_3 - \frac{f(n_3)(n_3 - n_2)}{f(n_3) - f(n_2)}$$

$$= 0.36785 - \frac{-0.46857(0.36785-1)}{-0.46857 - 1.71844}$$

$$\approx 0.50328$$

$$f(n_4) = -0.16746$$

∴ Third approximation to the root is

$$n_5 = n_4 - \frac{f(n_4)(n_4 - n_3)}{f(n_4) - f(n_3)}$$

$$= 0.50328 - \frac{-0.16746(0.50328 - 0.36785)}{-0.16746 + 0.46857}$$

$$\approx 0.57859$$

$$f(n_5) = 0.03197$$

∴ Fourth approximation to the root is:

$$n_6 = n_5 - \frac{f(n_5)(n_5 - n_4)}{f(n_5) - f(n_4)}$$

$$= 0.57859 - \frac{0.03197(0.57859 - 0.50328)}{0.03197 + 0.16746} \approx 0.56651$$

$$f(n_6) = -0.0017$$

\therefore Fifth approximation to the root is

$$n_7 = n_6 - \frac{f(n_6)}{f(n_7) - f(n_6)} (0.56651 - 0.57859)$$

$$= 0.56711 - \frac{-0.0017}{-0.0017} = 0.03197$$

$$\therefore f(n_7) = -0.000043$$

\therefore Sixth approximation to the root is

$$n_8 = n_7 - \frac{f(n_7)}{f(n_8) - f(n_7)} (n_7 - n_6)$$

$$= 0.56711 - \frac{-0.000043 (0.56711 - 0.56651)}{-0.000043 + 0.0017}$$

$$= 0.56712$$

$$\therefore f(n_8) = -0.000053$$

\therefore The approximation root of $e^x - ne^n = 0$ is

$$\approx 0.56712$$

Newton-Raphson Method:

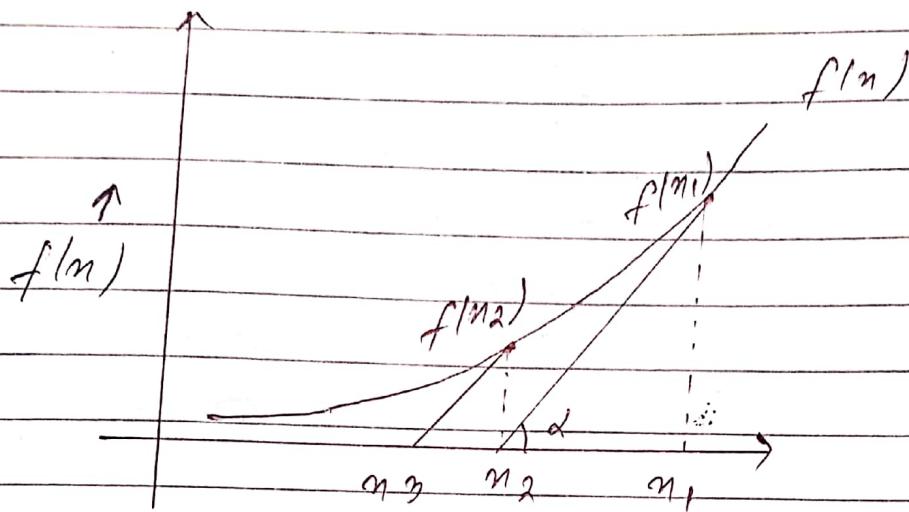


Fig. - Newton-Raphson method

Consider a graph $f(n)$, let us assume that n_1 is an approximate root of $f(n) = 0$. Now, draw a tangent at the curve $f(n)$ at $n = n_1$ as shown in figure. The point of intersection of this tangent with n -axis gives the second approximation to the root.

Let this point of the intersection be n_2 . The slope of tangent is given as,

$$\text{slope} = \frac{f(n_1)}{n_1 - n_2} = f'(n_1)$$

where $f'(n_1)$ is the slope of $f(n)$ at $n = n_1$.

Now, solving for n_2 , we get

$$\frac{f(n_1)}{n_1 - n_2} = f'(n_1)$$

$$\text{or, } n_2 = n_1 - \frac{f(n_1)}{f'(n_1)}$$

The next approximate would be

$$n_3 = n_2 - \frac{f(n_2)}{f'(n_2)}$$

In general,

$$n_{n+1} = n_n - \frac{f(n_n)}{f'(n_n)}$$

which is known as

Newton Iteration formula.

- Q. Find the root of the equation $f(n) = n^2 - 3n + 2$ in the vicinity of $n=0$ using Newton-Raphson method.

\Rightarrow Sol'n:

The given eqⁿ is : $f(n) = n^2 - 3n + 2$

$$f'(n) = 2n - 3$$

Initial point $n_0 = 0$

Now, $f(n_0) = 0 - 3 \times 0 + 2 = 2$

$$f'(n_0) = -3$$

. The first approximation to the root is

$$n_1 = n_0 - \frac{f(n_0)}{f'(n_0)}$$

$$= 0 - \frac{2}{-3}$$

$$\approx 0.6666$$

$$f(n_1) = 0.4445$$

$$f'(n_1) = -1.6668$$

: second approximation

$$n_2 = n_1 - \frac{f(n_1)}{f'(n_1)}$$

$$= 0.6666 - \frac{0.4445}{-1.6668}$$

$$= 0.9332$$

$$f(n_2) = -0.222 + 0.0712$$

$$f'(n_2) = -1.8886$$

Third approximation :-

$$n_3 = n_2 - \frac{f(n_2)}{f'(n_2)}$$

$$= 0.9332 - \frac{0.0712}{-1.1336}$$

$$= \text{true} 0.9960$$

$$f(n_3) = 0.0040$$

$$f'(n_3) = -1.008$$

: fourth approximation :-

$$n_4 = n_3 - \frac{f(n_3)}{f'(n_3)}$$

$$= 0.9960 + \frac{0.0040}{-1.008}$$

$$= 0.9999$$

$$f(n_4) = 0.0002$$

$$f'(n_4) = -1.0003$$

Fifth Approximation

$$\begin{aligned} f(n_5) &= n_4 - \frac{f(n_4)}{f'(n_4)} \\ &= 0.9999 + \frac{0.0001}{1.0002} \\ &= 0.9999 \end{aligned}$$

$$f(n_5) = 0.9999$$

$$f'(n_5) = -1.0002$$

∴ The approximation root of the eqⁿ is 0.9999.

Algorithm:

1. Start
 2. Assign an initial value to n_0 say n_0 , define error.
 3. Evaluate $f(n_0)$ and $f'(n_0)$
 4. Find the improved estimate of n_0 .
- $$n_1 = n_0 - \frac{f(n_0)}{f'(n_0)}$$
5. Check the accuracy of the latest estimate compare relative error to a predefined error value E .
If $\left| \frac{n_1 - n_0}{n_1} \right| \leq E$ stop,
otherwise continue.
 6. Replace n_0 by n_1 and repeat step 4 and step 5.

L8. Solve $\cos n = 2n$ using Newton-Raphson method, correct upto 5 decimal places with initial guess $n_0 = 0.5$

Q8. Find the root of the equation by Newton-Raphson's method $8n = \cos n$ with initial guess $n_0 = 0.6$ correct upto 4 decimal places.

Ans

$$f(n) = \cos n - 2n$$

$$f'(n) = -\sin n - 2$$

Initial point $n_0 = 0.5$

Now,

$$f(n_0) = \cos 0.5 - 2 \times 0.5 = -0.000038$$

$$f'(n_0) = -\sin 0.5 - 2 = -2.00872$$

∴ The first approximation to the root :-

$$n_1 = n_0 - \frac{f(n_0)}{f'(n_0)}$$

$$= 0.5 - \frac{0.000038}{-2.00872}$$

$$= 0.49998$$

$$f(n_1) = 0.0000019$$

$$f'(n_1) = -2.00872$$

∴ Second approximation to the root :-

$$n_2 = n_1 - \frac{f(n_1)}{f'(n_1)} = 0.49998 + \frac{0.0000019}{-2.00872}$$

$$= 0.49998$$

$$f(n_2) = 0.0000019$$

$$f'(n_1) = -2.00872$$

∴ The approximated root
is 0.49998

$$f(n_3) = -0.00013$$

$$f'(n_3) = 3.01163$$

∴ The approximated root is 0.6666

Derivation of Newton-Raphson method using Taylor series expansion.

⇒ Let n_0 be an approximate root of the equation $f(n)=0$. If $n_1 = n_0 + h$ be the exact root then $f(n_1) = 0$

Expanding $f(n_0+h)$ by Taylor series:

$$f(n_0) + h \frac{f'(n_0)}{1!} + h^2 \frac{f''(n_0)}{2!} + \dots$$

Since, h is small neglecting h^2 and higher power of h and equilize with zero.

we get,

$$f(n_0) + h f'(n_0) = 0$$

$$h = -\frac{f(n_0)}{f'(n_0)}$$

∴ A closer approximation to the root is given by

$$n_1 = n_0 - \frac{f(n_0)}{f'(n_0)}$$

Similarly, starting with n_1 , a still better approximation n_2 is given by

$$n_2 = n_1 - \frac{f(n_1)}{f'(n_1)}$$

In general,

$$n_{n+1} = n_n - \frac{f(n_n)}{f'(n_n)}$$

Fixed Point Method :-

Any function of the form $f(n)=0$ can be manipulated such that n is on the left hand side of the equation as,

$$n = \phi(n) \quad \text{--- (2)}$$

Equation (1) & (2) are equivalent and therefore, a root of eqⁿ (1) is also root of eqⁿ (2). The root of the eqⁿ (2) is given by the point of intersection of the curve.

$y = n$ and $y = \phi(n)$. This point of intersection is known as fixed point of $\phi(n)$.

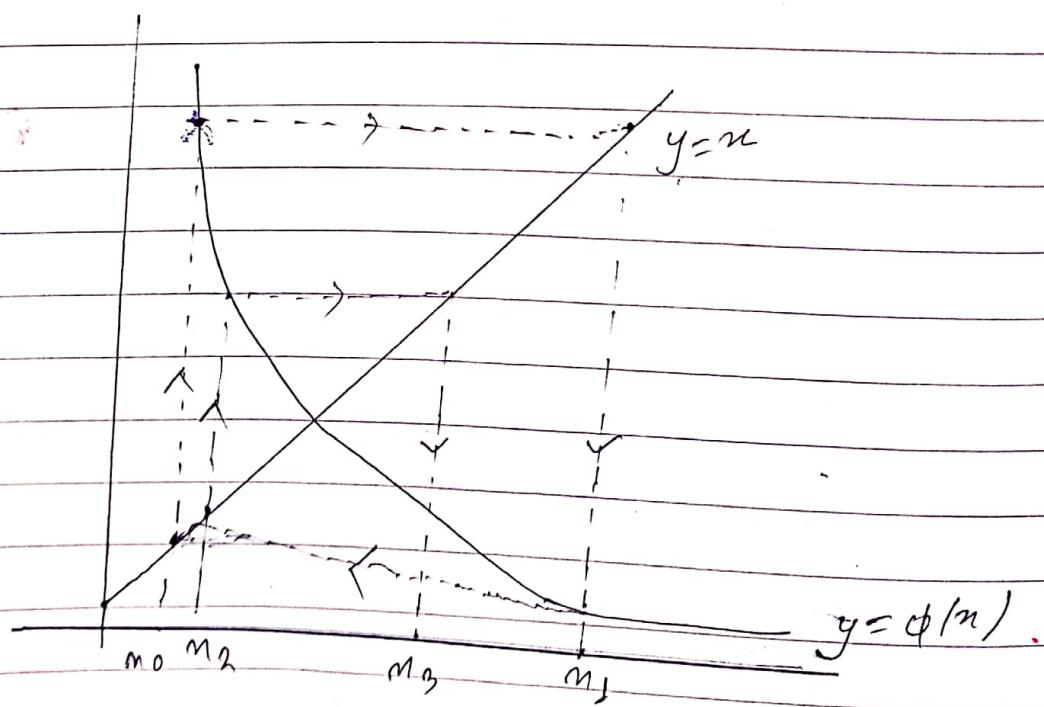


fig:- Geometrical interpretation of fixed point method

The equation $n = \phi(n)$ is known as the fixed point equation. It provides a convenient form for predicting the value of n as function of n . If n_0 is the initial guess to a root, then the next approximation is given by,

$$n_1 = \phi(n_0)$$

Further approximation is given by,

$$n_2 = \phi(n_1)$$

This iteration process can be expressed in general form as

$n_{i+1} = \phi(n_i)$, which is called fixed point iteration formula.

8. Find the root of the equation $\cos n = 3n - 1$, correct upto three decimal places using the fixed point method. Take initial guess as

$$n_0 = 0$$

\Rightarrow soln:

The given eqⁿ : $f(n) = \cos n - 3n + 1$

$$3n = \cos n + 1$$

$$n = \frac{\cos n + 1}{3} \text{ ie } \phi(n)$$

Initial guess : $n_0 = 0$

\therefore First approximation to the root :

$$n_1 = \phi(n_0) = \frac{\cos 0 + 1}{3}$$

$$= 0.666$$

Newton-Raphson method for system of non-linear equations.

The first order Taylor series is given as :

$$f(n_i+1) = f(n_i) + (n_i+1 - n_i) f'(n_i) \quad \text{--- (1)}$$

is used to derive the Newton's iteration formula

$$n_{i+1} = n_i - \frac{f(n_i)}{f'(n_i)} \quad \text{--- (2)}$$

For simplicity, let us again consider two equation of non-linear system.

$$f(n, y) = 0 \quad \text{and} \quad g(n, y) = 0$$

First order Taylor series of these eqⁿ can be written as :

$$f(n_i+1, y_i+1) = f(n_i, y_i) + (n_i+1 - n_i) \left[\frac{\partial f_i}{\partial n} \right] + (y_i+1 - y_i) \left[\frac{\partial f_i}{\partial y} \right] \quad \text{--- (3)}$$

$$g(n_i+1, y_i+1) = g(n_i, y_i) + (n_i+1 - n_i) \left[\frac{\partial g_i}{\partial n} \right] + (y_i+1 - y_i) \left[\frac{\partial g_i}{\partial y} \right] \quad \text{--- (4)}$$

If the root estimated are n_{i+1} and y_{i+1} then,
 $f(n_{i+1}, y_{i+1}) = g(n_{i+1}, y_{i+1}) = 0$

Now,

Substituting this in eqⁿ (3) & (4) respectively, we get the following two linear eqⁿ:

$$\Delta n f_1 + \Delta y f_2 + f = 0 \quad \text{--- (5)}$$

$$\Delta n g_1 + \Delta y g_2 + g = 0 \quad \text{--- (6)}$$

where, we denote,

$$\Delta n = n_i^o + 1 - n_i^o$$

$$\Delta y = y_{i+1}^o - y_i^o$$

$$f_1 = \left| \frac{\partial f_1}{\partial n} \right|, \quad f_2 = \left| \frac{\partial f_1}{\partial y} \right|, \quad g_1 = \left| \frac{\partial g_1}{\partial n} \right|, \quad g_2 = \left| \frac{\partial g_1}{\partial y} \right|$$

$$f = f(n_i^o, y_i^o), \quad g = g(n_i^o, y_i^o)$$

solving for n and y we get,

$$\Delta n = \frac{-f \cdot g_2 - g \cdot f_2}{f_1 g_2 - f_2 g_1} = -\frac{\Delta y}{D} \quad \text{--- (7)}$$

$$\Delta y = \frac{-g f_1 - f g_1}{f_1 g_2 - f_2 g_1} = -\frac{\Delta y}{D} \quad \text{--- (8)}$$

where,

$$D = \begin{vmatrix} f_1 & f_2 \\ g_1 & g_2 \end{vmatrix} = f_1 g_2 - f_2 g_1, \text{ is called}$$

Jacobian matrix

from eqn (7) & (8) we can establish the following relations.

$$n_{i+1}^o = n_i^o - \frac{\Delta n}{D} \quad \text{--- (9)}$$

$$y_{i+1}^o = y_i^o - \frac{\Delta y}{D} \quad \text{--- (10)}$$

eqn (9) and (10) are called the two eqn
Newton's formula.

Example: Determine the root of the eqⁿ

$$n^2 + ny = 6$$

$$n^2 - y^2 = 3 \text{ using Newton-Raphson method}$$

with Initial guess $n_0=1$ and $y_0=1$

\Rightarrow Solⁿ:

$$\text{Given eq}^n \text{ are: } f(n,y) = n^2 + ny - 6$$

$$g(n,y) = n^2 - y^2 - 3$$

Initial case: $n_0=1, y_0=1$

$$f(n_0, y_0)$$

$$f_1 = \frac{\partial f}{\partial n} = 2n + y, \quad g_1 = \frac{\partial g}{\partial n} = 2n$$

$$f_2 = \frac{\partial f}{\partial y} = n, \quad g_2 = \frac{\partial g}{\partial y} = -2y$$

First:

We have $n_0=1$, and $y_0=1$

in case of f_1 & f_2 Iteration 1:

$$g_1 = g_2 = f_1 = 2 + 1 = 3, \quad f_2 = 1$$

$$f_1 = 2, \quad f_2 = -2$$

$$\therefore D = \begin{vmatrix} f_1 & f_2 \\ g_1 & g_2 \end{vmatrix}$$

$$= \begin{vmatrix} 3 & 1 \\ 2 & -3 \end{vmatrix}$$

$$= -6 - 2$$

$$= -8$$

The value of functions at n_0 and y_0 ,

$$f = 1 + 1 - 6 = -4$$

$$g = 1 - 1 - 3 = -3$$

Algorithm

1. Start
 2. Define the functions f and g .
 3. Define the Jacobian elements f_1, f_2, g_1 and g_2 .
 4. Decide starting points m_0 and y_0 and error tolerance E .
 5. Evaluate $f(y), f_1, f_2, g_1, g_2$ at (m_0, y_0) .
Compute D_m, D_y and D .
 $m_1 = m_0 - D^{-1}D$
 $y_1 = y_0 - D_y^{-1}D$
 6. Test for accuracy
If $|m_1 - m_0| < E$ and
 $|y_1 - y_0| < E$, then
root obtained;
goto step 8
 7. Otherwise, set $m_0 = m_1$
 $y_0 = y_1$
goto step 5
 8. print the root
 9. Stop.
1. Start
2. $f(u) = 0, m_0$
3. convert $f(u) = 0$ into form $u = g(u)$
4. initial guess m_0
do 4. $m_1 = g(m_0)$
 $m_{i+1} = g(m_i)$
5. Error $|m_{i+1} - m_i| > E$ $|m_1 - m_0| > E$
 $m_{i+1} = m_i$ $m_0 = m_1$ goto step 4
6. Stop. otherwise root = m_1

Synthetic Division :-

Synthetic division is a method of performing polynomial division, with less writing and fewer calculations. It is mostly used for division by binomials of the form $(n-a)$. The most useful aspects of synthetic division is that, it allows to calculate without writing variables and uses fewer calculations.

Synthetic division is performed as follows:

Let $p(n)$ be a polynomial of degree n . If we divide $p(n)$ by $(n-a)$ we get another polynomial $q(n)$ of degree $n-1$.

Assume,

$$p(n) = \alpha_n n^n + \alpha_{n-1} n^{n-1} + \dots + \alpha_0$$

$$\text{and } q(n) = b_{n-1} n^{n-1} + b_{n-2} n^{n-2} + \dots + b_0$$

$$\Rightarrow p(n) = (n-a) \cdot q(n) [\because q(n) = \text{deflated polynomial}]$$

$$\alpha_n n^n + \alpha_{n-1} n^{n-1} + \dots + \alpha_0 = (n-a)(b_{n-1} n^{n-1} + b_{n-2} n^{n-2} + \dots + b_0)$$

Comparing the coefficient on both sides we get,

$$b_{n-1} = \alpha_n / [\alpha_n + a * b_{n-1}]$$

$$b_{n-2} = \alpha_{n-1} + a * b_{n-2}$$

!

$$b_0 = \alpha_1 + a * b_1$$

Thus,

$$b_{p-1} = \alpha_p + a * b_p, \text{ where } p=n, n-1, \dots, 1$$

$$\text{and } b_n = 0$$

This process is also known as deflation process.

Example:

The polynomial equation $p(n) = n^3 - 7n^2 + 15n - 9$ has root at $n=3$. Find the quotient polynomial $q(n)$ such that $p(n) = (n-3)q(n)$

for n :

The given eqn is:

$$P(n) = n^3 - 7n^2 + 15n - 9 = 0$$

We have,

$$\alpha_3 = 1, \alpha_2 = -7, \alpha_1 = 15, \alpha_0 = -9$$

$$\text{Now, } b_n = 0 \text{ i.e. } b_3 = 0 \quad [\text{general method}]$$

$$b_{n-1} = b_2 = \alpha_3 = 1 \quad [b_2 = \alpha_3 + \alpha_2 * b_3]$$

$$b_{n-2} = b_1 = \alpha_2 + \alpha_1 * b_2$$

$$= -7 + 3 * 1$$

$$= -4$$

$$b_{n-3} = b_0 = \alpha_1 + \alpha_0 * b_1$$

$$= 15 + 3 * (-4)$$

$$= 9$$

$$q(n) \Rightarrow b_2 n^2 + b_1 n + b_0 = 0$$

$$\Rightarrow n^2 - 4n + 3 = 0$$

\therefore The quotient polynomial $q(n)$ is
 $n^2 - 4n + 3 = 0$

Algorithm:

1. Start
2. Read coefficient of dividend polynomial say $\alpha_n, \alpha_{n-1}, \dots, \alpha_0$
3. Read magnitude of division polynomial by the form n^2 .
4. Set $b_n = 0$.
5. Repeat till $n > 0$
 - $b_{n-1} = \alpha_n + b_n * \alpha$
 - 6. remainder (R) = $\alpha_0 + b_{n-1} * \alpha$
 - 7. Display quotient polynomial & remainder
 - 8. STOP

18. Divide $2n^5 - 3n^4 + 4n^3 - 5n^2 + 6n - 9$ by $n^2 - n + 2$ synthetically.

2. Using Newton-Raphson method, find a root of the equation $n^4 + n^3 - 7n^2 - n + 5 = 0$ which lies between 2 and 3, correct upto 4 decimal places.

\Rightarrow b_n^n :

$$P(n) \Rightarrow 2n^5 - 3n^4 + 4n^3 - 5n^2 + 6n - 9 = 0 \text{ we get,}$$

$$\alpha_5 = 2, \alpha_4 = -3, \alpha_3 = 4, \alpha_2 = -5, \alpha_1 = 6, \alpha_0 = -9$$

$$\text{Now, } b_n = 0, i.e. b_5 = 0$$

$$\text{Then: } n^2 - n - 2 = 0$$

$$\Rightarrow n^2 - 2n + n - 2 = 0$$

$$\Rightarrow n(n-2) + 1(n-2) = 0$$

$$\Rightarrow (n+1)(n-2) = 0$$

Either, or

$$n+1=0 \quad n-2=0$$

$$n=-1 \quad n=2$$

$$\therefore n = \underline{-1}, 2$$

The root is 2, $b_5 = 0, \alpha = 2$

$$\begin{aligned} b_{n-1} &= b_{5-1} = b_4 = \alpha_5 + \alpha * b_5 \\ &= 2 + 2 * 0 \\ &= 2 \end{aligned}$$

$$\begin{aligned} b_{n-1} &= b_3 = \alpha_4 + \alpha * b_4 \\ &= -3 + 2 * 2 \\ &= -3 + 4 \\ &= 1 \end{aligned}$$

$$\begin{aligned} b_2 &= a_2 + \alpha \times b_1 \\ &= 4 + 2 \times 6 \\ &= 6 \end{aligned}$$

$$\begin{aligned} b_1 &= a_2 + \alpha \times b_2 \\ &= -5 + 2 \times 6 \\ &= -5 + 12 \\ &= 7 \end{aligned}$$

$$\begin{aligned} b_0 &= a_1 + \alpha \times b_1 \\ &= 6 + 2 \times 7 \\ &= 6 + 14 \\ &= 20 \end{aligned}$$

$$\begin{aligned} q(n) &\Rightarrow b_4 n^4 + b_3 n^3 + b_2 n^2 + b_1 n + b_0 = 0 \\ &\Rightarrow 2n^4 + n^3 + 6n^2 + 7n + 20 = 0 \end{aligned}$$

\therefore The quotient polynomial $q(n)$ is
 $2n^4 + n^3 + 6n^2 + 7n + 20 = 0$

~~2 \Rightarrow for n :~~

$$f(n) = n^4 + n^3 - 7n^2 - n + 5 = 0$$

$$f'(n) = 4n^3 + 3n^2 - 14n - 1 = 0$$

Initial guess: $n_0 = 3$

$$\text{Then, } f(n_0) = 2^4 + 2^3 - 7 \times 3^2 - 2 + 5 = -1$$

$$f'(n_0) = 15$$

\therefore 1st approximation to the root is:

$$n_1 = n_0 - \frac{f(n_0)}{f'(n_0)}$$

$$= 2 - \frac{1}{15}$$

$$= 2.0666$$

$$f(n_1) = (2.0666)^4 + (2.0666)^3 - 7x(2.0666) - 2.0666 + 5 \\ = 0.10478$$

$$f'(n_1) = 18.18753$$

\therefore Second approximation to the root \therefore

$$n_2 = n_1 - \frac{f(n_1)}{f'(n_1)}$$

$$= 2.0666 - \frac{0.10478}{18.18753}$$

$$= 2.06089$$

$$f(n_2) = 0.08067$$

$$f'(n_2) = 17.90194$$

\therefore Third approximation to the root:

$$n_3 = n_2 - \frac{f(n_2)}{f'(n_2)}$$

$$= 2.0608$$

\therefore The approximated root is 2.0608

Remainder Theorem:-

Remainder theorem states that "if the divisor polynomial $p(n)$ is divided by $n-c$, then remainder is $f(c)$ ". This means that we can apply synthetic division and the last number on the left right, which is the remainder, will tell us what the functional value of c is. Thus, the remainder theorem is useful for evaluating polynomials at a given value of n .

Example:

Let $f(n) = 3n^3 + n^2 + n - 5$, use the remainder theorem to find $f(-2)$.

\Rightarrow Sol'n:

We have, $c = -2$

$$f(n) = 3n^3 + n^2 + n - 5$$

Now,

$$\begin{array}{r|rrrr} -2 & 3 & 1 & 1 & -5 \\ \downarrow & -6 & 10 & -22 \\ 3 & -5 & 11 & -27 \end{array}$$

$$q(n) = 3n^2 - 5n + 11$$

$$\therefore f(-2) = -27$$

→ Remainder

[If remainder is zero then dividend is 000t]

II Horner's Method for polynomial Evaluation:

Horner's method is a way to evaluate a polynomial. The method splits the polynomial into its individual terms, and solve them incrementally. This method separates the lowest degree term in the polynomial from the rest, by grouping only terms with a higher degree into one unit, with degree 1, thus making any polynomial $K \times n + c$, where K is the group of all terms with degree higher than one and c is the term with degree 0. When the terms are grouped into K , their degree decrease by one as follows,

$$n^4 + 8n^3 + 6n^2 + 4n + 3 \\ = K \times n + 3, \text{ where } K = (n^3 + 8n^2 + 6n + 4)$$

This method can be applied recursively inside the K to simplify the polynomial that results after grouping, by repeating these steps, until there are no terms with degree higher than 1 and will yield an equation where there is no need to calculate any term of n^n . After applying this method to above polynomial we end up with,

$$\begin{aligned}
 & n^4 + 8n^3 + 6n^2 + 4n + 3 \\
 &= (n^3 + 8n^2 + 6n + 4) + n + 3 \\
 &= ((n^2 + 8n + 6) + n + 4)n + 3 \\
 &= (((n+8)n+6)n+4)n+3 \\
 &= (((((1)n+8)n+6)n+4)n+3)
 \end{aligned}$$

Generalizing this we can derive recursive relation for polynomial evaluation as

$$P(n) = \theta_0 + \theta_1 n + \theta_2 n^2 + \theta_3 n^3 + \dots + \theta_n n^n$$

where, $\theta_0, \theta_1, \dots, \theta_n$ are real numbers, we wish to evaluate a polynomial at a specific value say at n_0 .

To accomplish this we need to define a given polynomial in terms of nested multiplication as below:

$$P(n) = (\theta_0 + n(\theta_1 + \theta_2 n + \theta_3 n^2 + \dots + \theta_{n-1} n^{n-1}))$$

$$P(n) = (\theta_0 + n(\theta_1 + n(\theta_2 + \theta_3 n + \dots + \theta_{n-2} n^{n-2})))$$

$$P(n) = (\theta_0 + n(\theta_1 + n(\theta_2 + n(\theta_3 + n(\theta_4 + \dots + (\theta_{n-1} + n(\theta_n)))))))$$

Now, we can define a new sequence of constant as

$$b_n = \theta_n$$

$$b_{n-1} = \theta_{n-1} + b_n n,$$

Then b_0 is the value of $P(n_0)$.

Example:

Evaluate the polynomial

$$P(n) = 3n^3 - 4n^2 + 5n - 6 \text{ at } n_0 = 2$$

\Rightarrow for n_0

Given polynomial $a_n n^n$

$$P(n) = 3n^3 - 4n^2 + 5n - 6$$

Now,

$$\alpha_3 = 3, \alpha_2 = -4, \alpha_1 = 5, \alpha_0 = -6, n_0 = 2$$

Then, $b_n = \alpha_n$

$$b_3 = \alpha_3 = 3$$

$$b_2 = \alpha_2 + b_3 n_0 \quad [b_{n-1} = \alpha_{n-1} + b_n n_0] \\ = -4 + 3 \times 2$$

$$= 2$$

$$b_1 = \alpha_1 + b_2 n_0$$

$$= 5 + 2 \times 2 = 9$$

$$b_0 = \alpha_0 + b_1 n_0$$

$$= -6 + 9 \times 2 = 12$$

Hence, the functional value of given polynomial $P(n)$ at $n_0 = 2$ is 12.

Convergence ray

classmate

Date _____
Page _____

Algorithm:-

1. Start.
2. Read degree of polynomial say n .
3. Read the value at which polynomial
to be evaluated, x .
4. Initially set $b_0 = a_n$.
5. While ($n > 0$)
$$b_{n-1} = a_{n-1} + b_n * x$$
6. Print the value of b_0 , which is the
value of polynomial at x .
7. Stop.