

**BITS Pilani Hyderabad Campus**  
**CS F214 Logic in Computer Science,**  
**I Semester 2021-2022**  
**Lecture Notes**  
**Lecture 9**

*Example:* Prove the following sequent is valid.

$$p \rightarrow q, p \rightarrow \neg q \vdash \neg p$$

*Solution:*

1.	$p \rightarrow q$	premise												
2.	$p \rightarrow \neg q$	premise												
<table> <tr> <td>3.</td><td><math>p</math></td><td>assumption(3)</td></tr> <tr> <td>4.</td><td><math>q</math></td><td><math>\rightarrow e</math> 3,1</td></tr> <tr> <td>5.</td><td><math>\neg q</math></td><td><math>\rightarrow e</math> 3,2</td></tr> <tr> <td>6.</td><td><math>\perp</math></td><td><math>\neg e</math> 4,5</td></tr> </table>			3.	$p$	assumption(3)	4.	$q$	$\rightarrow e$ 3,1	5.	$\neg q$	$\rightarrow e$ 3,2	6.	$\perp$	$\neg e$ 4,5
3.	$p$	assumption(3)												
4.	$q$	$\rightarrow e$ 3,1												
5.	$\neg q$	$\rightarrow e$ 3,2												
6.	$\perp$	$\neg e$ 4,5												
7.	$\neg p$	$\neg i$ 3-6												

*Example:* Prove  $p \rightarrow (q \rightarrow r), p, \neg r \vdash \neg q$  is valid, without using MT.

*Solution:*

1.	$p \rightarrow (q \rightarrow r)$	premise												
2.	$p$	premise												
3.	$\neg r$	premise												
<table> <tr> <td>4.</td><td><math>q</math></td><td>assumption</td></tr> <tr> <td>5.</td><td><math>q \rightarrow r</math></td><td><math>\rightarrow e</math> 2,1</td></tr> <tr> <td>6.</td><td><math>r</math></td><td><math>\rightarrow e</math> 4,5</td></tr> <tr> <td>7.</td><td><math>\perp</math></td><td><math>\neg e</math> 3,6</td></tr> </table>			4.	$q$	assumption	5.	$q \rightarrow r$	$\rightarrow e$ 2,1	6.	$r$	$\rightarrow e$ 4,5	7.	$\perp$	$\neg e$ 3,6
4.	$q$	assumption												
5.	$q \rightarrow r$	$\rightarrow e$ 2,1												
6.	$r$	$\rightarrow e$ 4,5												
7.	$\perp$	$\neg e$ 3,6												
8.	$\neg q$	$\neg_i$ 4-7												

## 7 Derived Rules

MT can be derived from  $\rightarrow e$ ,  $\neg e$  and  $\neg_i$ .

1.	$\phi \rightarrow \psi$	premise									
2.	$\neg\psi$	premise									
<table> <tr> <td>3.</td><td><math>\neg\phi</math></td><td>assumption)</td></tr> <tr> <td>4.</td><td><math>\neg\psi</math></td><td><math>\rightarrow e</math> 3,1</td></tr> <tr> <td>5.</td><td><math>\perp</math></td><td><math>\neg e</math> 2,4</td></tr> </table>			3.	$\neg\phi$	assumption)	4.	$\neg\psi$	$\rightarrow e$ 3,1	5.	$\perp$	$\neg e$ 2,4
3.	$\neg\phi$	assumption)									
4.	$\neg\psi$	$\rightarrow e$ 3,1									
5.	$\perp$	$\neg e$ 2,4									
6.	$\neg\phi$	$\neg_i$ 3 – 5									

## 8 Law of the Excluded Middle (LEM)

- It is also called as Tertium non datur. ( There is no third possibility)
- Says that  $\phi \vee \neg\phi$  is always true.

$$\vdash \phi \vee \neg\phi$$

1.	$\neg(\phi \vee \neg\phi)$	assumption									
<table> <tr> <td>2.</td><td><math>\phi</math></td><td>assumption</td></tr> <tr> <td>3.</td><td><math>\phi \vee \neg\phi</math></td><td><math>\vee_i</math> 2</td></tr> <tr> <td>4.</td><td><math>\perp</math></td><td><math>\neg e</math> 3,1</td></tr> </table>			2.	$\phi$	assumption	3.	$\phi \vee \neg\phi$	$\vee_i$ 2	4.	$\perp$	$\neg e$ 3,1
2.	$\phi$	assumption									
3.	$\phi \vee \neg\phi$	$\vee_i$ 2									
4.	$\perp$	$\neg e$ 3,1									
5.	$\neg\phi$	$\neg_i$ 2-4									
6.	$\phi \vee \neg\phi$	$\vee_{i2}$ 5									
7.	$\perp$	$\neg e$ 6									
8.	$\phi \vee \neg\phi$	PBC 1-7									