BITS Pilani Hyderabad Campus CS F214 Logic in Computer Science, I Semester 2021-2022 Lecture Notes Lecture 21

Please read sec 1.5.2 from textbook regarding the conversion of any formula in CNF (Page 57).

The logical constants ('bottom') \perp and ('top') \top denote respectively unsatisfiable formula and tautology.

15 Horn Formula

A Horn formula is a formula ϕ in propositional logic, if it can be generated as instance of the following grammar.

$$P::=\bot|\top|P$$

$$A::=P|P\wedge A$$
 Horn Clause,
$$C::=A\rightarrow P$$
 Horn Formula,
$$H::=C|C\wedge H$$

Formula	Explanation
$1.(p \land q \land r \to p) \land (q \land s \to p) \land (\top \to s) \land (r \land s \to p) \land (\neg s \to$	Horn Formula
$(\bot) \land (\bot \land p \to r)$	
$2.(p \land q \land r \to \neg p) \land (q \land r \to q)$	Not Horn Formula due to $\neg p$
3. $(p \land r \land r \to \bot) \land (\neg q \land r \to p)$	Not Horn Formula due to $\neg q$
$4. (p_1 \land p_2 \land p_3 \to (p_4 \land p_5)) \land (\top \to p_5)$	Not Horn Formula due to $p_4 \wedge p_5$
$5. (p \land q \to r) \land (p \land q) \land (r \lor s \to p)$	Not Horn Formula due to \lor and
	$p \wedge q$

15.1 Deciding Satisfiability of Horn Formula

- Maintain a list of all occurrences of type P in your formula.
 - 1. It marks \top , if it occurs in that list.
 - 2. If there is a conjunct $p_1 \wedge p_2 \wedge ... \wedge p_k \to p'$ of ϕ such that all p_j with $1 \leq j < k$ is marked, then mark p' as well and go to step 2. Otherwise if there is no such conjunct go to step 3.
 - 3. If \perp is marked, print 'Unsatis fiable' and Stop. Otherwise go to step 4.
 - 4. Print 'Satisfiable'.

Example

$$1.(p \land q \land w \to \bot) \land (t \to \bot) \land (r \to p) \land (\top \to r) \land (\top \to q) \land (u \to s)$$

Solution:

- \bullet Mark all occurrences of $\top.$
- $\bullet\,$ Mark r, q, u.
- Mark p.
- Mark S.
- $\bullet\,$ Print Satisfiable.