An Introduction to Boolean Satisfiability

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Context & Introduction

Context: Solving (Very) Hard Combinatorial Problems

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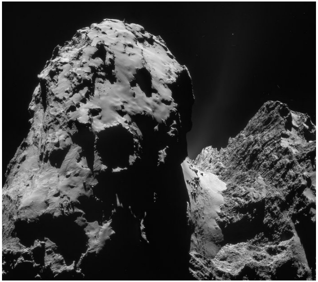

https://homepages.laas.fr/ehebrard/rosetta.html

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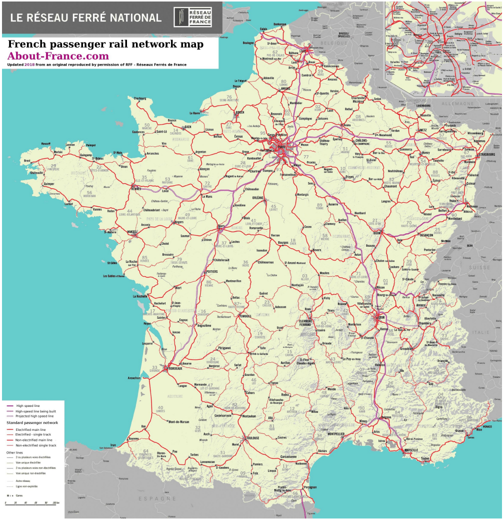

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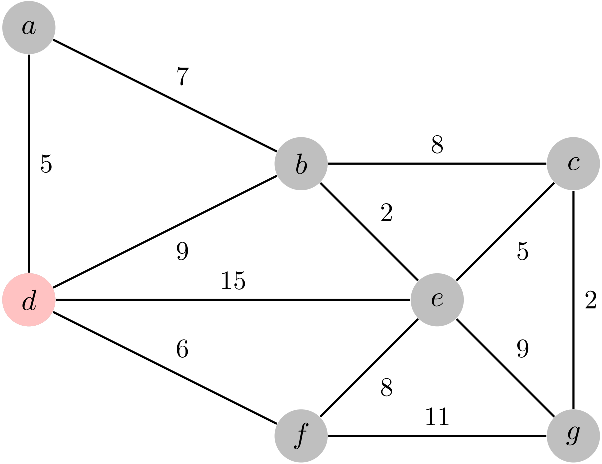
- SAT as a tool to tackle combinatorial problems
- We focus in this course on the modelling aspect
- Resources for combinatorial optimisation: Many! a good start would be the online course on discrete optimisation `https://www.coursera.org/learn/discrete-optimization`
- Handbook of Satisfiability - Second Edition - IOS Press, 2021

Exemple



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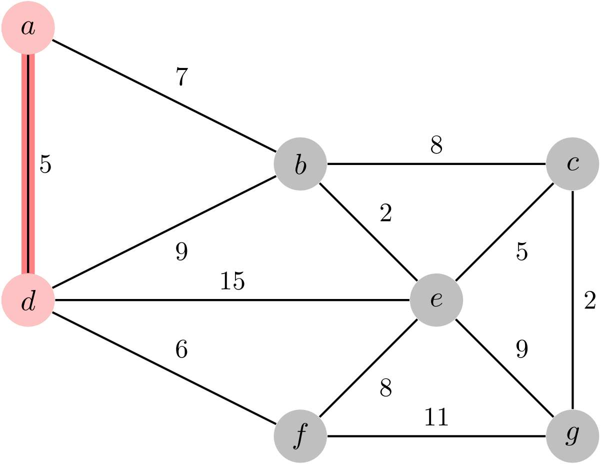

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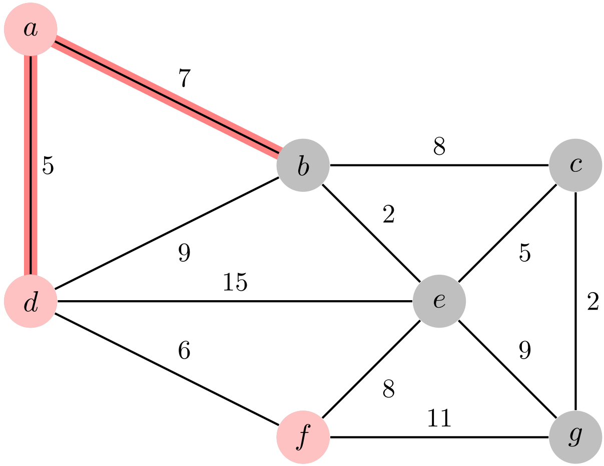

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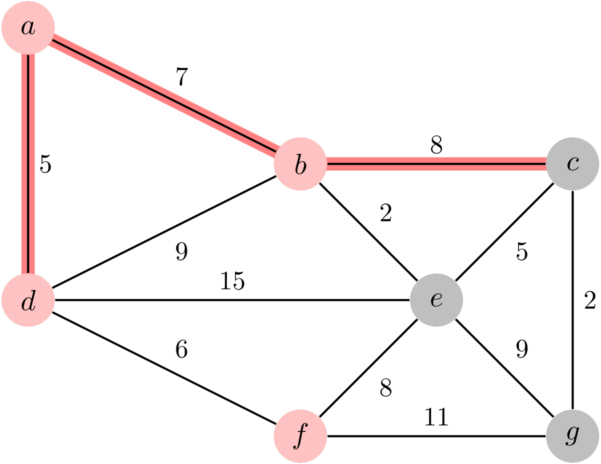

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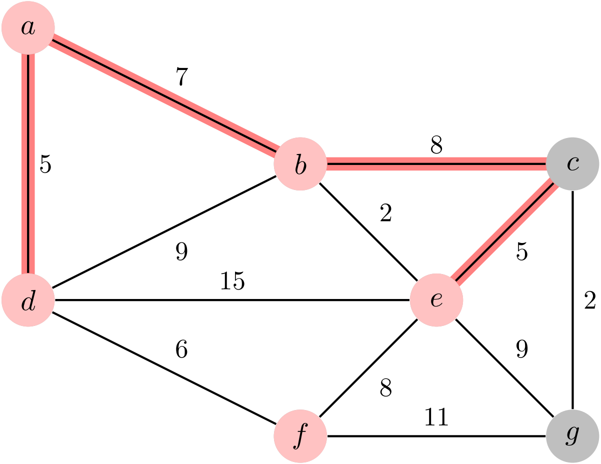

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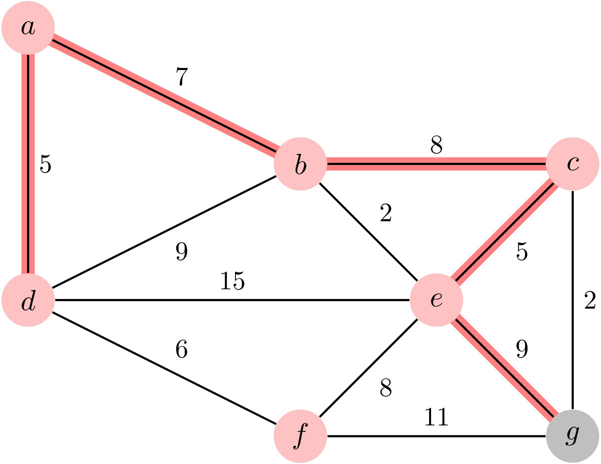

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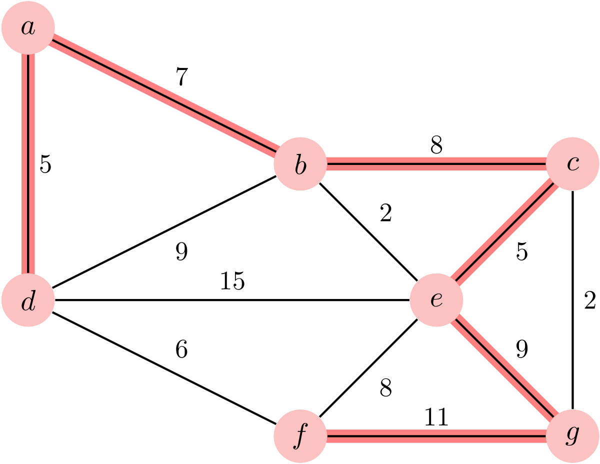

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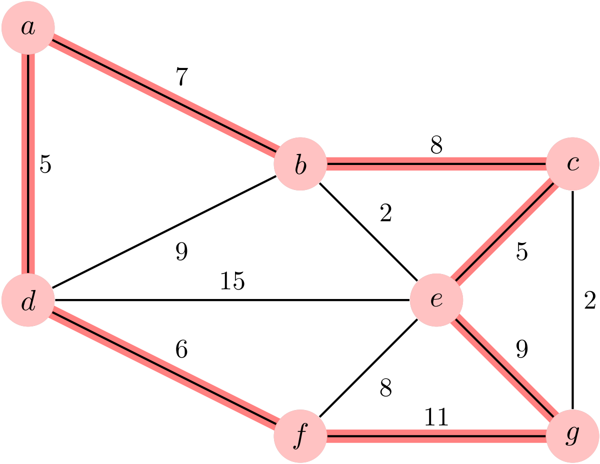

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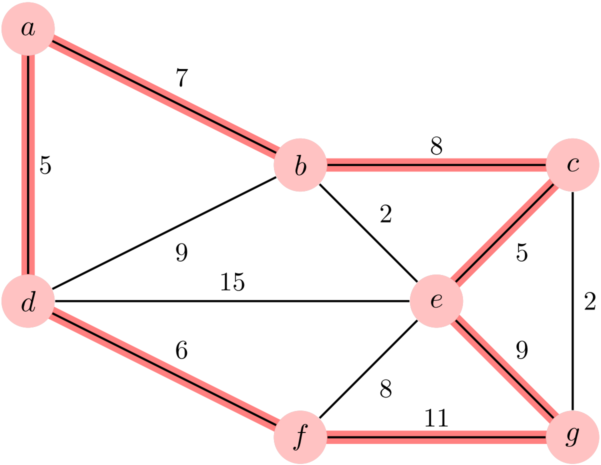

$--> Cost : 5 + 7 + 8 + 5 + 9 + 11 + 6 = 53Km$

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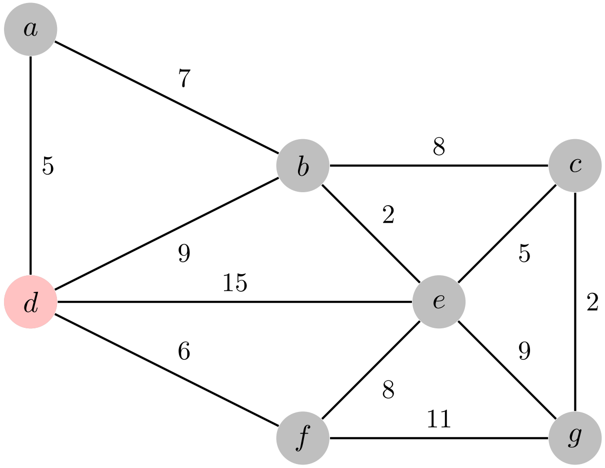

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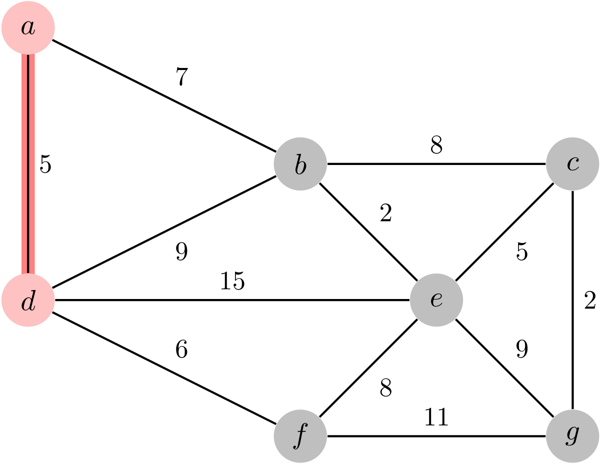

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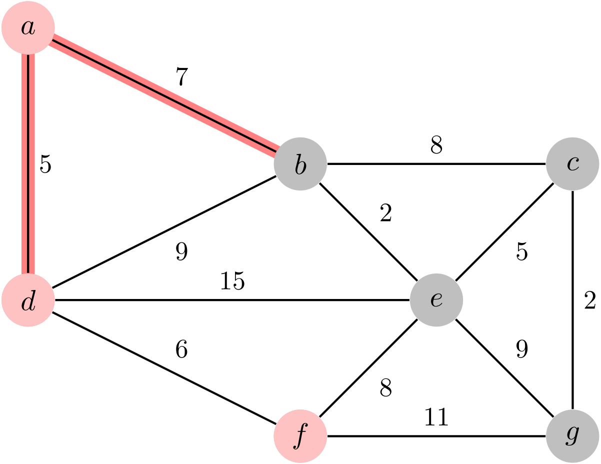

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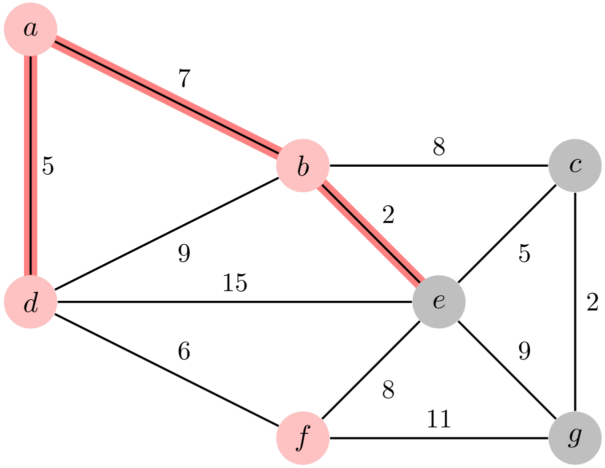

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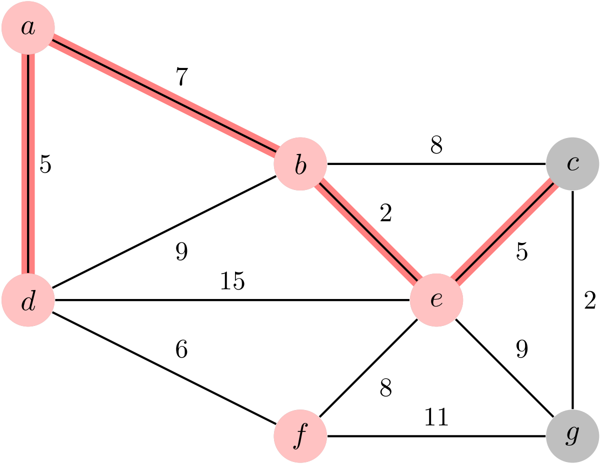

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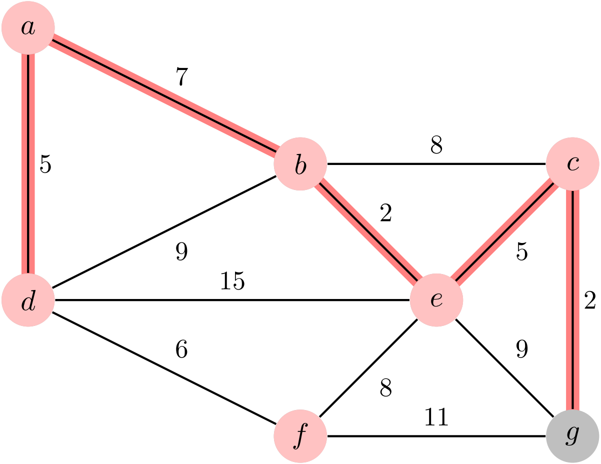

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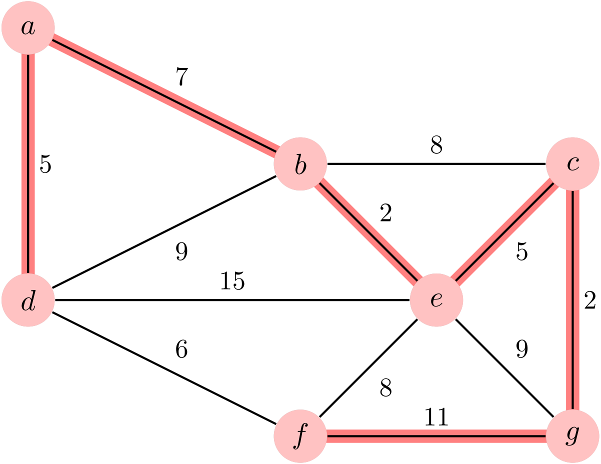

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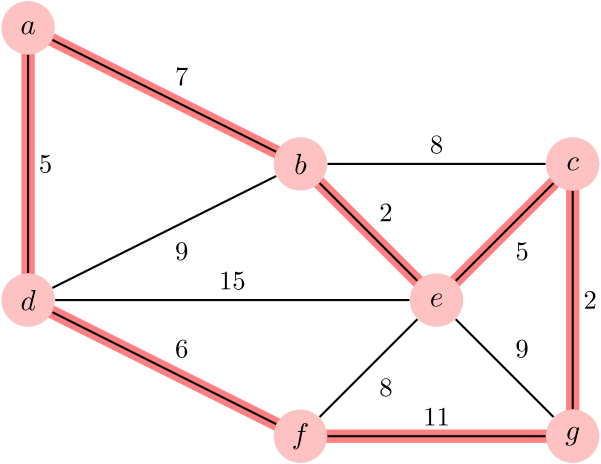

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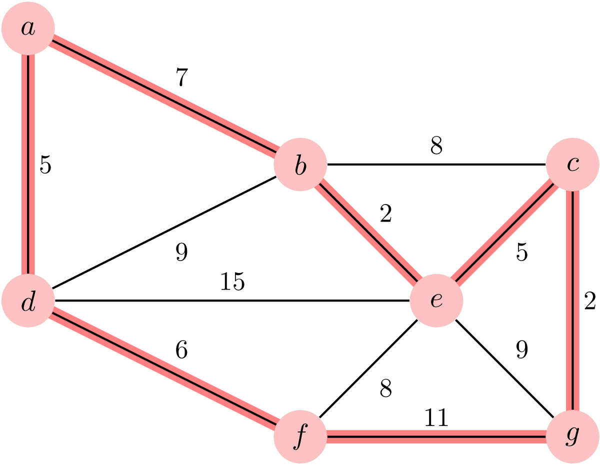

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The problem is inherently hard. However, the Concorde algorithm can solve instances up to 86 000 cities!

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Solving Methodologies

1. Adhoc methods
   1. Specific exact algorithm
   2. Heuristic method
   3. Meta-heuristic (genetic algorithms, ant colony, ..)
2. Declarative Approaches
   1. (Mixed) Integer Programming,
   2. Constraint Programming
   3. Boolean Satisfiability (SAT)
   4. ...

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- They are problem independent! The user models the problem in a specific language and the solver does the job!
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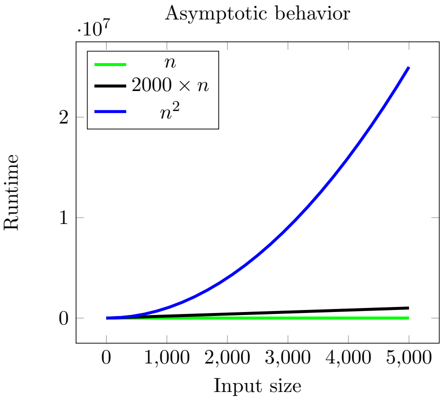

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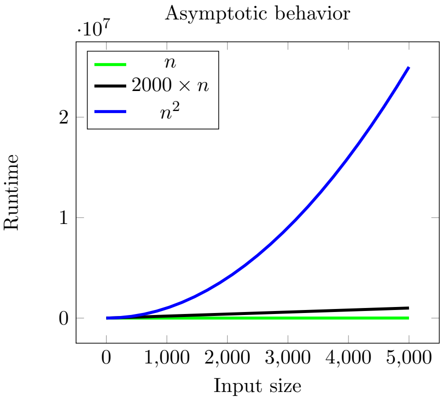

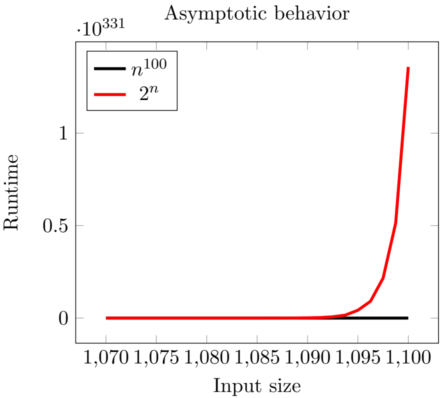

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- For many Problems in NP, we don't know if a polynomial time algorithm exists. Is P=NP?

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Given a set of Boolean variables $x_1, \ldots x_n$ and a CNF formula Φ over $x_1, \ldots x_n$, the Boolean Satisfiability problem (SAT) is to find an assignment of the variables that satisfies all the clauses.

Example

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$x \lor \neg y \lor z$

$\neg x \lor \neg z$

$y \lor w$

$\neg w \lor \neg x$

Example

$x \vee \neg y \vee z$

$\neg x \vee \neg z$

$y \vee w$

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A possible solution:

$$x \leftarrow 1; y \leftarrow 1; z \leftarrow 0; w \leftarrow 0$$

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- It is considered today as a powerful technology to solve computational problems
- Huge practical improvements in the past 2 decades or so

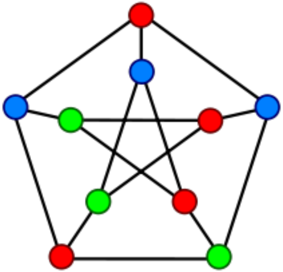
Examples of Applications

- AI Planning
- Scheduling
- Software verification
- Machine learning
  - Robustness
  - Synthesis
  - Verification
- Mathematical Proofs! `https://news.cnrs.fr/articles/the-longest-proof-in-the-history-of-mathematics`
- Timetabling
- ...

Modelling in SAT

The example of Graph Colouring

- Graph Coloring is a well known combinatorial problem that has many applications (in particular in scheduling problems)
- Let $G = (V, E)$ be an undirected graph where V is a set of n vertices and E is a set of m edges. Is it possible to colour the graph with k colours such that no two adjacent nodes share the same colour?



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- Forbid two nodes that share an edge to be coloured with the same colour

$$\forall \{i, j\} \in E, \forall a \in [1, k] : \neg x_i^a \vee \neg x_j^a$$

  (This is a translation of $x_i^a \rightarrow \neg x_j^a$)

The Example of Graph Coloring: The Model Size

What is the (space) size of the model?

- $n \times k$ Boolean variables
- Constraints form 1: n clauses with k literals each
- Constraints form 2: $n \times k^2$ binary clauses
- Constraints form 3: $m \times k$ binary clauses

The Example of Graph Coloring: The Minimization Version

- Propose a method that uses SAT for the minimisation version of the problem. That is, given $G = (V, E)$, we seek to find the minimum value of k to satisfy the colouring requirements.

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- Run iteratively the decision version until converging to the optimal value
- Let's call $SAT(V, E, K)$ the SAT model of the decision version of the problem (i.e., can we find a valid colouring of $G(V, E)$ with k colours). Use $SAT(V, E, K)$ as an oracle within an iterative search. For instance:

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A Straightforward Approach

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  - **Binary search:** Run iteratively $SAT(V, E, z)$ as long as $UB > LB$ where $z = \lceil (UB - LB)/2 \rceil$. If the result is satisfiable, then and $UB \leftarrow z$ otherwise $LB \leftarrow z$

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- An alternative approach is to look for valid theoretical bounds in the literature.

Modelling Cardinality Constraints

- A cardinality constraint takes as input a sequence of Boolean variables $[x_1, \ldots, x_n]$ and an integer k and enforces

$$\sum_{1}^{n} x_i \leq k$$

- Cardinality constraints are everywhere!
- There exist many ways in the literature to encode such constraints. See for instance `https://www.carstensinz.de/papers/CP-2005.pdf`

Quadratic encoding for $\sum_1^n x_i = 1$

- At least one constraint:

$$x_1 \vee x_2 \ldots \vee x_n$$

- at most one constraint:

$$\forall i, j : \neg x_i \vee \neg x_j$$

This generates one clause of size n and (n^2) binary clauses without introducing additional variables.

Linear encoding for $\sum_1^n x_i = 1$

A sequence of Boolean variables $[y_1, \ldots, y_n]$ is introduced such that $\forall i \in [1, n], y_i$ is true iff $\sum_{l=1}^{l=i} x_l = 1$. The set of clauses for the encoding is the following:

$$x_1 \vee x_2 \ldots \vee x_n$$

$$y_n^1$$

$$\forall i \in [1, n-1] : y_i \rightarrow y_{i+1}$$

$$\forall i \in [1, n-1] : y_i \rightarrow \neg x_{i+1}$$

$$\forall i \in [1, n] : x_i \rightarrow y_i$$

Size: n new variables, 1 $n-$ary clause and $3 \times n$ binary clauses,

Encoding for $\sum_1^n x_i \geq k$

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Encoding for $\sum_1^n x_i \geq k$

Size of the encoding:

- $\Theta(n \times k)$ variables
- $\Theta(n + k)$ unary clauses
- $\Theta(n \times k)$ binary clauses
- $\Theta(n \times k)$ ternary clauses

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- Encode $\sum_1^n x_i \geq k+1$
- Add y_n^k
- Replace y_n^{k+1} by $\neg y_n^{k+1}$
- The size of the encoding is the same as $\sum_1^n x_i \geq k$ (asymptotically)

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- Encode $\sum_1^n x_i \leq b$
- $\sum_1^n x_i \geq a$ with the same additional variables
- The size of the encoding is the same as $\sum_1^n x_i \geq k$ (asymptotically)

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- Check the MaxSAT competition

The Example of Graph Coloring: A Possible MaxSAT Model

Let $G = (V, E)$ be an undirected graph where V is the set of vertices and E is the set of edges. In the (decision version of the) graph colouring problem, we are given k colours to colour the graph such that no two adjacent nodes share the same colour.

- Propose a MaxSAT model for the minimisation version of the problem. That is, given $G = (V, E)$, we seek to find the minimum value of k to satisfy the colouring requirements.

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- We shall extend the previous model:
- Let u_a be a Boolean variable that is True iff. the colour $a \in [1, k]$ is used
- Consider the previous model $SAT(V, E, k)$ with k an upper bound.
- All the previous clauses are hard.
- Add the following hard clauses:

$$\forall i \in [1, n], \forall a \in [1, k] : \neg u_a \rightarrow \neg x_i^a$$

- Eventually we can add symmetry breaking constraints: $u_a \rightarrow u_{a-1}$
- Then add the soft clauses:

$$\forall a \in [1, k] : \neg u_a$$

- A MaxSAT Optimal solution satisfies all the hard coloring clauses (valid colouring) and maximizes the number of non used colours.

Extensions: Quantified Boolean Formula (QBF)

- A QBF has the form Q.F , where F is a CNF-SAT formulae, and Q is a sequence of quantified variables ($\forall x$ or $\exists x$).
- Example $\forall x, \exists y, \exists z, (x \vee \neg y) \wedge (\neg y \vee z)$
- QBF Solver Competition: `https://www.qbflib.org/solvers_list.php`

Extensions: Satisfiability Modulo Theories (SMT)

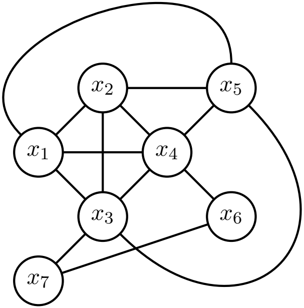
- SMT extends SAT by allowing higher level constraints
- Such constraints belong to certain theories
- Examples of theories include linear integer arithmetic, linear real arithmetic, difference logic, etc
- Check the SAT/SMT summer schools
  `http://satassociation.org/sat-smt-school.html`

Exercise: SAT for Machine Learning

- Let $F = [f_1, \dots f_k]$ be a set of k features and $E = [e_1, \ \dots e_n]$ a set of n examples.
- We want to build an undirected acyclic graph for prediction
- Task1: Propose a model for the topology of the graph
- Task 2: Extend the model to make sure that each example is well classified
- Task 3: Adapt the model to maximize the accuracy of the model

Exercise: Clique

A clique in a graph $G(V, E)$ (where V is the set of vertices and E is the set of edges) is a set of vertices $C \subseteq V$ such that $\forall a, b \in V, \{a, b\} \in E$.
For examples, in the example below: $\{x_1, x_2, x_3, x_4, x_5\}$ is a clique and $\{x_3, x_6, x_7\}$ is not a clique.



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- A possible solution:
  - x_i true iff v_i is in the clique
  - For each $\{i, j\} \notin E$:

$$\neg x_i \vee \neg x_j$$

  - Clique size:

$$\sum x_i \geq k$$

  - Implied constraints: If a vertex v_i has less than k edges it shouldn't be part of the clique:

$$\neg x_i$$

- Adapt your model into a MaxSAT formulae to find a clique with a maximum size

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  Same model without carnality constraints, without implies constraints, and each x_i is added as a soft clause

Conflict Driven Clause Learning

Modern SAT Solvers: Conflict Driven Clause Learning (CDCL)

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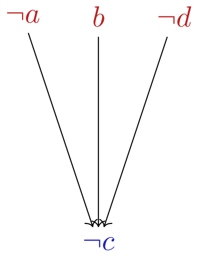
- But also :
  - Activity-based branching
  - Lazy data structures (2-Watched Literals)
  - Clause Database Reduction
  - Simplifications
  - Restarts
  - ...

Exercise: Propose a filtering algorithm to prune the variables domain in a given clause

Unit Propagation

Given a clause C of arity n. If $n-1$ literals are false then set the last one to be true.

Example: $(a \vee \neg b \vee \neg c \vee d)$



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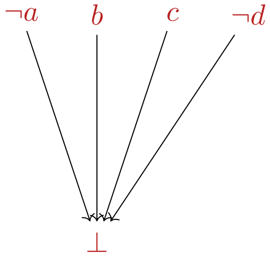

$\neg a \wedge b \wedge c \wedge \neg d \Rightarrow \perp$

Algorithm 1: Unit Propagation

Data: A clause C

if *All literals in* C *are false* **then**
    return Failure ;
else
    if *Only one literal* $l \in C$ *is unassigned and the rest are false* **then**
        Make l true ;
    end
end

Unit Propagation

- Observe first that propagation happens only in two cases:
  - The clause becomes unit (i.e., all variables except one is instantiated): Propagate the only uninstantiated literal to satisfy the clause
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- If a literal watching a clause C becomes $false$, look for replacement. If no replacement found, then perform propagation

Algorithm 2: Two watched Literals (decision d)

▷ Assuming initially that all variables are unassigned and that each clause contain at least 2 literals
▷ For each clause C, $W[C]$ is initialized with a set that contains two variables in C
▷ For each variable x, $B[x]$ is the set of clauses watched by x
▷ d is the latest decision ;

```
S ← {d} ;
while S ≠ ∅ do
    Let x ∈ S ;
    S ← S \ {x} ;
    while B[x] ≠ ∅ do
        Let C ∈ B[x] ;
        W[C] ← W[C] \ {x} ;
        if ∃x' ∈ C \ W[C] such that x' is unassigned then
            W[C] ← W[C] ∪ {x'} ;
            B[x'] ← B[x'] ∪ {C} ;
        else
            Let y ∈ W[C] ;
            if y is not assigned then
                assign y to a value that satisfies C ;
                S ← S ∪ {y} ;
                S ← ∅
            else
                if y does not satisfy C then
                    return FAILURE ;
                end
            end
        end
    end
end
```

Learning and Backjumping

Learning and Backjumping

- Definition: Explaining a failure: $l_1 \wedge \ldots \wedge l_n \rightarrow \perp$ where $\neg l_1 \vee \ldots \vee \neg l_n$ is the clause triggering the failure

Learning and Backjumping

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Learning and Backjumping

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- While there is more than one literal propagated in the last level in the current explanation, replace it with it's explanation from the triggering clause
- When there is only one literal uip propagated in the last level in the current explanation, learn the associated new clause C, backjump (to the last level of propagated literals in C), propagate $\neg uip$ via the new clause,

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- While there is more than one literal propagated in the last level in the current explanation, replace it with it's explanation from the triggering clause
- When there is only one literal *uip* propagated in the last level in the current explanation, learn the associated new clause C, backjump (to the last level of propagated literals in C), propagate $\neg uip$ via the new clause, and continue the exploration

Conflict Analysis

Algorithm 1: 1-UIP-with-Propagators

$\Psi \leftarrow explain(\bot)$;

while $|\{q \in \Psi \mid level(q) = current\ level\}| > 1$ **do**

  $p \leftarrow \arg\max_q(\{rank(q) \mid level(q) = current\ level\ \wedge\ q \in \Psi\})$;

  $\Psi \leftarrow \Psi \cup \{q \mid q \in explain(p) \wedge level(q) > 0\} \setminus \{p\}$;

return Ψ ;

Conflict Analysis

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while $|\{q \in \Psi \mid level(q) = current\ level\}| > 1$ **do**
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return Ψ ;

- Why stop with one literal l propagated at the last level ?
- **To make sure that when the algorithm backjumps, propagation takes place by making l true**
- When backjumping using a clause that contains more than one literal propagated at the last level, then no propagation can be performed.

Implication Graph

	f		

$\neg a \vee \neg f \vee g$

$\neg a \vee \neg b \vee \neg h$

$a \vee c$

$a \vee \neg i \vee \neg l$

$a \vee \neg k \vee \neg j$

$b \vee d$

$b \vee g \vee \neg n$

$b \vee \neg f \vee n \vee k$

$\neg c \vee k$

$\neg c \vee \neg k \vee \neg i \vee l$

$c \vee h \vee n \vee \neg m$

$c \vee l$

$d \vee \neg k \vee l$

$d \vee \neg g \vee l$

$\neg g \vee n \vee o$

$h \vee \neg o \vee \neg j \vee n$

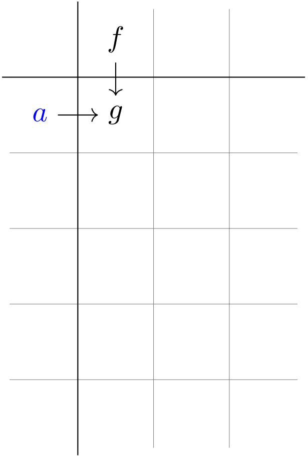
$\neg i \vee j$

$\neg d \vee \neg l \vee \neg m$

$\neg e \vee m \vee \neg n$

$\neg f \vee h \vee i$

Implication Graph



$\neg a \lor \neg f \lor g$

$\neg a \lor \neg b \lor \neg h$

$a \lor c$

$a \lor \neg i \lor \neg l$

$a \lor \neg k \lor \neg j$

$b \lor d$

$b \lor g \lor \neg n$

$b \lor \neg f \lor n \lor k$

$\neg c \lor k$

$\neg c \lor \neg k \lor \neg i \lor l$

$c \lor h \lor n \lor \neg m$

$c \lor l$

$d \lor \neg k \lor l$

$d \lor \neg g \lor l$

$\neg g \lor n \lor o$

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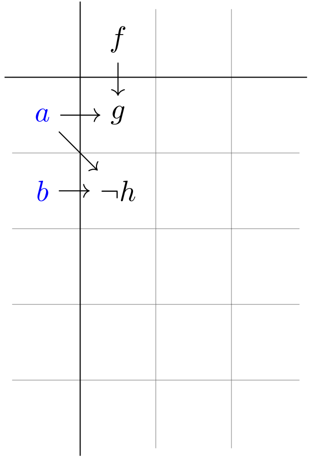
$\neg i \lor j$

$\neg d \lor \neg l \lor \neg m$

$\neg e \lor m \lor \neg n$

$\neg f \lor h \lor i$

Implication Graph



$\neg a \vee \neg f \vee g$

$\neg a \vee \neg b \vee \neg h$

$a \vee c$

$a \vee \neg i \vee \neg l$

$a \vee \neg k \vee \neg j$

$b \vee d$

$b \vee g \vee \neg n$

$b \vee \neg f \vee n \vee k$

$\neg c \vee k$

$\neg c \vee \neg k \vee \neg i \vee l$

$c \vee h \vee n \vee \neg m$

$c \vee l$

$d \vee \neg k \vee l$

$d \vee \neg g \vee l$

$\neg g \vee n \vee o$

$h \vee \neg o \vee \neg j \vee n$

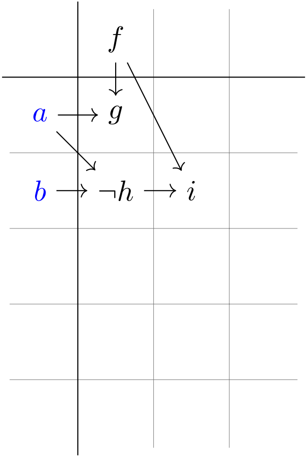
$\neg i \vee j$

$\neg d \vee \neg l \vee \neg m$

$\neg e \vee m \vee \neg n$

$\neg f \vee h \vee i$

Implication Graph



$\neg a \lor \neg f \lor g$

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$a \lor c$

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$\neg d \lor \neg l \lor \neg m$

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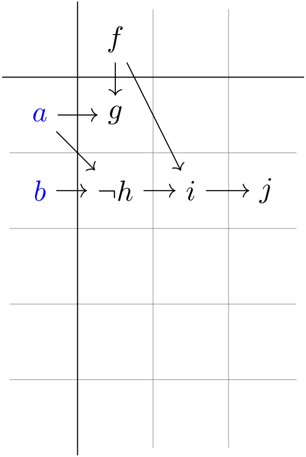
$\neg f \lor h \lor i$

Implication Graph



$\neg a \vee \neg f \vee g$

$\neg a \vee \neg b \vee \neg h$

$a \vee c$

$a \vee \neg i \vee \neg l$

$a \vee \neg k \vee \neg j$

$b \vee d$

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$\neg g \vee n \vee o$

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$\neg i \vee j$

$\neg d \vee \neg l \vee \neg m$

$\neg e \vee m \vee \neg n$

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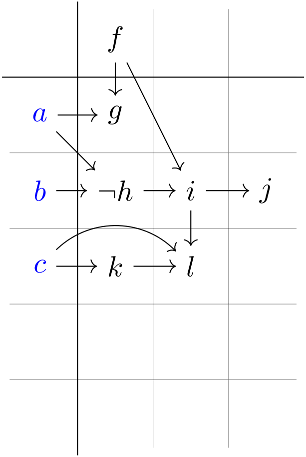
Implication Graph



$\neg a \vee \neg f \vee g$
$\neg a \vee \neg b \vee \neg h$
$a \vee c$
$a \vee \neg i \vee \neg l$
$a \vee \neg k \vee \neg j$
$b \vee d$
$b \vee g \vee \neg n$
$b \vee \neg f \vee n \vee k$
$\neg c \vee k$
$\neg c \vee \neg k \vee \neg i \vee l$

$c \vee h \vee n \vee \neg m$
$c \vee l$
$d \vee \neg k \vee l$
$d \vee \neg g \vee l$
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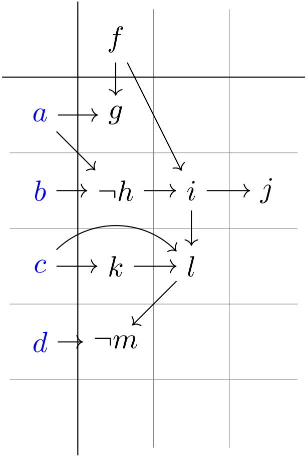
$\neg i \vee j$
$\neg d \vee \neg l \vee \neg m$
$\neg e \vee m \vee \neg n$
$\neg f \vee h \vee i$

Implication Graph



$\neg a \vee \neg f \vee g$
$\neg a \vee \neg b \vee \neg h$
$a \vee c$
$a \vee \neg i \vee \neg l$
$a \vee \neg k \vee \neg j$
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$\neg c \vee k$
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$c \vee h \vee n \vee \neg m$
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$d \vee \neg g \vee l$
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$\neg i \vee j$
$\neg d \vee \neg l \vee \neg m$
$\neg e \vee m \vee \neg n$
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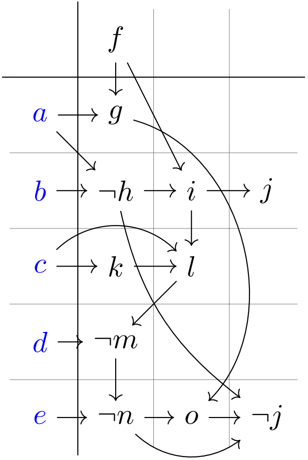
Implication Graph



$\neg a \lor \neg f \lor g$

$\neg a \lor \neg b \lor \neg h$

$a \lor c$

$a \lor \neg i \lor \neg l$

$a \lor \neg k \lor \neg j$

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$b \lor \neg f \lor n \lor k$

$\neg c \lor k$

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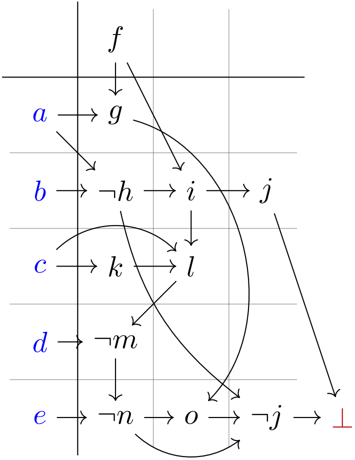
Implication Graph



$\neg a \vee \neg f \vee g$

$\neg a \vee \neg b \vee \neg h$

$a \vee c$

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$a \vee \neg k \vee \neg j$

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$d \vee \neg g \vee l$

$\neg g \vee n \vee o$

$\boxed{h \vee \neg o \vee \neg j \vee n}$

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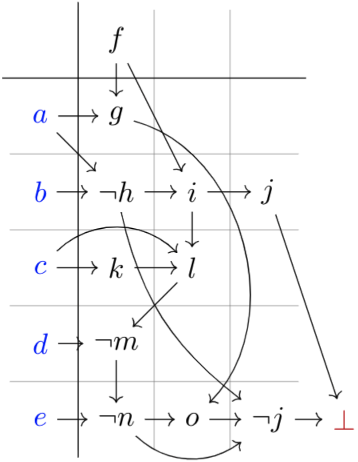
$\neg d \vee \neg l \vee \neg m$

$\neg e \vee m \vee \neg n$

$\neg f \vee h \vee i$

Implication Graph



$\neg a \lor \neg f \lor g$

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$a \lor c$

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$\neg i \lor j$

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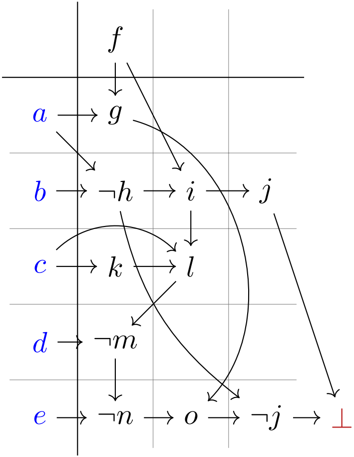
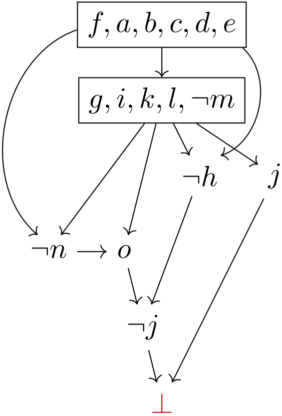
Conflict Analysis





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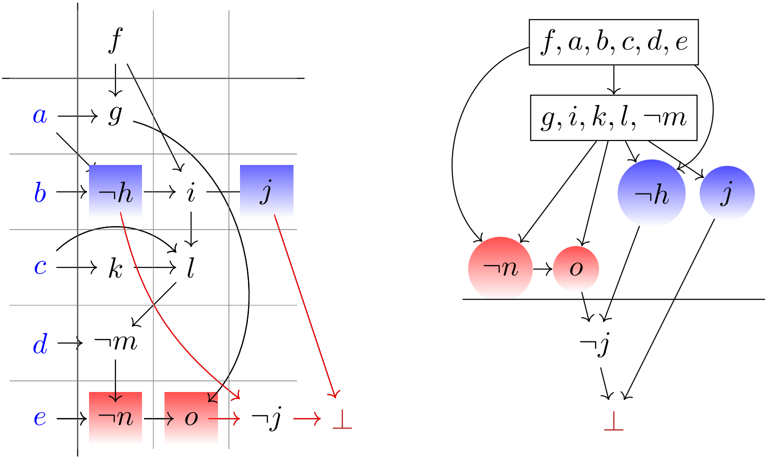
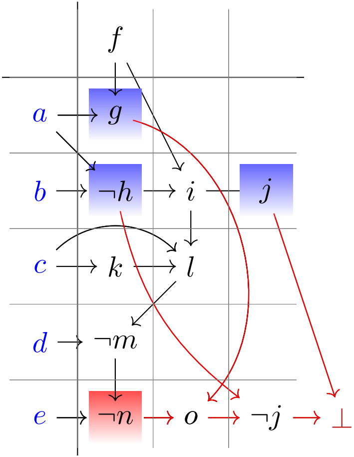
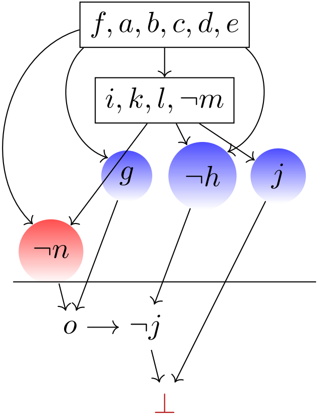
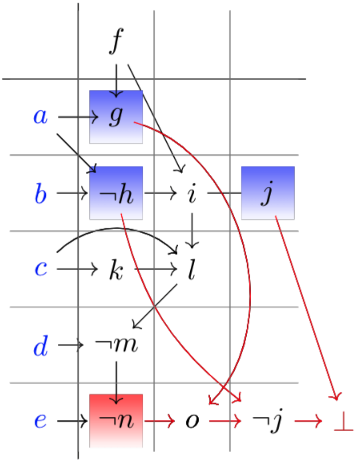


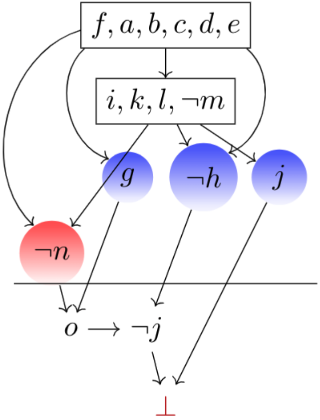
Conflict Analysis



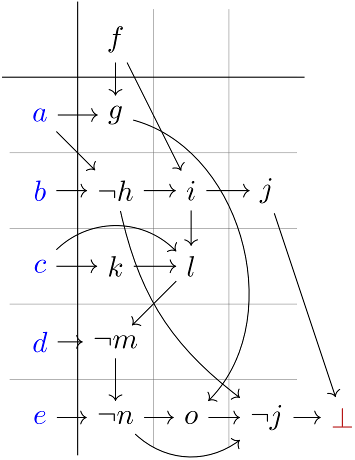


Conflict Analysis





Conflict analysis



$\neg a \vee \neg f \vee g$

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$a \vee c$

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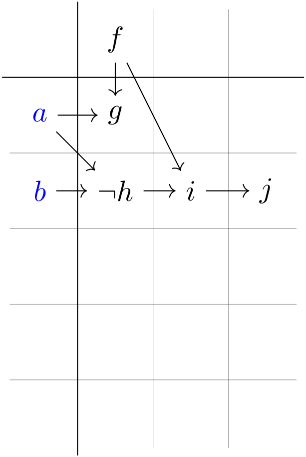
$\neg e \vee m \vee \neg n$

$\neg f \vee h \vee i$

Conflict analysis



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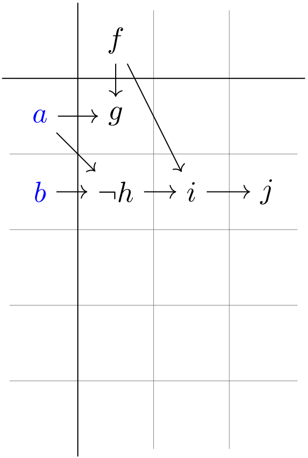
Conflict analysis



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$c \lor h \lor n \lor \neg m$

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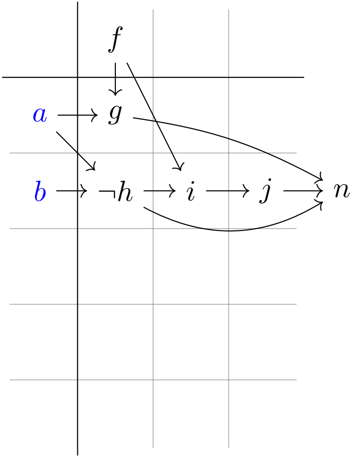
$\neg e \lor m \lor \neg n$

$\neg f \lor h \lor i$

$\boxed{\neg g \lor h \lor \neg j \lor n}$

Conflict analysis



$\neg a \vee \neg f \vee g$
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$\boxed{\neg g \vee h \vee \neg j \vee n}$

Boosting Search through Randomization and Restarts [Gomes et al., 1998]

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- Randomization: breaking ties, random decision between k best choices, ...
- Restarts: Geometric/Luby

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SAT Solvers (Few examples)

- MiniSat: http://minisat.se/
- Glucose: `http://www.labri.fr/perso/lsimon/glucose/`
- LingeLing `http://fmv.jku.at/lingeling`
- Any Solver by Armin Biere `http://fmv.jku.at/software/index.html`
- Any winner from past and future SAT competitions: https://www.satcompetition.org/

SAT vs CSP

Back to Constraint Programming

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- Mostly solvable by backtracking algorithms (Search and Filtering)

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Value Ordering

'Succeed-first' [Geelen, 1992]:
"Follow the best chances leading to a solution"

Filtering

- Filtering (propagation/pruning): inferences based on the current state

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Arc Consistency

Let C be a constraint and D be a list of domains for the variables in the scope of C.
C is Arc Consistent (AC) iff for every variable x in the scope of C, for every value $v \in D(x)$, there exists an assignment w in D satisfying C in which v is assigned to x

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- If each domain is a singleton, the propagator must be able to check if the assignment corresponds to a solution or not.

CP vs. SAT

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- CP vs. SAT: a fundamental difference is the presence of global reasoning win CP.

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- **Can we find something that takes advantage from both worlds?** $\rightarrow$ **Clause learning in CP**

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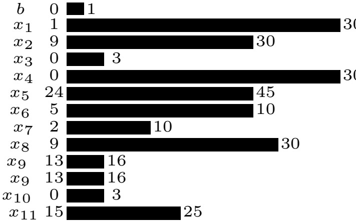

Learning in CP

$[\![x_1 = 1]\!]$

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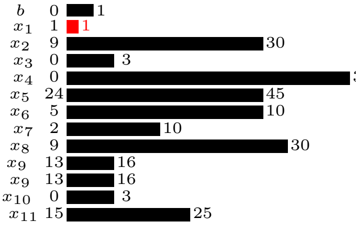

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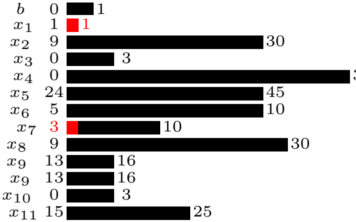

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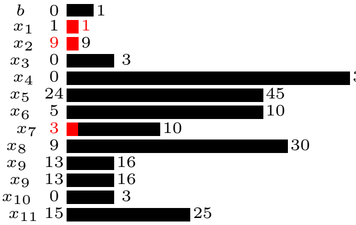

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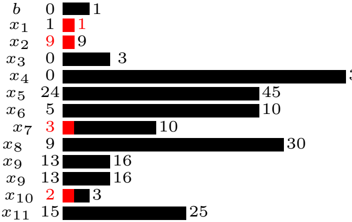

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b	0	1
x_1	1	1
x_2	9	9
x_3	2	2
x_4	0	3
x_5	24	45
x_6	5	10
x_7	3	10
x_8	9	30
x_9	13	16
x_9	13	16
x_{10}	2	3
x_{11}	15	25

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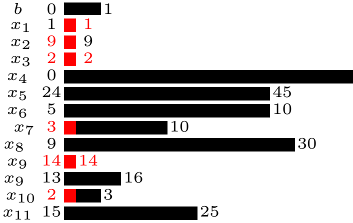

Learning in CP

$\llbracket x_1 = 1 \rrbracket \rightarrow \llbracket x_7 \geq 3 \rrbracket$

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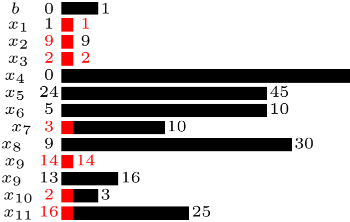

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$[\![x_1 = 1]\!] \rightarrow [\![x_7 \geq 3]\!]$

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$[\![x_4 = 0]\!]$

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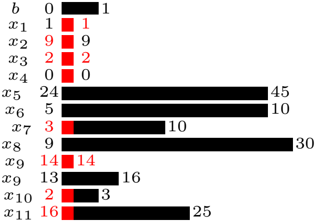

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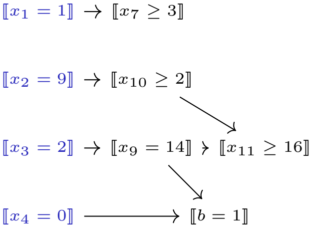

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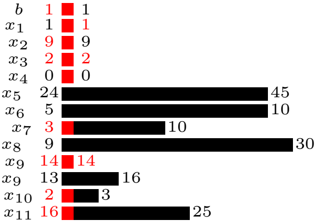

Learning in CP

$$
\begin{tikzcd}
\textcolor{blue}{[\![x_1 = 1]\!]} \arrow[r] & [\![x_7 \geq 3]\!] & & \\
\textcolor{blue}{[\![x_2 = 9]\!]} \arrow[r] & [\![x_{10} \geq 2]\!] \arrow[dr] & & \\
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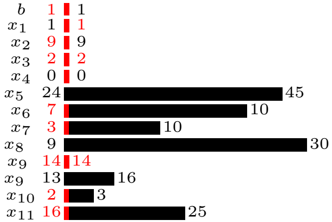

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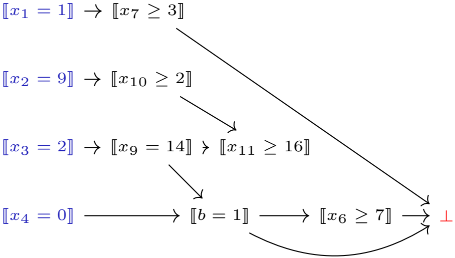

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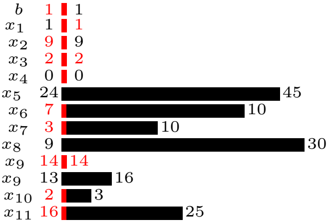

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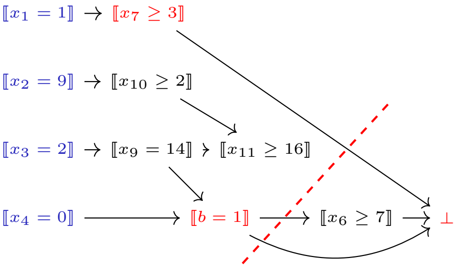

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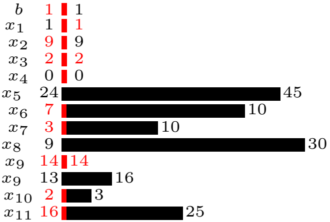
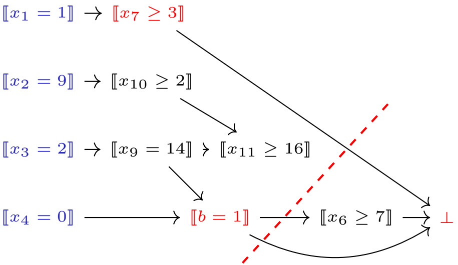


- Conflict analysis: $[\![b = 1]\!] \wedge [\![x_7 \geq 3]\!] \Rightarrow \perp$

Learning in CP



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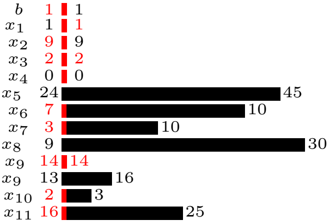
$b \rightarrow (x_6 \geq 7) \wedge$
$b \rightarrow (x_6 + x_7 \leq 9) \wedge$
$x_{11} \geq x_9 + x_{10}$



- Conflict analysis: $[\![b = 1]\!] \wedge [\![x_7 \geq 3]\!] \Rightarrow \perp$
- New clause: $[\![b \neq 1]\!] \vee [\![x_7 \leq 2]\!]$

Learning in CP

$x_1 + x_7 \geq 4 \wedge$
$x_2 + x_{10} \geq 11 \wedge$
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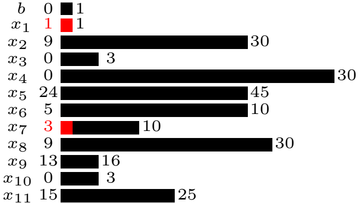

Learning in CP

$$[\![x_1 = 1]\!] \rightarrow \ [\![x_7 \geq 3]\!] \longrightarrow [\![b = 0]\!]$$

- Conflict analysis: $[\![b = 1]\!] \wedge [\![x_7 \geq 3]\!] \Rightarrow \bot$
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- Backtrack to level 1
- Propagate the learnt clause

$$\begin{aligned}
&x_1 + x_7 \geq 4 \wedge \\
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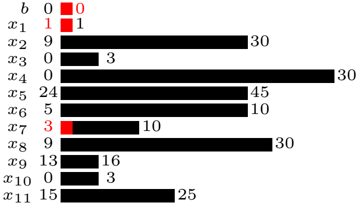

Learning in CP

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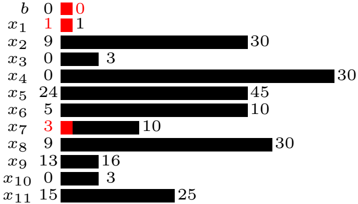
- Propagate the learnt clause
- Continue exploration



Conflict analysis

Algorithm 1: 1-UIP-with-Propagators

$\Psi \leftarrow explain(\perp)$;
while $|\{q \in \Psi \mid level(q) = current\ level\}| > 1$ **do**
    $p \leftarrow \arg\max_q(\{rank(q) \mid level(q) = current\ level\ \wedge\ q \in \Psi\})$;
    $\Psi \leftarrow \Psi \cup \{q \mid q \in explain(p) \wedge level(q) > 0\} \setminus \{p\}$;
return Ψ ;

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- To enable clause learning in CP, each propagator must be able to explain its filtering in the form of clauses ("Lazy Clause Generation").
- We distinguish two types of explanations:
  - Explaining Failure
  - Explaining Domain filtering
- Example: Explain the constraint $X \leq Y$ with two scenarios (failure and propagation).

- Let $(x_1, \ldots, x_n)$ be a sequence of Boolean variables, and let d be a positive integer.
- The CARDINALITY$(x_1, \ldots, x_n, d)$ constraint holds iff exactly d variables from the sequence $(x_1, \ldots, x_n)$ are true.
- Write a filtering algorithm for CARDINALITY.
- What is the time complexity?
- Does it enforce arc consistency?
- Explain the CARDINALITY filtering.

Correction

Algorithm 4: $\textsc{Cardinality}([x_1, \ldots, x_n], d)$

if $\left|\{x_j \mid \mathcal{D}(x_j) = \{1\}\}\right| > d$ **then**
    $\mathcal{D} \leftarrow \perp$;
if $\left|\{x_j \mid \mathcal{D}(x_j) = \{0\}\}\right| > n - d$ **then**
    $\mathcal{D} \leftarrow \perp$;
if $\left|\{x_j \mid \mathcal{D}(x_j) = \{1\}\}\right| = d$ **then**
    foreach $i \in \{1..n\}$ **do**
        if $\mathcal{D}(x_i) = \{0, 1\}$ **then**
            $\mathcal{D}(x_i) \leftarrow \{0\}$;
else
    if $\left|\{x_j \mid \mathcal{D}(x_j) = \{0\}\}\right| = n - d$ **then**
        foreach $i \in \{1..n\}$ **do**
            if $\mathcal{D}(x_i) = \{0, 1\}$ **then**
                $\mathcal{D}(x_i) \leftarrow \{1\}$;
return $\mathcal{D}$;

Explaining The Cardinality Constraint

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- Failure 1:

$$x^1 \wedge x^2 \wedge \ldots \wedge x^{d+1} \rightarrow \perp$$

Where $D(x^i) = \{1\}$

Explaining The Cardinality Constraint

- Failure 1:
$$x^1 \land x^2 \land \ldots \land x^{d+1} \rightarrow \bot$$
Where $D(x^i) = \{1\}$
- Failure 2:
$$\neg x^1 \land \neg x^2 \land \neg x^{n-d+1} \rightarrow \bot$$
Where $D(x^i) = \{0\}$
- Explaining the propagating the value 1: the conjunction of all the assigned variables

Explaining The Cardinality Constraint

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  Where $D(x^i) = \{1\}$

- Failure 2:

$$\neg x^1 \wedge \neg x^2 \wedge \neg x^{n-d+1} \rightarrow \perp$$

  Where $D(x^i) = \{0\}$

- Explaining the propagating the value 1: the conjunction of all the assigned variables
- Explaining the propagating the value 0: the conjunction of all the assigned variables

Encoding CSP into SAT

- How to encode the variables' domain ?
- How to encode each constraint into a set of clauses ?

Domain Encoding: Quadratic Encoding

- Suppose that $D(x) = \{v_1, \ldots, v_n\}$

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- As a clause: $\neg x_i \vee \neg x_j$
- The number of variables is linear
- The number of clauses is quadratic

Domain Encoding: Linear Encoding

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Domain Encoding: Linear Encoding

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- $y_j \rightarrow y_{j+1}$
- $x_i \rightarrow y_{v_i} \wedge \neg y_{v_i - 1}$
- The number of variables is linear in the size of the domain
- The number of clauses is linear. However, some clauses are of arity three

Exercise: Constraint encoding ?

- How to encode the AllDifferent constraint ?
- How to encode $\sum_i X_i \leq k$ (X_i is an integer variable)?
- How to encode $\sum_i a_i \times X_i \leq k$?

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