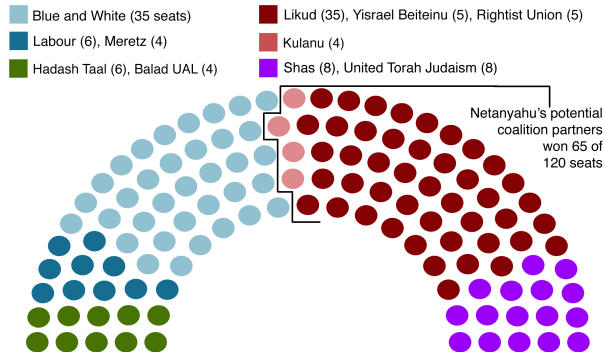


# Topic 5: Coalitional Games

## Israeli election 2019: Preliminary results



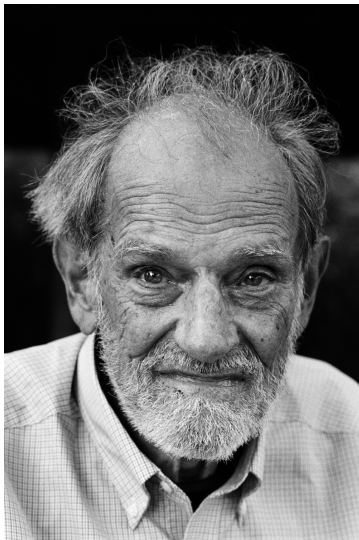
Source: Times of Israel, Jerusalem Post and Haaretz (95% of votes counted)



# Outcomes & Objectives

- ▶ Be proficient in solving coalitional games
  - ▶ Model player's rationality in forming coalitions via defining a value of a given coalition.
  - ▶ Identify some useful subclasses of games which produces some special coalitions.
  - ▶ Develop a solution concept called *Shapley value* to distribute a coalition's value in a fair manner.
  - ▶ Develop a solution concept called *core* that identifies a stable coalition structure in the game.

# Lloyd Shapley



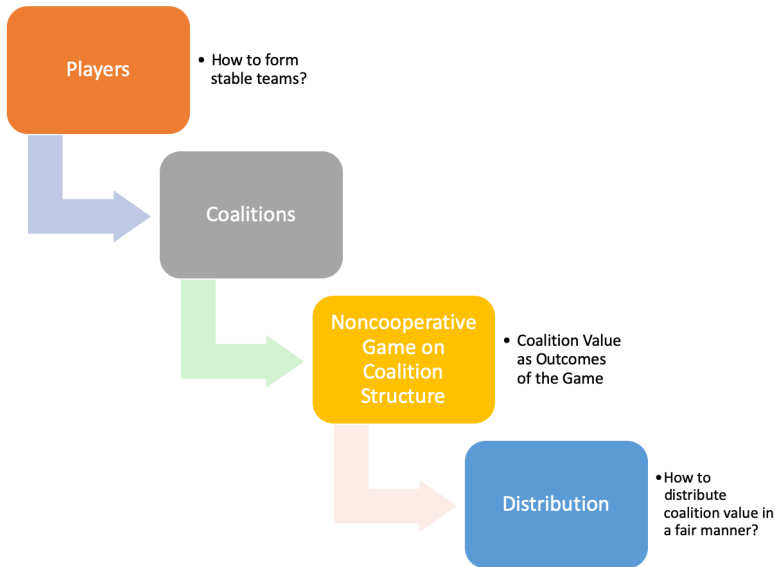
*Shapley was the greatest game theorist of all time.*

– Robert Aumann

# Applications of Coalitional Games

- ▶ **Political Coalitions:** Parties form coalitions if the elections did not result in one party with a majority votes. Coalitional governments resolve such concerns. However, the question is which coalitions form stable governments.
- ▶ **Cost Sharing for Network Design:** Users benefit from being connected to a server. So they have to build up a broadcast tree. However, it costs to maintain the server/network and the question is how to share the costs.
- ▶ **Queue Management:** Multiple users want to route traffic through a switch, which has a flow dependent delay (cost). The queueing delay cost has to be shared among the users.

# Coalitional Game: An Overview



# Coalitions and Transferable Utilities

## Definition

Given a set of players  $\mathcal{N} = \{1, \dots, N\}$ , a **coalition** is a subset of  $\mathcal{N}$ . Furthermore, a **grand coalition** is the set of all players  $\mathcal{N}$ .

## Definition

A **characteristic function game**  $\Gamma$  is a pair  $(\mathcal{N}, v)$ , where  $\mathcal{N}$  is the set of players, and  $v : 2^{\mathcal{N}} \rightarrow \mathbb{R}$  is a **characteristic function**, which assigns each coalition  $\mathcal{C} \subseteq \mathcal{N}$ , some real value  $v(\mathcal{C})$ .

## Definition

A characteristic function game  $\Gamma = (\mathcal{N}, v)$  is a **transferable utility game**, if the value of any coalition  $v(\mathcal{C})$  can be distributed amongst the members in  $\mathcal{C}$  in any way that the members of  $\mathcal{C}$  choose.

## Standard Assumptions:

- ▶ The value of a empty coalition is 0.
- ▶  $v(\mathcal{C}) \geq 0$ , for any  $\mathcal{C} \subseteq \mathcal{N}$ .

# Example

A fictional country  $X$  has a 101-member parliament, where each representative belongs to one of the three parties:

- ▶ Liberal ( $L$ ): 40 representatives
- ▶ Moderate ( $M$ ): 31 representatives
- ▶ Conservative ( $C$ ): 30 representatives

The parliament needs to decide how to allocate \$1bn of discretionary spending, and each party has its own preferred way of using this money. The decision is made by a simple **majority vote**, and we assume that all representatives vote along the party lines.

Parties can form **coalitions**; a coalition has value \$1bn if it can win the budget vote no matter what the other parties do, and value 0 otherwise.

This situation can be modeled as a three-player characteristic function game, where the set of players is  $\mathcal{N} = L, M, C$  and the characteristic function is given by

$$v(\mathcal{C}) = \begin{cases} 0, & \text{if } |\mathcal{C}| \leq 1, \\ 10^9, & \text{otherwise.} \end{cases}$$

# Coalition Structure

## Definition

Given a characteristic function game  $\Gamma = (\mathcal{N}, v)$ , a **coalition structure**  $\mathcal{C}$  is a partition of  $\mathcal{N}$ . In other words,  $\mathcal{C}$  is a collection of non-empty subsets  $\{\mathcal{C}_1, \dots, \mathcal{C}_K\}$  such that

- ▶  $\bigcup_{k \in \{1, \dots, K\}} \mathcal{C}_k = \mathcal{N}$ , and
- ▶  $\mathcal{C}_i \cap \mathcal{C}_j = \emptyset$ , for any  $i, j \in \{1, \dots, K\}$  such that  $i \neq j$ .

## Definition

A vector  $\mathbf{u} = \{u_1, \dots, u_N\} \in \mathbb{R}^N$  is the **utility profile** for a coalition structure  $\mathcal{C} = \{\mathcal{C}_1, \dots, \mathcal{C}_K\}$  over  $\mathcal{N}$  if

- ▶ **Non-Negativity:**  $u_i \geq 0$  for all  $i \in \mathcal{N}$ , and
- ▶ **Feasibility:**  $\sum_{i \in \mathcal{C}_k} u_i \leq v(\mathcal{C}_k)$  for any  $k \in \{1, \dots, K\}$ .



# Outcome, Efficiency and Social Welfare

## Definition

The **outcome** of a game  $\Gamma$  is a pair  $(\mathcal{C}, u)$ .

## Definition

An outcome  $(\mathcal{C}, u)$  is **efficient**, if all the utilities are distributed amongst the coalition members, i.e.

$$\sum_{i \in \mathcal{C}_k} u_i = v(\mathcal{C}_k), \text{ for all } k = 1, \dots, K.$$

## Definition

The social welfare of a coalition structure  $\mathcal{C}$  is

$$v(\mathcal{C}) = \sum_{k=1}^K v(\mathcal{C}_k)$$

# Individual Rationality and Imputation

## Definition

A player  $i$  is said to be **individually rational** in an outcome  $(\mathcal{C}, \mathbf{u})$ , if

$$u_i \geq v(\{i\}),$$

where  $v(\{i\})$  is the value of the coalition  $\{i\}$ , which only contains the  $i^{th}$  player.

## Definition

An outcome  $(\mathcal{C}, \mathbf{u})$  is said to be an **imputation**, if it is efficient, and if every player is individually rational within itself.

- ▶ Each player weakly prefers being in the coalition structure, than being on his/her own.
- ▶ Group deviations  $\Rightarrow$  Stability of Coalitions (covered later)

# Monotone Games

## Definition

A characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  is said to be **monotone** if it satisfies  $v(\mathcal{C}) \leq v(\mathcal{D})$ , for every pair of coalitions  $\mathcal{C}, \mathcal{D} \subseteq \mathcal{N}$ , such that  $\mathcal{C} \subseteq \mathcal{D}$ .

- ▶ Most games are monotone!
- ▶ However, non-monotonicity may arise because
  - ▶ some players intensely dislike each other, or
  - ▶ communication costs increase nonlinearly with coalition size.

**Example:** Three commuters can share a taxi. Individual journey costs:  $P_1 : 6$ ,  $P_2 : 12$ ,  $P_3 : 42$ . Then, the following characteristic function results in a monotone game:

$$v_1(\mathcal{C}) = \begin{cases} 6 & \text{if } \mathcal{C} = \{1\} \\ 12 & \text{if } \mathcal{C} = \{2\} \\ 42 & \text{if } \mathcal{C} = \{3\} \\ 12 & \text{if } \mathcal{C} = \{1, 2\} \\ 42 & \text{if } \mathcal{C} = \{1, 3\} \\ 42 & \text{if } \mathcal{C} = \{2, 3\} \\ 42 & \text{if } \mathcal{C} = \{1, 2, 3\}. \end{cases}$$

# Superadditive Games

## Definition

A characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  is said to be **superadditive** if it satisfies  $v(\mathcal{C} \cup \mathcal{D}) \geq v(\mathcal{C}) + v(\mathcal{D})$ , for every pair of disjoint coalitions  $\mathcal{C}, \mathcal{D} \subseteq \mathcal{N}$ .

## Proposition

If a superadditive game  $\Gamma = \{\mathcal{N}, v\}$  has a non-negative characteristic function  $v$ , then  $\Gamma$  is monotone.

*Proof:* For any pair of coalitions  $\mathcal{C} \subseteq \mathcal{D}$ , we have

$$v(\mathcal{C}) \leq v(\mathcal{D}) - v(\mathcal{D} - \mathcal{C}) \leq v(\mathcal{D}).$$

□

- ▶ Monotonicity  $\not\Rightarrow$  superadditivity. (Example:  $v(\mathcal{C}) = \log |\mathcal{C}|$ .)
- ▶ Always profitable for two groups to join forces  $\Rightarrow$  Grand Coalition.
- ▶ *Anti-trust* or *anti-monopoly* laws  $\Rightarrow$  Non-superadditive games.

# Superadditive Games: Example

Consider the same taxi example:

- ▶ Three commuters can share a taxi. Individual journey costs:  $P_1 : 6$ ,  $P_2 : 12$ ,  $P_3 : 42$ .
- ▶ Then,  $v_1(\mathcal{C})$  is not superadditive.
- ▶ However, the following characteristic function results in a superadditive game:

$$v_2(\mathcal{C}) = \begin{cases} 6 & \text{if } \mathcal{C} = \{1\} \\ 12 & \text{if } \mathcal{C} = \{2\} \\ 42 & \text{if } \mathcal{C} = \{3\} \\ 18 & \text{if } \mathcal{C} = \{1, 2\} \\ 48 & \text{if } \mathcal{C} = \{1, 3\} \\ 55 & \text{if } \mathcal{C} = \{2, 3\} \\ 80 & \text{if } \mathcal{C} = \{1, 2, 3\}. \end{cases}$$

# Convex Games

## Definition

A characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  is said to be **convex** if the characteristic function  $v$  is supermodular, i.e., it satisfies  $v(\mathcal{C} \cup \mathcal{D}) + v(\mathcal{C} \cap \mathcal{D}) \geq v(\mathcal{C}) + v(\mathcal{D})$  for every pair of coalitions  $\mathcal{C}, \mathcal{D} \subseteq \mathcal{N}$ .

## Proposition

A characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  is convex, if and only if, for every pair of coalitions  $\mathcal{C}, \mathcal{D}$  such that  $\mathcal{C} \subset \mathcal{D}$ , and for every player  $i \in \mathcal{N} - \mathcal{D}$ , we have

$$v(\mathcal{C} \cup \{i\}) - v(\mathcal{C}) \leq v(\mathcal{D} \cup \{i\}) - v(\mathcal{D})$$

- ▶ Players become more useful if they join bigger coalitions.
- ▶ Convexity  $\Rightarrow$  Superadditivity.
- ▶ However, the converse may not be true!

**3-player majority game:** Consider a game  $\Gamma = (\mathcal{N}, v)$ , where  $\mathcal{N} = \{1, 2, 3\}$ , and  $v(\mathcal{C}) = 1$  if  $|\mathcal{C}| \geq 2$ , and  $v(\mathcal{C}) = 0$  otherwise. This game is superadditive. On the other hand, for  $\mathcal{C} = \{1, 2\}$  and  $\mathcal{D} = \{2, 3\}$ , we have  $v(\mathcal{C}) = v(\mathcal{D}) = 1$ ,  $v(\mathcal{C} \cup \mathcal{D}) = 1$ ,  $v(\mathcal{C} \cap \mathcal{D}) = 0$ .

# Simple Games

## Definition

A characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  is said to be **simple** if it is monotone and its characteristic function only takes values 0 and 1, i.e.  $v(\mathcal{C}) \in \{0, 1\}$ , for any  $\mathcal{C} \subseteq \mathcal{N}$ .

- ▶  $v(\mathcal{C}) = 1 \Rightarrow$  Winning Coalition.
- ▶  $v(\mathcal{C}) = 0 \Rightarrow$  Loosing Coalition.

## Claim

A simple game  $\Gamma = \{\mathcal{N}, v\}$  is superadditive, only if the complement of every winning coalition loses.

# Solution Concepts

Outcomes can be evaluated based on two sets of criteria:

- ▶ **Fair Distribution:** How well each agent's payoff reflects his/her contribution?
  - ▶ **Shapley Value**
  - ▶ Banzhaf Index
- ▶ **Coalition Stability:** What are the incentives for the agents to stay in the coalition structure?
  - ▶ Stable Set
  - ▶ **Core**
  - ▶ Nucleolus
  - ▶ Bargaining Set



# Fair Distribution: Shapley's Axioms

Let  $u_i^\Gamma$  denote the allocation (utility) to the  $i^{th}$  player in a game  $\Gamma = (\mathcal{N}, v)$ . Then, we desire the following four properties:

- **Efficiency:** Distribute the value of grand coalition to all agents, i.e.

$$\sum_{i \in \mathcal{N}} u_i^\Gamma = v(\mathcal{N}).$$

- **Dummy Player:** If a player  $i$  does not contribute to any coalition in  $\Gamma$ , then

$$u_i^\Gamma = 0.$$

- **Symmetry:** If two players  $i$  and  $j$  contribute equally to each coalition in  $\Gamma$ , then

$$u_i^\Gamma = u_j^\Gamma.$$

- **Additivity:** If the same set of players are involved in two coalitional games  $\Gamma_1 = (\mathcal{N}, v_1)$  and  $\Gamma_2 = (\mathcal{N}, v_2)$ , if we define  $\Gamma = \Gamma_1 + \Gamma_2 = (\mathcal{N}, v_1 + v_2)$ , then for every player  $i$ , we have

$$u_i^\Gamma = u_i^{\Gamma_1} + u_i^{\Gamma_2}.$$

# Finding a Fair Distribution...

*Assume we have a superadditive game, which results in a grand coalition!*

- ▶ Agent's allocation is proportional to his/her contribution in  $v(\mathcal{N})$ .
- ▶ Idea: As each agent joins to form the grand coalition, compute how much the value of the coalition increases, i.e., allocate  $u_i = v(\mathcal{N}) - v(\mathcal{N} - \{i\})$  to player  $i$ .

*This contribution is evaluated when the player is the last inclusion in  $\mathcal{N}$ .*

*But, what about players who joined the coalition before the last player?*

Let  $\Pi_{\mathcal{N}}$  denote the set of all permutations of  $\mathcal{N}$ , i.e., one-to-one mappings from  $\mathcal{N}$  to itself. Given a permutation  $\pi \in \Pi_{\mathcal{N}}$ , we denote by  $S_{\pi}(i)$  the set of all predecessors of  $i$  in  $\pi$ , i.e., we set

$$S_{\pi}(i) = \{ j \in \mathcal{N} \mid \pi(j) < \pi(i) \}.$$

*Example:* If  $\mathcal{N} = \{1, 2, 3\}$ , we have

$$\Pi_{\mathcal{N}} = \{ \{1, 2, 3\}, \{1, 3, 2\}, \{2, 1, 3\}, \{2, 3, 1\}, \{3, 1, 2\}, \{3, 2, 1\} \}.$$

Then, if  $\pi = \{2, 1, 3\}$ , we have

$$S_{\pi}(2) = \emptyset \quad S_{\pi}(1) = \{2\} \quad S_{\pi}(3) = \{1, 2\}$$

# Shapley Value

## Definition

The **marginal contribution** of an agent  $i$  with respect to a permutation  $\pi$  in a game  $\Gamma = (\mathcal{N}, v)$  is given by

$$\Delta_{\pi}^{\Gamma}(i) = v[S_{\pi}(i) \cup \{i\}] - v[S_{\pi}(i)].$$

## Definition

Given a characteristic function game  $\Gamma = (\mathcal{N}, v)$  with  $|\mathcal{N}| = N$ , the **Shapley value** of an agent  $i \in \mathcal{N}$  is given by

$$u_i(\Gamma) = \frac{1}{N!} \sum_{\pi \in \Pi_{\mathcal{N}}} \Delta_{\pi}^{\Gamma}(i).$$

## Theorem

Shapley's axioms *uniquely* characterize Shapley value. In other words, Shapley value is the only fair distribution scheme that satisfies all the Shapley's axioms.

# Shapley Value: Example

Consider the same ridesharing example, as stated earlier.

- ▶ Three commuters can share a taxi.
- ▶ Individual journey costs:  $P_1 : 6$ ,  $P_2 : 12$ ,  $P_3 : 42$ .
- ▶ The characteristic function is

$$v_1(C) = \begin{cases} 6 & \text{if } C = \{1\} \\ 12 & \text{if } C = \{2\} \\ 42 & \text{if } C = \{3\} \\ 12 & \text{if } C = \{1, 2\} \\ 42 & \text{if } C = \{1, 3\} \\ 42 & \text{if } C = \{2, 3\} \\ 42 & \text{if } C = \{1, 2, 3\}. \end{cases}$$

Permutation set  $\Pi_{\mathcal{N}} = \{\pi_1, \pi_2, \pi_3, \pi_4, \pi_5, \pi_6\}$ , where

$$\begin{aligned} \pi_1 &= \{1, 2, 3\}, & \pi_2 &= \{1, 3, 2\}, & \pi_3 &= \{2, 1, 3\}, \\ \pi_4 &= \{2, 3, 1\}, & \pi_5 &= \{3, 1, 2\}, & \pi_6 &= \{3, 2, 1\}. \end{aligned}$$

# Shapley Value: Example (cont...)

Given

$$\begin{aligned}\pi_1 &= \{1, 2, 3\}, & \pi_2 &= \{1, 3, 2\}, \\ \pi_3 &= \{2, 1, 3\}, & \pi_4 &= \{2, 3, 1\}, \\ \pi_5 &= \{3, 1, 2\}, & \pi_6 &= \{3, 2, 1\},\end{aligned}$$

$$\text{and } v_1(C) = \begin{cases} 6 & \text{if } C = \{1\} \\ 12 & \text{if } C = \{2\} \\ 42 & \text{if } C = \{3\} \\ 12 & \text{if } C = \{1, 2\} \\ 42 & \text{if } C = \{1, 3\} \\ 42 & \text{if } C = \{2, 3\} \\ 42 & \text{if } C = \{1, 2, 3\}. \end{cases}$$

Marginal contributions:

- ▶  $\pi_1: \Delta_1^\Gamma(1) = 6, \Delta_1^\Gamma(2) = 6, \Delta_1^\Gamma(3) = 30$
- ▶  $\pi_2: \Delta_2^\Gamma(1) = 6, \Delta_2^\Gamma(2) = 0, \Delta_2^\Gamma(3) = 36$
- ▶  $\pi_3: \Delta_3^\Gamma(1) = 0, \Delta_3^\Gamma(2) = 12, \Delta_3^\Gamma(3) = 30$
- ▶  $\pi_4: \Delta_4^\Gamma(1) = 0, \Delta_4^\Gamma(2) = 12, \Delta_4^\Gamma(3) = 30$
- ▶  $\pi_5: \Delta_5^\Gamma(1) = 0, \Delta_5^\Gamma(2) = 0, \Delta_5^\Gamma(3) = 42$
- ▶  $\pi_6: \Delta_6^\Gamma(1) = 0, \Delta_6^\Gamma(2) = 0, \Delta_6^\Gamma(3) = 42$

Shapley value:

- ▶  $u_1(\Gamma) = \frac{1}{6} \sum_{i=1}^6 \Delta_i^\Gamma(1) = 2$
- ▶  $u_2(\Gamma) = \frac{1}{6} \sum_{i=1}^6 \Delta_i^\Gamma(2) = 5$
- ▶  $u_3(\Gamma) = \frac{1}{6} \sum_{i=1}^6 \Delta_i^\Gamma(3) = 35$

# Stability of Coalitions: Core

- ▶ Consider a characteristic function game  $\Gamma = \{\mathcal{N}, v\}$  with an outcome  $(\mathcal{C}, u)$ .
- ▶ Let  $u(\mathcal{C})$  denote the total payoff of a coalition  $\mathcal{C}$  under  $u$ .
- ▶ Given a coalition  $\mathcal{C}$ , if  $u(\mathcal{C}) < v(\mathcal{C})$ , some agents can abandon  $\mathcal{C}$  and form their own coalition.

## Definition

A utility profile  $u$  is **stable** through a coalition  $\mathcal{C}$  if

$$\sum_{i \in \mathcal{C}} u_i \geq v(\mathcal{C}).$$

## Definition

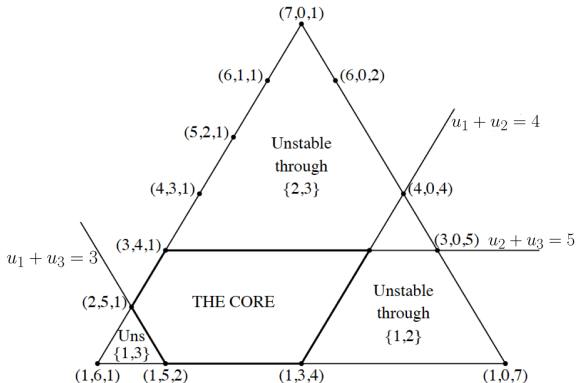
**Core** is defined as the set of all stable utility profiles, which is denoted as

$$\mathbb{C} = \left\{ u \in \mathbb{R}_+^N \mid \sum_{i \in \mathcal{C}} u_i \geq v(\mathcal{C}), \text{ for all } \mathcal{C} \subset \mathcal{N} \right\}.$$

# Core: An Example

Consider a characteristic function game  $\Gamma = \{\mathcal{N}, v\}$ , where  $\mathcal{N} = \{1, 2, 3\}$  and

$$v(C) = \begin{cases} 1 & \text{if } C = \{1\} \\ 0 & \text{if } C = \{2\} \\ 1 & \text{if } C = \{3\} \\ 4 & \text{if } C = \{1, 2\} \\ 3 & \text{if } C = \{1, 3\} \\ 5 & \text{if } C = \{2, 3\} \\ 8 & \text{if } C = \{1, 2, 3\}. \end{cases}$$



- Then, the utility profiles are those such that  $u_1 + u_2 + u_3 = 8$  such that  $u_1 \geq 1$ ,  $u_2 \geq 0$  and  $u_3 \geq 1$ .
- This is a hyperplane with vertices  $(7, 0, 1)$ ,  $(1, 0, 7)$ , and  $(1, 6, 1)$ .

*Is core always non-empty?*

# Core in Convex and Simple Games

## Theorem

Any convex game  $\Gamma = (\mathcal{N}, v)$  has a non-empty core.

## Definition

In a characteristic function game  $\Gamma = (\mathcal{N}, v)$ , a player  $i$  is called a **veto player**, if  $v(\mathcal{C}) = 0$  for all  $\mathcal{C} \subseteq \mathcal{N} - \{i\}$ .

## Theorem

A simple game  $\Gamma = (\mathcal{N}, v)$  has a non-empty core, if and only if there is a veto player in  $\mathcal{N}$ . Moreover, a utility profile  $u$  is in the core of  $\Gamma$  if and only if  $u_i = 0$  for every player  $i$ , who is not a veto player in  $\Gamma$ .



# Core and Superadditive Covers

## Definition

$\Gamma^* = (\mathcal{N}, v^*)$  is called a **superadditive cover** of  $\Gamma = (\mathcal{N}, v)$  if, for every coalition  $\mathcal{C} \subseteq \mathcal{N}$ ,

$$v^*(\mathcal{C}) = \max_{\mathcal{P}_{\mathcal{C}}} \sum_{\mathcal{C}_i \in \mathcal{P}_{\mathcal{C}}} v(\mathcal{C}_i),$$

where  $\mathcal{P}_{\mathcal{C}}$  denotes a partition of the coalition  $\mathcal{C}$ .

Consider  $\Gamma = (\mathcal{N}, v)$ :  $\mathcal{N} = \{1, 2, 3\}$  and

$$v(\mathcal{C}) = \begin{cases} 5 & \text{if } \mathcal{C} = \{1\} \\ 0 & \text{if } \mathcal{C} = \{2\} \\ 0 & \text{if } \mathcal{C} = \{3\} \\ 1 & \text{if } \mathcal{C} = \{1, 2\} \\ 1 & \text{if } \mathcal{C} = \{1, 3\} \\ 1 & \text{if } \mathcal{C} = \{2, 3\} \\ 1 & \text{if } \mathcal{C} = \{1, 2, 3\}. \end{cases}$$

Its superadditive cover  $\Gamma^* = (\mathcal{N}, v^*)$  is

$$v^*(\mathcal{C}) = \begin{cases} 5 & \text{if } \mathcal{C} = \{1\} \\ 0 & \text{if } \mathcal{C} = \{2\} \\ 0 & \text{if } \mathcal{C} = \{3\} \\ 5 & \text{if } \mathcal{C} = \{1, 2\} \\ 5 & \text{if } \mathcal{C} = \{1, 3\} \\ 1 & \text{if } \mathcal{C} = \{2, 3\} \\ 6 & \text{if } \mathcal{C} = \{1, 2, 3\}. \end{cases}$$

## Theorem

A characteristic function game  $\Gamma = (\mathcal{N}, v)$  has a non-empty core if and only if its superadditive cover  $\Gamma^* = (\mathcal{N}, v^*)$  has a non-empty core.

# Summary

- ▶ **Characteristic function game:** How to model players' rationality in coalitional games?
- ▶ **Subclasses:** Are there any special games that result in some specific coalitions?
- ▶ **Shapley value:** How to distribute a coalition's value in a fair manner amongst its members?
- ▶ **Core:** What is a stable coalition?