

Part A:  $i = 2^{2^k}$   
 $2^{2^k} < n$

runtime:  
 $\Theta(\log \log(n))$

$$\log(2^{2^k}) < \log(n)$$

$$2^k < \log(n)$$

$$k < \log \log(n)$$

Part B:

$$\sum_{j=0}^{\sqrt{n}} \sum_{i=1}^{j^3-1} \Theta(1) = n^3$$

$$\sum \Theta(1) + \sum_{j=0}^{\sqrt{n}} n^3 = \Theta(n^{7/2})$$

Part C:

iterating  $g$

$n$	$g$	$n$
1	0	$n$
2	1	
1	2	
0	3	

finding  $m$   
 $m = 2$

$$\log(m) = g \log(2)$$

$$g = \log(m)$$

$$\log(m) \log(m)$$

since its going for  $n$  many

if loop is true  $n$  times

first case: when if loop is true

$$\sum_{k=1}^n \log(n) = n \log n$$

second case: when if loop is not true  
 $\Theta(1)$  will execute  $n^2$  times

$$\Theta(n \log n) + \Theta(n^2)$$

have higher term

$$\Theta(n^2)$$

Part d:

inner for loop

$$\sum_{i=1}^{\text{size}} 1 = \text{size} = S$$

if if-statement is true

0 1 2 3 4 5 6 7 8 9 10

if(10 <= 12)  $\rightarrow$  3.10/2 = 15

6

15

15

5:70=15

Geometric Series

runs till  $\infty$   
if  $0 < r < 1$

$$\sum_{i=0}^n c^i = \sum_{i=0}^n \frac{c^{n+1} - 1}{c - 1} = \Theta(c^n)$$

factor is  $\frac{3}{2}$

$$\sum_{i=0}^n \left(\frac{3}{2}\right)^i = \sum_{i=0}^n \frac{3/2^{n+1} - 1}{(-1)} = \Theta(3/2^n)$$

if statement is not true

$\Theta(1)$  will execute  $n$  times

$$i < n \quad K = f(n)$$

$$\log\left(\frac{3}{2}\right)^K = n$$

$$\left(\frac{3}{2}\right)^K = \frac{n}{10}$$

$$\log\left(\frac{3}{2}\right)^{n/10} = K$$

Therefore the answer  
is  $\Theta(n)$