

ARM Instructions Worksheet #1

Function Call and Return

And their effect on registers PC, LR, and SP.

Prerequisite Reading: Chapter 3: Sections 3.1 and 3.2

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Objectives: To become acquainted with the web-based simulator ("CPUlator") and to use it to better understand how the ...

- 1. Program Counter (PC) is used to fetch an instruction,
- 2. Branch and Link (BL) instruction is used to call a function,
- 3. Branch Indirect (BX) instruction is used to return from a function,
- 4. Link Register (LR) is used to hold the return address, and
- 5. PUSH and POP instructions use the Stack Pointer (SP) to preserve and restore register content.

To do offline: Answer the questions that follow the listing below. (Numbers at far left are memory addresses.)

```
unified
                 .syntax
                 .global
                              start
00000000
          _stack_end:
                                     // Reserve memory for stack
                 .skip 100
00000064
          _tos:
00000064
          start:
                                     // *** EXECUTION STARTS HERE ***
00000064
                LDR
                       SP,=_tos
                BL
                       f1
                                     // Simple function call
00000068
                BL
                       f2
                                     // Nested function call
000006C
00000070
                BL
                       f3
                                     // Optimized nested function
                                     // End of demo
00000074
                       done
                                     // Simply returns
00000078 f1:
                BX
                       LR
0000007c
         f2:
                PUSH
                       {LR}
                                     // Preserve LR
                                     // Call f1 (changes LR)
00000080
                BL
                       f1
00000084
                POP
                       {LR}
                                     // Restore LR
8800000
                BX
                       LR
                                     // Return (Copies LR into PC)
0000008C
         f3:
                PUSH
                       {LR}
                                     // Preserve LR
                                     // Call f1 (Changes LR)
00000090
                BL
                       f1
00000094
                POP
                       {PC}
                                     // Return
00000098 done: B
                       done
                                     // infinite loop
                 .end
```

What is left in SP after executing the LDR instruction at 00000064 ₁₆ ?	Address of_tos
What is left in PC after executing the LDR instruction at 00000064 ₁₆ ?	00000068 base 16
What instruction is at the address that's now in the PC? (Include any referenced label)	BL f1
What address is left in register PC after executing the BL f1 instruction?	0000006C base 16
What instruction is at the address that's now in the PC? (Include any referenced label)	BL f2
What address is left in register LR after executing the BL f1 instruction?	00000068
What instruction is at the address that's now in the LR? (Include any referenced label)	BL f1
What value is in register PC after executing the BX $$ LR instruction at 00000078_{16} ?	0000070
What instruction is at the address that's now in the PC? (Include any referenced label)	0000007C

Getting ready: Now use the simulator to collect the following information and compare to your earlier answers.

- 1. Click here to open a browser for the ARM instruction simulator with pre-loaded code.
- 2. Press Ctrl-E to open the "Editor" window and notice the LDR pseudo-instruction.
- 3. Press Ctrl-D to replace the editor by the "Disassembly" window. Notice how the LDR pseudo-instruction has been replaced by a *real* LDR instruction that loads SP from a word in memory whose content is the address of label "tos" (top of stack).

Step 1: Executing the first instruction

The CPU registers are shown in the "Registers" window. Note that the PC value is 00000064₁₆. This is the starting address of the program. At that address is the LDR instruction that initializes the stack pointer (SP), highlighted in yellow to indicate that it is the next instruction to be executed. Press F2 once on the to execute that LDR instruction.

What is left in SP after executing the LDR instruction at 00000064₁₆?

What is left in PC after executing the LDR instruction at 00000064₁₆?

What instruction is at the address that's now in the PC? (Include any referenced label)

b1 0*78 (0*78:f1)

Step 2: Call function f1

The PC should contain the address of the instruction, "BL f1". Press F2 once to execute the instruction.

What address is left in register PC after executing the BL f1 instruction?

What instruction is at the address that's now in the PC? (Include any referenced label)

What address is left in register LR after executing the BL f1 instruction?

What instruction is at the address that's now in the LR? (Include any referenced label)

b1 0*7c (0*7c:f2)

Step 3: Return from function f1

The PC should contain the address of the instruction, "BX LR". Press F2 once to execute the instruction.

What is in register PC after executing the BX LR instruction at 00000078₁₆?

What instruction is at the address that's now in the PC? (Include any referenced label)

b1 0*7c (0*7c:f2)

Step 4: Continue exploring

Continue pressing F2 to step through the program, noting changes to registers PC, LR and SP at each step. Function f2 contains a call to function f1 that overwrites the return address of f2 in LR. In order for f2 to return properly, we use a PUSH {LR} at the entry of f2 to copy the return address onto the stack and then restore it with a POP {LR} before the return. Function f3 does the same, but eliminates the BX LR by popping directly into the PC.