## (1) Definition:

- · Synchronization: Threads share the same memory space ie they can also share objects
- (b) There might arise some situations (oritical) where it's desirable that only one thread at a time can access to a shared resources. (Let's deep dive to it by an example)
  - · If there is a licket booking applications and there's a field called remaining seats.
  - · There are multiple threads here responsible to book remaining seats for multiple users.
    - · If two (or multiple threads) by to book the remaining seats simultaneously.
    - · Assume, Remaining seat (0)
    - when 1 St thread books the seat and perform some other operations before updating seats.
    - . At the same time other threads also come Iseal visibility | Remaining seat >0 | as not updated yet. This is called . scace condition not recommended.

understand reace condition practically (Hands-on) and how to

Package: Synchronization STACK > LIFO (Last In, Frost Out)

el synchronized Blocks)

File: Stack Java Implementing Stack using an Array with threads / Synchronization Demo with Package Synchronization; Stacks (Synchronized met Stacks (Synchronized methods

public class Stack of

111) Attributes

private int [ Javay; Ustack Storage

private int et

. At the same time other threads also come into the actions and book the seat as seal visibility Remaining seal >0 as not updated yet. This is called race condition not recommended. understand peace condition practically (Hands-on) and how to prevent it' Package: Synchronization STACK -> LIFO (Last In, First Out) File: Stack java Implementing Stack using an Array with threads Synchronization Demo with Stacks (Synchronized methods Package Synchronization; El synchronized Blacks) public class Stack of 111) Attributes private int [ Javray; 11 Stack Storage private int start Top 112) Constantor with capacity of stack as an argument Public Stack(int capacity) { away = new int [capacity]; 110 bject for away Stack Top Index = -1; Il No element in the stack initially 1/3) methods () (113.1) check if stack is empty Public boolean isEmpty () { return stack Top Index <0; 113.2) check it stock is full public boolean is Full () of seetwer stack top Index >= avocay. length -1; Pockage -> Synchronization File - Stackk. java

other operations before updating seats.

. When 1 St th

```
113.3) PUSH() - Insert clement into STACKED
 public boolean push (int element) {
                 if (is Full ()) {
                       return false; 11 STACK is full, cannot insert
                     ++ stack TopIndex; 11 increment stack topindex
                     toy of
          Thread. sleep (1000); [I simulate delay (optional)

make thread wait (1 sec = 1000ms)
        Je Catch (Exception e) {
        System out println ("Thread interrupted: "+ e-getnessage ());
        avray[stack TopIndex] = clement; [Insert element into stack return true; [Push successful
     Joseph September 1960 00 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 000 1960 
113.4) POP() - Delete element from STACK
             Public int popl) {
                         if (is Emptyl)) of
                           seetwas Integer. MIN-VALUE; 11 Stack underflow case
                           int obj = away [ stack TopIndex ]; [ 1 store the top element
                           toy &
```

113.4) POP() -> Delete element from STACK (0) public int popl) q if (is Empty ()) of sections Integer. MIN\_VALUE; [[Stack underflow case] int obj = away [ stack topIndex ]; [ latore the top element Thread. Sleep (1000); // Stimulate delay (optional)
make thread wait (isec=1000ms) Jo catch (Exception e) of System out pointly ("Thread interrupted: "+ e.get Message W. Stack Tophodex --: Il correctly decrement stack printer return obj; // Return popped element (Moreale class Thread Testere) class Thread Tester of public Static void main (String I J args) { System out point in ("main is starting"); Mobject for class Stack Stack Stack = new Stack (copocity: 5); new Thread (1) > 9 int counter=0; While (++ counter(10) System out · println ( "Pushed: " + Black · push (climat: 1000 )); f, name: "pusher"). Start (); new Thread (() > & int counter = 0; While (++ counter (10) Systemout. Pointln ("Popped. 11+ 6tack. POP ()); 3. name: "popper"). Stort (); System. out . printly (Main is exiting");

(output Synchronization. Thread Tester (Package. Class Main is Starting Main is exiting "pusher" java lang krayInderOutOfBound Exception Exception in thread Popped:0 Popped: -2147483648 Popped: -2147483648 Topic - Synchronized Block [ Package: Synchronization. Synchronized Block File: Stackk-java a particular resource / particular piece of code me need to use synchronized keymord (Synchronized block with lock). - There are 2 ways to make code block synchronize (a) Apply synchronized on the method itself. (b) To make particular piece of code synchronized (c) Use synchronized block with explicit tock Synchronized ( Il holds critical code Il Allows just one thread to access a particular block of code at a time so called synchronized block. out of many within the task to get executed by performing it is task -> Now question arises how will I know which

time so called synchronized block.

Now we made just access to one thread out of many within the task to get executed by

-> Now paultin arised how will I finance which the able to access this Porticular synchronized alock.

- · So thread need to acquire lock.
- · Whichever thread will have access to the lock it will be able to enter this critical section (inside synchronized block) of execute that piece of code.
  - . In Java every object can we lock (All wrapper class can use it easily).
  - · Object Instance of wropper can use it s

    Synchronized (lock ) the synchronized

    black black

· In java we can use any objects as a lock.

Questions VVVI

To a multiple threads (ti,t2,t3) buying to gain access to push () to execute it.

The multiple threads (ti,tu,t5) trying to gain access to pop () to execute it.

Package: - Synchronization · Synchronization Block - multiple Throught

File: -> Stackk · java

Access same fine

## Answer

- (a) Now thread(ti) buying to gain access to push () method with explicit lock to get access to it -
- (b) and at the same time thread (t) trips to gets access to pop() with lock as well.
- (c) But if thread (t2) gets a lock and gets access to push() and perform it's execution in the meantime thread (t,) is waiting for lock.
  - (d) But thread (t) eart go and execute popl) because popls of push() are bounded by the same lock object, since this lock object can only be with one thread at a time.
  - (PWHI) El popi) are bounded by the same lock Objects.
- (f) Therefore, where threads gets access to the explicit lock he/she will only be able to access any of this methods () other of threads need to writ.
- (9) These two methods() pushing popil might be completely different from each other (Based on their functions) but since both are bounded by some lock therefore,

Gusp of clear answer for Interview

If a particular synchronized method has a

lock and that lock also restricts access to

ather synchronized methods as well all the

access to synchronized methods() will

be blocked for all other threads that

these have that particular lock.

(Pues) If there are 2 separate lock 1 + for push ()
adher for popl) then ?

Pockage -> Synchronization. Synchronized Block. Separate Locks

## Answer

- (a) If we assign 2 different explicit lock for two different methods () . i.e. Pash() & Pop().
- (b) Ideally we allow two threads to simultaneously sun the push () (1) pop () (at the same time).
- (C) This will not solve problem for pushes of pop() -> nace condition arises.
- (d) So it's strongly recommended, for pushes and Popl) in the stack we need to have same explicit lock so that we if pushes getting executed other threads won't be able to call popls that is fine but it its also
- (d) so it's strongly recommended to have same explicit lock for both the methods push() and pop() so that if one method() getting executed, other threads can't call pop() methods(other method) to avoid space condition (as both methods() are bounded by the same lock)

Synchronized Block ends here )

## Define race condition (WWVI) XXXXX

Race condition occurs when faw (or) more threads simultaneously update the same value and as a consequence, leave the value in an undefined (or) in consistent state.

(9) How can we make entire method () synchronized?

-> By applying synchronized on the method() itself

Package: Synchronization. Synchronized method Tile: Stackk. java

(92) When making method () synchronized what is the lock being used?

Define race condition (WIVI) XXXXX

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Package: Synchronization. Synchronized method The: Stackk. java

(92) When making method () synchronized what is the lock being used?

Behind the scenes/At background the entire Riece of code is awapped bising Synchronized (this) keyword.

-> compiler uses instance of the coverent object as the lock i.e. this keywood is a bock.

- for non-static synchronized method this keyword is used as lock.



Topic - Static synchronized methods		
Fockage File-	-> Synchronization Stack k. java	Static Synchronized methods Static Synchronized methods
200		
Prustim -	lunchumized sto	Air methods?

lock itself.

- we synchronized static methods on the class

> In Static Synchronized methods ()

Stockk. Class
is med as lock

(Class-name Class)

Topic > Thread safety (VVVVI)

Intervious

Thread safety means that a program (or) code segment can be safely accessed by multiple threads without leading to race conditions, data corruption (or) unexpected behavior.

Eg -> String Buffer