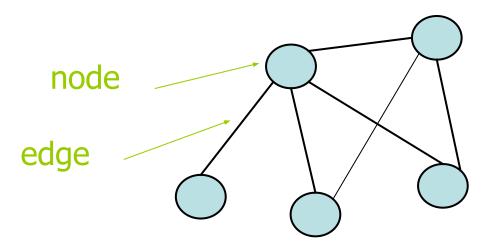
Graphs and Spanning Trees

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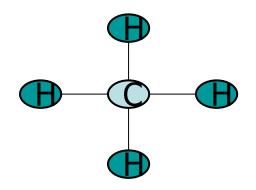
What is a graph?

- Graph represents relationship among the data items
- A graph G consists of
 - a set V of nodes (vertices)
 - a set E of edges (each edge connects two nodes)

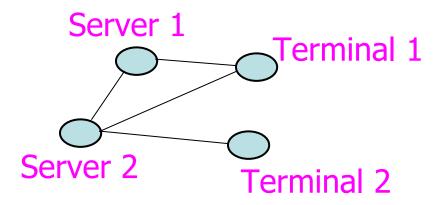


Examples of graphs

Molecular Structure



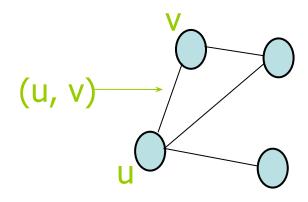
Computer Network



Other examples: electrical and communication networks, airline routes, flow chart, etc.

Formal Definition of graph

- The set of nodes is denoted as V
- For any nodes u and v, if u and v are connected by an edge, such edge is denoted as (u, v)
- The set of edges is denoted as E
- A graph G is defined as a pair (V, E)

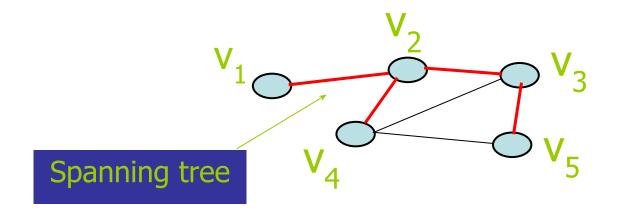


Problems related to Graph

- Spanning Tree
- Minimum Spanning Tree
- Shortest Path

Spanning Tree

 Given a connected undirected graph G, a spanning tree of G is a subgraph of G that contains all of G's nodes and enough of its edges to form a tree.



Spanning tree is not unique!

$$E=\{(v_1,v_2),(v_2,v_3),(v_2,v_4),(v_3,v_5)\}$$

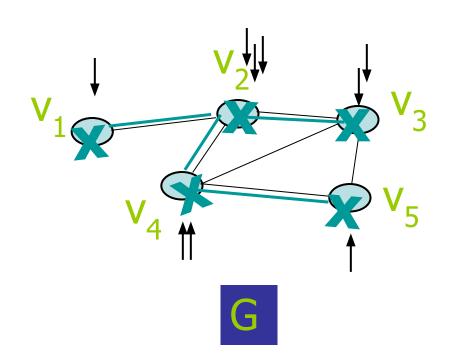
DFS spanning tree

```
Algorithm dfsSpanningTree(v)
mark v as visited;
for (each unvisited node u adjacent to v) {
   mark the edge from u to v;
   dfsSpanningTree(u);
```

 Similar to DFS, the spanning tree edges can be generated based on BFS traversal.

Generating SP based on DFS

		stack
→	V ₃	V ₃
	V ₂	V ₃ , V ₂
→	v ₁	v ₃ , v ₂ , v ₁
	backtrack	V ₃ , V ₂
→	V ₄	V ₃ , V ₂ , V ₄
	V ₅	V_3, V_2, V_4, V_5
→	backtrack	V ₃ , V ₂ , V ₄
	backtrack	V ₃ , V ₂
	backtrack	V ₃
	backtrack	empty

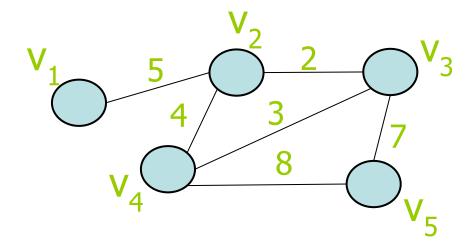


Minimum Spanning Tree (MST)

- Consider a connected undirected graph where
 - Each node x represents a country x
 - Each edge (x, y) has a number which measures the cost of placing telephone line between country x and country y
- Problem: connecting all countries while minimizing the total cost
- Solution: find a spanning tree with minimum total weight, that is, minimum spanning tree

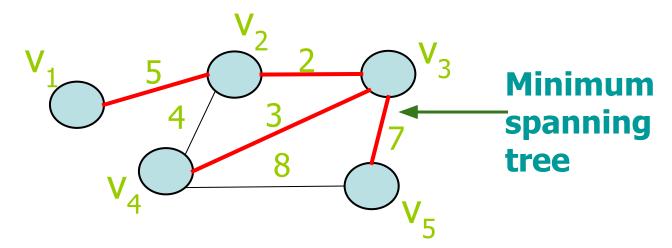
Formal definition of MST

- Given a connected undirected graph G.
- Let T be a spanning tree of G.
- $cost(T) = \sum_{e \in T} weight(e)$
- MST is a spanning tree T which minimizes cost(T)



Formal definition of MST

- Given a connected undirected graph G.
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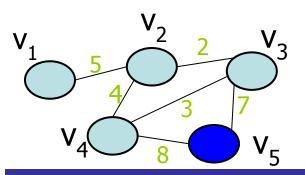


Prim's algorithm

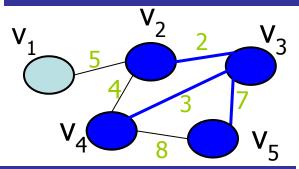
Algorithm PrimAlgorithm(v)

- Mark node v as visited and include it in the minimum spanning tree;
- while (there are unvisited nodes)
 - find the minimum edge (v, u) between a visited node v and an unvisited node u;
 - mark u as visited;
 - add both v and (v, u) to the minimum spanning tree;

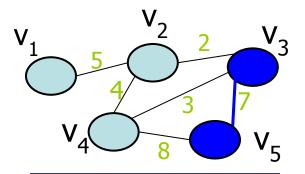
Prim's algorithm



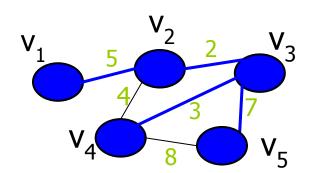
Start from v₅, find the minimum edge attach to v₅

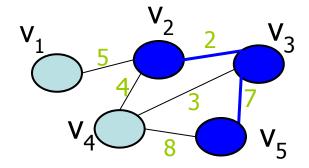


Find the minimum edge attach to v_2 , v_3 , v_4 and v_5



Find the minimum edge attach to v_3 and v_5





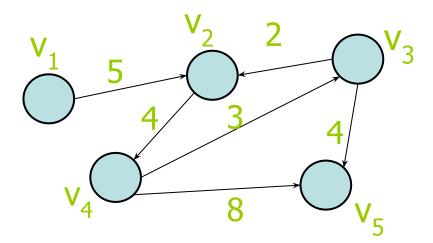
Find the minimum edge attach to v_2 , v_3 and v_5

Shortest path

- Consider a weighted directed graph
 - Each node x represents a city x
 - Each edge (x, y) has a number which represent the cost of traveling from city x to city y
- Problem: find the minimum cost to travel from city x to city y
- Solution: find the shortest path from x to y

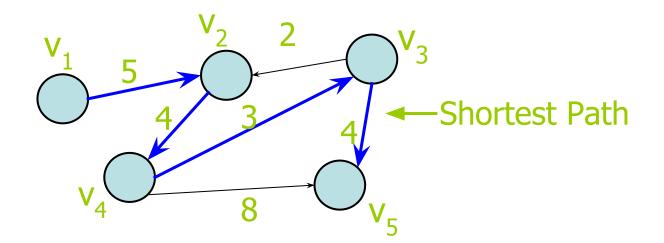
Formal definition of shortest path

- Given a weighted directed graph G.
- Let P be a path of G from x to y.
- $cost(P) = \sum_{e \in P} weight(e)$
- The shortest path is a path P which minimizes cost(P)



Formal definition of shortest path

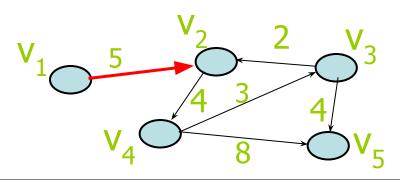
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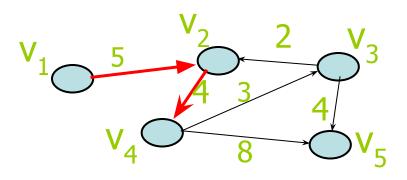
- Let G be a graph
- Each edge (u, v) has a weight w(u, v) > 0
- Find the shortest path starting from v₁ to any node v_i
- Let VS be a subset of nodes in G
- Let cost[v_i] be the weight of the shortest path from v₁ to v_i that passes through nodes in VS only

Algorithm shortestPath()

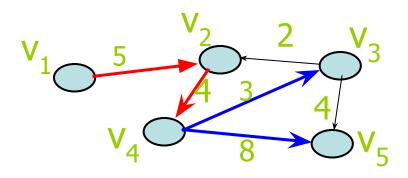
```
n = number of nodes in the graph;
for i = 1 to n
     cost[v_i] = w(v_1, v_i);
VS = \{ v_1 \};
for step = 2 to n {
     find the smallest cost[v<sub>i</sub>] s.t. v<sub>i</sub> is not in VS;
     include v<sub>i</sub> to VS;
     for (all nodes v<sub>i</sub> not in VS) {
           if (cost[v_i] > cost[v_i] + w(v_i, v_i))
                cost[v_i] = cost[v_i] + w(v_i, v_i);
```



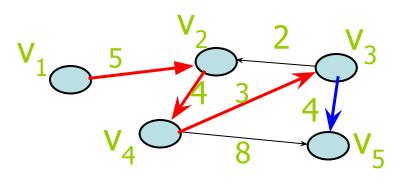
	V	VS	cost[v ₁]	cost[v ₂]	cost[v ₃]	cost[v ₄]	cost[v ₅]
1		[V ₁]	0	5	∞	8	8



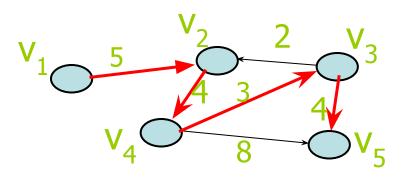
	V	VS	cost[v ₁]	cost[v ₂]	cost[v ₃]	cost[v ₄]	cost[v ₅]
1		[V ₁]	0	5	∞	8	8
2	V ₂	[V ₁ , V ₂]	0	5	∞	9	∞



	V	VS	cost[v ₁]	cost[v ₂]	cost[v ₃]	cost[v ₄]	cost[v ₅]
1		[V ₁]	0	5	∞	8	∞
2	V ₂	[V ₁ , V ₂]	0	5	∞	9	∞
3	V ₄	$[v_1, v_2, v_4]$	0	5	12	9	17



	V	VS	cost[v ₁]	cost[v ₂]	cost[v ₃]	cost[v ₄]	cost[v ₅]
1		[V ₁]	0	5	∞	∞	8
2	V ₂	[V ₁ , V ₂]	0	5	∞	9	∞
3	V_4	$[V_1, V_2, V_4]$	0	5	12	9	17
4	V ₃	$[V_1, V_2, V_4, V_3]$	0	5	12	9	16



	V	VS	cost[v ₁]	cost[v ₂]	cost[v ₃]	cost[v ₄]	cost[v ₅]
1		[v ₁]	0	5	8	8	∞
2	V ₂	[V ₁ , V ₂]	0	5	∞	9	∞
3	V_4	$[V_1, V_2, V_4]$	0	5	12	9	17
4	V_3	$[V_1, V_2, V_4, V_3]$	0	5	12	9	16
5	V ₅	$[V_1, V_2, V_4, V_3, V_5]$	0	5	12	9	16

Summary

- We have studied some basic concepts and algorithms
 - Spanning Tree
 - Minimum Spanning Tree
 - Shortest Path