

DAY
49

Recursion with ArrayList

1. Date → 25 October 2021.

2. Day → Monday.

3. Agenda →

1. Get Subsequence = CW
 2. Get KPC = CW
 3. Get stairs Path = CW
 4. Get Maze Path
 5. First Index
 6. Last Index
 7. All Indices of array
- Done in day 46 HW

Ques. 1 Get Subsequence

→ हमें एक string मिलता है जो इनput है, हमें उसे subsequence के subsequences return करने हैं जो output in a ArrayList है।

→ यहाँ परी question strings का module में आया है।
लेकिन iteratively किया है।

→ हमें अब यह question Recursively करना है और solve.
तो पहले जानते हैं What is Subsequence?

A subsequence is a part derived from a given string which is obtained by deleting some or no elements of the given string without changing the order of the remaining elements.

For a String of n length,
we have 2^n subsequences

eg → String = abc.

Subsequences of String \Rightarrow

Step 1) → Take a blank string

अब इसे भी लो (इस String को)
और एक इसकी copy लगाया
परियों c को उसकी correction
position पर रखें

Step 2)

→ Put last character
at correct position
same

— — c.

→ copied & put **c** (last char.)
at its pos.

इसे लो और

Copy करो तक (इन दोनों String की)

उत्तर कोड में b को उसकी correct
Position पर रखो

Step 3)

→ Copy it
→ Put last second
character at correct
position

— — c

→ copied same as
above & put
b at its correct
position

b —

Take these

String & Copy all these
& string and put a at correct posn

→ last second character

Step 4

— — —
— — c
— b —
— b c.

a — —
a — c
a b —
a b c.

Copied as Previous
Strings and
have put
a (last 3rd) character
at its correct position

These are the Subsequences
of abc, because

character in abc = 3

Subsequences of abc = $(2)^n = (2)^3 = 8$

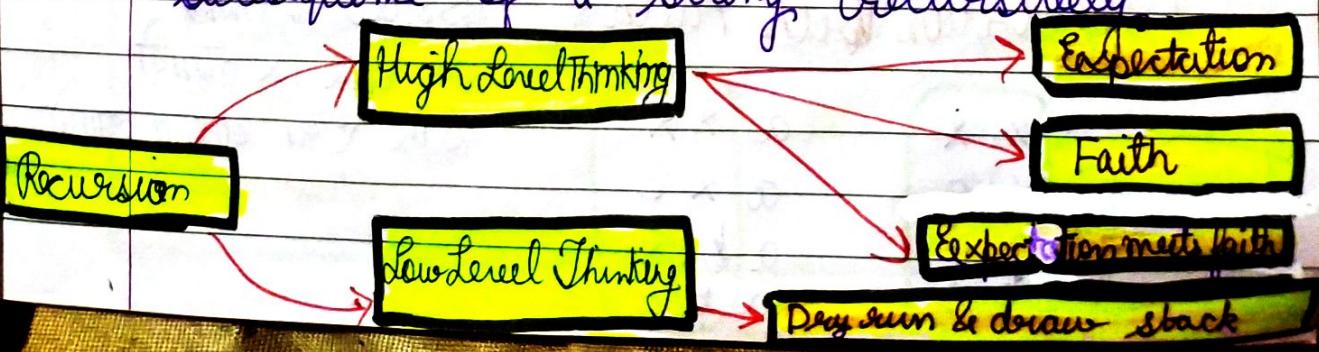
& these
are total

8 Subsequences.

Hence, we are done
with finding all
the subsequences of
abc

* अब हमने ये समझा किया है कि Subsequence का एक एक व्यापकीय है।

Now we need its code to find all the
subsequence of a string, recursively.



1. High Level Thinking

a) Expectation

$$gss(abc) = \boxed{x \ x \ x}$$

get subsequence.

$$x \ x \ c.$$

$$x \ b \ x$$

$$x \ b \ c.$$

$$a \ x \ x.$$

$$a \ x \ c.$$

$$a \ b \ x$$

$$a \ b \ c.$$

Expectation

of gss(abc)

function

of 8 subsequences

देता है

$x \ x \ x.$	$a \ x \ x.$
$x \ x \ c.$	$a \ x \ c.$
$x \ b \ x.$	$a \ b \ x$
$x \ b \ c.$	$a \ b \ c.$

b) Faith → faith, expectation का दोनों part

Same.

$$\begin{matrix} x & x \\ x & c \\ b & x \\ b & c \end{matrix}$$

होता है

$$gss(bc) \rightarrow$$

) faith की पैरे output हो
अपने आप ही मिलेगा
अगर हमने bc pass
किया तो.

c) Expectation meets Faith

$$\begin{matrix} x & x & x \\ x & x & c \\ x & b & x \\ x & b & c \end{matrix}$$

$$\begin{matrix} a & x & x \\ a & x & c \\ a & b & x \\ a & b & c \end{matrix}$$

एक बार x लगाए
उसके बार a लगाए

High Level Thinking about E; so now, we can write the basic code.

Get Subsequence <code>

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```
→ import java.util.*;  
→ import java.io.*;  
→ public class Main {  
    → public static void main (String [] args)  
        → Scanner scan = new Scanner (System.in);  
        → String str = scan.nextLine();  
        → ArrayList <String> Subsequence = gss(str);  
        → System.out.println (Subsequence);  
    }  
    → getsubsequence
```

Expectation // $gss(abc) = [\dots, \dots, \dots, abc] + [a\dots; a\cdot c, ab\dots, abc]$

Truth // $gss(abc) = [\dots, \dots, \dots, bc] + [a\dots; a\cdot c, b\dots, bc]$

Expectation meets truth // $= \dots + gss(bc) + a + gss(bc)$

```
→ public static ArrayList <String> gss (String str)
```

```
This is  
the  
base  
case  
& it  
is  
explained  
in Low  
level  
Thinking  
[ if (str.length () == 0) → get subsequence  
    → { ArrayList <String> bres = new ArrayList <> ();  
        → String theSubsequence = ""; } Blank etajit subsequen  
        → bres.add (theSubsequence);  
    } return bres;  
} first character of given string  
char ch = str.charAt(0);  
String ros = str.substring(1); → recursion tip result  
ArrayList <String> rres = gss(ros);  
ArrayList <String> mres = new ArrayList <> ();  
for (String rstr: rres) → my result  
    → { mres.add ("_" + rstr); }  
    → for each  
        → for (String rstr: rres) → for each element  
            → { mres.add (" " + ch + rstr); }  
    }  
return mres;
```

2. Low Level Thinking

→ Dry run & draw stack.

String at
0th character
P(Σ)
P(Σ)

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Line 1 → char ch = str.charAt(0); // abc → a
 Line 2 → String res = str.substring(1); // abc → bc
 Line 3 → ArrayList<String> reses = gss(res); // gss(bc)
 Line 4 → ArrayList<String> mres = newArrayList();

rest of
the string
(0th character
of abc)

→ for (String rstr : reses)

Line 5 } { mres.add(ch + rstr);

{ }

{ }

// mstr = [---, -c, -b-, -bc]

for (String mstr : mres)

Line 6 } { mres.add(ch + rstr);

{ }

// mstr = [---, -c, -b-, -bc]

Line 7 } { mres.add(ch + rstr);

return mres;

→ This is the Base Case.

$$\rightarrow x^0 = 1 \quad (\text{Multiplication Identity 1})$$

$$\rightarrow x^2 = x * x \quad (\text{Sum of Identity})$$

$$\rightarrow x^1 = x \quad (\text{Identity})$$

→ In how many ways can we select 0 apple out of 5 apples?

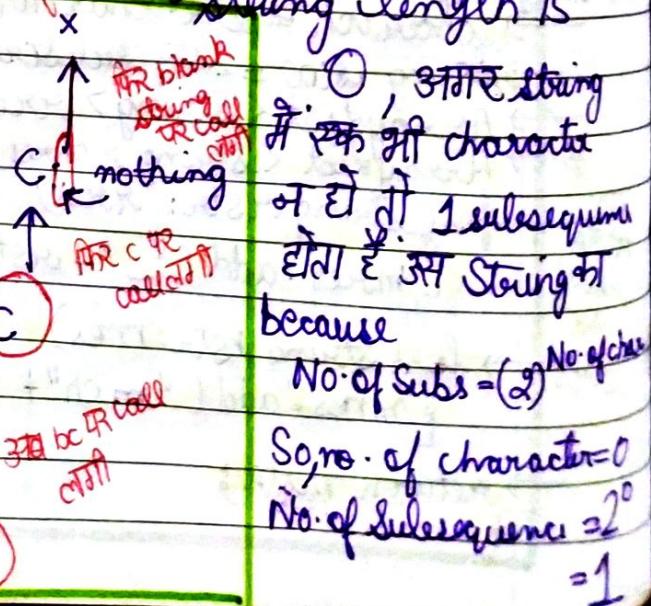
→ 0 apple can be

selected in 1 way

i.e. Don't select any apple

1 2 3 right side call

a b c



If old String $x = [""]$, एक blank string with 0 characters
 String $\in x$, तो उसके Subsequence.
 हमें 1

$$\text{bcz no. of character} = 0 \\ \text{no. of subsequence} = (2)^{\text{no. of character}} \\ = (2)^0 \\ = 1$$

Time Complexity of get Subsequence
 $O(n) \rightarrow$ Linear bcz a recursion call is made along with using for loop.

Space Complexity of get Subsequence $\rightarrow O(1) \rightarrow$ No extra space required.: constant.

अब क्या हमें तो पता चलगया कि 1 Subsequence होगा लेकिन उस Subsequence में value क्या होगी?

[]

कोई भी value नहीं होगी

Blank Subsequence

तो Base Case Code कैसे लिखें?

if (str.length() == 0)

base result

ArrayList<String> bres = new ArrayList<>();
 bres.add("");
 return bres;

प्रारंभिक blank subsequence add करे arraylist में, even without a space

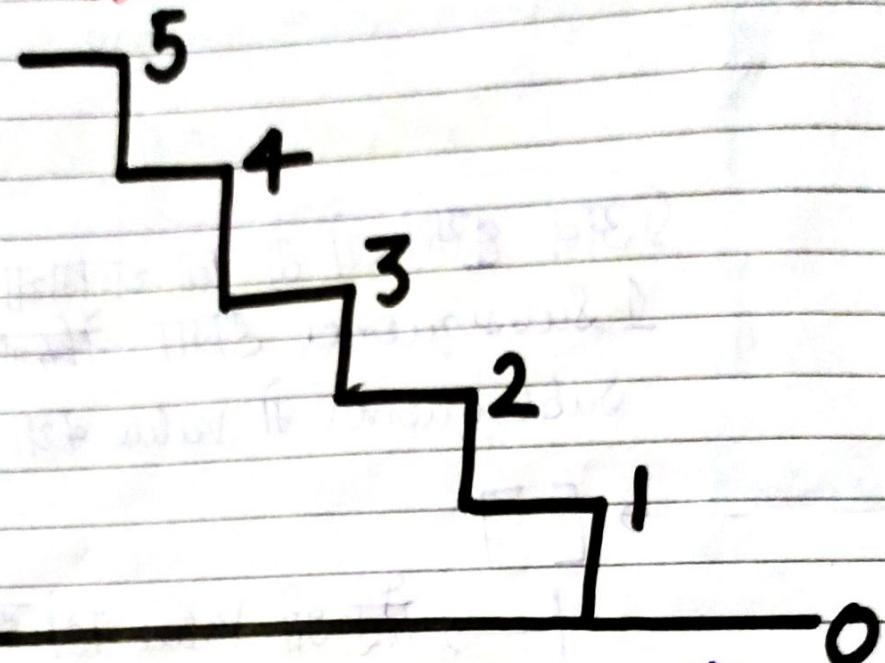
[" "]

ArrayList का Size 1 है अब

bres.get(0).length = 0 लेकिन इस ArrayList में भी String है उसका Size 0 है।

Get Stair Paths

L एवं यूक्टि की स्केच (स्टीपली) गेम है।



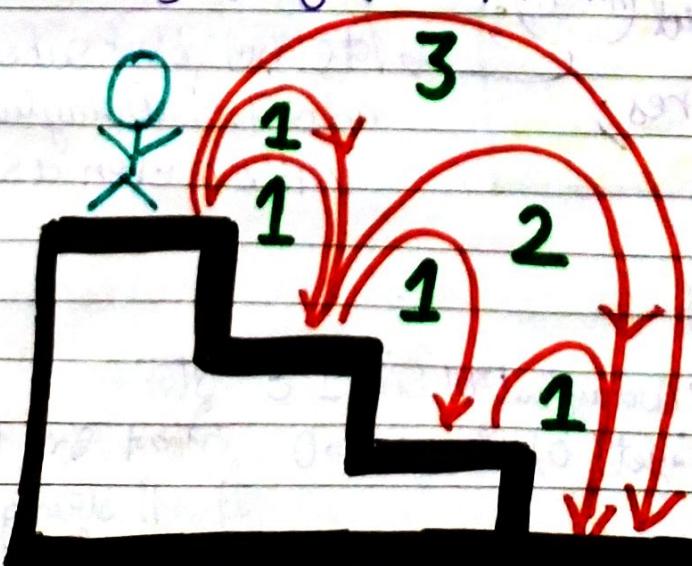
इस स्केच पर यह कदम रख सकते हैं।

या को कदम रख सकते हैं।

या तीन कदम रख सकते हैं।

स्केच में 0 से 5 तक जाने के सारे तरीके कानून हैं।
0 से 5 तक पहुँचने के 16 तरीके हो सकते हैं।

या फिर 5 से 0 तक जाने के सारे तरीके बहुमध्य हैं।
5 से 0 तक जाने के भी 16 तरीके हो सकते हैं।

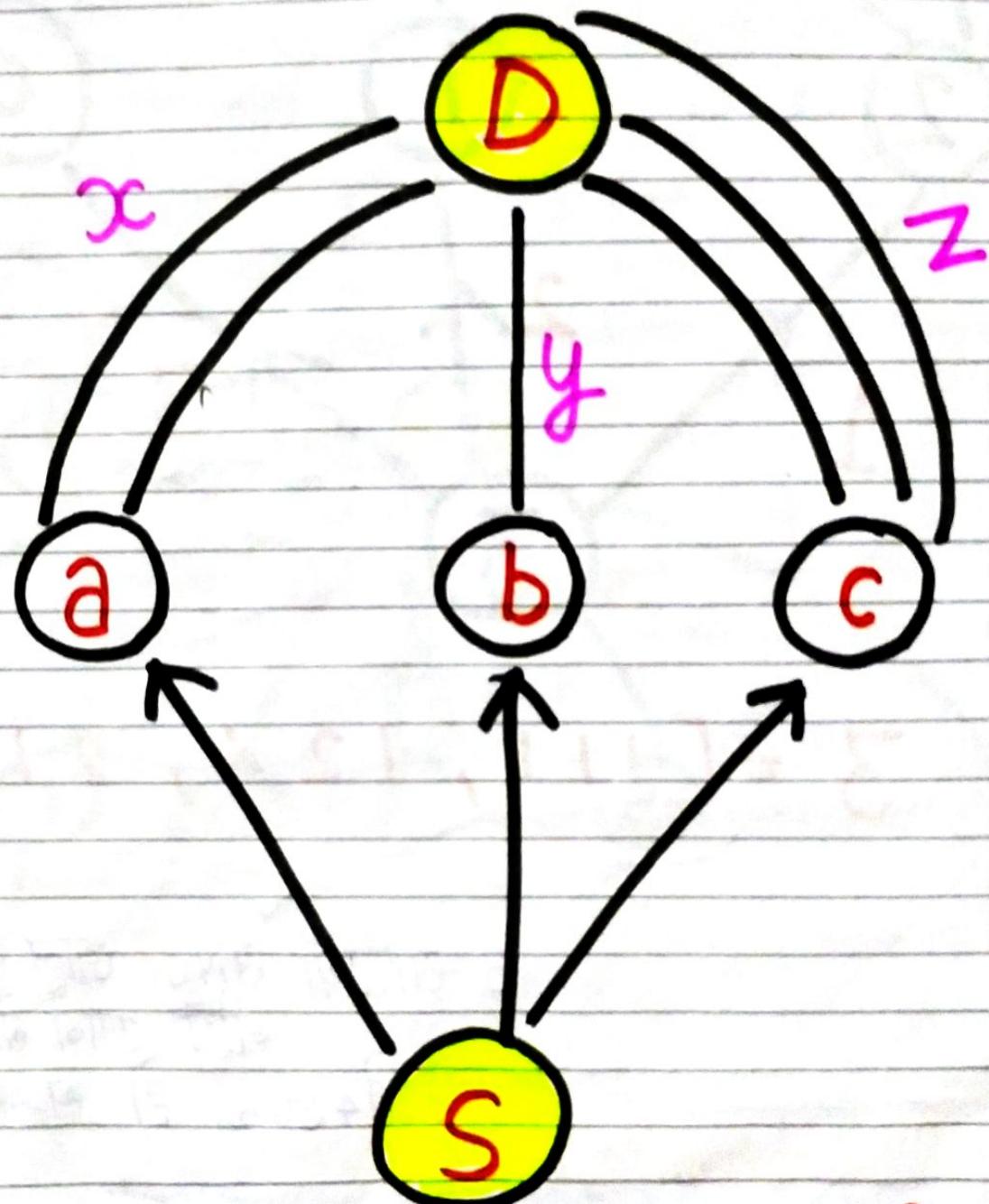


S=Source

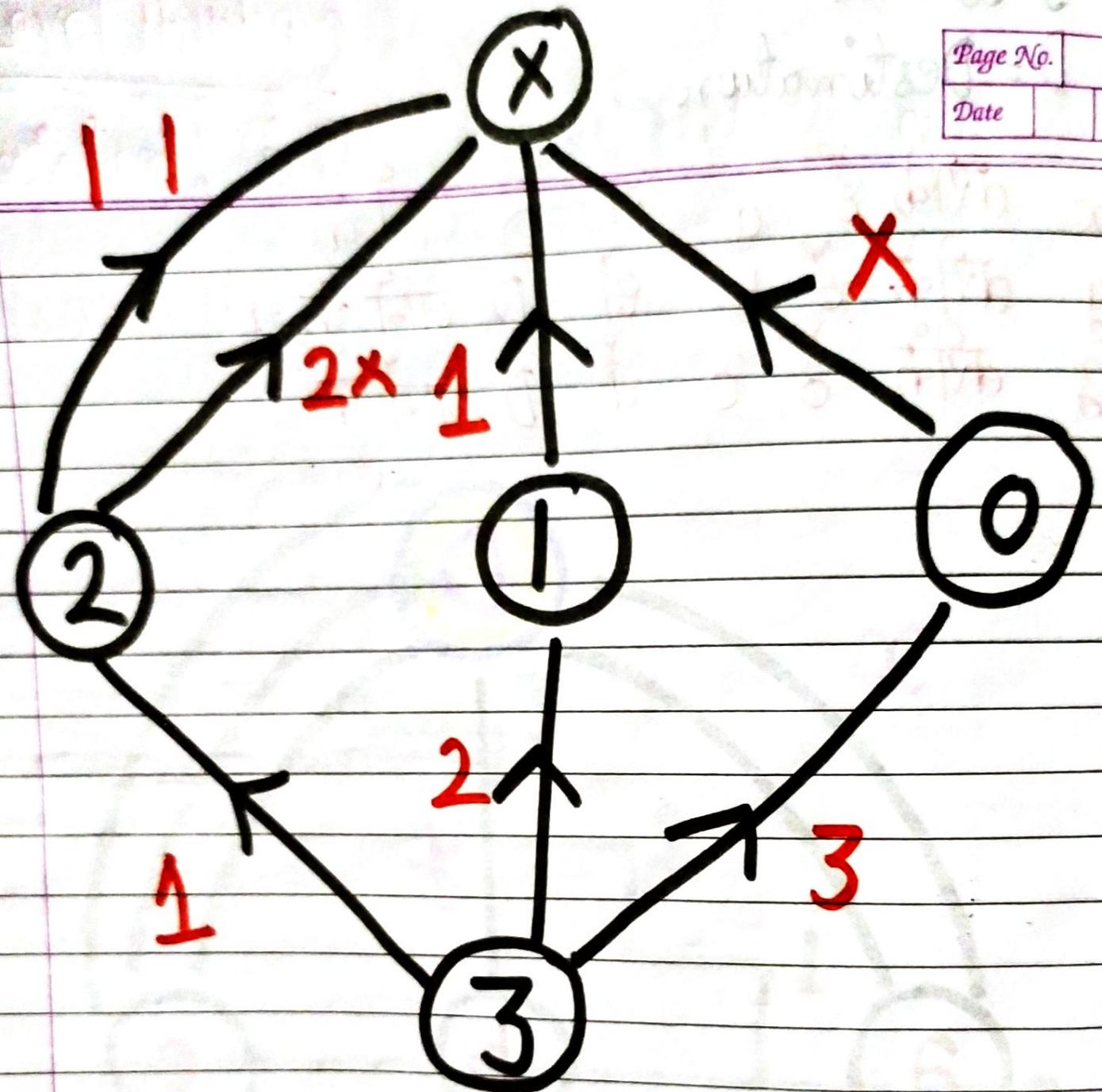
D= Destination

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x तरीके से a से D जाने के
y तरीके से b से D जाने के
z तरीके से c से D जाने के

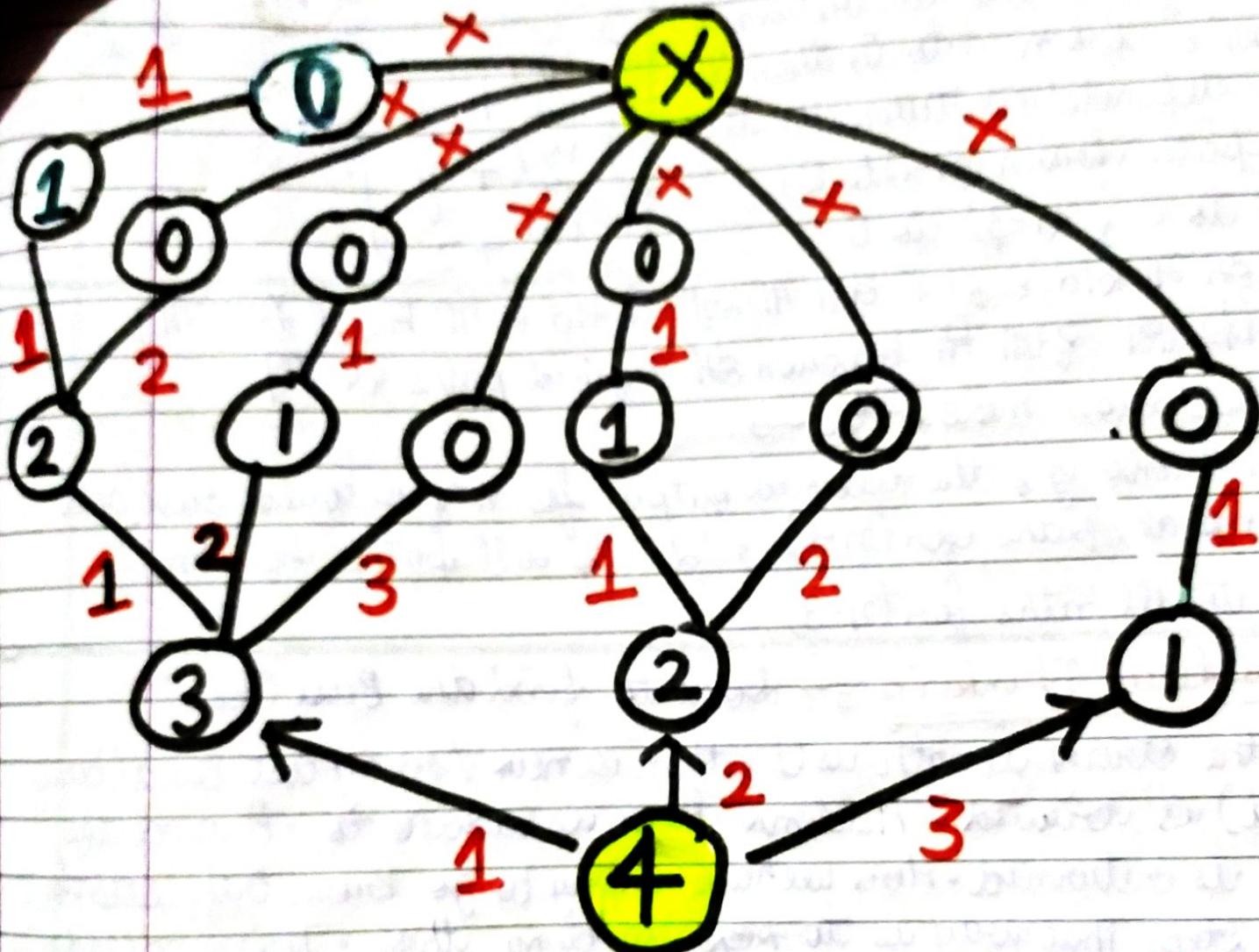


तो S से D तक जाने के
 $x + y + z$ तरीके होंगे।



$$3 = [111, 12x, 21, 3x]$$

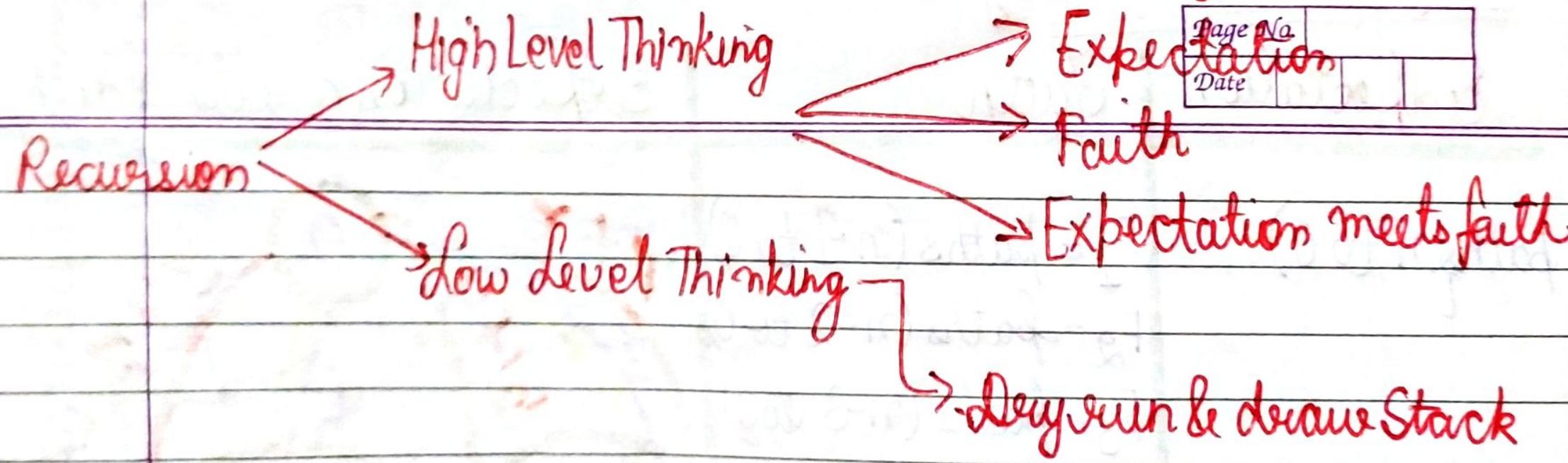
3 सीढ़ियाँ ऊपर जाने के/
नीचे जाने के
पैर 4 Steps हो सकते हैं।



4. [1111, 112, 121, 13, 211, 22, 31]

4 सीटियों ऊपर जाने के / नीचे आने के थे 7
तरीके हो सकते हैं।

We need to solve this Problem Recursively



1. High Level Thinking :-

Expectation | Faith

paths(n to 0).

$$F_1 = \text{paths}(n-1 \text{ to } 0)$$

$$F_2 = \text{paths}(n-2 \text{ to } 0)$$

$$F_3 = \text{paths}(n-3 \text{ to } 0)$$

1. Set Expectation → We expect that we will get all the paths from n to 0.

2. Build Faith → यह faith है कि हमारा code सब path return कर सकता है. from n to 0, then it can definitely return us all the path from (n-1) to 0, (n-2) to 0, (n-3) to 0.

ऐसा हम मानलेते हैं, ये बस मानलो, ये मत सोचो कि ये कैसे होगा।
बस ऐसा दर्शक कि function द्वारा required paths होंगा।

3. Expectation meets faith →

→ for printing the desired output for n, we could just add "1" to all paths for (n-1), "2" to all paths for (n-2), "3" to all paths for (n-3).

2. Low Level Thinking :- Done to find the Base Case

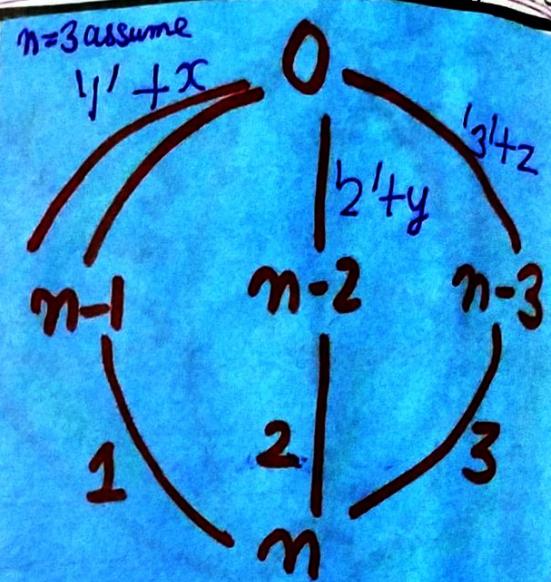
a) If the string length is 0, then a new ArrayList (base result) is created. A blank path is added to it and the list is returned. Here we have 1 way to go from 0th step to 0th step. That way is to keep standing there. This is called the Positive Base Case.

b) When n < 0, it means that there are no stairs left and the person will be going into the basement which isn't practically possible. Here, we don't add any path to the ArrayList.

This is called the Negative Base Case.

Expectation meets faith

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$$\left\{ \begin{array}{l} 1' + x \\ 2' + y \\ 3' + z \end{array} \right\}$$

Time Complexity:- $O(3^n)$

This time

complexity is exponential because for each state , 3 recursion calls are made .

Space Complexity: $O(1)$

~~As no extra space is required . So Space complexity is constant . If we include the space in recursive stack , then space complexity is $O(n)$.~~

Get Stair Path <code>

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```
import java.util.*;  
import java.io.*;  
public class Main {  
    public static void main (String [] args) {  
        Scanner scan = new Scanner (System.in);  
        int n = scan.nextInt();  
        ArrayList < String > paths = getStairPaths (n);  
        System.out.println (paths);  
    }  
    public static ArrayList < String > getStairPath (int n) {  
        if (n == 0) {  
            ArrayList < String > leres = new ArrayList <> ();  
            leres.add ("");  
            return leres; // [ ] base result  
        } else if (n < 0) {  
            return null; // [ ] base result  
        }  
        ArrayList < String > leres = new ArrayList <> ();  
        for (String le : leres) {  
            leres.add (le + "1");  
            leres.add (le + "2");  
        }  
        return leres; // [ ] base result  
    }  
}
```

These `ArrayList < String > pathfromnm1 = getStairPath (n-1);`
are `ArrayList < String > pathfromnm2 = getStairPath (n-2);`
Recursive `ArrayList < String > pathfromnm3 = getStairPath (n-3);`,
Tells `ArrayList < String > pathfromn = new ArrayList <>();`

String
path from n to m
for each string path in ArrayList paths from nm1

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→ for (String pathfromnm1 : pathsfromnm1)

{

String pathfromn = 1 + pathfromnm1;
pathsfromn.add(pathfromn);

}

→ for (String pathfromnm2 : pathsfromnm2)

{

String pathfromn = 2 + pathfromnm2;
pathsfromn.add(pathfromn);

}

→ for (String pathfromnm3 : pathsfromnm3)

{

String pathfromn = 3 + pathfromnm2;
pathsfromn.add(pathfromn);

}

return pathsfromn;

}