

addition
minimum

K query
P query

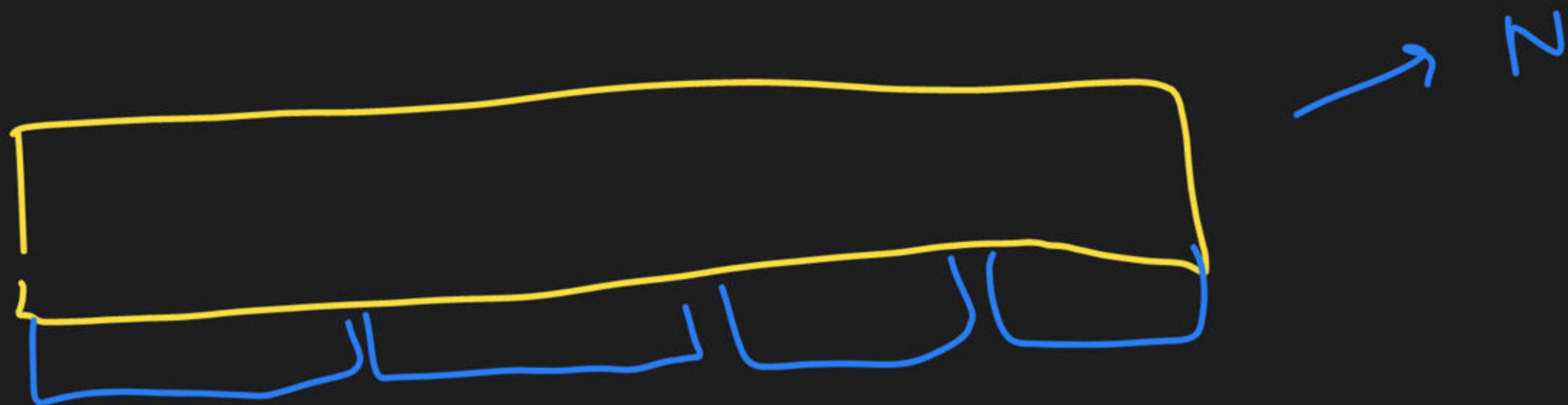
Square Root Decomposition on Arrays

Course on Square Root Decomposition

Queries

Tree (Bonus)

MO's algorithm



B → we choose the best B

$$O\left(B + \frac{N}{B}\right)$$

Best value of B , if $T = O\left(B + \frac{N}{B}\right)$

A. $N^{1/2}$ ✓✓←

B. $N^{1/3}$

C. $N^{1/4}$

D. N

→ Calculus

→ AM = GM

$$B = \frac{N}{B}$$

$$\Rightarrow B = \sqrt{N}$$

Best value of B , if $T = O\left(B^2 + \frac{N}{B}\right)$

A. $N^{1/2}$

B. $N^{1/3}$ ✓

C. $N^{1/4}$

D. $(N \log N)^{1/2}$

E. $N^{1/2} (\log N)$

$$B^2 = \frac{N}{B}$$

$$B = \sqrt[3]{N}$$

Best value of B, if $T = O\left(B + \frac{N \log N}{B}\right)$

A. $N^{1/2}$

B. $N^{1/3}$

C. $N^{1/4}$

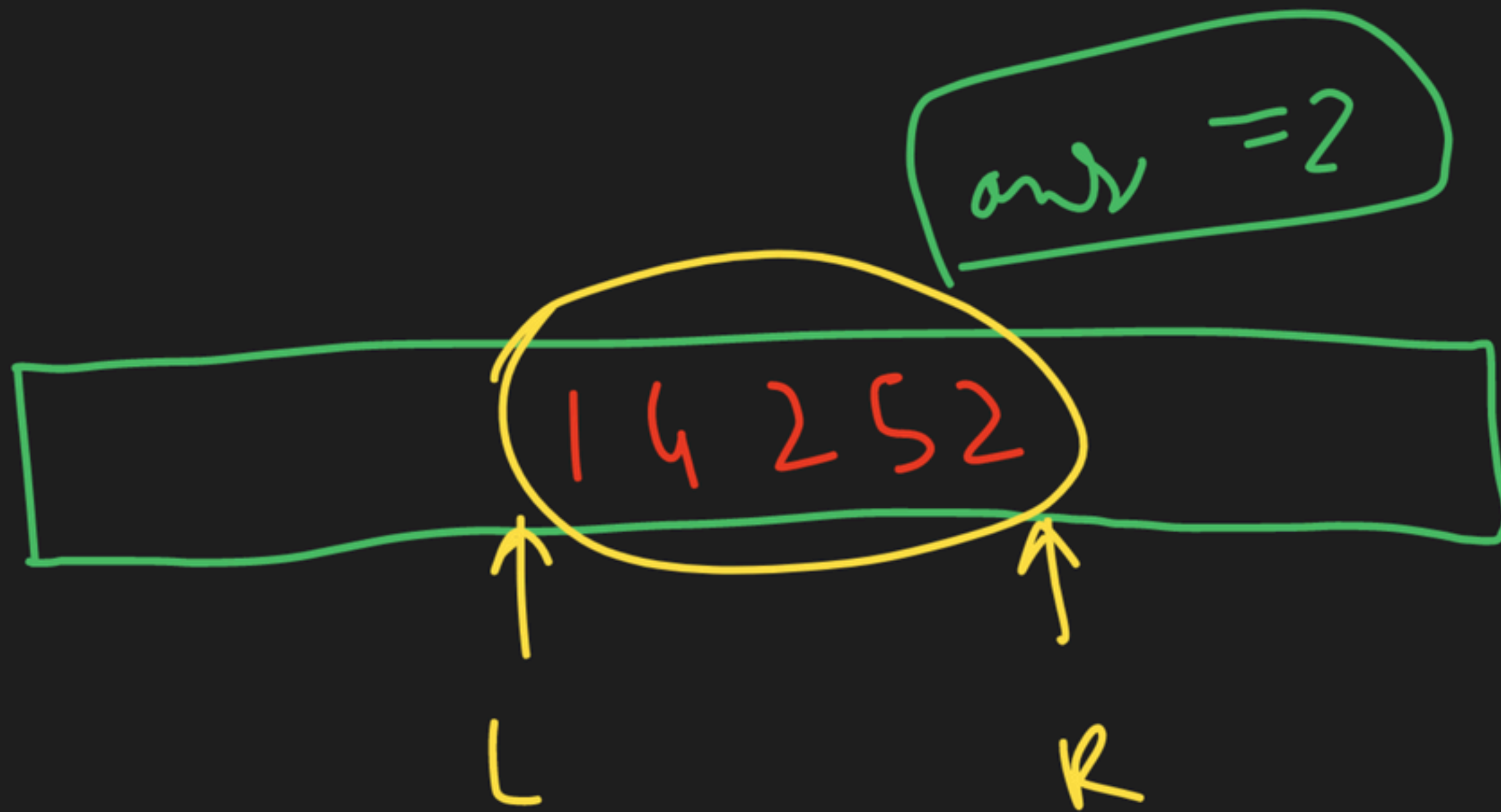
D. $(N \log N)^{1/2}$ ✓

E. $N^{1/2}(\log N)$

$$K = N \log N$$

$$O\left(B + \frac{K}{B}\right)$$

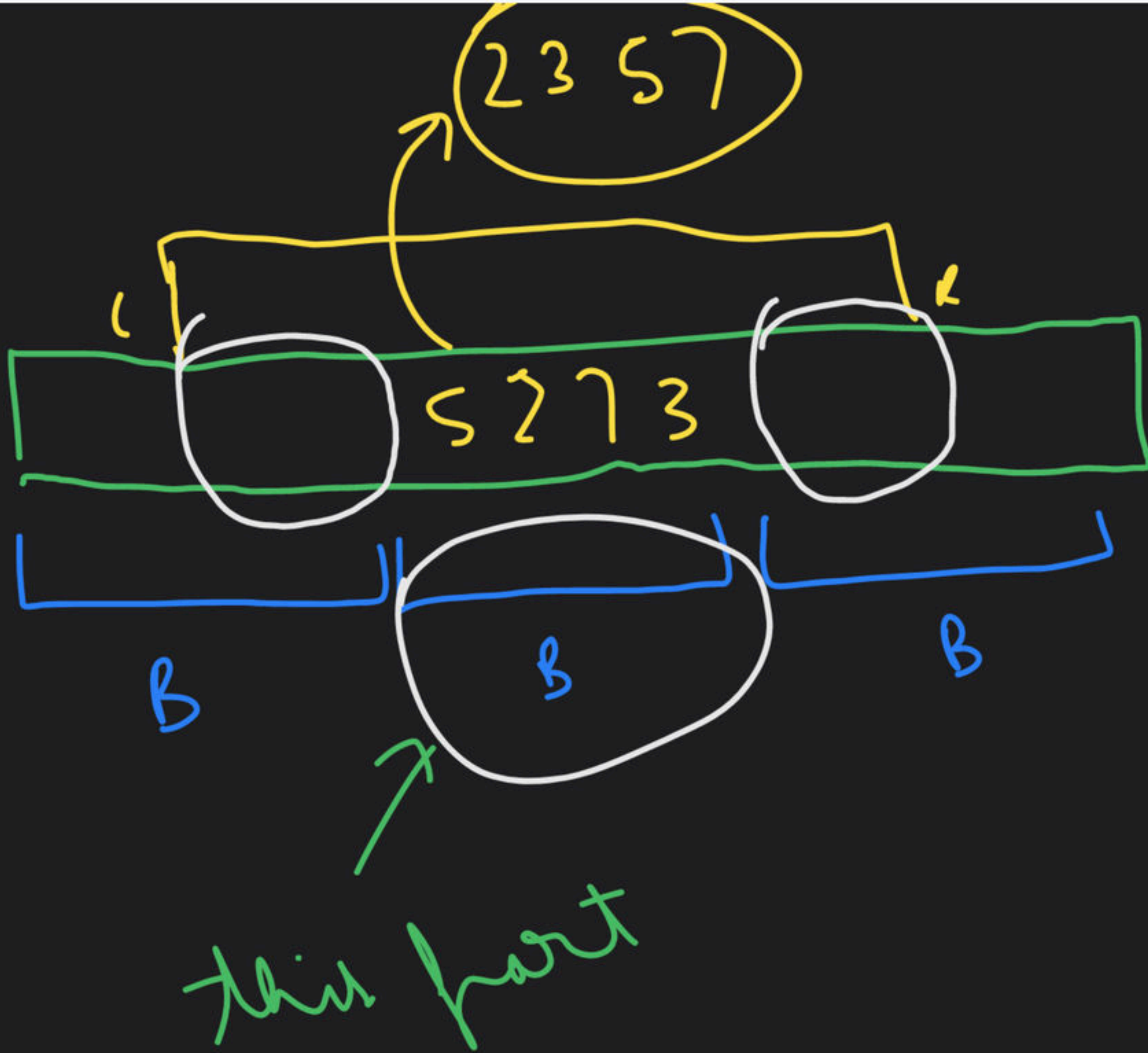
$$\rightarrow B = \sqrt{K} \\ = \sqrt{N \log N}$$



(L, R, x)

A red arrow points from the 'x' in the tuple to a red curly brace '}' below it.

update \rightarrow (id, val)



$>$ $<$
 $||$
 \wedge

unknown
 $O(B)$ constant

$O\left(\frac{N}{B} \cdot \log B\right)$

→ worst

$$O\left(B + \frac{N \log B}{B}\right)$$

something

$$O\left(\frac{N}{B} \log B\right)$$

wrong

→ $\log N$

~~$\frac{N}{B} \log N$~~

binary
search
no find

2 3 5 7
↑

$O(B \log B)$

$O(B)$

5 2 7 3

$$O(B) \left(B + \overset{\text{SES}}{\frac{N}{B} \log B} \right)$$

queries only !!

$$O\left(B + \frac{N \log B}{B}\right) \approx O\left(B + \frac{N \log \sqrt{N}}{B}\right)$$

Approx best value of B !

$$B = \text{slightly} > \sqrt{N}$$

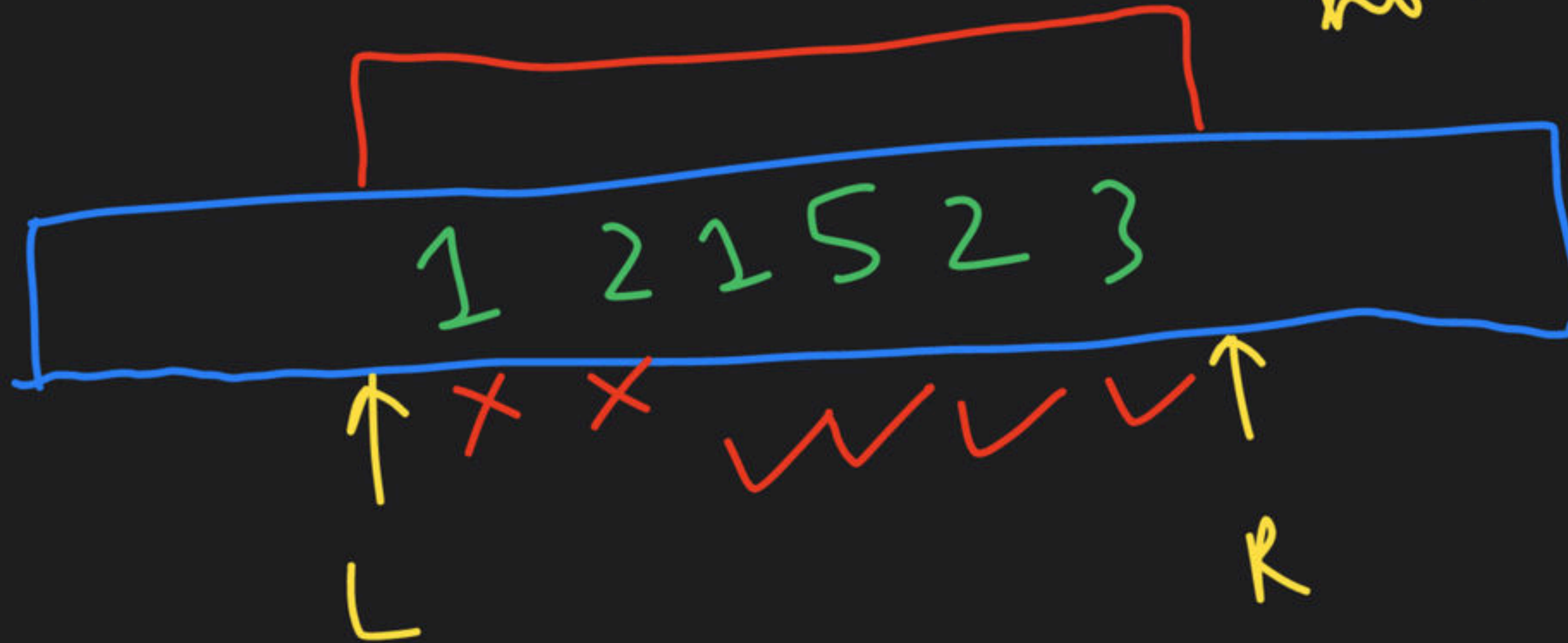
$$B \approx \sqrt{N}$$

$$O\left(B + \underbrace{N \log \sqrt{N}}_B\right)$$

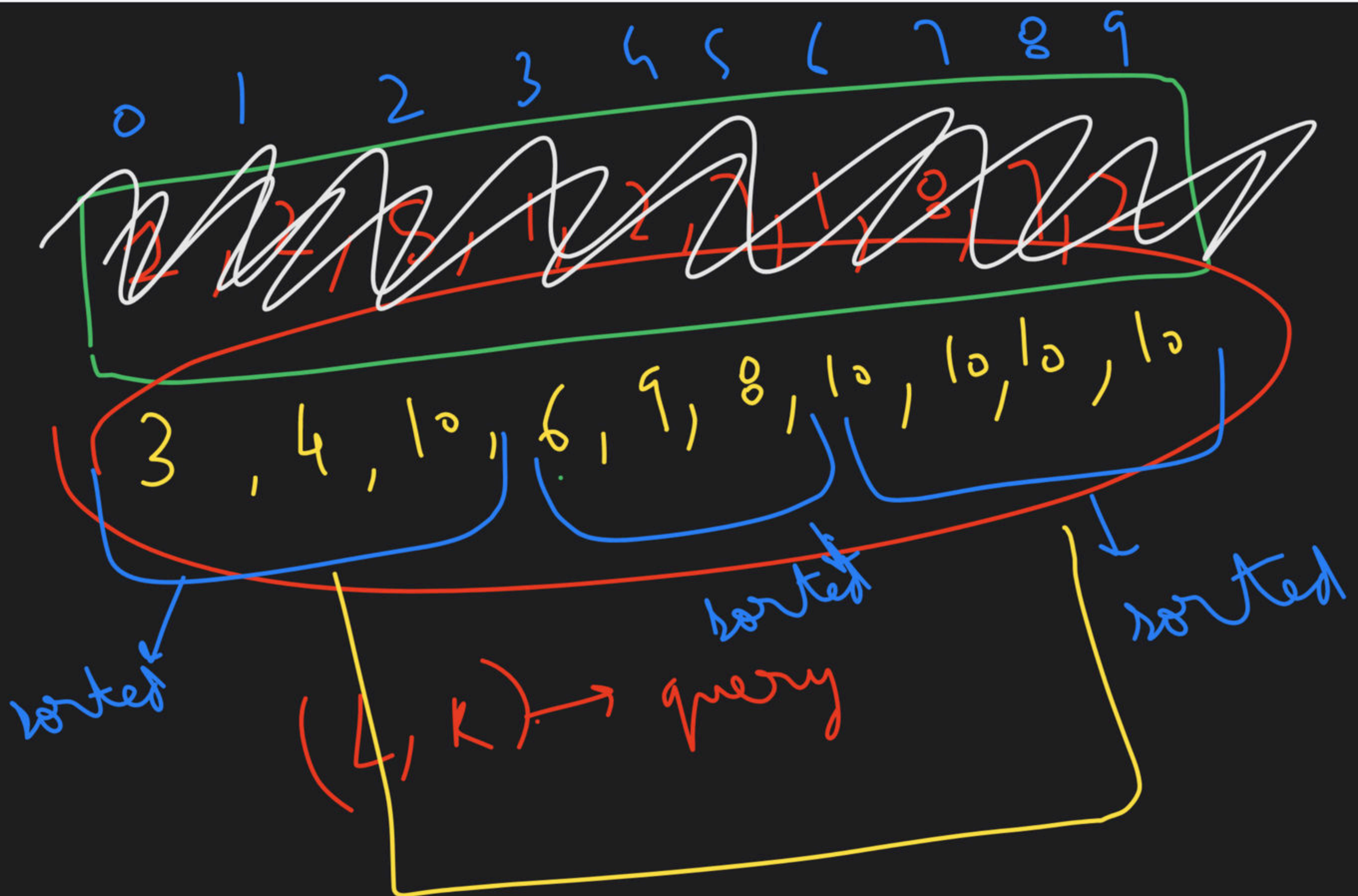
$$B \approx \sqrt{N \log \sqrt{N}} \approx \sqrt{\frac{N \log N}{2}}$$

D2 way →

Shoj



how many
distinct
elements



K Query problem

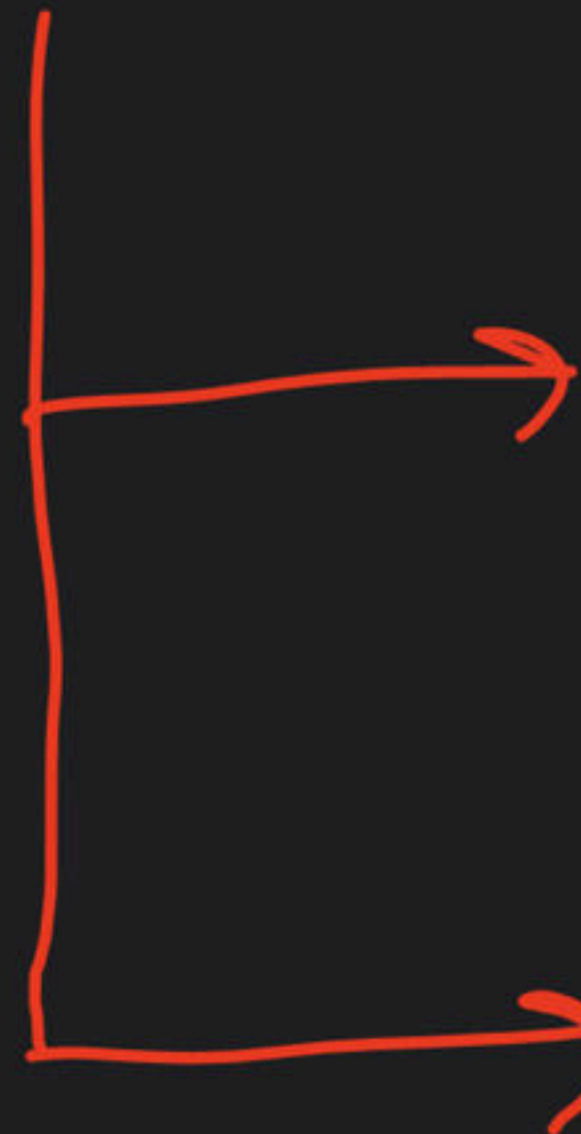
no. of elements in range $[L, R]$

which are $> R$

discussed
earlier
today



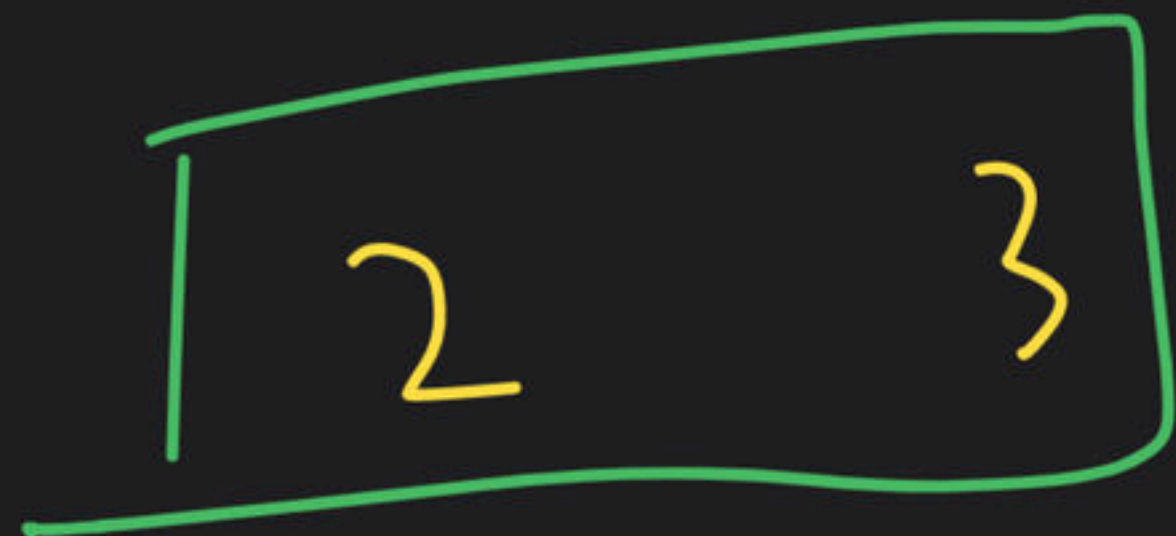
prefix sum



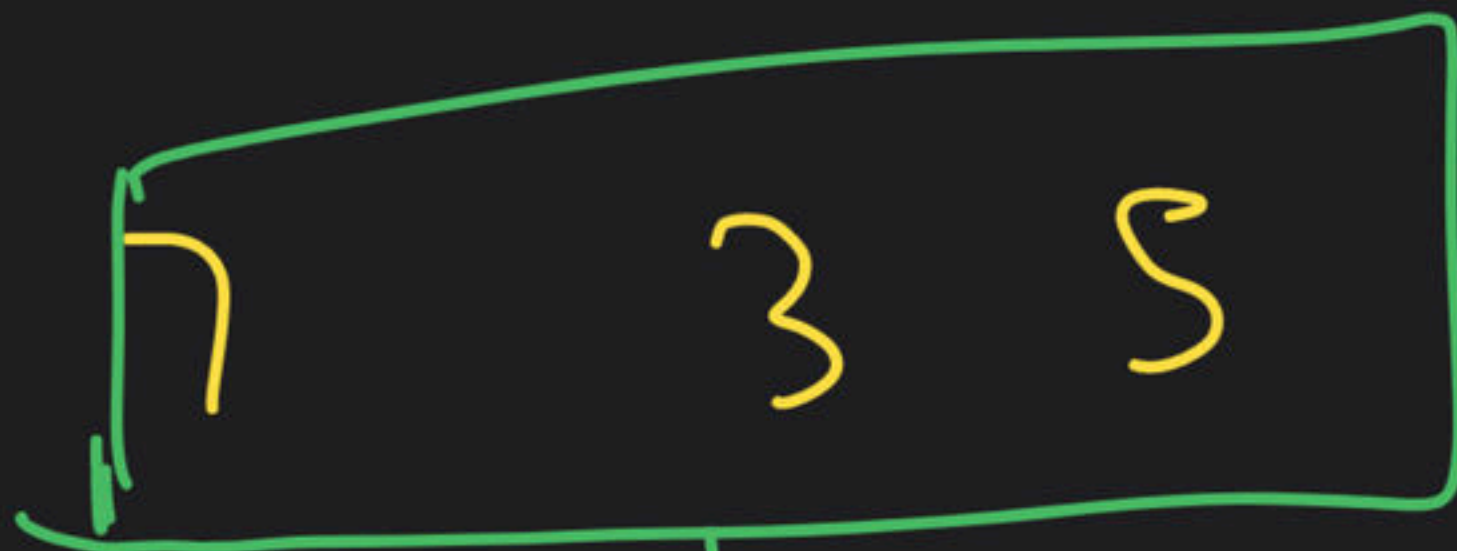
combinable
(made up term)

inversion

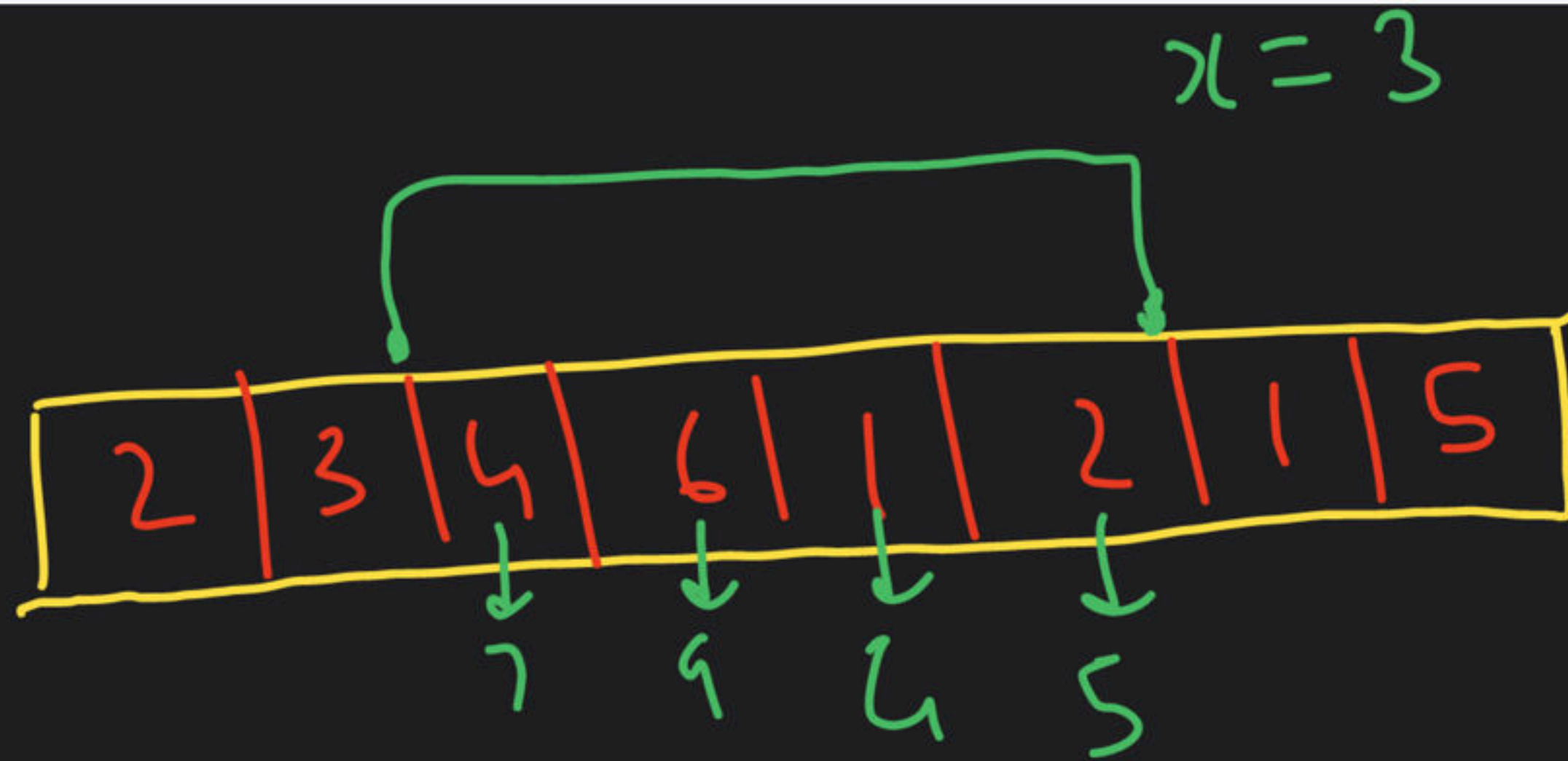
(add \leftrightarrow subtract)



↓
2 distinct



↓
3 distinct



using
sqrt decomposition

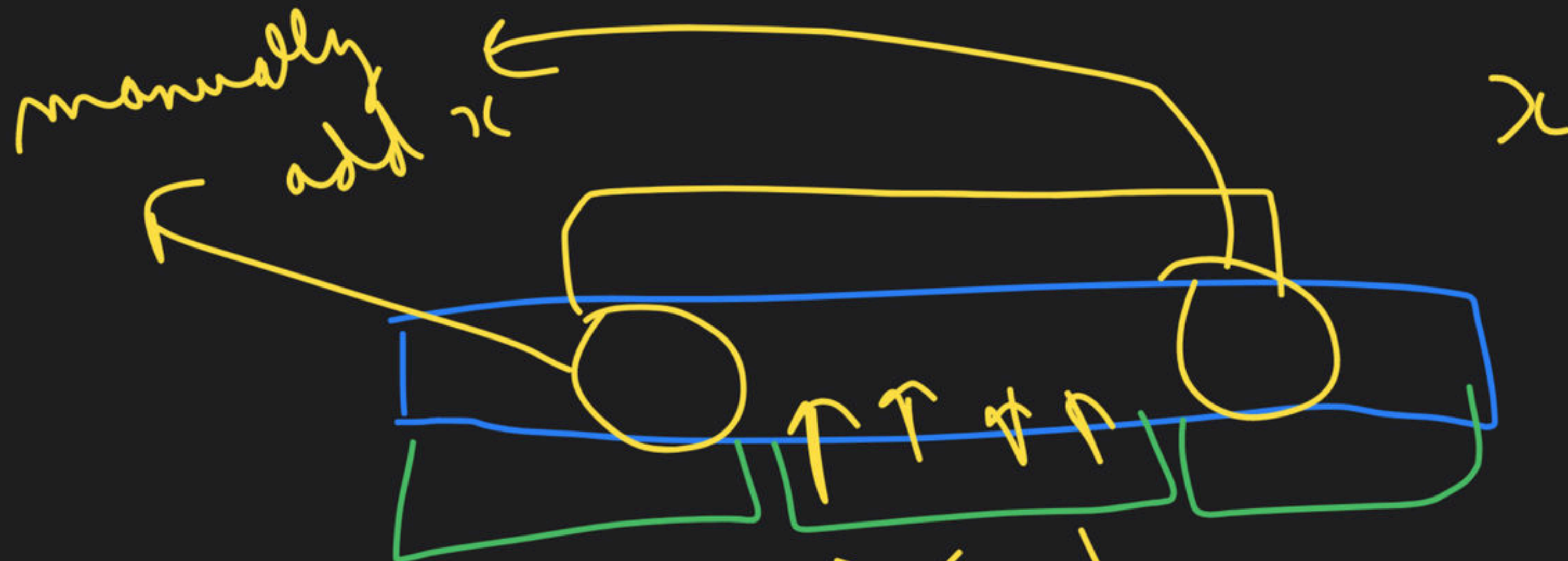
Q

1. $[L, R]$

→ sum of elements

2. $[L, R, x]$

→ add x to all
elements in range
 $[L \dots R]$



sum

~~sum~~

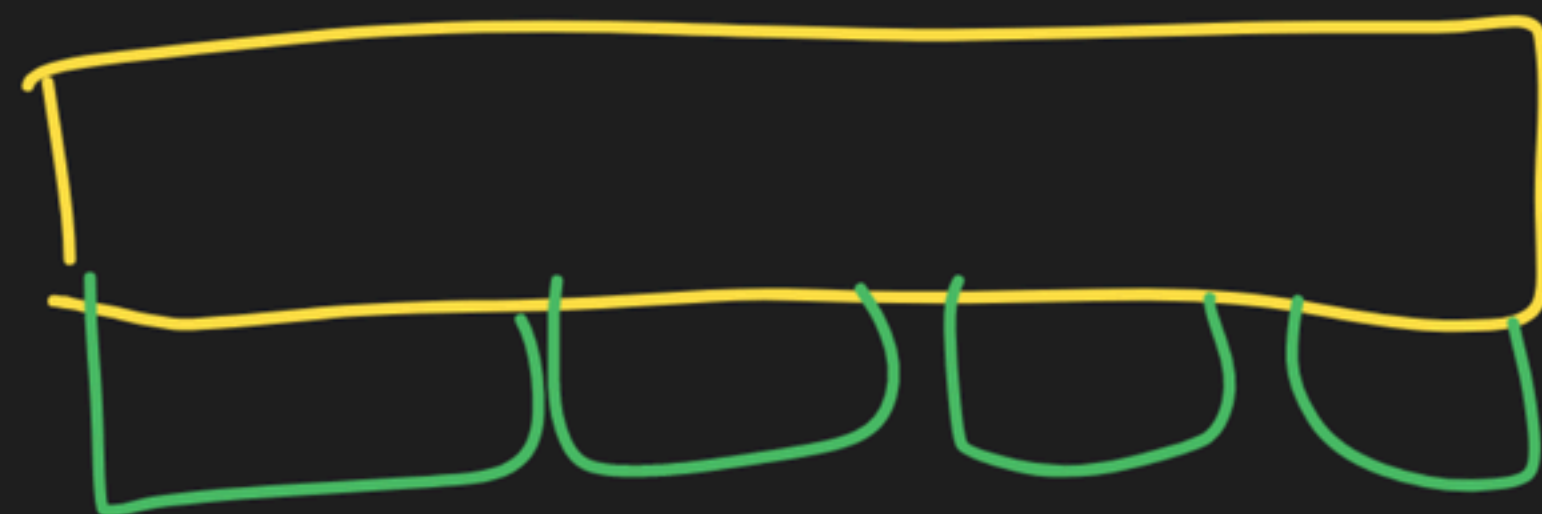
sum

watchman

sum
"

sum + Bx x

lazy-value += x
!!
0



Lazy



person

Propagation



only when
needed



greater \rightarrow upper bound

\geq \rightarrow lower bound

sorted
version

