Contents

Section outline

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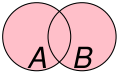

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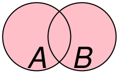

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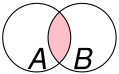

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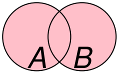

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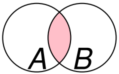

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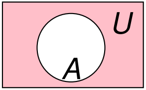

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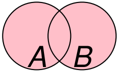

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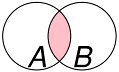

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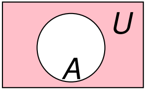
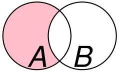


- Set difference: $A - B = A \cap \overline{B}$



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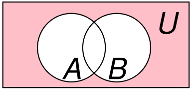

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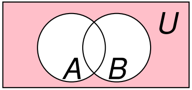

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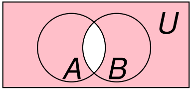

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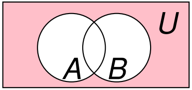

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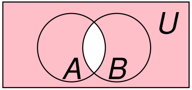

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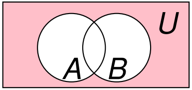

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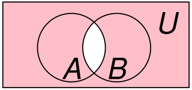

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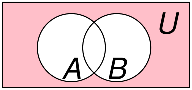

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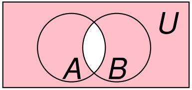

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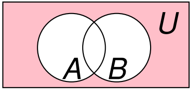

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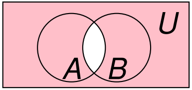

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  $A \cap \overline{B}$, $B \cap \overline{A}$, $A \cap B$ and $\overline{A \cup B}$ constitute a partition of U

Set algebra

Idempotence	$A \cup A = A$	$A \cap A = A$
Associativity	$(A \cup B) \cup C = A \cup (B \cup C)$	$(A \cap B) \cup C = A \cap (B \cup C)$
Commutativity	$A \cup B = B \cup A$	$A \cap B = B \cap A$
Distributivity	$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
Identity	$A \cup \{\} = A, A \cup U = U$	$A \cap \{\} = \{\}, A \cup U = A$
Involution	$\bar{\bar{A}} = A$	
Complements	$\bar{U} = \{\}, A \cup \bar{A} = U$	$\overline{\{\}} = U, A \cap \bar{A} = \{\}$
DeMorgan	$\overline{A \cup B} = \bar{A} \cap \bar{B}$	$\overline{A \cap B} = \bar{A} \cup \bar{B}$

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- Equivalence relation: $\mathcal{R}$ is reflexive, symmetric and transitive
- An equivalence relation induces a partition and vice versa

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- Binary relation $\mathcal{R}$ on sets A and B: $R \subseteq A \times B$
- Characteristic function of $\mathcal{R}$: $\chi_{\mathcal{R}}(a, b) = \begin{cases} 1 & \text{if } \langle a, b\rangle \in \mathcal{R} \\ 0 & \text{if } \langle a, b\rangle \notin \mathcal{R} \end{cases}$
- $\mathcal{R} \subseteq A \times A$ is reflexive if $\forall x \in A.\ x\mathcal{R}x$
- $\mathcal{R} \subseteq A \times A$ is symmetric if $\forall x, y \in A.\ x\mathcal{R}y \Rightarrow y\mathcal{R}x$
- $\mathcal{R} \subseteq A \times A$ is transitive if $\forall x, y, z \in A.\ x\mathcal{R}y \wedge y\mathcal{R}z \Rightarrow x\mathcal{R}z$
- $\mathcal{R} \subseteq A \times A$ is antisymmetric if $\forall x, y \in A.\ x\mathcal{R}y \wedge y\mathcal{R}x \Rightarrow x = y$
- Equivalence relation: $\mathcal{R}$ is reflexive, symmetric and transitive
- An equivalence relation induces a partition and vice versa
- Partial order: $\mathcal{R}$ is reflexive, antisymmetric and transitive

Relations (contd.)

- Connected relation: $\forall x, y \in A$, either $x\mathcal{R}y$ or $y\mathcal{R}x$

Relations (contd.)

- Connected relation: $\forall x, y \in A$, either $x\mathcal{R}y$ or $y\mathcal{R}x$
- Total order: Connected partial order (eg $\preceq$ on $\mathbb{R}$)

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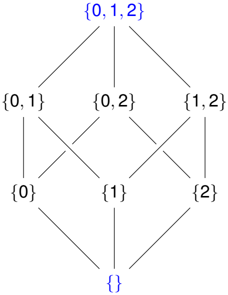

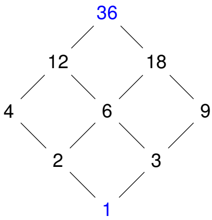

Divisors of 36 (60?)

Suppose $\langle A, \preceq\rangle$ is a poset, $M \in A$ $(m \in A)$, $S \subseteq A$

M (m) is a maximal (minimal) element of S iff $M \in S$ $(m \in S)$ and $\nexists x \in S$ st $M < x$ $(x < m)$

M (m) is a maximum (minimum) of S iff $M \in S$ $(m \in S)$ and $\forall x \in S,\ x \preceq M$ $(m \preceq x)$

A PO can be embedded in a total order

Relations (contd.)

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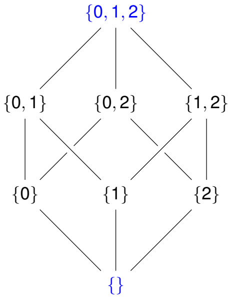

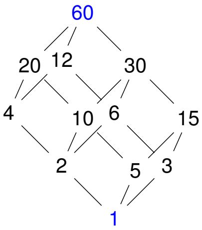

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A PO can be embedded in a total order

Lattices

Let $\langle A, \preceq \rangle$ be a poset, let $x, y \in A$

- The *meet* of x and y ($x \wedge y$), is the maximum of all lower bounds for x and y: $x \wedge y = \max\{w \in A : w \preceq x, w \preceq y\}$, *glb* for x and y
- The *join* of x and y ($x \vee y$), is the minimum of all upper bounds for x and y; $x \vee y = \min\{z \in A : x \preceq z, y \preceq z\}$, *lub* for x and y

A poset $\langle A, \preceq \rangle$ is a lattice iff every pair of elements in A have both a meet and a join

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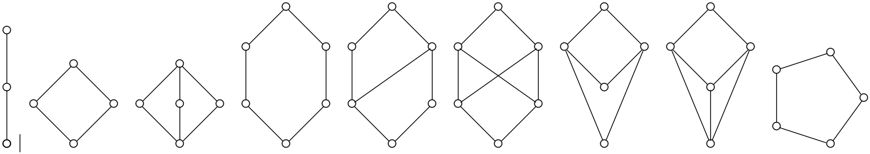


Lattices (contd.)

Basic order properties of meet and join

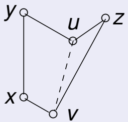
- $x \wedge y \preceq \{x, y\} \preceq x \vee y$
- $x \preceq y$ iff $x \wedge y = x$
- $x \preceq y$ iff $x \vee y = y$
- If $x \preceq y$, then $x \wedge z \preceq y \wedge z$ and $x \vee z \preceq y \vee z$
- If $x \preceq y$ and $z \preceq w$, then $x \wedge z \preceq y \wedge w$ and $x \vee z \preceq y \vee w$

Theorem

If $x \preceq y$, then $x \wedge z \preceq y \wedge z$ and $x \vee z \preceq y \vee z$

Proof.

- Let $v = x \wedge z$ and $u = y \wedge z$
- By transitivity, v is a lb for y and z
- By definition of $\wedge$, $v \preceq u$ (as u is the maximum among all lbs)



Similarly, the other clause may be proven □

Lattices (contd.)

Commutativity $x \wedge y = y \wedge x, \quad x \vee y = y \vee x$

Associativity $(x \wedge y) \wedge z = x \wedge (y \wedge z), \quad (x \vee y) \vee z = x \vee (y \vee z)$

Absorption $x \wedge (x \vee y) = x, \quad x \vee (x \wedge y) = x$

Idempotence $x \wedge x = x, \quad x \vee x = x$

Associativity of meet.

- $(x \wedge y) \wedge z \preceq x \wedge y \preceq x$ [$x \wedge y \preceq \{x, y\}$ applied twice]
- $(x \wedge y) \wedge z \preceq x$ [transitivity of $\preceq$]

Lattices (contd.)

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Lattices (contd.)

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- Thus $(x \wedge y) \wedge z$ is a lb of both x and $y \wedge z$

$\therefore$ $(x \wedge y) \wedge z \preceq x \wedge (y \wedge z)$ [glb of x and $y \wedge z$]

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- Also, $x \wedge (y \wedge z) \preceq (x \wedge y) \wedge z$ [on similar lines]

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$\therefore$ $(x \wedge y) \wedge z = x \wedge (y \wedge z)$ [if $a \preceq b$ and $b \preceq a$ then $a = b$] $\square$

Lattices (contd.)

Absorbtion.

- $x \preceq x \vee y$ [$\{x, y\} \preceq x \vee y$]

$\therefore$ $x \wedge (x \vee y) = x$ [$x \preceq y$ iff $x \wedge y = x$] $\square$

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Idempotence.

- $x \wedge x = x \wedge (x \vee (x \wedge y)) = x$ [Absorbtion, applied twice] $\square$

Principle of Duality

The dual of any theorem in a lattice is also a theorem.

Lattices (contd.)

Bounded lattice: It has a maximum element (1) and a minimum element (0), in which case identity properties hold:

- $0 \vee x = x = x \vee 0, \quad 1 \wedge x = x = x \wedge 1$
- $0 \wedge x = 0 = x \wedge 0, \quad 1 \vee x = 1 = x \vee 1$

Every finite lattice is bounded

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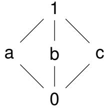
Distributive lattice: If $\forall x, y, z \in A$,

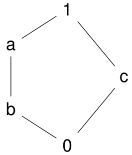
- $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$ and
- $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$

Are these lattices distributive?



$a \wedge (b \vee c) = a$

$(a \wedge b) \vee (a \wedge c) = 0$



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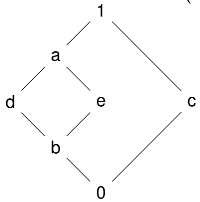

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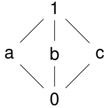

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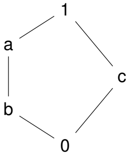

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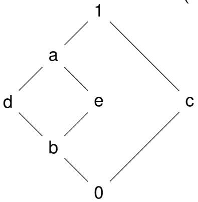

$a \wedge (b \vee c) = a \quad (a \wedge b) \vee (a \wedge c) = b$

Is $\mathcal{P}(A)$ for set A distributive?

? $x \cap (y \cup z) = (x \cap y) \cup (x \cap z)$ and

? $x \cup (y \cap z) = (x \cup y) \cap (x \cup z)$

Complemented lattice

- *Complement in a bounded lattice:* z is the complement of x iff
  - $x \wedge z = 0$ and
  - $x \vee z = 1$
- **Bounded complemented lattice:** every element has a complement
- In a bounded distributive lattice with minimum 0 and maximum 1, the complements of elements are unique, provided they exist – let $\bar{x}$ and z be complements of x ...

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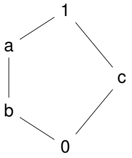
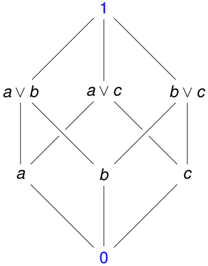
- $\bar{x} = \bar{x} \wedge 1 = \underbrace{\bar{x} \wedge (x \vee z)}_{D_f(\wedge,\vee)} =$
- $(\bar{x} \wedge x) \vee (\bar{x} \wedge z) =$
- $0 \vee (\bar{x} \wedge z) =$
- $\underbrace{(x \wedge z) \vee (\bar{x} \wedge z)}_{D_r(\wedge,\vee)} =$
- $(x \vee \bar{x}) \wedge z = 1 \wedge z = z$





Boolean lattice

Boolean lattice: Bounded complemented distributive lattice

- Extreme elements: Max: 1, Min: 0
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De Morgan's laws in a Boolean lattice $\langle \mathcal{A}, \preceq, \bar{\ }, 0, 1 \rangle$

1. $\overline{x \wedge y} = \overline{x} \vee \overline{y}$
2. $\overline{x \vee y} = \overline{x} \wedge \overline{y}$

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Meet of complements is 0

- $(x \wedge y) \wedge (\overline{x} \vee \overline{y}) =$
  $(x \wedge y \wedge \overline{x}) \vee (x \wedge y \wedge \overline{y})$
- $= 0 \vee 0 = 0$

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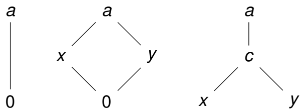
- $= 0 \vee 0 = 0$

Join of complements is 1

- $(x \wedge y) \vee (\overline{x} \vee \overline{y}) = ((x \wedge y) \vee \overline{x}) \vee \overline{y}$
- $= ((x \vee \overline{x}) \wedge (y \vee \overline{x})) \vee \overline{y}$
- $= (1 \wedge (y \vee \overline{x})) \vee \overline{y}$
- $= (y \vee \overline{x}) \vee \overline{y}$
- $= \overline{x} \vee (y \vee \overline{y})$
- $= \overline{x} \vee 1 = 1$

Boolean lattice structure

- Let $\mathcal{A}$ be a lattice with min 0
- $a \neq 0 \in \mathcal{A}$ is join irreducible if $x, y \preceq a \Rightarrow x \vee y \preceq a$, alternatively $a = x \vee y$ implies $a = x$ or $a = y$
- 0 is join irreducible?



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Boolean lattice structure

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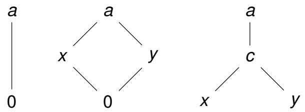

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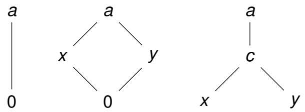

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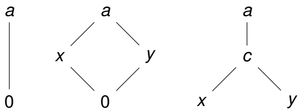

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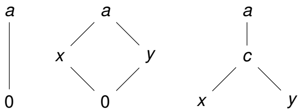

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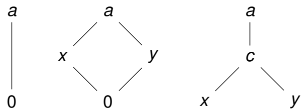

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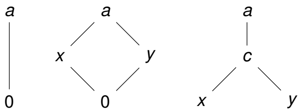

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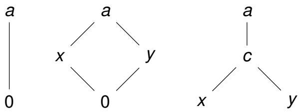

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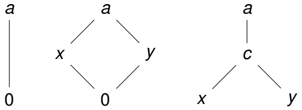

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- Unique (up to permutation) for distributive lattice

Boolean lattice representation (contd.)

Unique irredundant irreducible sum representation

- Let $a = c_1 \vee c_2 \vee \ldots \vee c_m = d_1 \vee d_2 \vee \ldots \vee d_n$

Boolean lattice representation (contd.)

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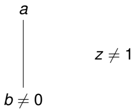
  – making the representation unique (up to permutation)

Boolean lattice structure (contd.)

- Let z be the complement of a in a lattice as shown
- So, $a \vee z = 1$ and $a \wedge z = 0$
- Suppose a has b as a unique predecessor
- Now, $b \vee z = a \vee z = 1$ and $b \wedge z = a \wedge z = 0$ as b is the immediate predecessor of a
- So, b is also a complement of z



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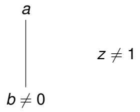
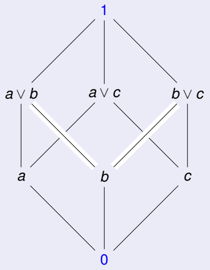
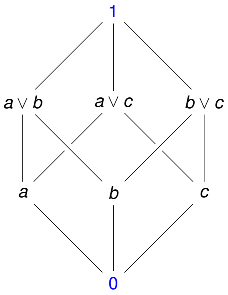


Join irreducible elements in a Boolean lattice

- A lattice with an element having a non-zero join irreducible element as a predecessor will not have unique complements
- In a Boolean lattice all non-zero join irreducible elements are atoms



Stone representation of Boolean lattices



- Atom of a Boolean lattice: Non-trivial minimal element of $A \setminus \{0\}$
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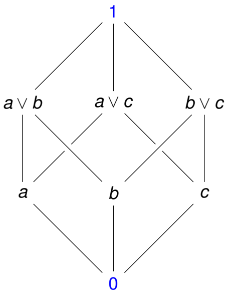

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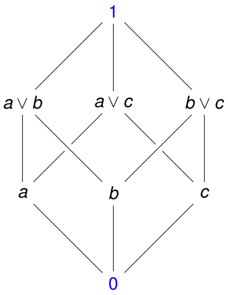

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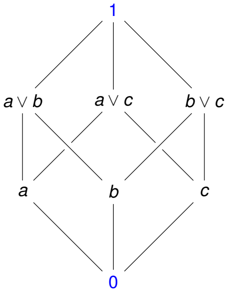

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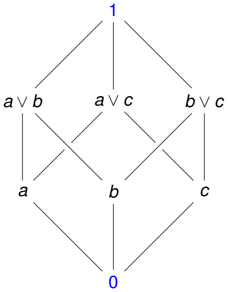

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- There will be $\binom{n}{\imath}$ such sets in level $\imath$, totaling to $\sum_{\imath=0}^{\imath=n} \binom{n}{\imath} = 2^n$

Boolean algebra from Boolean lattice

- For the Boolean lattice $\langle \mathcal{A}, \preceq, 0, 1 \rangle$ consider the algebraic system $\langle \mathcal{A}, +, \cdot, \bar{}, 0, 1 \rangle$ where $\vee \mapsto +$, $\wedge \mapsto \cdot$ and $\forall x \in \mathcal{A}$, $\overline{x} \mapsto z | x + z = 1, x \cdot z = 0$
- This system satisfies the Huntington's postulates for a Boolean algebra

B1: Commutative Laws

1. $x + y = y + x$
2. $x \cdot y = y \cdot x$

B2: Distributive Laws

1. $x \cdot (y + z) = x \cdot y + x \cdot z$
2. $x + (y \cdot z) = (x + y) \cdot (x + z)$

B3: Identity Laws

1. $x + 0 = x = 0 + x$
2. $x \cdot 1 = x = 1 \cdot x$

B4: Complementation Laws

1. $x + \bar{x} = 1 = \bar{x} + x$
2. $x \cdot \bar{x} = 0 = \bar{x} \cdot x$

Additional Boolean algebra properties

- These properties carry over from the Boolean lattice
- May be proven independently from the Huntington's postulates

Idempotence:

1. $x + x = x$
2. $x \cdot x = x$

Axiomatic proof

- $x + x = (x + x) \cdot 1$
- $= (x + x) \cdot (x + \bar{x})$
- $= x + (x \cdot \bar{x})$
- $= x + 0 = x$

Absorption:

1. $x + xy = x$
2. $x \cdot (x + y) = x$
3. $x + \bar{x}y = x + y$
4. $x \cdot (\bar{x} + y) = xy$

Axiomatic proof

- $x + xy = (x \cdot 1) + xy$
- $= x(1 + y) = x(y + 1)$
- $= x \cdot 1 = x$

Boolean algebra (contd.)

Boundedness/annihilation:

1. $x + 1 = 1$
2. $x \cdot 0 = 0$

Axiomatic proof

- $x + 1 = 1 \cdot (x + 1)$
- $= (x + \bar{x}) \cdot (x + 1)$
- $= x + (\bar{x} \cdot 1)$
- $= x + \bar{x} = 1$

Truth table for Boolean AND, OR, NOT:

x	y	$\overline{x}$	$x \cdot y$	$x + y$
0	0	1	0	0
0	1	1	0	1
1	0	0	0	1
1	1	0	1	1

Associativity:

1. $(x + y) + z = x + (y + z)$
2. $(x \cdot y) \cdot z = x \cdot (y \cdot z)$

Boolean algebra (contd.)

Axiomatic proof of associativity of Boolean $+$

- Let $x = a + (b + c)$ and $y = (a + b) + c$
- $ax = aa + a(b + c) = a + a(b + c) = a$
- $bx = ba + b(b + c) = ba + (bb + bc) = ba + (b + bc) = ba + b = b$
- Similarly, $cx = c$ and $ay = a$, $by = b$ and $cy = c$
- $yx = ((a+b)+c)x = (a+b)x+cx = (ax+bx)+cx = (a+b)+c = y$
- $xy = (a+(b+c))y = ay+(b+c)y = ay+(by+cy) = a+(b+c) = x$
- Thus, $x = xy = yx = y$

Additional Boolean algebra properties (contd.)

Uniqueness of Complement:

If $(a + x) = 1$ and $(a \cdot x) = 0$, then $x = \overline{a}$.

Involution:

$$\overline{(\overline{a})} = a$$

Complements of extreme elements:

- $\overline{0} = 1$
- $\overline{1} = 0$

Axiomatic proof

- $1 + 0 = 1$ [identity]
- $1 \cdot 0 = 0$ [boundedness]

$\therefore$ 0 is the complement of 1

DeMorgan's laws:

- $\overline{(x + y)} = \overline{x} \cdot \overline{y}$
- $\overline{(x \cdot y)} = \overline{x} + \overline{y}$

- Let $a \preceq b$ if $a \cdot b = a$ or $a + b = b$ then $\cdot \mapsto \wedge$ and $+ \mapsto \vee$
- Properties from axiomatic proofs allow Boolean algebras to be expressed as Boolean lattices – they are equivalent