

Recitation 5

CSCI 2270
July 12, 2018

Outline

Announcement: Moving office hours from Monday to Tuesday 11-12:30 next week

Homework common mistakes

Depth First Search

Breadth First Search

Dijkstra's Algorithm

Recitation Quiz

Homework

Counting the number of movies (642 not 50)

maze

K — L

Works on directed and undirected graphs

Keep track of:

- pre/post visit
- visited (property of node)
- stack

Depth First Search Pseudocode

```
function explore(G,v)
```

```
    v.visited = true
```

```
    v.previsit = clock
```

```
    clock += 1
```

```
    for nei in v.neighbors
```

```
        If nei.visited = false
```

```
            explore(G, nei)
```

```
    v.postvisit = clock, clock += 1
```

```
function DFS(G)
```

```
    for all v in G:
```

```
        v.visited = false
```

```
    for all v in G:
```

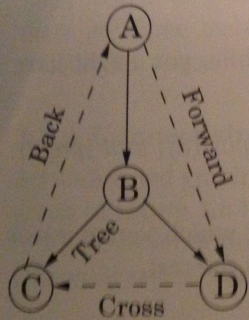
```
        if v.visited = false
```

```
            explore(G, v)
```

```
            distirctNum += 1
```

Why keep track of visit time?

DFS tree



Tree edges are actually part of the DFS forest.

Forward edges lead from a node to a *nonchild* descendant in the DFS tree.

Back edges lead to an ancestor in the DFS tree.

Cross edges lead to neither descendant nor ancestor; they therefore lead to a node that has already been completely explored (that is, already postvisited).

pre/post ordering for (u, v)

Edge type

$\begin{bmatrix} & & & \end{bmatrix}$
 $u \quad v \quad v \quad u$

Tree/forward

$\begin{bmatrix} & & & \end{bmatrix}$
 $v \quad u \quad u \quad v$

Back

$\begin{bmatrix} & & & \end{bmatrix}$
 $v \quad v \quad u \quad u$

Cross

Try it out

Breadth First Search

`BFS(G, starting_node)`

Used to find *connected components*

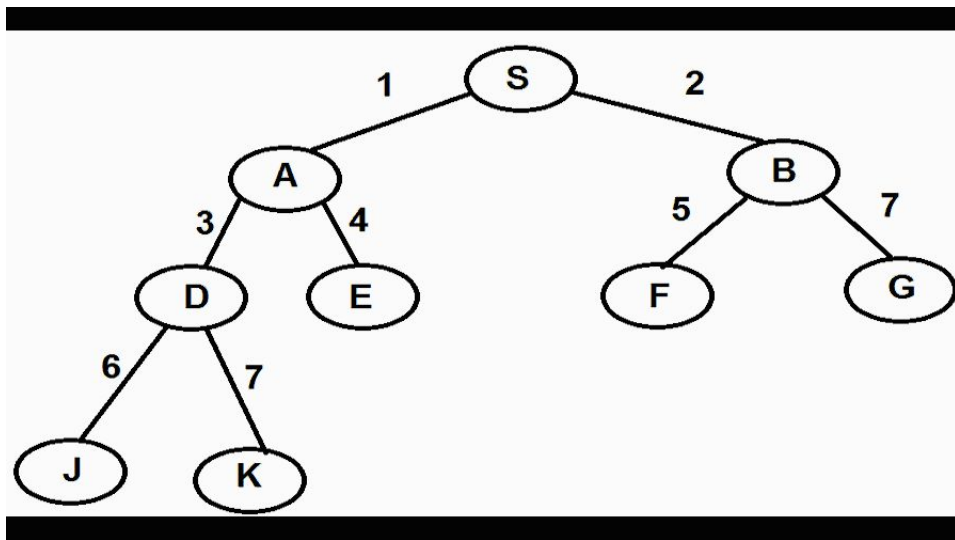
Keep track of:

- Queue (vector)
- Visited (boolean property of node)

$O(|V| + |E|)$ algorithm

On directed or undirected, unweighted graphs

Can find shortest path on *unweighted* graphs



Breadth First Search Pseudocode (without shortest distance)

BFS(starting_node)

 current = starting_node

 queue.enqueue(starting_node)

 while (queue is not empty)

 for neighbor in current.neighbors

 if neighbor.visited = false

 queue.enqueue(neighbor)

 neighbor.visited = true

 current = queue.pop_front()

return visited neighbors

Try it out

Dijkstra's Algorithm

Like BFS, but keeping track of weighted distances

Shortest distance from `starting_node` to every other node in graph

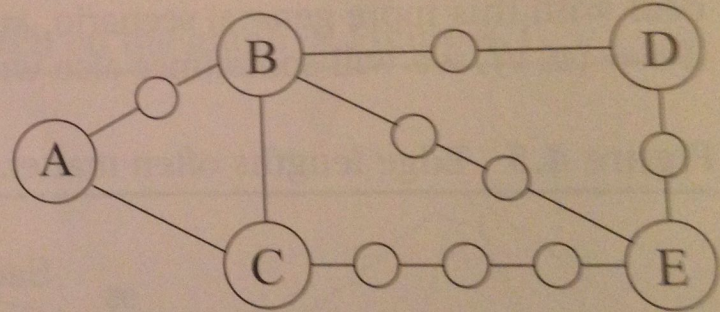
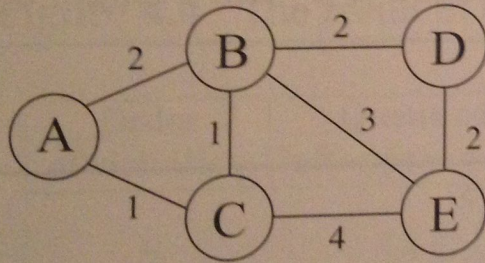
Relies on the property that the shortest path from s to t is also the shortest path to any of the vertices along the path.

This is exactly what BFS does.

Nonnegative edge weights

Couldn't we just tweak the graph and run BFS?

Figure 4.6 Breaking edges into unit-length pieces.



Well, yeah, but it gets really inefficient.

Dijkstra's pseudocode

dijkstra(G, w, s)

for all u in V : $u.dist = \text{infinity}$, $u.prev = \text{NULL}$

$s.dist = 0$

$H = \text{Makequeue}(V)$

while H is not empty:

$u = H.pop_front()$

 for all v in $u.neighbors$:

 if $v.dist > u.dist + w(u, v)$:

$v.dist = u.dist + w(u, v)$

$v.prev = u$

Quiz Password

Quiz Password

FirstWatch

(Does that exist here???)