# CSE110A: Compilers

#### **Topics**:

Scope

```
int main() {
  printf("");
  return 0;
}
```

# Homework questions?

#### Scope

What is scope?

 Can it be determined at compile time? Can it be determined at runtime?

• C vs. Python

Anyone have any interesting scoping rules they know of?

#### One consideration: Scope

Lexical scope example

```
int x = 0;
int y = 0;
{
  int y = 0;
  x+=1;
  y+=1;
}
x+=1;
y+=1;
What are the final values in x and y?
```

Symbol table object

- Two methods:
  - lookup(id): lookup an id in the symbol table.
    Returns None if the id is not in the symbol table.
  - insert(id,info): insert a new id (or overwrite an existing id) into the symbol table along with a set of information about the id.

#### a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"
INTYPE = "int"
LBRAC = "{"
RBRAC = "}"
SEMI = ";"
```

statements are either a declaration or an increment

#### a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"

TYPE = "int"

LBRAC = "{"

RBRAC = "}"

SEMI = ";"
int x;

int y;

x++;

y++;

y++;
```

statements are either a declaration or an increment

#### a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"

TYPE = "int"

LBRAC = "{"

RBRAC = "}"

SEMI = ";"

int x;

{
    int y;
    int y;
```

statements are either a declaration or an increment

• SymbolTable ST;

Say we are matched the statement: int x;

```
declare_statement ::= TYPE ID SEMI
{ }
```

lookup(id) : lookup an id in the symbol table. Returns None if the
id is not in the symbol table.

insert(id,info) : insert a new id (or overwrite an existing id) into
the symbol table along with a set of information about the id.

• SymbolTable ST;

```
Say we are matched the statement: int x;
```

```
declare_statement ::= TYPE ID SEMI
  self.eat(TYPE)
  variable name = self.to match[1] # lexeme value
  self.eat(ID)
  ST.insert(variable name, None)
  self.eat(SEMI)
```

• SymbolTable ST;

Say we are matched string: x++;

```
inc_statement ::= ID INCREMENT SEMI
{ }
```

lookup(id) : lookup an id in the symbol table. Returns None if the
id is not in the symbol table.

insert(id,info) : insert a new id (or overwrite an existing id) into
the symbol table along with a set of information about the id.

• SymbolTable ST; inc\_statement ::= ID INCREMENT SEMI variable name = self.to match[1] # lexeme value if ST.lookup(variable name) is None: raise SymbolTableException(variable name) self.eat(ID) self.eat(INCREMENT) self.eat(SEMI)

Say we are matched string: x++;

• SymbolTable ST;

statement : LBRAC statement\_list RBRAC

```
int x;
{
    int y;
    x++;
    y++;
}
```

• SymbolTable ST;

statement : LBRAC statement\_list RBRAC

start a new scope S

remove the scope S

```
int x;
{
    int y;
    x++;
    y++;
}
```

- Symbol table
- four methods:
  - lookup(id) : lookup an id in the symbol table.
    Returns None if the id is not in the symbol table.
  - insert(id, info): insert a new id into the symbol table along with a set of information about the id.
  - push\_scope() : push a new scope to the symbol table
  - pop\_scope() : pop a scope from the symbol table

• SymbolTable ST;

statement : LBRAC statement\_list RBRAC

You will be adding the functions to push and pop scopes in your homework

- Thoughts? What data structures are good at mapping strings?
- Symbol table
- four methods:
  - lookup(id): lookup an id in the symbol table.
    Returns None if the id is not in the symbol table.
  - insert(id,info): insert a new id into the symbol table along with a set of information about the id.
  - push\_scope() : push a new scope to the symbol table
  - pop\_scope() : pop a scope from the symbol table

Many ways to implement:

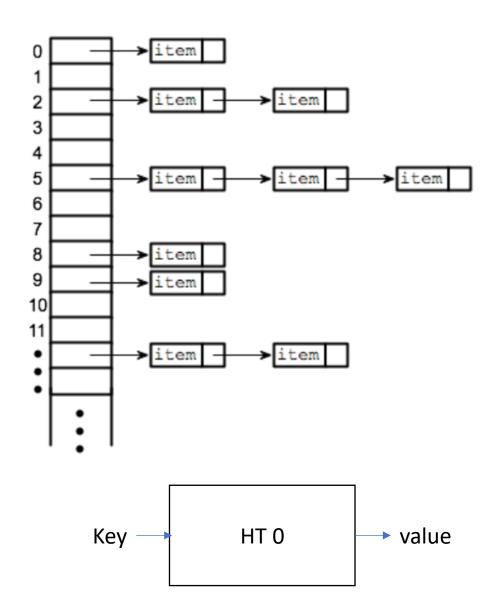
A good way is a stack of hash tables:

base scope

HT 0

#### What is a Hash Table?

- A HASH Table starts with a hashing function that maps a string (key) to an index.
- The Index goes into a table of buckets.
   The more indexes (bucket lists) the less likelihood of a collision.
- Bucket lists entries store each key, value pair.
- Dicts in Python are implemented using hash tables.



Many ways to implement:

A good way is a stack of hash tables:

push\_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

adds a new
Hash Table
to the top of the stack

HT 1

push\_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

insert(id,data)

HT 0

Many ways to implement:

A good way is a stack of hash tables:

insert(id,data)

insert (id -> data) at
top hash table

HT 1

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

lookup(id)

HT 0

lookup(id)

Many ways to implement:

A good way is a stack of hash tables:

check here first HT 1

Many ways to implement:

A good way is a stack of hash tables:

lookup(id) then check here HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

pop\_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 0

Example

```
int x = 0;
int y = 0;
{
  int y = 0;
  x++;
  y++;
}
x++;
y++;
```

HT 0

#### WHAT ELSE?

- GLOBAL VARIABLES
- NAMED SPACES
- DYNAMIC SCOPES

#### What about Dynamic Scope?

#### Example from Perl

```
use strict;
use warnings;
our $x = "global"; # Global variable
sub a {
  print "In a: x = x\n";
  b();
sub b {
  print "In b: x = x\n";
sub main {
  local $x = "local in main"; # Temporarily override
$x for dynamic scope
  print "In main: x = x\n";
  a();
main();
print "After main: x = x\n";
```

Some languages support dynamic scope, e.g. Perl and Lisp.

In the example x is a global variable (our \$x) and a dynamic variable in main (local x) when main calls a() it masks out the global variable x and passes down the local variable to be used in function a and b. X in those functions is derived from the call chain on the stack.

Try it, play with it.



Stack of hash tables
Is dynamically chained
Based on the call stack

Learn more click **HERE** 

## Next Topic:

• We will discuss parser generators