Quiz-6-CFG-AMBIG-PREC

8 questions

There are certain patterns that regular expressions can express that context-free grammars cannot express. But it is not an issue because those patterns do not show up in practice

O True

O False

- We just need to show fundamental operators
 - concat, choice, star

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```
add_expr ::= NUM '+' NUM
```

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How to express "a*" in BNF?

- We just need to show fundamental operators
 - concat, choice, star

How to express "a*" in BNF?

A production rule consists of:
☐ Terminals
☐ Regular Expressions
☐ Non-terminals
☐ function calls

Context-free grammar

We will use *Backus–Naur form* (BNF) form

- Production rules contain a sequence of either non-terminals or terminals
- In our class, terminals will either be string constants or tokens
- Traditionally tokens will be all caps.

Examples:

```
add_expr ::= NUM '+' NUM
```

```
mult expr ::= NUM '*' NUM
```

```
joint_expr ::= add_expr '*' add_expr
```

a left derivation will always produce the same parse tree as a right derivation
○ True
○ False

Not discussed yet. This is TRUE if the Grammar is unambiguous. The trees will create different sentential forms at different steps, but the tree will always be unique if the grammar is unambiguous. To be discussed in class.

Different programming languages make structure more or less explicit, e.g. using ()s and {}s.

Write a few sentences on any programming language experience you have w.r.t. structure and how you use it. For example do you use {}s when you write if statements, even if they contain a single statement? Why or Why not? Do you think Python's use of whitespace is a good construct for structure? Have you ever used <u>S expressions</u> in a Lisp language?

Programming language structure

Should conditionals require braces?

VS.
$$5 + (6 * 3)$$

should expressions require parenthesis?

S expressions (lisp) require explicit structure

Quiz: Top Down Parsing

What is an example of input recognized by the following grammar?		
$a \rightarrow a X$		
$a \rightarrow Y$		
○ XXXXXXXXY		
\bigcirc XYYYYYYYY		
○ YXXXXXXXX		

 \bigcirc YYYYYYYX

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one?

XXXXXXXY

RULE	Sentential Form
start	а

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
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How about this one?

XXXXXXXY

RULE	Sentential Form
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Applying either rule gives us a sentential form that won't create the string

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one? XYYYY

RULE	Sentential Form
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- 1 $a \rightarrow a X$
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How about this one? XYYYY

RULE	Sentential Form
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Similar reason: strings that are longer than 1 character cannot end in y

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one?

YXXXXXX

RULE	Sentential Form
start	Α

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one?

YXXXXXX

RULE	Sentential Form
start	Α
1	aX
1	aXX
•••	••••
2	YXXXXX

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one?

YYYYYX

RULE	Sentential Form
start	А

What is an example of input recognized by the following grammar?

- 1 $a \rightarrow a X$
- 2 $a \rightarrow Y$

How about this one?

YYYYYYX

RULE	Sentential Form
start	Α

We can only produce 1 y, so we cannot derive this string

Which grammar is ambiguous?

```
(a)
```

 $e \rightarrow e PLUS e$

 $e \rightarrow ID$

(b)

 $e \rightarrow e$ PLUS ID

 $e \to \mathsf{ID}$

(c)

 $e \rightarrow ID PLUS e$

 $e \rightarrow ID$

(d)

 $e \rightarrow ID PLUS ID$

 $\mathsf{e}\to\mathsf{ID}$

Which grammar is ambiguous?

(a)

 $e \rightarrow e PLUS e$

 $e \rightarrow ID$

(b)

 $e \rightarrow e PLUS ID$

 $e \rightarrow ID$

(c)

 $e \rightarrow ID PLUS e$

 $e \rightarrow ID$

(d)

 $e \rightarrow ID PLUS ID$

 $e \rightarrow ID$

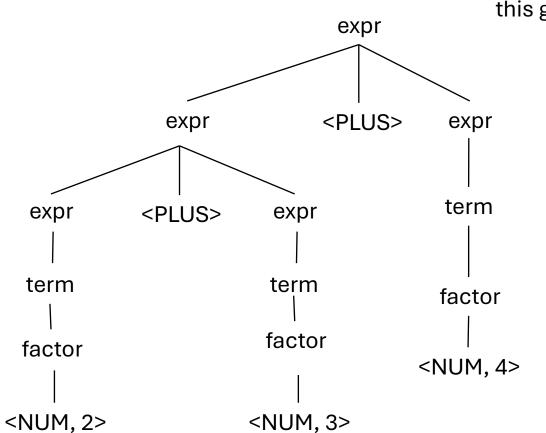
Let's look at some examples.

Let's assume that E is an "expr" and x is a number

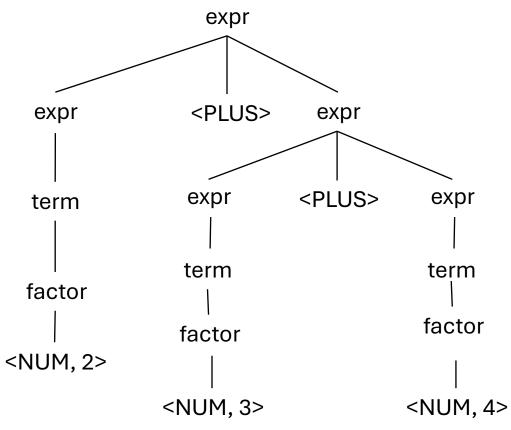
$$e \rightarrow e PLUS e$$

 $e \rightarrow ID$

input: 2+3+4



Both parse trees are valid, this grammar is ambiguous



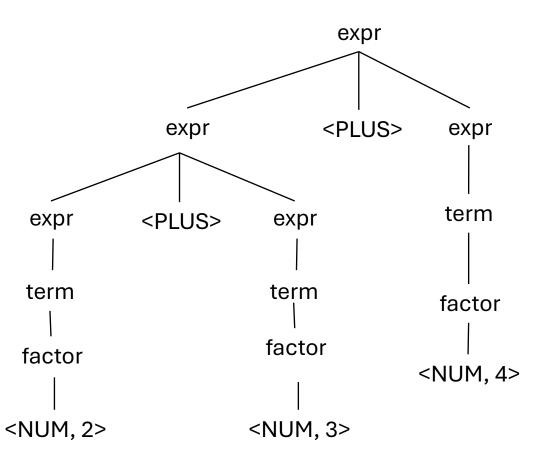
$$e \rightarrow e$$
 PLUS ID

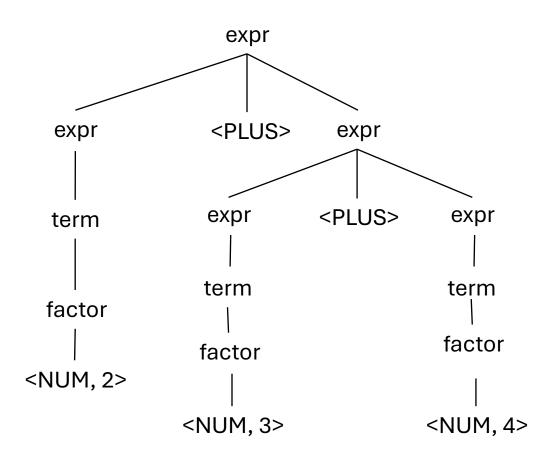
 $\mathsf{e}\to\mathsf{ID}$

input: 2+3+4

See next slide

What about this one?



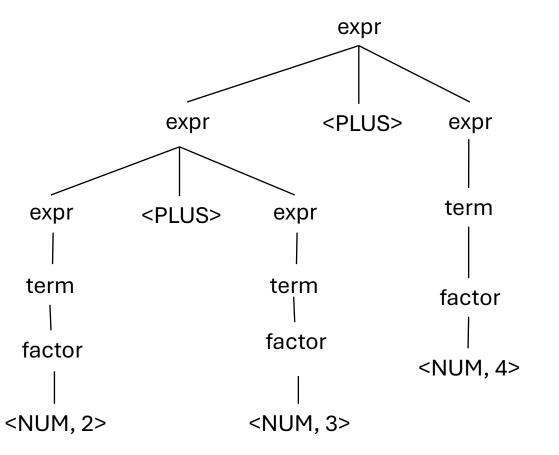


 $e \rightarrow e PLUS ID$

 $e \rightarrow ID$

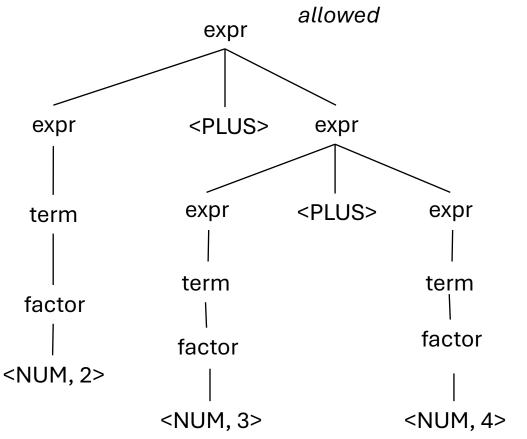
input: 2+3+4

What about this one?



UNAMBIGUOUS

Doesn't allow an expression on the RHS. This parse tree is not allowed



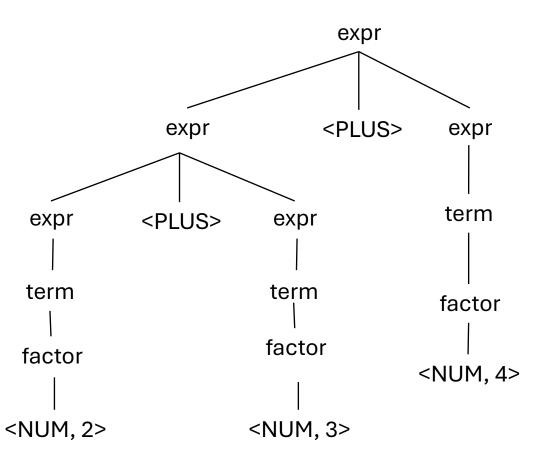
 $e \rightarrow ID PLUS e$

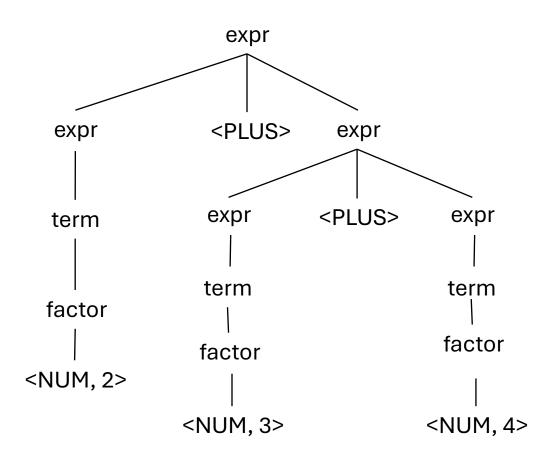
 $e \rightarrow ID$

input: 2+3+4

See next slide

What about this one?



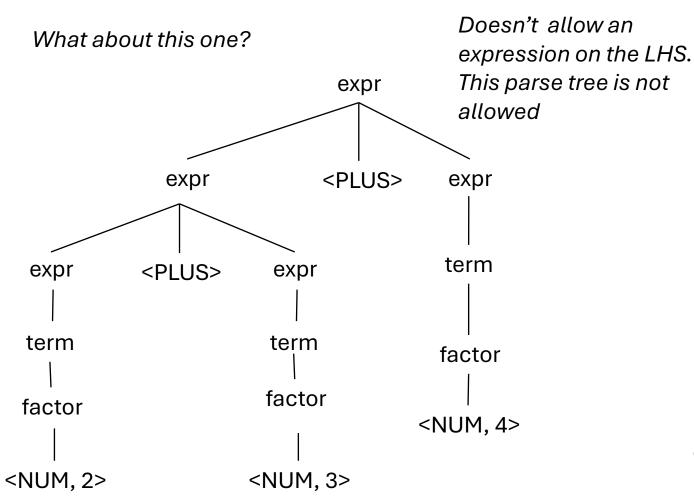


e → ID PLUS e

 $e \rightarrow ID$

input: 2+3+4

UNAMBIGUOUS



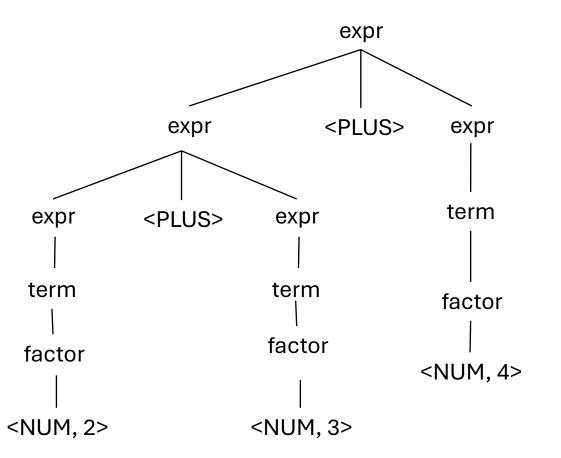
expr <PLUS> expr expr <PLUS> expr expr term term term factor factor factor <NUM, 2> <NUM, 3> <NUM, 4> $e \rightarrow ID PLUS ID$

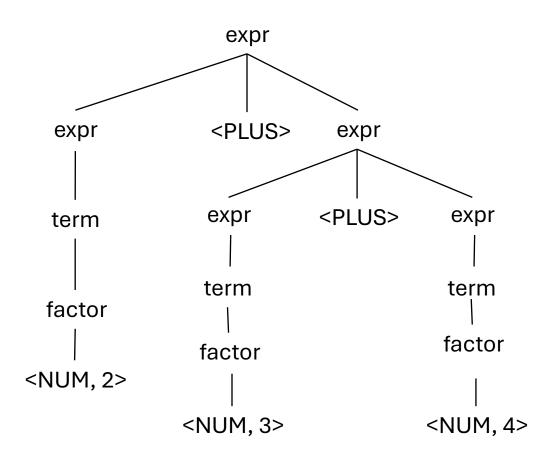
 $e \rightarrow ID$

input: 2+3+4

See next slide

What about this one?





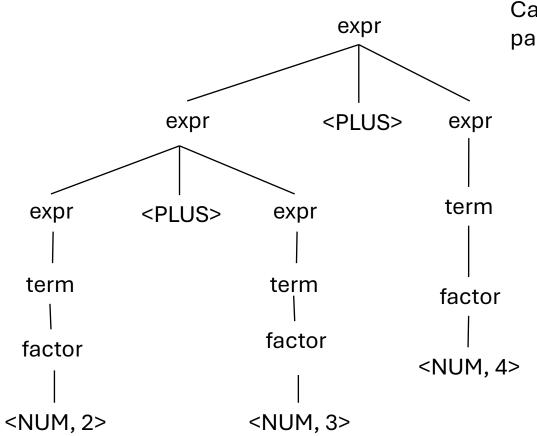
e → ID PLUS ID

 $e \rightarrow ID$

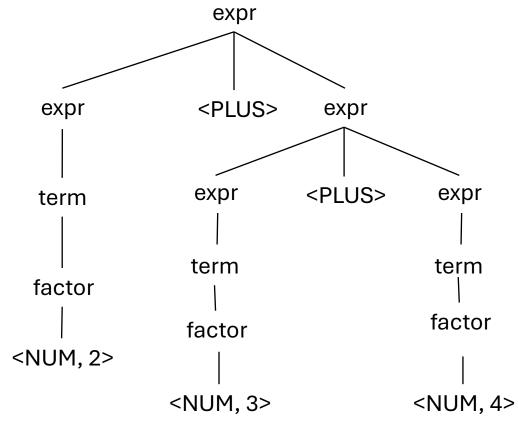
input: 2+3+4

UNAMBIGUOUS
But not expressive
No recursion

What about this one?



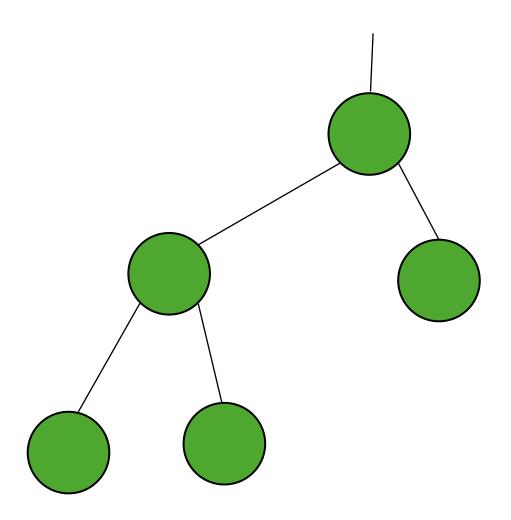
Cannot produce either parse tree!



operators with higher precedence should appear in production rules that appear higher in the parse tree

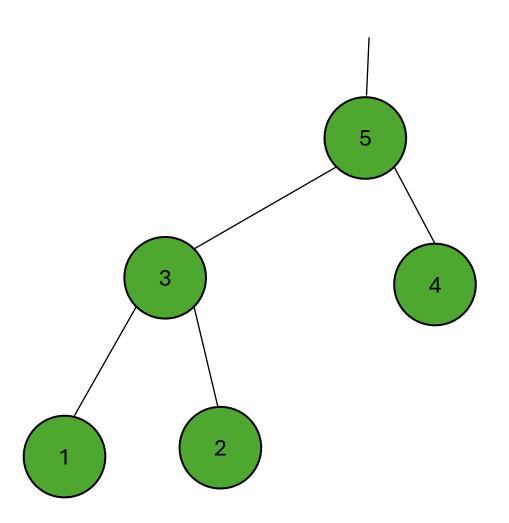
- \bigcirc True
- False

Post order traversal



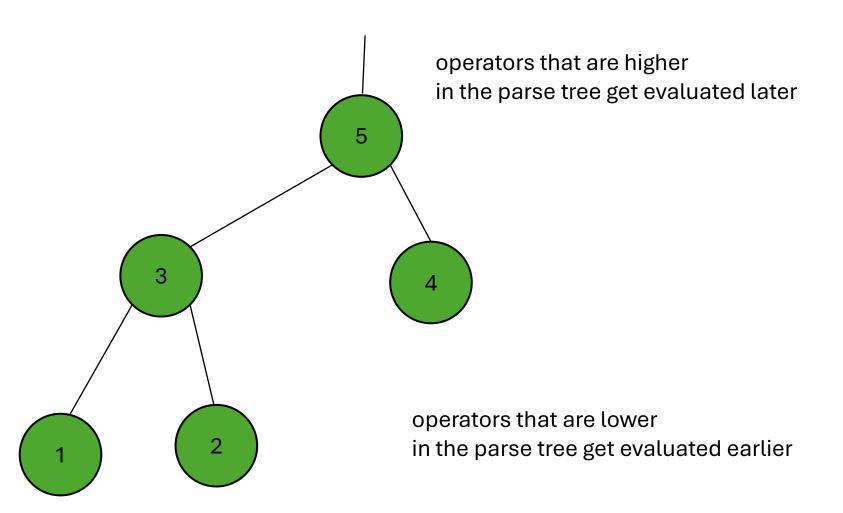
What is the post order traversal of this tree?

Post order traversal



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Post order traversal

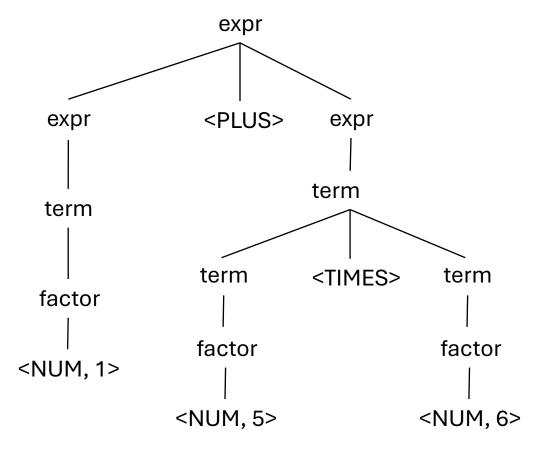


What is the post order traversal of this tree?

Evaluating a parse tree

Operator	Name	Productions
+	expr	: expr PLUS expr
*	term	: term TIMES term factor
()	factor	: LPAREN expr RPAREN NUM

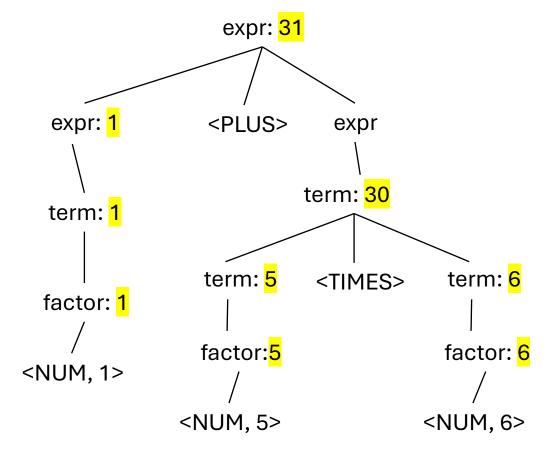
input: 1+5*6



Evaluating a parse tree

Operator	Name	Productions
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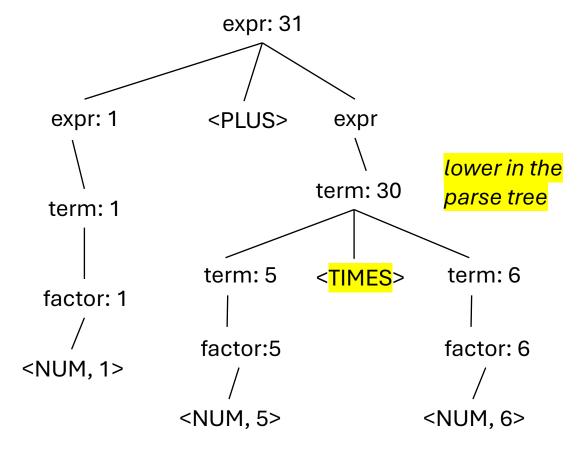
input: 1+5*6



Evaluating a parse tree

Operator	Name	Productions
+	expr	: expr PLUS expr
*	term	: term <mark>TIMES</mark> term factor
()	factor	: LPAREN expr RPAREN NUM

input: 1+5*6



Avoiding Ambiguity

- new production rules
 - One non-terminal for each level of precedence
 - lowest precedence at the top
 - highest precedence at the bottom
- How would we add power? ^

See this: See this

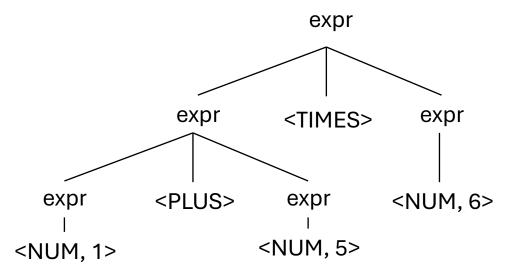
Precedence increases going down

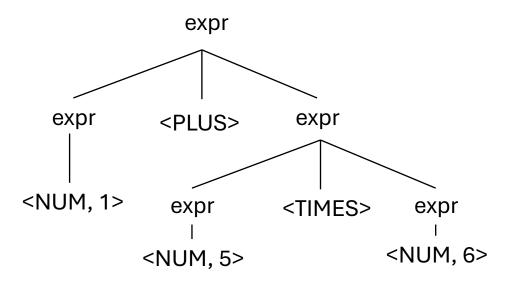
Operator	Name	Productions
+	expr	: expr PLUS expr
*	term	: term TIMES term factor
()	factor	: LPAREN expr RPAREN NUM

Write a few sentences about why it might be bad to have an ambiguous grammars

Ambiguous grammars

•input: 1 + 5 * 6





Ambiguous grammars

•input: 1 + 5 * 6

Evaluations are different!

