CSE110A: Compilers

Topics:

Scope

```
int main() {
  printf("");
  return 0;
}
```

Homework questions?

Scope

What is scope?

 Can it be determined at compile time? Can it be determined at runtime?

• C vs. Python

Anyone have any interesting scoping rules they know of?

One consideration: Scope

Lexical scope example

```
int x = 0;
int y = 0;
{
  int y = 0;
  x+=1;
  y+=1;
}
x+=1;
y+=1;
What are the final values in x and y?
```

Symbol table object

- Two methods:
 - lookup(id): lookup an id in the symbol table.
 Returns None if the id is not in the symbol table.
 - insert(id,info): insert a new id (or overwrite an existing id) into the symbol table along with a set of information about the id.

a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"
INTYPE = "int"
LBRAC = "{"
RBRAC = "}"
SEMI = ";"
```

statements are either a declaration or an increment

a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"

TYPE = "int"

LBRAC = "{"

RBRAC = "}"

SEMI = ";"
int x;

int y;

x++;

y++;

y++;
```

statements are either a declaration or an increment

a very simple programming language

```
ID = [a-z]+
INCREMENT = "\+\+"

TYPE = "int"

LBRAC = "{"

RBRAC = "}"

SEMI = ";"

int x;

{
    int y;
    int y;
```

statements are either a declaration or an increment

• SymbolTable ST;

Say we are matched the statement: int x;

```
declare_statement ::= TYPE ID SEMI
{ }
```

lookup(id) : lookup an id in the symbol table. Returns None if the
id is not in the symbol table.

insert(id,info) : insert a new id (or overwrite an existing id) into
the symbol table along with a set of information about the id.

• SymbolTable ST;

```
Say we are matched the statement: int x;
```

```
declare_statement ::= TYPE ID SEMI
  self.eat(TYPE)
  variable name = self.to match[1] # lexeme value
  self.eat(ID)
  ST.insert(variable name, None)
  self.eat(SEMI)
```

• SymbolTable ST;

Say we are matched string: x++;

```
inc_statement ::= ID INCREMENT SEMI
{ }
```

lookup(id) : lookup an id in the symbol table. Returns None if the
id is not in the symbol table.

insert(id,info) : insert a new id (or overwrite an existing id) into
the symbol table along with a set of information about the id.

• SymbolTable ST; inc_statement ::= ID INCREMENT SEMI variable name = self.to match[1] # lexeme value if ST.lookup(variable name) is None: raise SymbolTableException(variable name) self.eat(ID) self.eat(INCREMENT) self.eat(SEMI)

Say we are matched string: x++;

• SymbolTable ST;

statement : LBRAC statement_list RBRAC

```
int x;
{
    int y;
    x++;
    y++;
}
```

• SymbolTable ST;

statement : LBRAC statement_list RBRAC

start a new scope S

remove the scope S

```
int x;
{
    int y;
    x++;
    y++;
}
```

- Symbol table
- four methods:
 - lookup(id) : lookup an id in the symbol table.
 Returns None if the id is not in the symbol table.
 - insert(id, info): insert a new id into the symbol table along with a set of information about the id.
 - push_scope() : push a new scope to the symbol table
 - pop_scope() : pop a scope from the symbol table

• SymbolTable ST;

statement : LBRAC statement_list RBRAC

You will be adding the functions to push and pop scopes in your homework

- Thoughts? What data structures are good at mapping strings?
- Symbol table
- four methods:
 - lookup(id): lookup an id in the symbol table.
 Returns None if the id is not in the symbol table.
 - insert(id,info): insert a new id into the symbol table along with a set of information about the id.
 - push_scope() : push a new scope to the symbol table
 - pop_scope() : pop a scope from the symbol table

Many ways to implement:

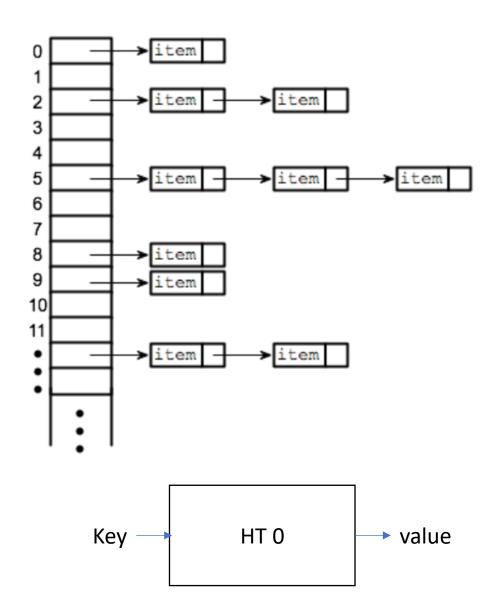
A good way is a stack of hash tables:

base scope

HT 0

What is a Hash Table?

- A HASH Table starts with a hashing function that maps a string (key) to an index.
- The Index goes into a table of buckets.
 The more indexes (bucket lists) the less likelihood of a collision.
- Bucket lists entries store each key, value pair.
- Dicts in Python are implemented using hash tables.



Many ways to implement:

A good way is a stack of hash tables:

push_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

adds a new
Hash Table
to the top of the stack

HT 1

push_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

insert(id,data)

HT 0

Many ways to implement:

A good way is a stack of hash tables:

insert(id,data)

insert (id -> data) at
top hash table

HT 1

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

lookup(id)

HT 0

lookup(id)

Many ways to implement:

A good way is a stack of hash tables:

check here first HT 1

Many ways to implement:

A good way is a stack of hash tables:

lookup(id) then check here HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 1

pop_scope()

HT 0

Many ways to implement:

A good way is a stack of hash tables:

HT 0

Example

```
int x = 0;
int y = 0;
{
  int y = 0;
  x++;
  y++;
}
x++;
y++;
```

HT 0

WHAT ELSE?

- GLOBAL VARIABLES
- NAMED SPACES
- DYNAMIC SCOPES

What about Dynamic Scope?

Example from Perl

```
use strict;
use warnings;
our $x = "global"; # Global variable
sub a {
  print "In a: x = xn;
  b();
sub b {
  print "In b: x = x n;
sub main {
  local $x = "local in main"; # Temporarily
override $x for dynamic scope
  print "In main: x = x\n";
  a();
main();
print "After main: x = x\n";
```

Some languages support dynamic scope, e.g. Perl and Lisp.

In the example x is a global variable (our \$x) and a dynamic variable in main (local x) when main calls a() it masks out the global variable x and passes down the local variable to be used in function a and b. X in those functions is derived from the call chain on the stack.

Try it, play with it.



Stack of hash tables
Is dynamically chained
Based on the call stack

Next Topic:

• We will discuss parser generators