

# How to type-set Fitch natural deductions using GNU troff, pic and eqn

*Sigrid Lundberg*  
*sigfrid (at) sigfrid-lundberg.se*

Fitch is a notation for natural deduction (Pelletier and Hazen, 2024), and `troff` is a software system for type-setting using Unix™ and related operating systems (Ossanna and Kernighan, 1994). Brian W. Kernighan was one of the creators of Unix and the C programming language. `pic` is a system for typesetting graphs, also created by Kernighan (1982). GROFF AKA GNU troff is the implementation I am using (FSF, 1990). There are other competitors, but this is the version I use.

The Fitch notations has got its name after its inventor, Fredric Fitch. This notation seems to be a de facto standard: It is used in all the text books I have been able to find electronically, and seems to be taught at logics courses in mathematics as well as philosophy. I wrote this note while learning Fitch; I used the writing was a method for learning. My intention is to demonstrate how to format natural deduction on this platform. I cannot teach you how to format scientific text, neither can I give an introduction to natural deduction.

Table 1. Unicode characters for logical signs and operators. On some operating systems you can type them by pressing `ctrl-shift-u` and then the four character code (following `u`). The Groff name is usually better to use than the Unicode character, but I tend to use the latter.

Unicode	Character	Groff name
u+00AC	¬	\(no
u+2227	∧	
u+2228	∨	
u+2200	∀	
u+2203	∃	
u+2192	→	
u+2194	↔	
u+22A5	⊥	
u+22A2	⊢	
u+2261	≡	
u+25A1	□	
u+25C7	∴	
u+2234	∴	
U+2208	∈	
U+2209	∉	

1		$A \vee B$	
2		$\neg A$	
3			$A$
4			$\perp$ $\perp$ Intro: 3,2
5			$B$ $\perp$ Elim: 4
6			$B$
7			$B$ Reit: 6
8		$B$	$\vee$ Elim: 6-7,3-5,1

Figure 1. Proof that  $A \vee B, \neg A \therefore B$ . The line numbering is in the left-most margin. Then there is a vertical line, as long as the proof. The step 1-2 in the proof is where the premisses lives. The horizontal line after step 2 is usually referred to as the *fitch line*. The two groups, 3-5 and 3-6 are sub-proofs, with their own premisses, vertical lines and fitch lines

## References

- FSF, Free Software Foundation, *Groff* (1990).
- Kernighan, Brian W., "PIC — A language for typesetting graphics," *Software: Practice and Experience* **12** (1982).
- Ossanna, Joseph F. and Kernighan, Brian W., "Troff User's Manual," *Computing Science Technical Report* **54** (1994).
- Pelletier, Francis Jeffry and Hazen, Allen, "Natural Deduction Systems in Logic" in *The Stanford Encyclopedia of Philosophy (Spring 2024 Edition)*, ed. Zalta, E. N. and U. Nodelman (2024).

```
#
# The proof is initialized by calling this macro, which
# informs the scripts on the number of steps in the proof
# and its maximum depth, i.e., how many proofs we have
# inside proofs.
#

set_steps_and_depths(8,3)

#
# Any proof (the root proof or any sub-proof) starts
# with the start_proof() which also names that proof.
#
# after started we add its premises, and end it with
# premis_end()

start_proof(START);
add_premis(START, "A $\vee$ B");
add_premis(START, " $\neg$ A");
premis_end(START);

#
# Here comes the sub-proofs
#

start_proof(SUB1);
add_premis(SUB1, "A");
premis_end(SUB1);

#
# The add_step() macro has three arguments, the name of the
# current proof, the result of the step, and finally the
# references to the steps needed for reaching the step.
#

add_step(SUB1, " $\perp$ ", " $\perp$  Intro: 3,2");
add_step(SUB1, "B", " $\perp$  Elim: 4");
end_proof(SUB1);

start_proof(SUB2);
add_premis(SUB2, "B");
premis_end(SUB2);
add_step(SUB2, "B", "Reit: 6");
end_proof(SUB2);

add_step(START, "B", " $\vee$  Elim: 6-7, 3-5, 1");
end_proof(START)
```

Figure 2. The PIC code needed to generate Figure 1.

1	$A \vee B$	
2	$\neg B \vee C$	
3	$A$	
4	$A \vee C$	$\vee$ Intro:3
5	$B$	
6	$\neg B$	
7	$\perp$	$\perp$ Intro:6,5
8	$A \vee C$	$\perp$ Elim:7
9	$C$	
10	$A \vee C$	$\vee$ Intro:9
11	$A \vee C$	$\vee$ Elim:6-8,9-10,2
12	$A \vee C$	$\vee$ Elim:1,3-4,5-11

Figure 3. A longer example, " $A \vee B, \neg B \vee C \therefore A \vee C$ ".