## TERMWORK

## Title :- Leaky Bucket Algorithm

Objective :

The primary objective is to successfully implement the Leaky

Bucket algorithm in order to understand functioning and

practical application in controlling data flow within a network

Truestigating the influence of varying parameters such as

bucket capacity and tolen addition rate, on the algorithm's

behaviour.

Theory : The Leaky Bucket algorithm is willely used traffic shoping mechanism in computer networks designed to control the rate at which data is teansmitted from a source to a desknoton. The concept is analogous to a bucket that can hold a limited amount of water, with a leak of the bottom In the context of data teansmission, the bucket represents a buffer with a fixed capacity and the leak symbolizes the allowed output rates. Tokens which symbolize units of datas, are added to the bucket of a specific rate. If the bucket is full, excess totans are discorded. When packets are to be transmitted, the algorithm checks of there are to hers available In the buckel. It tokens are present, packets are transmitted and the tookens are then leaked from the bucket. This Process is fundamental for managing network toother, providing a balance between resource ultalization and preventing excessive data rates that could degrade network performance.

Algorithm : In Haloze leaky bucket with 'copacity', 'rate', 'tokens' and 'last-time' 2 In the add-tokens method, calculate elopsed time since the last update. Add tokens based on elapsed time and rate ensuring copacity or not exceeded update last-time In the packet - process method, call add-tokens to update tokens based on theme elapsed. Check if there are enough and point packet sent successfully is no drop the packet and print packet dropped due to inssufficeent tokens. Process each packet using process-packet method. leaky bucket regulates packet transmission rate, avoiding network congertion Brodson : class leaky bucket: def -- Port -- (self, capacity, rate): self. capacity = capacity self. rate = rate self. tokens = 0 self. last-time = 0 def odd-tokens (self): current-those = those tene () time-elopsed = current-time = self. last-time self. tokens = min (self. capacity, self. tokens + time-elopse self-rate) Self. last-time = current-time

```
del process-packet (self, packet-size):
                 self. add tokens ()
                 If packetsize <= self. tokas :
                     self. tokens = packet-size
                      print (1" Packet of size spacket-size ] sent
                                       successfully")
                 else:
                     prent (f" Packet of size & packet-size & dropped due
                              to sosufficent tokens")
  import time
  def man ():
       capacity = 1000
       rate=1
       packets = [200, 500, 600, 700, 450, 400, 200]
       bucket = leaky Bucket (copacity, rate)
       for packet-size in packets:
             bucket, process-packet (packet-size)
et -- name -- = " -- main -- "
      magn ()
```

Outromes :

\* Learn to analyze the simpact of varying parameters (e.g. bucket capacity, token addition rater) on algorithm behaviour.

\* Recognizing the significance of the leaky bucket to instigating network issues and ensuring the reliability of data transfer.

Conclusion :

The leaky bucket algorithm is an effective tool for controlling the flow of data in a network. By regulating the rate of which data is transmitted, it helps prevent congestion and ensures a smoother data transfer expresence. Through, this experiment, a deeper understanding of traffic shaping and algorithm control has been achieved.

## TERM WORK 6

Tottle : To show and demonstrate the use of Cyclec Redundancy Check

Objective: To demonstrate the use of cyclic redundancy check and to understand the back concepts of error detection.

Theory of Cyclic Redundancy Check ((RC) is an error detecting code commonly used on disgotal metworks and storage derives to detect accordental changes to row data. Here's the theory behind it,

- \* Polynomial Codes CRC ?s based on the alegebraic structure of polynomial codes one finite fields, specifically the field of two elements.
- \* Devision The sender treats the date as a polynomeal and developed by by a specific polynomial known as the generator polynomial
- \* treas Detection The receiver performs the same disson aperation and if the remainder is non-zero, an error is detected

Algorithm:

- \* When the receiver gets the message, they perform the same calculations to get their own CRC chroksom
- \* If their checksum marknes the one we sett if means the message was toonsmotted correctly.
- \* If the checksins don't moter, it mean there was an

```
det crc_remoinder (Popul bitstong, polynomeal, inital-foll):
         displand= Poput - bits tong +' 0' + (len (polynomeal) -1)
         derpland = lest (deredend)
          to fromgal - 13t (to fromgal)
          for in in range (len (apput-bitstang)):
                ? 1 derodand[0] = = 11 3
                      for gin range (len(polynomial)):
                              Kindend [it] = sto (in (dindend (it]))
        remograder = ' ' goen (dividend) [-len (polynomial) + 1:]
        codeward = Popul - bits hang + remainder
         seturn remoinder, codeword.
det crc_check (codeward, polynomial):
       semonder, - = cxc semonder (radeword, polynomial, 'O'* (len (polynomial)
       return remainder == '6' & (len (polynomial) - 1)
        remarder, codeward = crc - ormarder (message, polynomed), 'O'
        print ("Original Message:", message)
        prot! "Polynomal ", polynomial)
        pant ("Transmostled (oderward) transmost)
        print ("Remainder:", remainder)
         chooce = eval (? put())
         94 chare = = 1 :
               ?s_valed = crc_check (codeword, polynomeal)
         else :
              error- poston = eval (supatt" Enter the bit poston:")
              Cadevard = cadevard [: error-poston]
               print ("Received Codewood with Eoros:", cadewad)
              95 va 19d= arc-check (codeward, polynama)
```

Program :

else.

Prent (" No error detected!")

Prent (" Firor detected!")

Course Outcome ?

\* To understand error detection and error mitigation in computer networks

\* To understand the basic concepts of CRC.

Conclusion:

In conclusion, the (RC program effectiently detects errors in data teansmission. The code demonstrates key concepts such as polynomial codes, disperson and checksum calculation.

Reference &

Larry L Peterson and Bruce S Darke, Computer Metabols, fith edition, EISENIER