

FA2022 Week 07

Crypto I

Anakin



Outline

Basics

XOR

Diffie-Hellman

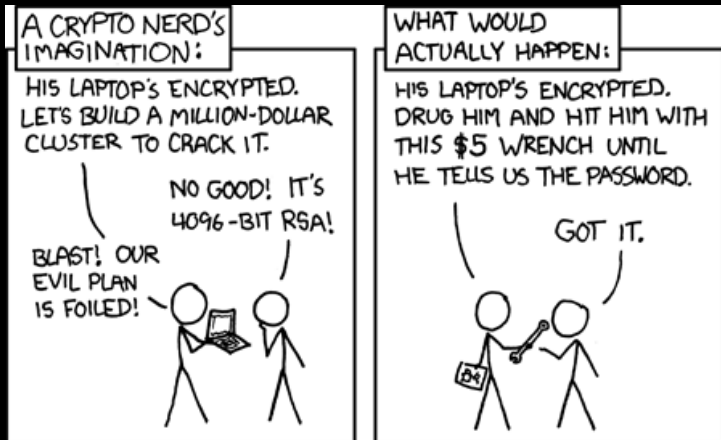


Announcements

- ACM Cleanup Sunday after Crypto II



sigpwny{n0t_that_crypt0_but_th3_0th3r_0n3}
ctf.sigpwny.com



Section 1

Basics



What is Crypto Anyways?



Alice



Bob



Eve



Why Do We Care?



Crypto in Ye Olden Days

- Relied on simple patterns
- Hard / annoying to break by hand, **easy to break by computer**
- Examples:
 - Caesar Cipher (rot k)
 - $a \rightarrow c, b \rightarrow d, \dots, y \rightarrow a, z \rightarrow b$ (rot 2)
 - Substitution
 - Create a table mapping each letter to another
 - Generalization of Caesar Cipher
 - Many More
 - **All insecure!!**



Data Representation

- TL;DR: computers store things in binary (0s and 1s), and we have different ways of representing this
- Tip: In Python, always work with bytes / bytestrings, never with normal strings (Python 3.8+)
- Look at the challenge source if given and mimic what they do



Conversion Cheatsheet

This is hard to read, download the slides!!

Format	Description	From Bytes	To Bytes
base64	uses printable letters to encode more complex binary	<code>base64.b64encode</code>	<code>base64.b64decode</code>
hex	uses symbols 0-9, A-F	<code>bytearray.hex()</code> , <code>binascii.hexlify()</code>	<code>bytes.fromhex()</code> , <code>binascii.unhexlify()</code>
integer	normal integers	<code>Crypto.Util.number.bytes_to_long</code> (PyCryptoDome), <code>int.from_bytes</code>	<code>Crypto.Util.number.long_to_bytes</code> (PyCryptoDome), <code>int.to_bytes</code>

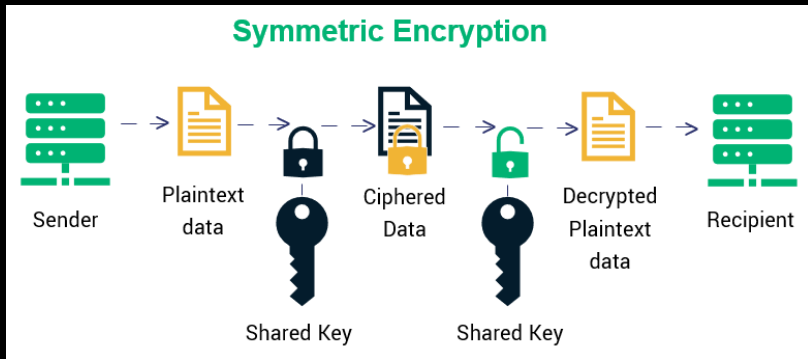


Section 2

XOR



Symmetric Encryption



XOR

A	B	$A \oplus B$
0	0	1
0	1	0
1	0	1
1	1	0



XOR

- XOR has some really nice properties that make it perfect for symmetric encryption
- Say M is some message as a bitstring, K is some key
- Then let $C = M \oplus K$ be a ciphertext
- Properties:
 - **Order doesn't matter:** $M \oplus K = K \oplus M$
 - **Group as needed:** $M \oplus (K \oplus K) = (M \oplus K) \oplus K$
 - **0 is the identity:** $M \oplus 0 = M$
 - **Self Inverse:** $K \oplus K = 0$
- All of this means $C \oplus K = M \oplus K \oplus K = M \oplus 0 = M$



Overview of Some Attacks

- For certain reasons, in general XOR is really really hard to break
 - Without more information, you need to try 2^λ guesses to break a bitstring of length λ
- Usually you need to know some information about the plaintext
 - Known plaintext
 - Properties like language
- You may need to know some information about the key
 - Really short keys are able to be brute forced
 - Check if code leaks information about key length



Section 3

Diffie-Hellman



Modular Arithmetic

- In cryptography, numbers can get really big really fast
- We use **modular arithmetic** to deal with this
- Modular arithmetic is **arithmetic with remainders after division**



Remainders

- Assume we have some number n . We are going to do some computation **mod** n
- For now, say $n = 101$

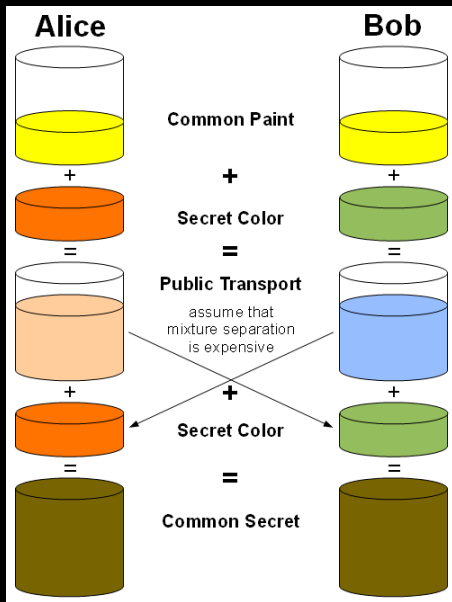
$$\begin{aligned} 131 + 140 * (102)^{2000} &\equiv 131 + 39 * (102)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 * (102)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 * (1)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 && (\text{mod } 101) \\ &\equiv 69 && (\text{mod } 101) \end{aligned}$$



Discrete Log

- If $a^b \equiv X \pmod{p}$, b is the **discrete log of X with base a** .
- Given some random X and a , finding b is really hard to compute for large primes p
- This **Discrete Log Problem** is the basis for many modern cryptography standards





Painting with Numbers

- Let g be a public number we call a generator and p be some public prime
- Alice generates secret a and computes $A \equiv g^a \pmod{p}$
- Bob generates secret b and computes $B \equiv g^b \pmod{p}$
- Alice sends Bob A and Bob sends Alice B
- Alice computes $B^a \pmod{p}$
- Alice computes $A^b \pmod{p}$

$$A^b \equiv (g^a)^b \equiv g^{ab} \equiv (g^b)^a \equiv B^a \pmod{p}$$



Overview of Some Attacks

Remember, discrete log in general is **hard**

- Small Primes
- Primes are generated in specific ways (“smooth primes”)
- “Oracle” attacks (access to a special machine that leaks information)

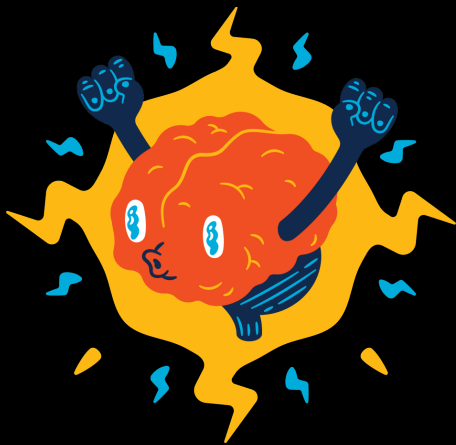


Tools!

- Python + [SageMath](#) is your friend
- [PyCryptodome](#) is an extremely useful Python crypto library
- [PwnTools](#) will allow you to automate parts of your attacks
- Google + StackOverflow (“how to crack DH with ...”)
- Installation is **annoying**, use the [CryptoHack Docker](#)



Practice @ CryptoHack



Next Meetings

2022-10-16 – This Sunday

- Crypto II
- More Diffie-Hellman + RSA
- ACM Cleanup afterwards

2022-10-20 – Next Thursday

- Rev II with Richard
- angr + Z3

2022-10-23 – Next Sunday

- Research Presentation from Mingjia
- Stealing Hospital Information





SIGPwny